## Exercise 1

#### Solution

a)

It is clear that:

$$T_1 \in O(n)$$

$$T_2 \in O(n^2)$$

Given the proper definition of the function class O(g) ("Big O") of a function  $g: \mathbb{N} \to \mathbb{R}$  as follows;

$$O(g) \coloneqq \{f: \mathbb{N} \to \mathbb{R} \mid \exists c \in \mathbb{R}, n_0 \in \mathbb{N}, f(n) \leq cg(n), \forall n > n_0\}$$

b)

We assume  $n \neq 0$  , for obviously  $T_1(0) = T_2(0)$ , so  $T_2$  is more efficient if:

$$T_1(n) > T_2(n) \Leftrightarrow 625n > n^2 \Leftrightarrow 625 > n$$

And the opposite is true for 625 < n, for n = 625 the algorithms perform equally.

# Exercise 2

### **Solution**

Algorithms  $T_1,...,T_5,T_7,...,T_9$  are trivial. for  $T_6$ , if n>1 then  $T_6\in O(n\log^2(n))$ . For  $T_{10}$ , notice that:

$$T(2) = 2T(1) + 2 \in O(1),$$
 
$$T(3) = 2T(2) + 2 = 2[2T(1) + 2] + 2$$
 
$$\cdot$$
 
$$\cdot$$
 
$$\cdot$$
 
$$T(n) = 2^n[T(1) + 1] + 2 \in O(2^n).$$

# Exercise 3

### **Solution**

a)

```
1 int a = 0, b = 0;
2 for (i = 0; i < n; i++) {
3     a = a + i;
4 }
5 for (j = 0; j < m; j++) {
6     b = b + j;
7 }</pre>
```

The algorithm operates in O(n) + O(m), the first for grows linearly with n and the second one grows with m.

b)

```
G C++
1
   float what2(int *arr, int n) {
2
        int a = 0;
3
       for (int i = 0; i < n; i++) {
4
           if(arr[i] > 10) {
5
               for (int j = 0; j < n; j++) {
6
                   a += n / 2;
7
               }
8
           } else {
               printf("ok :(")
9
10
           }
11
       }
12 }
```

The algorithm has 2 for 's' with n operations each, so it is at least  $\in O(n^2)$  , it also does check