Algortimos de streaming

Como elegir un número al azar de un stream infinito con memoria constante y otros algoritmos online.

Que es un stream?

$$X = \langle x_1, x_2, x_3, x_4, ..., x_n, ... \rangle$$

n desconocido se lee serialmente

Algoritmo de streaming?

Algoritmos que leen streams y calculan algo usando menos memoria que todo el stream¹.

Disclaimer

Para que le agarren el gustito, no para que lo sepan.

En criollo

```
class Procesador(object):
    def __init__(self, parametros):
        pass

def update(self, value):
        pass

def value(self):
    pass
```

Motivación

vuela coco

- algortimos probabilisticos
- soluciones piolas
- problemas interesantes

util

- base de datos, machine learning, optimizacion, metricas, y mil cosas mas.
- el procesamiento secuencial de datos es normalmente mucho mas rapido

Promedio

Dado un stream de números mantener el promedio de los números vistos

Promedio

```
class Promedio(Procesador):
   def __init__(self):
       self.n = 0
       self.suma = 0.0
   def update(self, value):
       self.n += 1
       self.suma += value
   def value(self):
       return self.suma / self.n
```

Ojo!

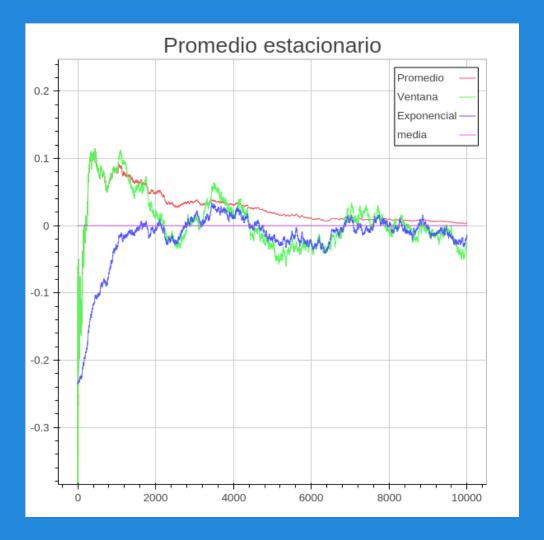
```
>>> sum([1, 1e100, 1, -1e100])
0.0
>>> math.fsum([1, 1e100, 1, -1e100])
2.0
```

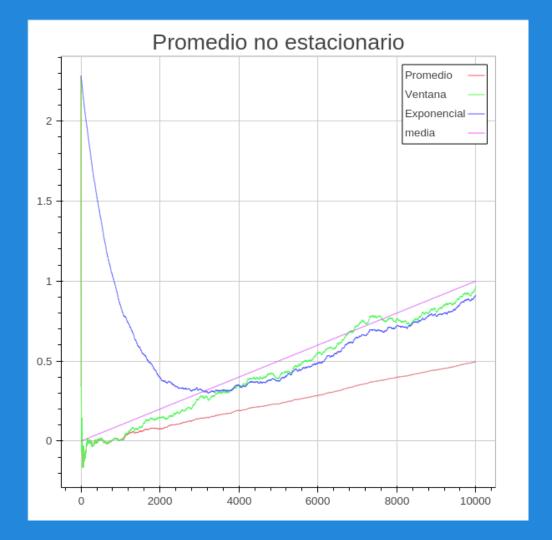
Sliding Window

```
class SlidingWindow(Procesador):
   def __init__(self, size):
       self.size = size
       self.window = []
   def update(self, value):
       self.window.append(value)
       if len(self.window) > self.size:
           self.window = self.window[1:]
   def value(self):
       return sum(self.window) / len(self.window)
```

Exponencial

```
class Weighted(Procesador):
    def __init__(self, alpha):
        assert(alpha < 1)</pre>
        self.alpha = alpha
        self._value = None
    def update(self, value):
        if self._value is None:
            self._value = value
        else:
            self._value = self._value * (1 - self.alpha) + value * self.alpha
    def value(self):
        return self._value
```



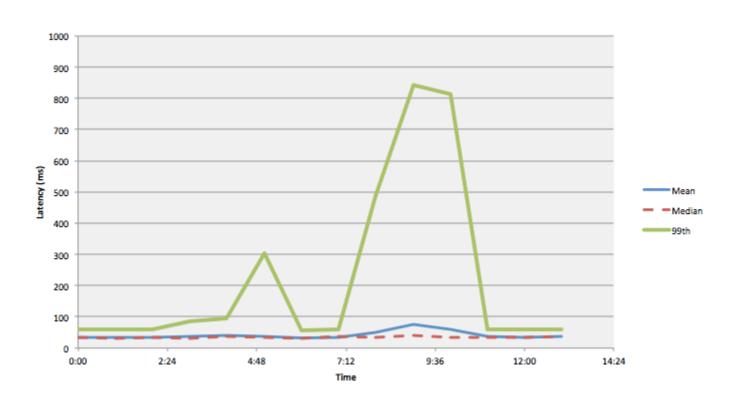


Usos

Monitoreo!

(y un millón de cosas más)

Ojo como medimos latencia



Número al azar

Como elegir un elemento al azar de un stream

Número al azar

$$\langle x_1 \rangle \to p(x_1) = 1$$

 $\langle x_1, x_2 \rangle \to p(x_1) = \frac{1}{2}$
 $\langle x_1, x_2, x_3 \rangle \to p(x_1) = \frac{1}{3}$
:
 $\langle x_1, x_2, x_3, ..., x_n \rangle \to p(x_1) = \frac{1}{n}$

Inducción

$$f(1) \rightarrow x_1$$

$$f(n) \rightarrow x_n$$
 $p = 1/n$
 $f(n-1)$ $p = 1 - 1/n$

Implementación

```
class EleccionUniforme(Procesador):
    def __init__(self):
        self.n = 0
        self._value = None
    def update(self, value):
        self.n += 1
        # range of random is [0.0, 1.0)
        if random.random() < 1. / self.n:</pre>
            self._value = value
    def value(self):
        return self._value
```

Usos

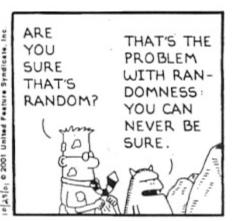
Muestreo!

Se puede extender a tomar múltiples elementos y priorizando los más recientes.

Testing

DILBERT By Scott Adams





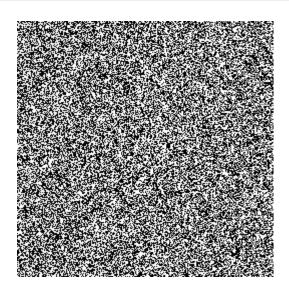
Seed fijo

Prueba que se ejecuta el algoritmo correctamente, no que elige bien al azar

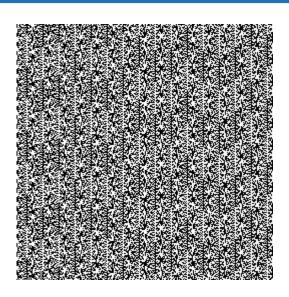
```
def test_uniforme():
    p = EleccionUniforme()
    random.seed(1)
    map(p.update, range(10))
    assert p.value() == 9
```

```
int getRandomNumber()
{
return 4; // chosen by fair dice roll.
// guaranteed to be random.
}
```

wat



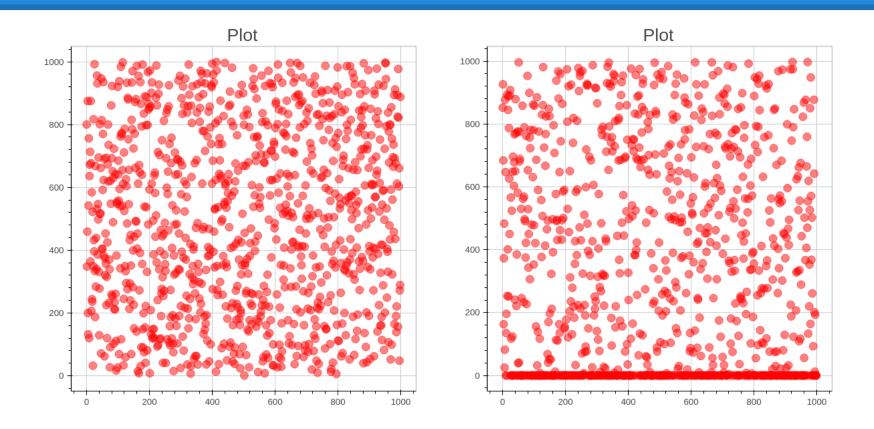
random.org



PHP rand() on Windows

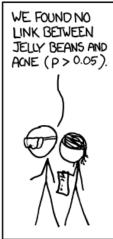
source: http://www.random.org/analysis/

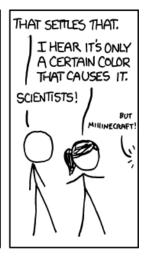
Visual

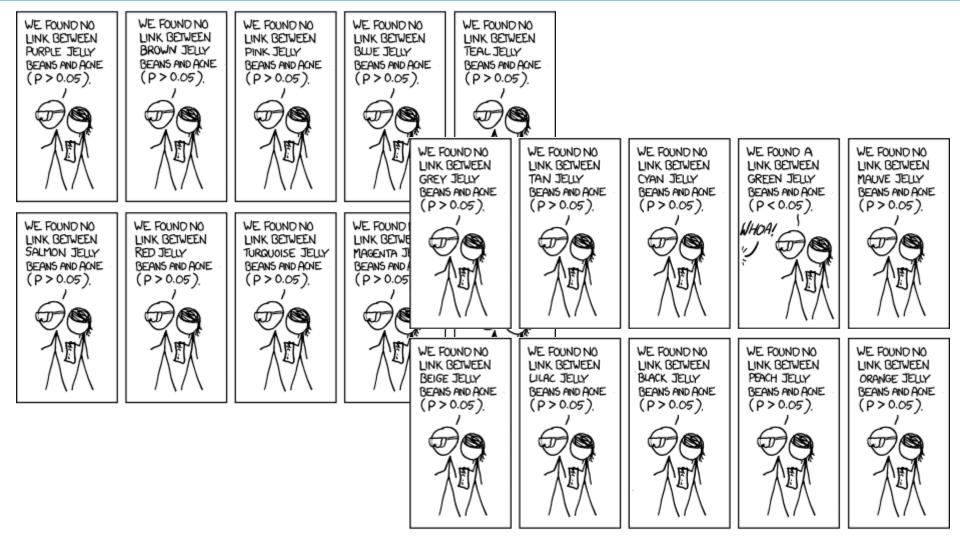


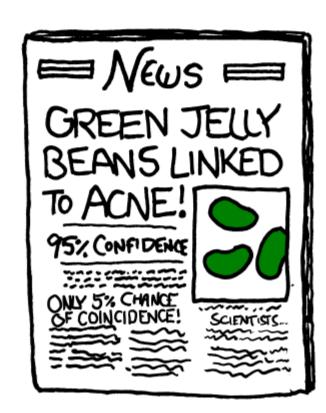
Estadística!











<pre>\$ nosetests statistical_tests.pyF</pre>
FAIL: test_fails.test_15
<pre>Traceback (most recent call last): File "/usr/lib/python2.7/dist-packages/nose/case.py", line 197, in runTest self.test(*self.arg) File "statistical_tests.py", line 46, in test_15 assert(p < 0.05) AssertionError</pre>
Ran 20 tests in 0.002s
FAILED (failures=1)

Sketches

Representamos el stream en espacio menor al que ocupa.

Tradeoff entre espacio y precisión.

Las respuestas suelen tener un error $< \varepsilon$ con probabilidad $< 1-\delta$.

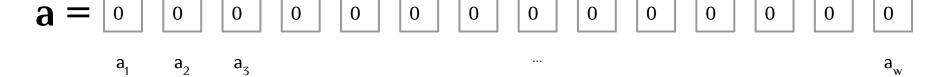
 ϵ y δ suelen ser los parámetros para construir el sketch.

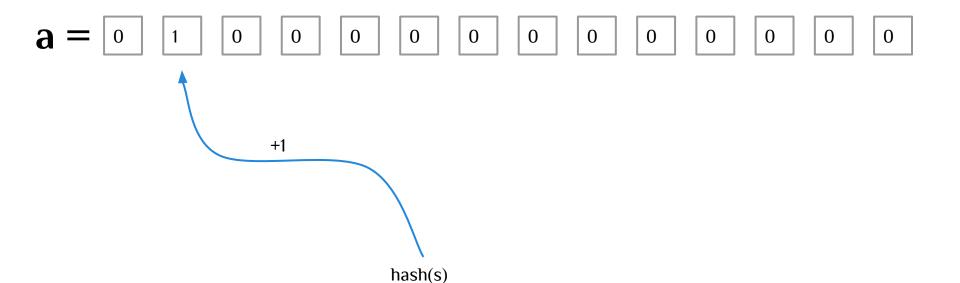
Usos

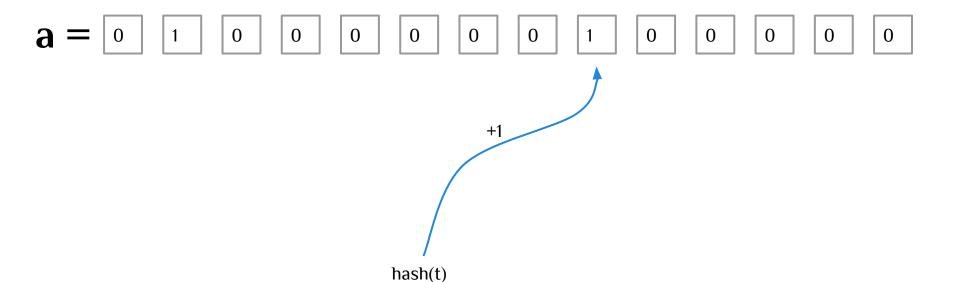
- Agregacion
 - o escalabilidad
 - o redes de sensores
- Queries mas interesantes

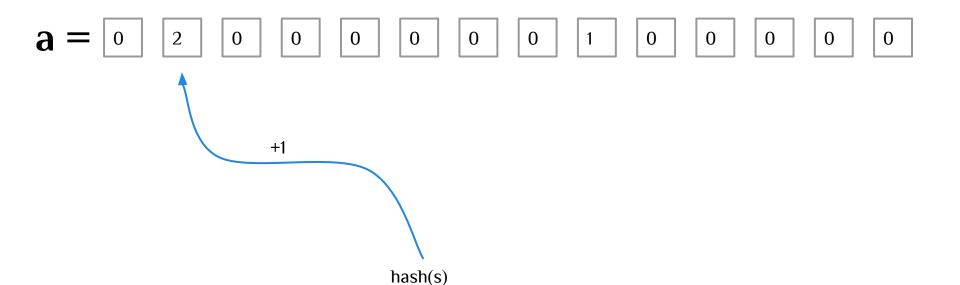
Un sketch que nos permita (entre otras cosas) dado un stream de elementos, estimar cuántas veces vimos cada uno.

stream =
$$<$$
s, t, s, u, s, v, s $>$
Point(s) = 4

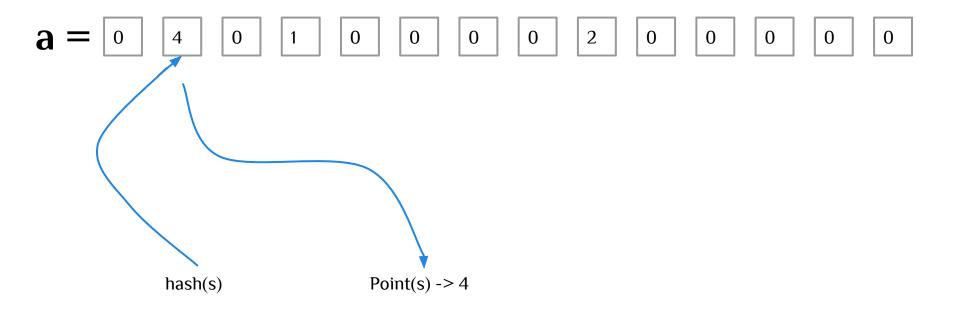


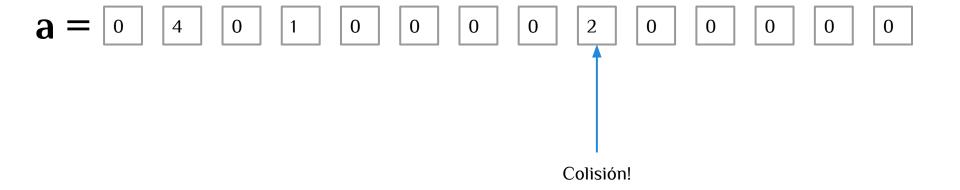


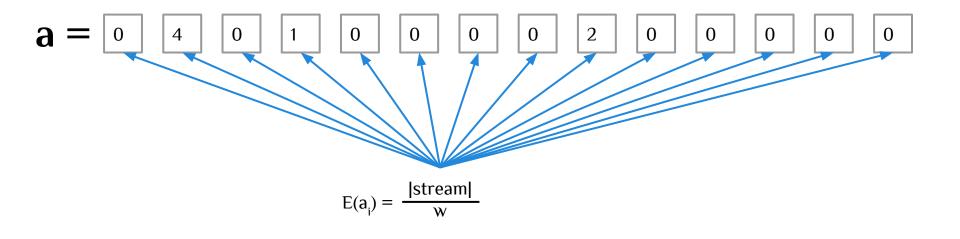




$$\mathbf{a} = |0| |4| |0| |1| |0| |0| |0| |2| |0| |0| |0| |0|$$

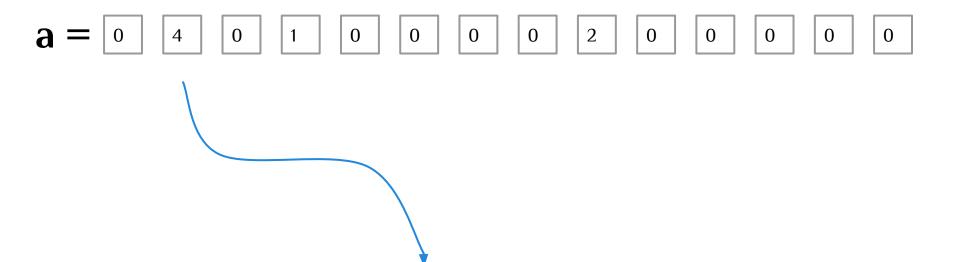






 $Count(s) \le Point(s) \le Count(s) + e|stream|/w$

(con alguna probabilidad)



Count(s) -> $4 \pm 7/14 = 4 \pm \frac{1}{2}$

Mejorando la probabilidad de error

count

count₁= count₂= $count_z =$ count_d = hash₂(s) hash₁(s) hash_z(s) hash_d(s)

Point(s)

$$Point(s) = min_j count_j[hash(s)]$$

Elegimos el mejor caso!

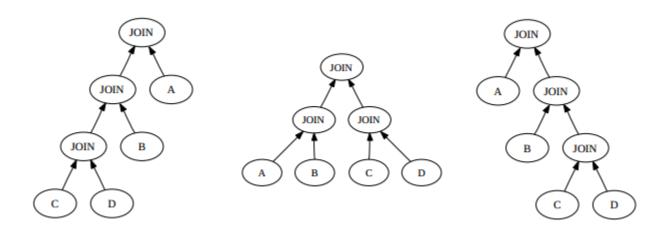
Garantias

$$1 - e^{-d}$$

Ejemplos

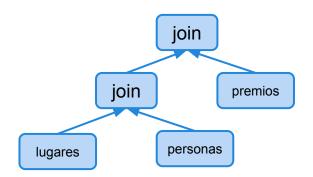
w	d	size	error	probabilidad
		w * d	e / w * 100%	1-e ^{-d}
6	3	18	45%	0.95
272	6	1632	1%	0.995
2719	8	21752	0.1%	0.9995

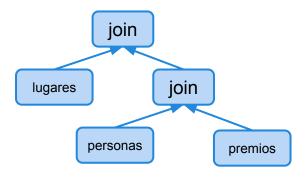
Usos: Query planning



Usos: Query planning

Donde nacieron las personas que ganaron premios nobel?





pg_stats

Name	Туре	Description
schemaname	name	Name of schema containing table
tablename	name	Name of table
attname	name	Name of the column described by this row
n_distinct	real	If greater than zero, the estimated number of distinct values in the column. If less than zero, the negative of the number of distinct values divided by the number of rows. (The negated form is used when ANALYZE believes that the number of distinct values is likely to increase as the table grows; the positive form is used when the column seems to have a fixed number of possible values.) For example, -1 indicates a unique column in which the number of distinct values is the same as the number of rows.
most_common_vals	anyarra y	A list of the most common values in the column. (Null if no values seem to be more common than any others.)
most_common_freqs	real[]	A list of the frequencies of the most common values, i.e., number of occurrences of each divided by total number of rows. (Null when most_common_vals is.)
histogram_bounds	anyarra y	A list of values that divide the column's values into groups of approximately equal population. The values in most_common_vals, if present, are omitted from this histogram calculation. (This column is null if the column data type does not have a < operator or if the most_common_vals list accounts for the entire population.)
correlation	real	Statistical correlation between physical row ordering and logical ordering of the column values. This ranges from -1 to +1. When the value is near -1 or +1, an index scan on the column will be estimated to be cheaper than when it is near zero, due to reduction of random access to the disk. (This column is null if the column data type does not have a < operator.)
most_common_elems	anyarra y	A list of non-null element values most often appearing within values of the column. (Null for scalar types.)
most_common_elem_f reqs	real[]	A list of the frequencies of the most common element values, i.e., the fraction of rows containing at least one instance of the given value.
elem_count_histogram	real[]	A histogram of the counts of distinct non-null element values within the values of the column, followed by the average

number of distinct non-null elements. (Null for scalar types.)

The end.

Preguntas?
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