

Session 4

White-Box Testing

3.5 Data Flow Testing

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Dataflow Coverage

- Based on the idea that program paths along which variables are **defined** and then **used** should be covered.
- A **family of path selection criteria** has been defined, each providing a different degree of coverage.

Data Flow Testing

- definition-use path
- definition-clear path

Variable Definitions and Uses

- A program variable is **DEFINED** when it appears:
 - on the *left* hand side of an assignment statement (e.g., `Y := 17`)
 - in an input statement (e.g., `input(Y)`)
 - as an OUT parameter in a subroutine call (e.g., `DOIT(X:IN,Y:OUT)`)

Variable Definitions and Uses (cont'd)

- A program variable is **USED** when it appears:
 - on the *right* hand side of an assignment statement (e.g., $Y := X + 17$)
 - as an IN parameter in a subroutine or function call (e.g., $Y := \text{SQRT}(X)$)
 - in the predicate of a branch statement (e.g., if $X > 0$ then...)

Variable Definitions and Uses (cont'd)

- Use of a variable in the predicate of a branch statement is called a ***p-use*** ("p-use"). Any other use is called a ***c-computation-use*** ("c-use").
- For example, in the program statement:

```
If (X>0) then  
    print(Y)  
end_if_then
```

there is a p-use of X and a c-use of Y.

Variable Definitions and Uses (cont'd)

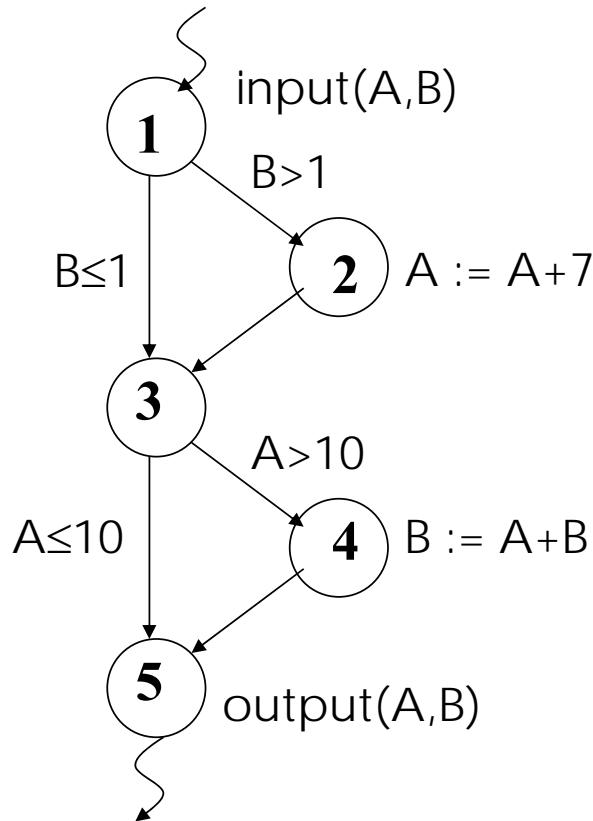
- A variable can also be used and then re-defined in a single statement when it appears:
 - on *both* sides of an assignment statement (e.g., $\textcolor{orange}{Y} := \textcolor{orange}{Y} + X$)
 - as an IN/OUT parameter in a subroutine call (e.g., `INCREMENT($\textcolor{orange}{Y}$:IN/OUT)`)

Other Dataflow Terms and Definitions

- A *definition-use pair (“du-pair”)* with respect to a variable v is a double (d,u) such that d is a node in the program’s flow graph at which v is defined, u is a node or edge at which v is used.
- Note: Broad definition.

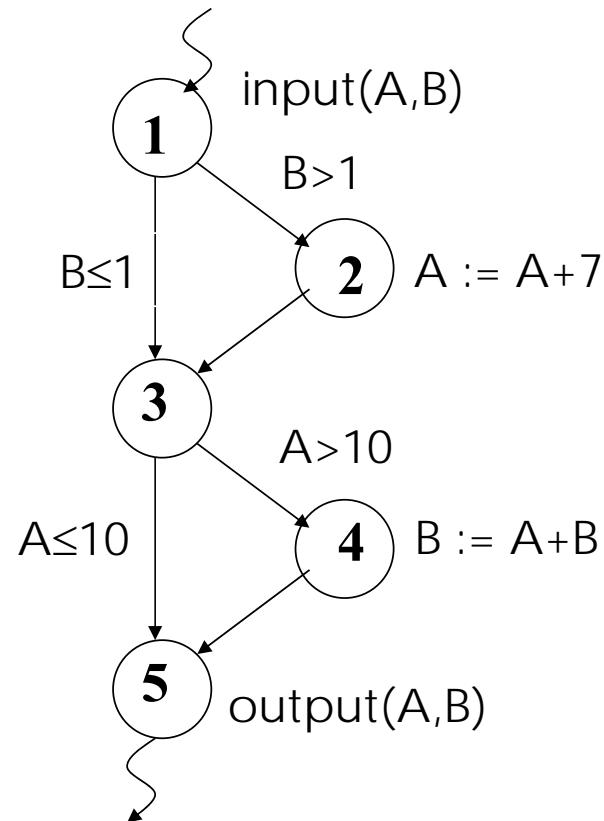
Example 1

1. `input(A,B)`
- `if (B>1) then`
2. `A := A+7`
- `end_if`
3. `if (A>10) then`
4. `B := A+B`
- `end_if`
5. `output(A,B)`



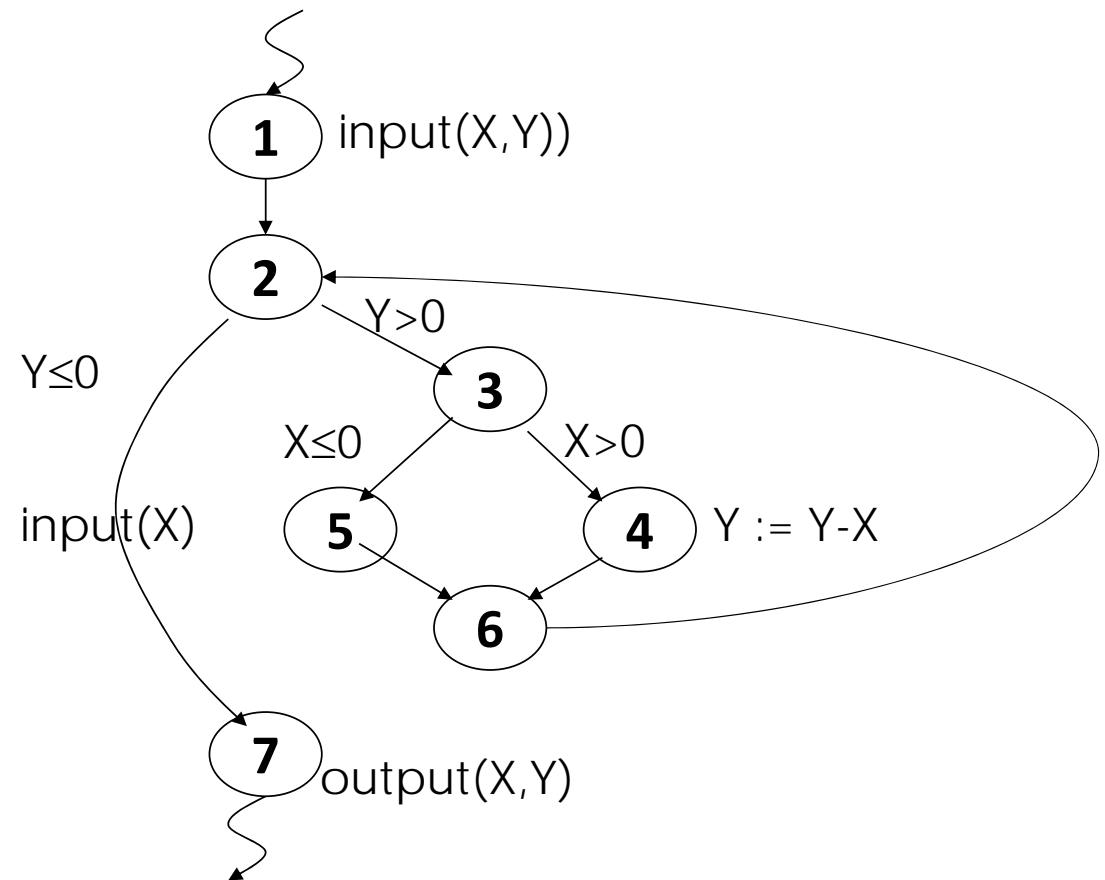
Identifying DU-Pairs – Variable A

<u>du-pair</u>	<u>path(s)</u>
(1,2)	<1,2>
(1,4)	<1,3,4>
(1,5)	<1,3,4,5>
	<1,3,5>
(1,<3,4>)	<1,3,4>
(1,<3,5>)	<1,3,5>
(2,4)	<2,3,4>
(2,5)	<2,3,4,5>
	<2,3,5>
(2,<3,4>)	<2,3,4>
(2,<3,5>)	<2,3,5>



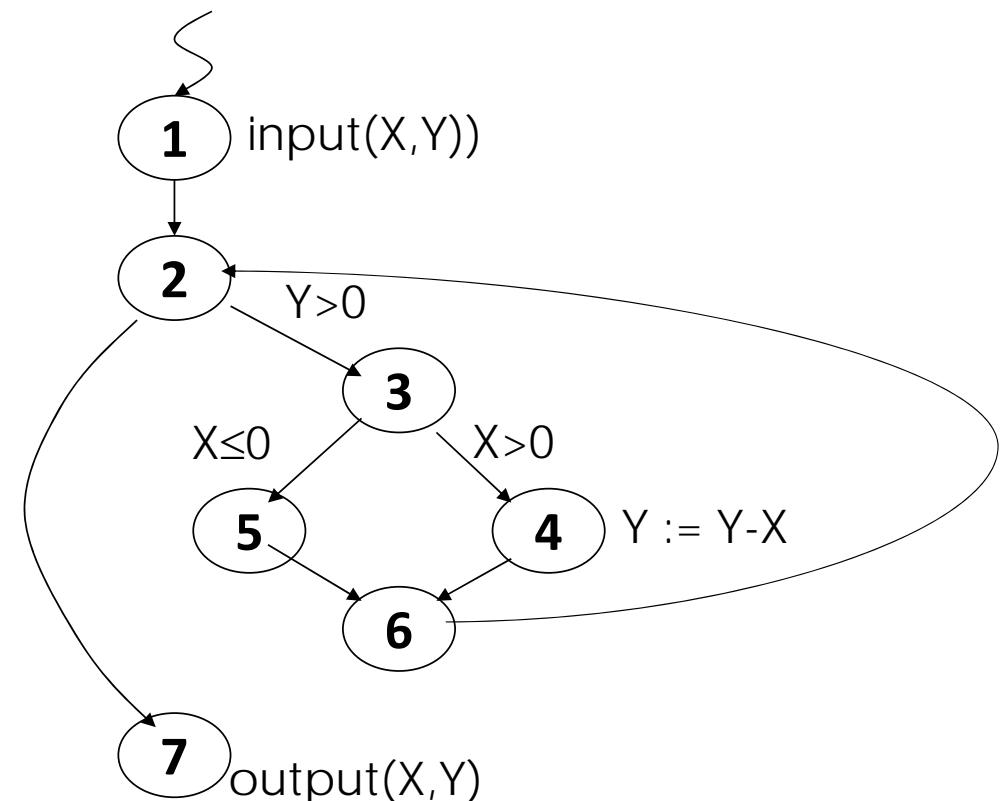
Exercise

1. `input(X,Y)`
2. `while (Y>0) do`
3. `if (X>0) then`
4. `Y := Y-X`
- `else`
5. `input(X)`
- `end_if_then_else`
6. `end_while`
7. `output(X,Y)`



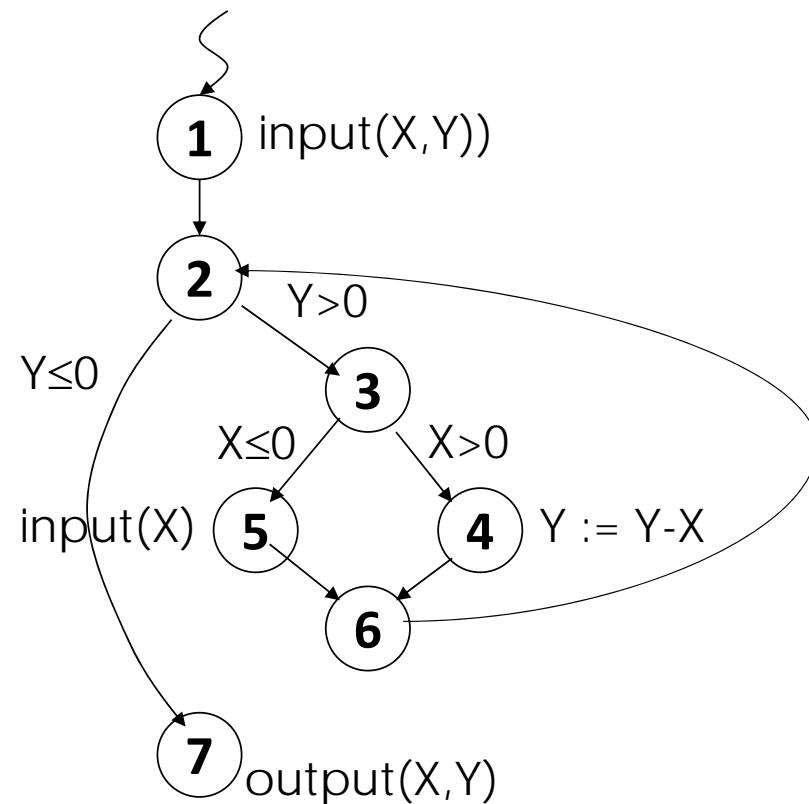
Identifying DU-Pairs – Variable X

<u>du-pair</u>	<u>path(s)</u>
(1,4)	<1,2,3,4>
	<1,2,3,4,(6,2,3,4)*>
(1,7)	<1,2,7>
	<1,2,3,4,6,2,7>
	<1,2,3,4,6,(2,3,4,6)*,2,7>
(1,<3,4>)	<1,2,3,4>
	<1,2,3,4,(6,2,3,4)*>
(1,<3,5>)	<1,2,3,5>
(5,4)	<5,6,2,3,4>
	<5,6,2,3,4,(6,2,3,4)*>



Identifying DU-Pairs – Variable X

<u>du-pair</u>	<u>path(s)</u>
(5,7)	<5,6,2,7> [†]
	<5,6,2,3,4,6,2,7>
	<5,6,2,3,4,6,(2,3,4,6)*,2,7>
(5,<3,4>)	<5,6,2,3,4>
	<5,6,2,3,4,(6,2,3,4)*>
(5,<3,5>)	<5,6,2,3,5>
	<5,6,2,3,4,6,2,3,5> [†]
	<5,6,2,3,4,6,(2,3,4,6)*,2,3,5> †
[†] infeasible	



Data Flow Testing

- definition-use path
- definition-clear path

def-clear path

- A path is *definition clear (“def-clear”)* with respect to a variable v if there is *no re-definition* of v within the path.
- A *definition-use pair (“du-pair”)* with respect to a variable v is a double (d, u) such that d is a node in the program’s flow graph at which v is defined, u is a node or edge at which v is used, and there is a *def-clear path* with respect to v from d to u .
 - Rigorous definition 严谨版本定义

More Dataflow Terms and Definitions

- A path (either partial or complete) is **simple** if all edges within the path are distinct (i.e., different).
- A path is **loop-free** if all nodes within the path are distinct (i.e., different).

Simple and Loop-Free Paths

path	Simple?	Loop-free?
<1,3,4,2>	Yes	Yes
<1,2,3,2>	Yes	No
<1,2,3,1,2>	No	No
<1,2,3,2,4>	Yes	No

Simple and Loop-Free Paths (cont'd)

Which is *stronger*, **simple** or **loop-free**?

环路为1

无环路

More Dataflow Terms and Definitions

A path $\langle n_1, n_2, \dots, n_j, n_k \rangle$ is a **du-path** with respect to a variable v if v is defined at node n_1 and either:

1. there is a **c-use** of v at node n_k and $\langle n_1, n_2, \dots, n_j, n_k \rangle$ is a def-clear **simple** path,

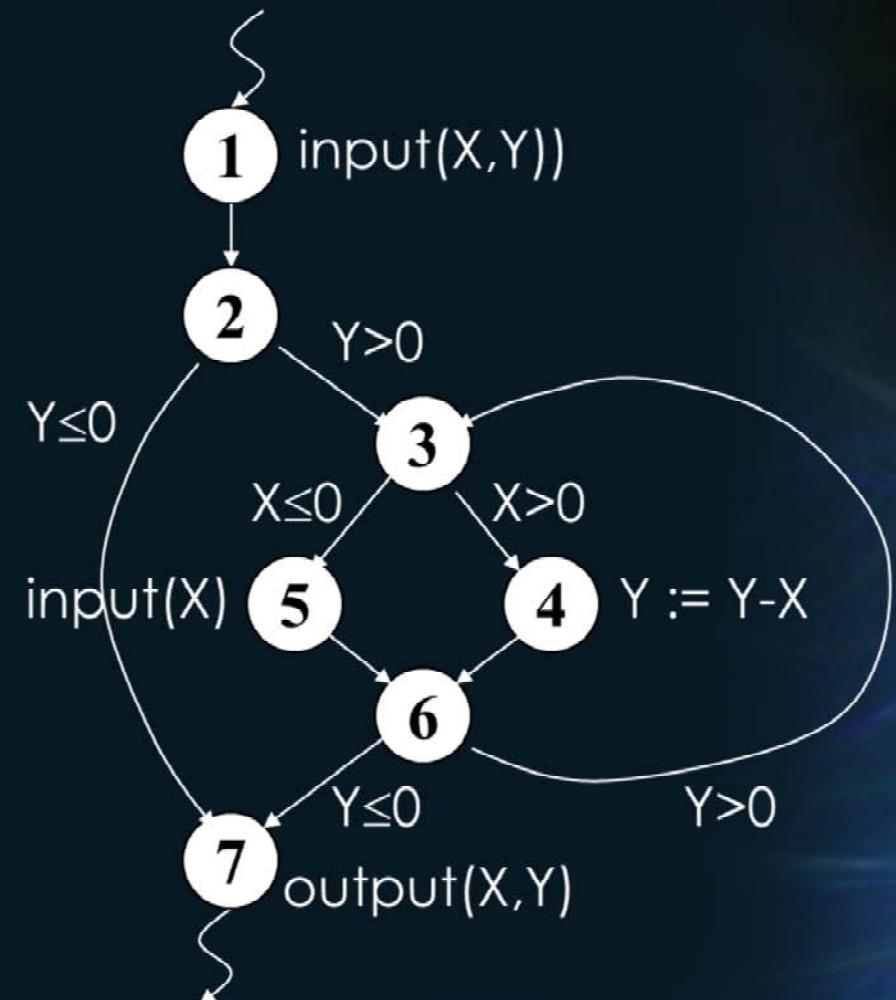
or

1. there is a **p-use** of v at edge $\langle n_j, n_k \rangle$ and $\langle n_1, n_2, \dots, n_j \rangle$ is a def-clear **loop-free** path.

NOTE!

Identifying du-paths

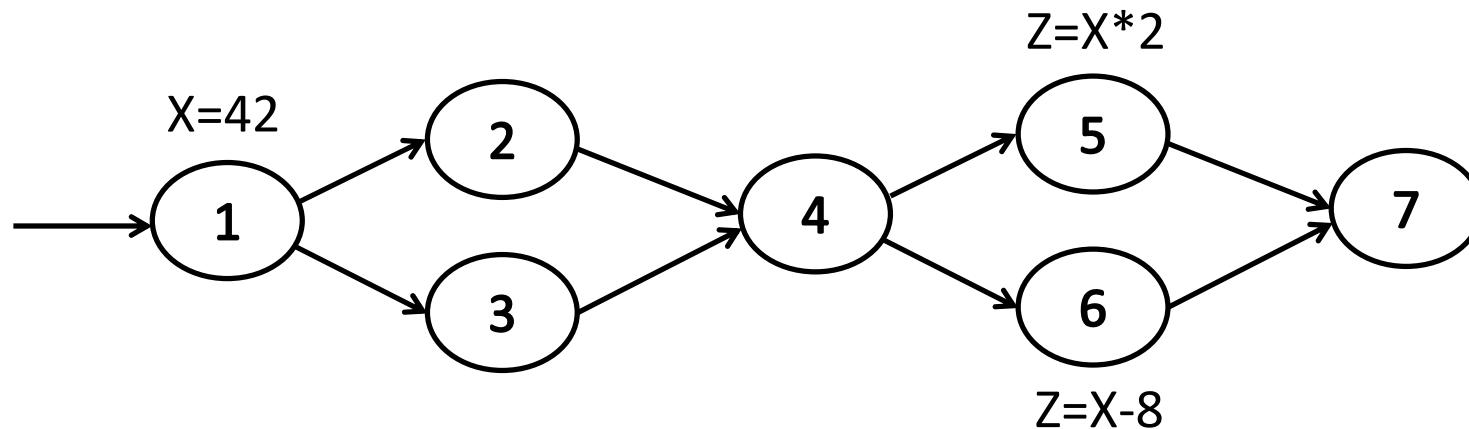
<u>du-pair</u>	<u>path(s)</u>	<u>du-path?</u>
X: (5,7)	<5,6,7> †	Yes
	<5,6,3,4,6,7>	Yes
	<5,6,3,4,6,(3,4,6)*,7>	No
X: (5,<3,4>)	<5,6,3,4>	Yes
	<5,6,3,4,(6,3,4)*>	No
X: (5,<3,5>)	<5,6,3,5>	Yes
	<5,6,3,4,6,3,5> †	No
	<5,6,3,4,6,(3,4,6)*,3,5> †	No
† infeasible		



Dataflow Test Coverage Criteria

- *All-Defs*: for every program variable v , at least one def-clear path from **every definition** of v to **at least one c-use or one p-use** of v must be covered.

Example



All-defs for X

$\langle 1, 2, 4, 5 \rangle$

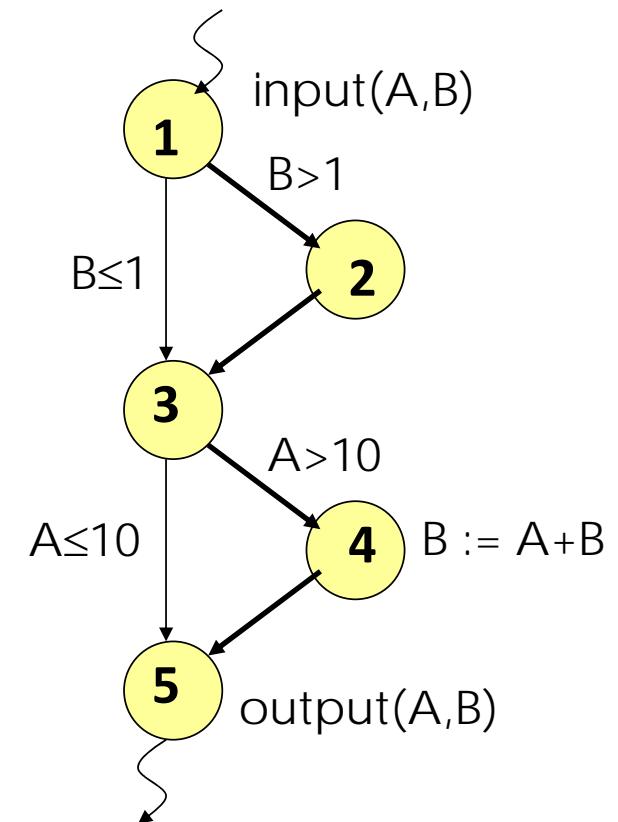
Dataflow Test Coverage Criteria (cont'd)

- Consider a test case executing path:

1. **<1,2,3,4,5>**

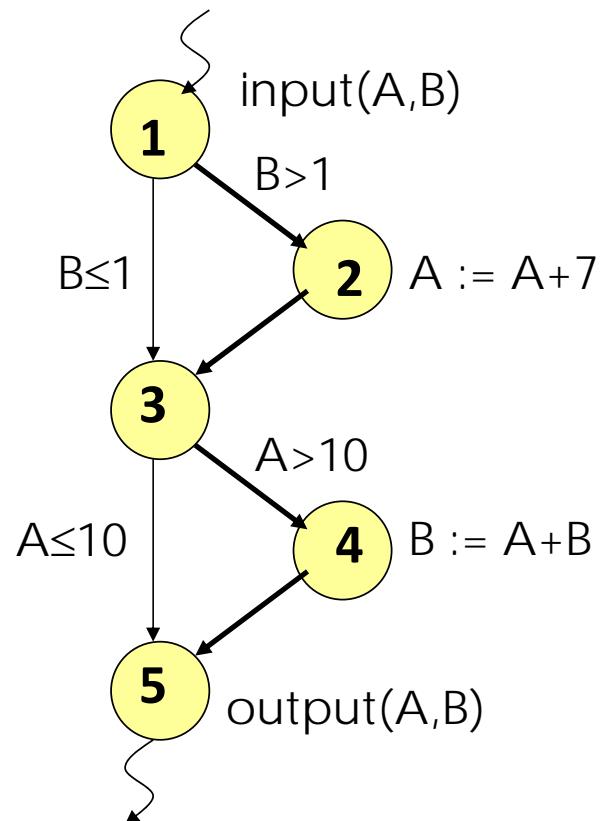
- Identify all def-clear paths covered (i.e., subsumed) by this path for each variable.

- Are all definitions for each variable associated with at least one of the subsumed def-clear paths?



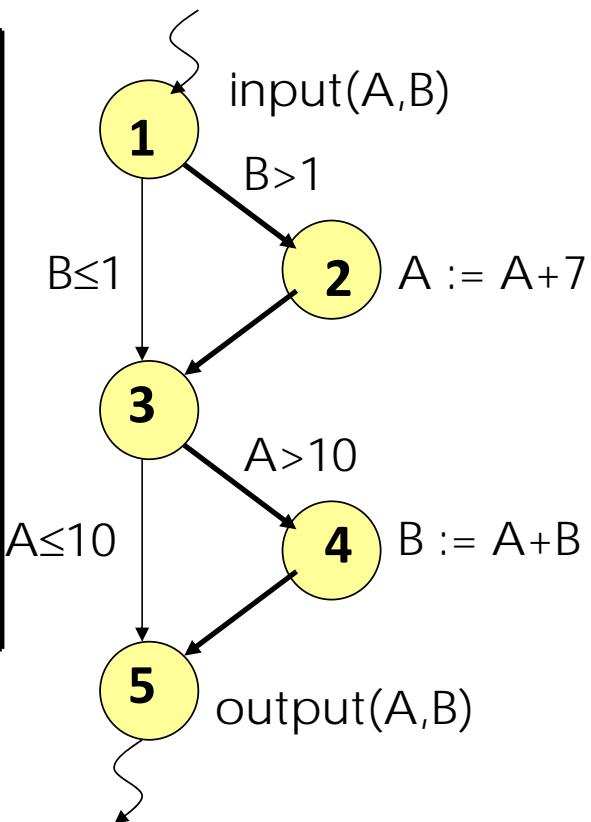
Def-Clear Paths Subsumed by $\langle 1,2,3,4,5 \rangle$ for Variable A

<u>du-pair</u>	<u>path(s)</u>
(1,2)	$\langle 1,2 \rangle$ ✓
(1,4)	$\langle 1,3,4 \rangle$
(1,5)	$\langle 1,3,4,5 \rangle$
	$\langle 1,3,5 \rangle$
(1,<3,4>)	$\langle 1,3,4 \rangle$
(1,<3,5>)	$\langle 1,3,5 \rangle$
(2,4)	$\langle 2,3,4 \rangle$ ✓
(2,5)	$\langle 2,3,4,5 \rangle$ ✓
	$\langle 2,3,5 \rangle$
(2,<3,4>)	$\langle 2,3,4 \rangle$ ✓
(2,<3,5>)	$\langle 2,3,5 \rangle$



Def-Clear Paths Subsumed by $\langle 1,2,3,4,5 \rangle$ for Variable B

<u>du-pair</u>	<u>path(s)</u>
(1,4)	$\langle 1,2,3,4 \rangle$ ✓
	$\langle 1,3,4 \rangle$
(1,5)	$\langle 1,2,3,5 \rangle$
	$\langle 1,3,5 \rangle$
(4,5)	$\langle 4,5 \rangle$ ✓
(1, $\langle 1,2 \rangle$)	$\langle 1,2 \rangle$ ✓
(1, $\langle 1,3 \rangle$)	$\langle 1,3 \rangle$



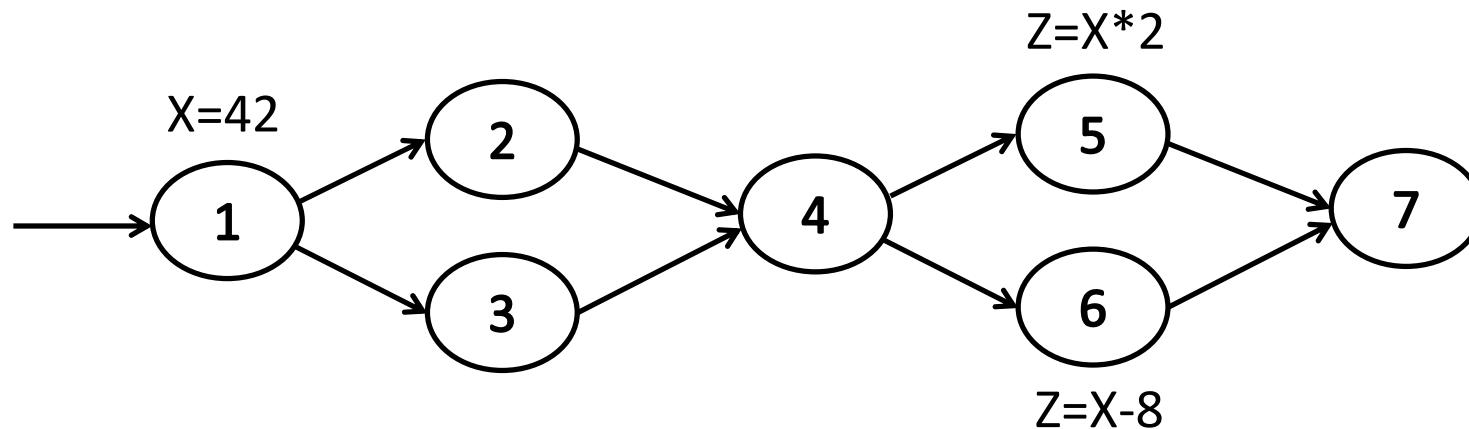
Dataflow Test Coverage Criteria (cont'd)

- Since **<1,2,3,4,5>** covers at least one def-clear path from every definition of A/B to at least one c-use or p-use of A/B,
All-Defs coverage is achieved.

Dataflow Test Coverage Criteria (cont'd)

- **All-Uses:** for **every program variable** v , at least one def-clear path from **every definition** of v to **every c-use** and **every p-use** of v must be covered.

Example



All-defs for X

$\langle 1, 2, 4, 5 \rangle$

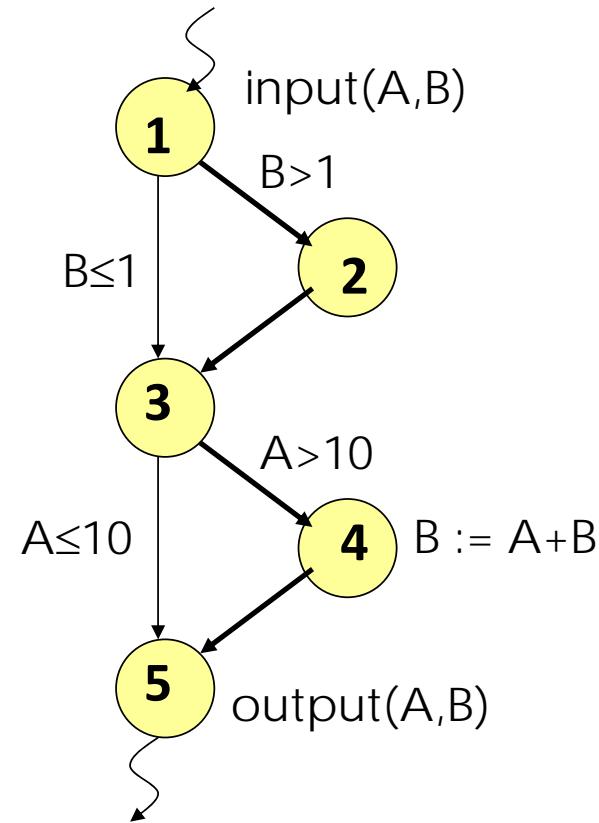
All-uses for X

$\langle 1, 2, 4, 5 \rangle$

$\langle 1, 2, 4, 6 \rangle$

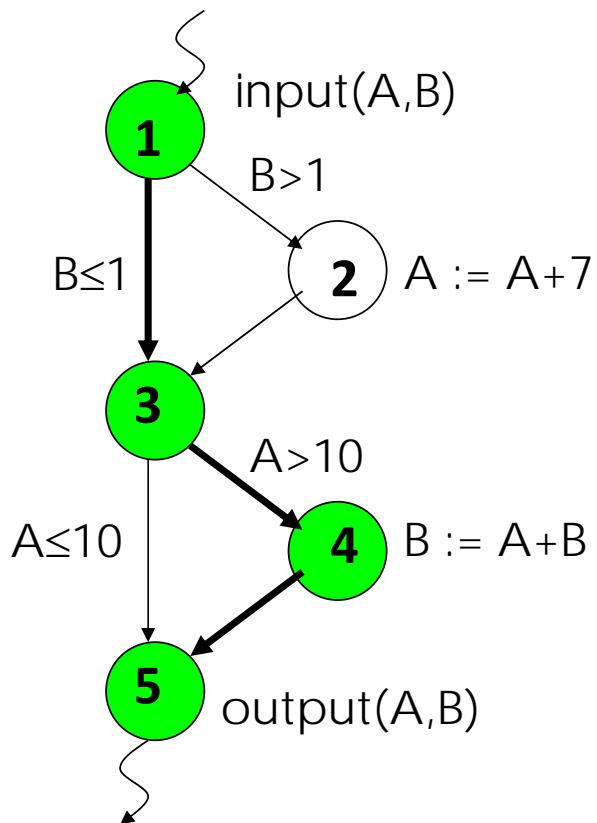
Dataflow Test Coverage Criteria (cont'd)

- Consider additional test cases executing paths:
 1. **<1,2,3,4,5>**
 2. **<1,3,4,5>**
 3. **<1,2,3,5>**
- Do all three test cases provide All-Uses coverage?



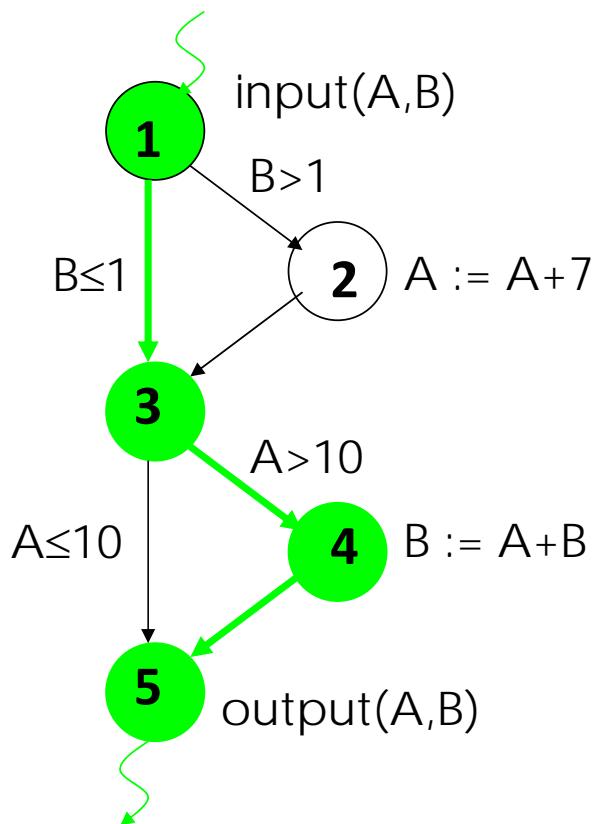
Def-Clear Paths Subsumed by $\langle 1,3,4,5 \rangle$ for Variable A

<u>du-pair</u>	<u>path(s)</u>
(1,2)	$\langle 1,2 \rangle$ ✓
(1,4)	$\langle 1,3,4 \rangle$ ✓
(1,5)	$\langle 1,3,4,5 \rangle$ ✓
	$\langle 1,3,5 \rangle$
(1,<3,4>)	$\langle 1,3,4 \rangle$ ✓
(1,<3,5>)	$\langle 1,3,5 \rangle$
(2,4)	$\langle 2,3,4 \rangle$ ✓
(2,5)	$\langle 2,3,4,5 \rangle$ ✓
	$\langle 2,3,5 \rangle$
(2,<3,4>)	$\langle 2,3,4 \rangle$ ✓
(2,<3,5>)	$\langle 2,3,5 \rangle$



Def-Clear Paths Subsumed by $\langle 1,3,4,5 \rangle$ for Variable B

<u>du-pair</u>	<u>path(s)</u>
(1,4)	$\langle 1,2,3,4 \rangle$ ✓
	$\langle 1,3,4 \rangle$ ✓
(1,5)	$\langle 1,2,3,5 \rangle$
	$\langle 1,3,5 \rangle$
(4,5)	$\langle 4,5 \rangle$ ✓ ✓
(1, $\langle 1,2 \rangle$)	$\langle 1,2 \rangle$ ✓
(1, $\langle 1,3 \rangle$)	$\langle 1,3 \rangle$ ✓



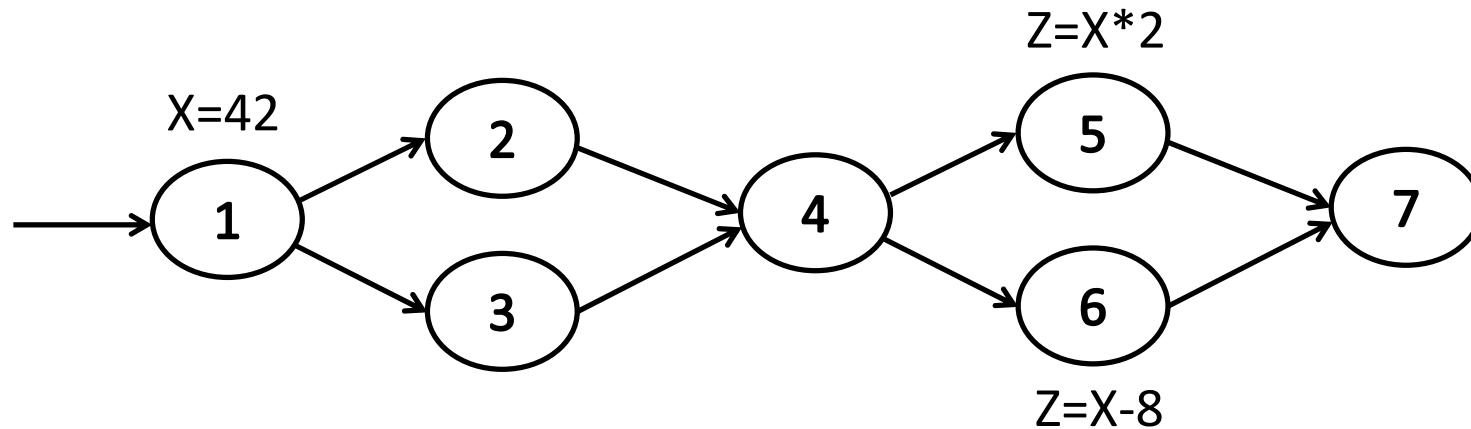
Dataflow Test Coverage Criteria (cont'd)

- Since none of the three test cases covers the du-pair $(1, <3, 5>)$ for variable A, All-Uses Coverage is not provided.

Another Dataflow Test Coverage Criterion

- **All-DU-Paths:** for **every program variable v**, **every du-path** from **every definition** of v to **every c-use** and **every p-use** of v must be covered.

Example



All-defs for X

$\langle 1, 2, 4, 5 \rangle$

All-uses for X

$\langle 1, 2, 4, 5 \rangle$

$\langle 1, 2, 4, 6 \rangle$

All-DU-Paths for X

$\langle 1, 2, 4, 5 \rangle$

$\langle 1, 3, 4, 5 \rangle$

$\langle 1, 2, 4, 6 \rangle$

$\langle 1, 3, 4, 6 \rangle$

Summary of White-Box Coverage Relationships

