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# 1 Contest

## 1.1 Header

```
#pragma once

#include <bits/stdc++.h>
using namespace std;

template <class TH>
void _dbg(const char *sdbg, TH h) { cerr << sdbg << '=' << h << endl; }

template <class TH, class... TA>
void _dbg(const char *sdbg, TH h, TA... a)
{
    while (*sdbg != ',')
        cerr << *sdbg++;
    cerr << '=' << h << ', ';
    _dbg(sdbg + 1, a...);
}

template <class L, class R>
ostream &operator<<(ostream &os, pair<L, R> p)
{
    return os << "(" << p.first << ", " << p.second << ")";
}

template <class Iterable, class = typename enable_if<!is_same<string,
    Iterable>::value>::type>
auto operator<<(ostream &os, Iterable v) -> decltype(os << *begin(v))
{
    os << "[";
    for (auto vv : v)
        os << vv << ", ";
    return os << "]";
}

#define debug(...) _dbg(__VA_ARGS__, __VA_ARGS__)

typedef pair<int, int> pii;
typedef long long ll;
typedef long double ld;

const int inf = 0x3f3f3f3f;
const long long infll = 0x3f3f3f3f3f3f3f3fll;

#define sz(x) ((int)(x).size())

// Return 1 if x > 0, 0 if x == 0 and -1 if x < 0.
template <class T>
```

```

int sign(T x) { return (x > 0) - (x < 0); }

template <class T>
T abs(const T &x) { return (x < T(0)) ? -x : x; }

// Pretty good compilation command:
// g++ -g a.cpp --std=c++14 -Wall -Wextra -Wno-unused-result -Wconversion -
// Wfatal-errors -fsanitize=undefined,address -o a.out

// int main()
// {
//   cin.sync_with_stdio(0);
//   cin.tie(0);
//   cin.exceptions(cin.failbit);
// }

```

## 1.2 Sample Debug

```

32cfcc  #include "header.hpp"
d41d8c
13a4b1  int main(void)
f95b70  {
3e8410      int a = 11, b = 12, c = 13;
b9ee34      vector<vector<int>> v = {{a, b}, {c}, {0, 1}};
2a803c      set<int> s = {a, b};
6b3b13      map<double, int> m;
af2bd1      m[2.5] = 2;
37a428      m[-3.1] = 3;
d41d8c
632de9      map<string, int> tab;
34aeb8      tab["abc"] = (int) '\t' + 'b' + '\r';
cb5ab5      tab["abz"] = (int) 'a' + '\n' + 'z';
bd6def      int array[3] = {1, 2, 5};
d41d8c
5939a6      debug(a, b, c);
fb9ee1      debug(v);
b45aab      debug(s, m);
3cf91d      debug(tab);
d95ee2      debug(array); // This one does not work.
cbb184  }
Full file hash: 6d218a

```

## 1.3 Hash Code

```

#!/bin/bash
# Hashes each line of a file, ignoring all whitespace and comments (multi-
# line comments will be bugged).

while IFS= read -r line; do # Loops lines of stdin.
    echo "$line" | cpp -dD -P -fpreprocessed | tr -d '[:space:]' | md5sum | cut
    -c-6 | tr -d '[:space:]';

```

```
    echo "  $line"; # Before $line is a tab.  
done
```

## 2 Data Structures

### 2.1 BIT

```

5d1131 #include "../contest/header.hpp"
d41d8c
d41d8c /*
ccd9cd     BIT: element update, range sum query and sum lower_bound in O(log(N
    )),
2716d1     Represents an array of elements in range [1, N].
c4c9bd */
d41d8c
4fce64 template <class T>
2d55ba struct bit
f95b70 {
b9a249     int n, LOGN;
2262a2     vector<T> val;
695052     bit(int _n) : n(_n), LOGN(log2(n + 1)), val(_n + 1, 0) {}
d41d8c
d41d8c     // val[pos] += x
b29da0     void update(int pos, T x)
f95b70     {
259a9f         for (int i = pos; i <= n; i += -i & i)
ac55c8             val[i] += x;
cbb184     }
d41d8c
d41d8c     // sum of range [1, pos]
8a835d     T query(int pos)
f95b70     {
56622d         T retv = 0;
ac430c         for (int i = pos; i > 0; i -= -i & i)
106953             retv += val[i];
6272cf         return retv;
cbb184     }
d41d8c
d41d8c     // min pos such that sum of [1, pos] >= sum, or n + 1 if none
exists.
79d23b     int lower_bound(T x)
f95b70     {
501ce1         T sum = 0;
bec7a6         int pos = 0;
d41d8c
51d707         for (int i = LOGN; i >= 0; i--)
0328f7             if (pos + (1 << i) <= n && sum + val[pos + (1 << i)] < x)
420193                 sum += val[pos += (1 << i)];
d41d8c
7e21de         return pos + 1; // pos will have position of largest value less
than x.
cbb184     }
2145c1 };

```

Full file hash: 6acfcc

## 2.2 BIT2D

```

5d1131 #include "../contest/header.hpp"
d41d8c
d41d8c /*
caf843     BIT: element update, range sum query in  $O(\log(n) * \log(m))$ . This
          can also be generalized for 3d.
a6cfe6     Represents a matrix of elements in range [1 ... n][1 ... m].
c4c9bd */
d41d8c
4fce64 template <class T>
f6f3a7 struct bit2d
f95b70 {
14e0a7     int n, m;
f7ea55     vector<vector<T>> val;
9c8214     bit2d(int _n, int _m) : n(_n), m(_m), val(_n + 1, vector<T>(_m + 1,
          0)) {}
d41d8c
d41d8c     // val[i][j] += x
4460cb     void update(int r, int c, T x)
f95b70     {
9e45d9         for (int i = r; i <= n; i += -i & i)
13d333             for (int j = c; j <= m; j += -j & j)
ff237f                 val[i][j] += x;
cbb184     }
d41d8c
d41d8c     // sum of positions (1 ... r, 1 ... c)
450f85     T query(int r, int c)
f95b70     {
56622d         T retv = 0;
bc7409         for (int i = r; i > 0; i -= -i & i)
d53722             for (int j = c; j > 0; j -= -j & j)
86df71                 retv += val[i][j];
6272cf         return retv;
cbb184     }
d41d8c
d41d8c     // sum of positions (ri ... rf, ci ... cf). (1 <= ri <= rf <= n)
          and (1 <= ci <= cf <= m). TODO: test me.
bdc664     T query_rect(int ri, int ci, int rf, int cf)
f95b70     {
6072bc         return query(rf, cf) - query(rf, ci - 1) - query(ri - 1, cf) +
          query(ri - 1, ci - 1);
cbb184     }
2145c1 };

```

Full file hash: a4a33c

## 2.3 Dynamic Seg

```

5d1131 #include "../contest/header.hpp"

```

```

d41d8c
d41d8c  /*
811629     Segment tree with dynamic memory allocation and arbitrary interval.
91eb69     Every operation is O(log(r-l))
b5a53f     Uses O(min(r-l, n*log(r-l))) memory, where n is the number of
           insertions.
d41d8c
ca2095     Constraints:
3dcfba     Segment tree range [l, r] must be such that 0 <= l <= r.
d41d8c
3db72f     Author: Arthur Pratti Dadalto
c4c9bd  */
d41d8c
4fce64  template<class T>
e4accb  struct node
f95b70  {
f48ea0      T val;
af32d9      node *left, *right;
d41d8c
995125      T get(int l, int r, int a, int b)
f95b70      {
47234b          if (l == a && r == b)
d943f4              return val;
814ad2          int mid = (l + 0ll + r) / 2;
f890f2          if (b <= mid)
ac57ce              return left ? left->get(l, mid, a, b) : 0;
a54f0c          else if (a > mid)
1c7837              return right ? right->get(mid + 1, r, a, b) : 0;
2954e9          else
9b1cb1              return (left ? left->get(l, mid, a, mid) : 0) + (right ? right
->get(mid + 1, r, mid + 1, b) : 0);
cbb184      }
d41d8c
14d5ea      void update(int l, int r, int a, T x)
f95b70      {
bd3398          if (l == r)
c43fe0              val = x;
2954e9          else
f95b70              {
814ad2                  int mid = (l + 0ll + r) / 2;
a49729                  if (a <= mid)
1ec55a                      (left ? left : (left = new node()))->update(l, mid, a, x);
2954e9                  else
92fe63                      (right ? right : (right = new node()))->update(mid + 1, r, a,
x);
d41d8c
dd51dd                      val = (left ? left->val : 0) + (right ? right->val : 0);
cbb184              }
cbb184      }
2145c1  };

```



d41d8c

Full file hash: eef0d8

## 2.4 Linear Container

5d1131 #include "../contest/header.hpp"

d41d8c

d41d8c /\*

1ee0c3 Line Container (most common for convex hull trick). Amortized  $O(\log N)$  per operation.48cf95 Container where you can add lines of the form  $kx+m$ , and query maximum values at points  $x$ .

dc45cd Useful for dynamic programming.

d41d8c

1d1558 Source: <https://github.com/kth-competitive-programming/kactl/blob/master/content/contest/template.cpp>

c4c9bd \*/

d41d8c

3fe318 struct line

f95b70 {

3e2604 mutable ll k, m, p;

889941 bool operator&lt;(const line &amp;o) const { return k &lt; o.k; }

abfd1f bool operator&lt;(ll x) const { return p &lt; x; }

2145c1 };

d41d8c

0c8ce5 struct line\_container : multiset&lt;line, less&lt;&gt;&gt;

f95b70 {

d41d8c // (for doubles, use inf = 1/.0, div(a,b) = a/b)

f5e3e7 const ll inf = LLONG\_MAX;

d41d8c

9608c5 ll div(ll a, ll b)

f95b70 { // floored division

353cf0 return a / b - ((a ^ b) &lt; 0 &amp;&amp; a % b);

cbb184 }

d41d8c

9c092f bool isect(iterator x, iterator y)

f95b70 {

f959d1 if (y == end())

f95b70 {

09a75e x-&gt;p = inf;

d1fe4d return false;

cbb184 }

3cca77 if (x-&gt;k == y-&gt;k)

83e301 x-&gt;p = x-&gt;m &gt; y-&gt;m ? inf : -inf;

2954e9 else

b4284e x-&gt;p = div(y-&gt;m - x-&gt;m, x-&gt;k - y-&gt;k);

870ec6 return x-&gt;p &gt;= y-&gt;p;

cbb184 }

d41d8c

928f4b void add(ll k, ll m)

```

f95b70      {
116e6c      auto z = insert({k, m, 0}), y = z++, x = y;
2d9d80      while (isect(y, z))
96cee5      z = erase(z);
d94b4e      if (x != begin() && isect(--x, y))
c07d21      isect(x, y = erase(y));
57dd20      while ((y = x) != begin() && (--x)->p >= y->p)
77462a      isect(x, erase(y));
cbb184      }
d41d8c
e8b5c2      ll query(ll x)
f95b70      {
229883      assert(!empty());
7d13b8      auto l = *lower_bound(x);
96a2bc      return l.k * x + l.m;
cbb184      }
2145c1  };
d41d8c
Full file hash: 66b35a

```

## 2.5 Min Queue

```

d41d8c  /*
958401   max(min) queue with O(1) get_max(min).
d41d8c
f67dcb   Tips:
c53808   - Useful for sliding window 1D and 2D.
af9dc1   - For 2D problems, you will need to pre-compute another matrix,
55e3e9   by making a row-wise traversal, and calculating the min/max value
79c288   beginning in each cell. Then you just make a column-wise traverse
b21db2   as they were each an independent array.
c4c9bd  */
d41d8c
8f0a66  struct max_queue
f95b70  {
84841a   queue<ll> q;
889d23   deque<ll> s;
d41d8c
dbb27b   int size()
f95b70   {
593f12   return (int)q.size();
cbb184   }
d41d8c
a1fe24   void push(ll val)
f95b70   {
d41d8c   // while (!s.empty() && s.back() > val) -> for a min_queue
1cb658   while (!s.empty() && s.back() < val) // for a max_queue
342ca4   s.pop_back();
fcc849   s.push_back(val);
d41d8c

```

```

380c99     q.push(val);
cbb184   }
d41d8c
d99fc4   void pop()
f95b70   {
7a8432     ll u = q.front();
833270     q.pop();
d41d8c
de7036     if (!s.empty() && s.front() == u)
784c93       s.pop_front();
cbb184   }
d41d8c
ba28bf   ll get_max()
f95b70   {
eccd4b     return s.front(); // same for min and max queue
cbb184   }
d41d8c
2145c1   };
Full file hash: 82549d

```

## 2.6 Persistent Seg

```

5d1131   #include "../contest/header.hpp"
d41d8c
d41d8c   /*
a032aa     Persistent Segment Tree:
115cd2       Segment tree that stores all previous versions of itself.
91eb69       Every operation is  $O(\log(r-l))$ 
ad671a       Uses  $O(n \cdot \log(r-l))$  memory, where  $n$  is the number of updates.
d41d8c
b95cae     Usage:
aca1a7       A new root is created for every persistent update (p_update) and
             returned.
0d5abd       Queries can be performed on any root as if it were a usual
             segment tree.
61a20d       You should keep a list of roots. Something like:
072987         vector<node*> roots = {new node()};
bf8bc0         roots.push_back(p_update(roots.back(), 0, 2*MAXV, a[i] + MAXV,
             v + 1));
d41d8c
ca2095     Constraints:
3dcfba       Segment tree range  $[l, r]$  must be such that  $0 \leq l \leq r$ .
d41d8c
3db72f     Author: Arthur Pratti Dadalto
c4c9bd   */
d41d8c
e4accb   struct node
f95b70   {
97f03f     int val;
af32d9     node *left, *right;

```

```

d41d8c
1f6b0f    node(int x=0) : val(x), left(NULL), right(NULL) {}
2f77b9    node(node *l, node *r) : left(l), right(r) { val = (left ? left->
    val : 0) + (right ? right->val : 0); }
d41d8c
f219f1    int get(int l, int r, int a, int b)
f95b70    {
47234b        if (l == a && r == b)
d943f4            return val;
814ad2        int mid = (l + 0ll + r) / 2;
f890f2        if (b <= mid)
ac57ce            return left ? left->get(l, mid, a, b) : 0;
a54f0c        else if (a > mid)
1c7837            return right ? right->get(mid + 1, r, a, b) : 0;
2954e9        else
9b1cb1            return (left ? left->get(l, mid, a, mid) : 0) + (right ? right
    ->get(mid + 1, r, mid + 1, b) : 0);
cbb184    }
2145c1    };
d41d8c
63f202    node *p_update(node *prev, int l, int r, int a, int x)
f95b70    {
bd3398        if (l == r)
13478f            return new node(x);
d41d8c
814ad2        int mid = (l + 0ll + r) / 2;
a49729        if (a <= mid)
b73799            return new node(p_update(prev ? prev->left : NULL, l, mid, a, x),
    prev ? prev->right : NULL);
2954e9        else
460332            return new node(p_update(prev ? prev->left : NULL, p_update(prev ? prev->
    right : NULL, mid + 1, r, a, x)));
cbb184    }
d41d8c
Full file hash: 707f69

```

## 2.7 Lazy Seg

```

2b74fa    #include <bits/stdc++.h>
ca417d    using namespace std;
d41d8c
d41d8c    /*
6f561b        Segment Tree with Lazy updates:
d8b1dc            Range update and range query in O(log(MAX_RANGE))
c329b0            Binary search on tree in O(log(MAX_RANGE))
05382c            Given as an example since it is not worth it to copy a generic
    tree during a contest.
d41d8c
e3c955        Solves: https://codeforces.com/contest/1179/problem/C
c4c9bd    */

```

```

d41d8c
ab0dbf #define MAX_RANGE 1123456
d41d8c
fd87fe int val[4 * MAX_RANGE];
802d92 int delta[4 * MAX_RANGE];
d41d8c
4ee394 #define left(i) ((i) << 1)
56e5cf #define right(i) (((i) << 1) + 1)
d41d8c
0379af void prop(int id, int l, int r)
f95b70 {
cfd4b4     if (l != r)
f95b70     {
d41d8c         // Updates need to be numerically stackable (e.g. not valid to
have a list of updates).
df541b         delta[left(id)] += delta[id];
966351         delta[right(id)] += delta[id];
cbb184     }
d41d8c
21c2c8     val[id] += delta[id]; // Node value needs to be obtainable without
propagating all the way to root.
0a8860     delta[id] = 0;
cbb184 }
d41d8c
d41d8c // Sum x in all elements in range [a, b].
f2b4f2 void update(int id, int l, int r, int a, int b, int x)
f95b70 {
addc1f     if (a == l && b == r)
f95b70     {
d50197         delta[id] += x;
b62cfe         prop(id, l, r);
cbb184     }
2954e9     else
f95b70     {
b62cfe         prop(id, l, r);
ae007b         int mid = (l + r) / 2;
f890f2         if (b <= mid)
f95b70         {
6dbd37             update(left(id), l, mid, a, b, x);
384ec5             prop(right(id), mid + 1, r);
cbb184         }
a54f0c         else if (a > mid)
f95b70         {
859d13             update(right(id), mid + 1, r, a, b, x);
221ad0             prop(left(id), l, mid);
cbb184         }
2954e9         else
f95b70         {
fc79c7             update(left(id), l, mid, a, mid, x);
04c83e             update(right(id), mid + 1, r, mid + 1, b, x);

```

```

cbb184     }
d41d8c
caf644     val[id] = min(val[left(id)], val[right(id)]);
cbb184     }
cbb184 }
d41d8c
d41d8c // Get the minimum value in range [a, b].
9fed20 int get(int id, int l, int r, int a, int b)
f95b70 {
b62cfe     prop(id, l, r);
addc1f     if (a == l && b == r)
a0328b         return val[id];
2954e9     else
f95b70     {
ae007b         int mid = (l + r) / 2;
f890f2         if (b <= mid)
c55f80             return get(left(id), l, mid, a, b);
a54f0c         else if (a > mid)
26dd34             return get(right(id), mid + 1, r, a, b);
2954e9         else
5e3fad             return min(get(left(id), l, mid, a, mid), get(right(id), mid +
1, r, mid + 1, b));
cbb184     }
cbb184 }
d41d8c
d41d8c // Find index of rightmost element which is less than x. (works
because this is a seg of min)
0529b3 int bsearch(int id, int l, int r, int x)
f95b70 {
b62cfe     prop(id, l, r);
d41d8c
bd3398     if (l == r)
f7d2ed         return (val[id] < x) ? l : -1;
2954e9     else
f95b70     {
ae007b         int mid = (l + r) / 2;
221ad0         prop(left(id), l, mid);
384ec5         prop(right(id), mid + 1, r);
f01b35         if (val[right(id)] < x)
018a94             return bsearch(right(id), mid + 1, r, x);
2954e9         else
bad725             return bsearch(left(id), l, mid, x);
cbb184     }
cbb184 }
d41d8c
1037bf #define MAXN 312345
d41d8c
a58cd5 int a[MAXN];
c4b25f int b[MAXN];
d41d8c

```

```

13a4b1  int main(void)
f95b70  {
b067b3      int n, m, q, tp, x, y;
d69917      scanf("%d %d", &n, &m);
5359f3      for (int i = 1; i <= n; i++)
f95b70      {
9376f3          scanf("%d", &a[i]);
49e934          update(1, 1, 1000000, 1, a[i], -1);
cbb184      }
d41d8c
8eae24      for (int i = 1; i <= m; i++)
f95b70      {
264aeb          scanf("%d", &b[i]);
472fcc          update(1, 1, 1000000, 1, b[i], 1);
cbb184      }
d41d8c
4aaeab      scanf("%d", &q);
a953ae      while (q--)
f95b70      {
960099          scanf("%d %d %d", &tp, &x, &y);
abc772          if (tp == 1)
f95b70          {
996a9b              update(1, 1, 1000000, 1, a[x], 1);
e603e6              a[x] = y;
28cfa0              update(1, 1, 1000000, 1, a[x], -1);
cbb184          }
2954e9          else
f95b70          {
8dbabe              update(1, 1, 1000000, 1, b[x], -1);
0464a9              b[x] = y;
bc18aa              update(1, 1, 1000000, 1, b[x], 1);
cbb184          }
d41d8c
584906          int tmp = bsearch(1, 1, 1000000, 0);
d41d8c
d41d8c          // Test of get and bsearch. Make sure all to the right are non-
5a5bec          if (tmp != 1000000)
5df0f6              assert(get(1, 1, 1000000, tmp == -1 ? 1 : (tmp + 1), 1000000)
>= 0);
c3e568          if (tmp != -1)
1d95f2              assert(get(1, 1, 1000000, tmp, tmp) < 0);
d41d8c
b03a7a          printf("%d\n", tmp);
cbb184      }
cbb184  }
Full file hash: 90a905

```

## 2.8 Key Treap

```
5d1131 #include "../contest/header.hpp"
d41d8c
d41d8c /*
1977a5     Treap:
3ca64f     This treap implements something like a c++ set with additional
           operations: find the k-th element and count elements less than a given
           value.
d41d8c
4c88cf     Time: O(log N) per operation.
d41d8c
3db72f     Author: Arthur Pratti Dadalto
c4c9bd */
d41d8c
41c55a namespace treap
f95b70 {
e4accb struct node
f95b70 {
97f03f     int val; // node key.
ee1179     int p;   // node heap priority.
59afd1     int num; // node subtree size.
af32d9     node *left, *right;
d41d8c
71091e     node(int _val) : val(_val), p(rand()), num(1), left(NULL), right(
NULL) {}
2145c1 };
d41d8c
48f3b4 int get_num(node *root)
f95b70 {
424a36     return (root == NULL) ? 0 : root->num;
cbb184 }
d41d8c
68f1eb void update_num(node *root)
f95b70 {
47a6f1     root->num = get_num(root->left) + get_num(root->right) + 1;
cbb184 }
d41d8c
afdba0 node *rotate_left(node *root)
f95b70 {
d25f1b     node *a = root;
a95379     node *b = root->right;
d41d8c
b51426     a->right = b->left;
e7e30a     b->left = a;
a5e0c3     update_num(a);
2b11db     update_num(b);
73f89f     return b;
cbb184 }
d41d8c
f17a34 node *rotate_right(node *root)
f95b70 {
```



```

d25f1b     node *a = root;
eb0328     node *b = root->left;
d41d8c
a09684     a->left = b->right;
7352c4     b->right = a;
a5e0c3     update_num(a);
2b11db     update_num(b);
73f89f     return b;
cbb184 }
d41d8c
d41d8c     // Insert new node with key x in treap rooted at root if not already
there.
960bce     node *insert(node *root, int x)
f95b70     {
0edbc9         if (root == NULL)
13478f             return new node(x);
6b2a0b         if (x > root->val)
34c9df             root->right = insert(root->right, x);
ba0dc8         else if (x < root->val)
12f5b5             root->left = insert(root->left, x);
d41d8c
622638         update_num(root);
d41d8c
4f4bcf         if (root->right && root->right->p > root->p)
04107a             root = rotate_left(root);
c93ea7         if (root->left && root->left->p > root->p)
3f3108             root = rotate_right(root);
e2fc54         return root;
cbb184     }
d41d8c
d41d8c     // Remove node with key x in treap rooted at root if present.
d0ba77     node *remove(node *root, int x)
f95b70     {
0edbc9         if (root == NULL)
ea9b0a             return NULL;
6b2a0b         if (x > root->val)
fed39a             root->right = remove(root->right, x);
ba0dc8         else if (x < root->val)
6cf773             root->left = remove(root->left, x);
fb8e77         else if (root->left == NULL)
4de2d2             root = root->right;
a15580         else if (root->right == NULL)
2d4ff4             root = root->left;
386129         else if (root->left->p > root->right->p)
f95b70             {
3f3108                 root = rotate_right(root);
fed39a                 root->right = remove(root->right, x);
cbb184             }
2954e9         else
f95b70             {

```

```

04107a     root = rotate_left(root);
6cf773     root->left = remove(root->left, x);
cbb184     }
e6a2b0     if (root)
622638         update_num(root);
e2fc54     return root;
cbb184     }
d41d8c
d41d8c     // Return the k-th smallest element in tree rooted at root.
3576ec     int kth(node *root, int k)
f95b70     {
f9e30a         if (get_num(root->left) >= k)
7473ee             return kth(root->left, k);
f3e79f         else if (get_num(root->left) + 1 == k)
ae0ddc             return root->val;
2954e9         else
235aa0             return kth(root->right, k - get_num(root->left) - 1);
cbb184     }
d41d8c
d41d8c     // Return the number of elements smaller than x in tree rooted at
    root.
194e12     int count(node *root, int x)
f95b70     {
0edbc9         if (root == NULL)
bb30ba             return 0;
83010a         if (x < root->val)
da7c4c             return count(root->left, x);
08e5c0         else if (x == root->val)
140f45             return get_num(root->left);
2954e9         else
b73a02             return get_num(root->left) + 1 + count(root->right, x);
cbb184     }
cbb184     } // namespace treap
Full file hash: 85f362

```

## 2.9 Sequential Treap

```

5d1131     #include "../contest/header.hpp"
d41d8c
d41d8c     /*
1977a5         Treap:
5c39c7         A short self-balancing tree. It acts as a sequential container
    with log-time splits/joins, and
df7261         is easy to augment with additional data.
d41d8c
4c88cf         Time:  $O(\log N)$  per operation.
d41d8c
ca2095         Constraints:
c1b810         Acts as a vector of size  $N$ , with positions in range  $[0, N-1]$ .
d41d8c

```

```

1d1558 Source: https://github.com/kth-competitive-programming/kactl/blob/
      master/content/data-structures/Treap.h
d41d8c
b95cae Usage:
24eb84 To insert elements, create one node treaps. (e.g. treap::ins(root
      , new treap::node(x), i))
acfc60 To augment with extra data you should mostly add stuff to the
      recalc function. (e.g. to make it work like a seg tree)
03bb33 See applications for more usage examples.
c4c9bd */
d41d8c
41c55a namespace treap
f95b70 {
e4accb struct node
f95b70 {
8f5901     node *l = 0, *r = 0;
97f03f     int val;    // Any value associated with node.
ee1179     int p;    // Node heap priority.
c6aff2     int c = 1; // Node subtree size.
674490     node(int val) : val(val), p(rand()) {}
86d631     void recalc();
2145c1 };
d41d8c
853943 int cnt(node *n) { return n ? n->c : 0; }
9af082 void node::recalc() { c = cnt(l) + cnt(r) + 1; }
d41d8c
d41d8c // Apply function f on each tree node in order.
044d82 template <class F>
d5442c void each(node *n, F f)
f95b70 {
f63660     if (n)
f95b70     {
cbc351         each(n->l, f);
ed31a5         f(n->val);
f5ab50         each(n->r, f);
cbb184     }
cbb184 }
d41d8c
d41d8c // Split treap rooted at n in two treaps containing positions [0, k)
      and [k, ...)
de9c69 pair<node *, node *> split(node *n, int k)
f95b70 {
a020ba     if (!n)
e70a07         return {NULL, NULL};
9416bd     if (cnt(n->l) >= k) // "n->val >= k" for lower_bound(k)
f95b70     {
215a80         auto pa = split(n->l, k);
f3cfa7         n->l = pa.second;
2f09c0         n->recalc();
c05937         return {pa.first, n};

```

```

cbb184     }
2954e9     else
f95b70     {
7c23f0         auto pa = split(n->r, k - cnt(n->l) - 1); // and just "k"
d37e77         n->r = pa.first;
2f09c0         n->recalc();
7af31a         return {n, pa.second};
cbb184     }
cbb184 }
d41d8c
d41d8c // Merge treaps l and r keeping order (l first).
7f5419 node *merge(node *l, node *r)
f95b70 {
0c92a8     if (!l)
4c1f3c         return r;
6bf95d     if (!r)
792fd4         return l;
a0ade2     if (l->p > r->p)
f95b70     {
ed7b68         l->r = merge(l->r, r);
bf6a1f         l->recalc();
792fd4         return l;
cbb184     }
2954e9     else
f95b70     {
654f23         r->l = merge(l, r->l);
cda92d         r->recalc();
4c1f3c         return r;
cbb184     }
cbb184 }
d41d8c
d41d8c // Insert treap rooted at n into position pos of treap rooted at t.
3fc637 node *ins(node *t, node *n, int pos)
f95b70 {
ca9a9f     auto pa = split(t, pos);
cc8215     return merge(merge(pa.first, n), pa.second);
cbb184 }
d41d8c
d41d8c // Remove node at position pos from treap rooted at t.
1e0b32 node *rem(node *t, int pos)
f95b70 {
abdf75     node *a, *b, *c;
cf9546     tie(a, b) = split(t, pos);
0052e9     tie(b, c) = split(b, 1);
d41d8c
625cf2     delete b;
a300e4     return merge(a, c);
cbb184 }
d41d8c
d41d8c // Example application: do a query in range [l, r].

```

```
0475c8 node *query(node *t, int l, int r)
f95b70 {
abdf75     node *a, *b, *c;
a8341d     tie(a, b) = split(t, l);
89f194     tie(b, c) = split(b, r - l + 1);
d41d8c
d41d8c     // printf("%lld\n", b->tab);
d41d8c
53aa0f     return merge(merge(a, b), c);
cbb184 }
d41d8c
d41d8c // Example application: move the range [l, r) to index k.
b51124 void move(node *&t, int l, int r, int k)
f95b70 {
abdf75     node *a, *b, *c;
a8341d     tie(a, b) = split(t, l);
e81a2b     tie(b, c) = split(b, r - l);
1527bb     if (k <= l)
eeb6c2         t = merge(ins(a, b, k), c);
2954e9     else
646d6a         t = merge(a, ins(c, b, k - r));
cbb184 }
cbb184 } // namespace treap
Full file hash: 02c35c
```

## 3 Geometry

### 3.1 2D

```

5d1131 #include "../..//contest/header.hpp"
d41d8c
d41d8c // 2D geometry operations. This file should not have algorithms.
d41d8c // Author: some of it by Arthur Pratti Dadalto.
d41d8c // Source: some of it from https://github.com/kth-competitive-
    programming/kactl/blob/master/content/geometry/.
d41d8c // Usage: avoid int unless necessary.
d41d8c
d41d8c // When increasing EPS, keep in mind that  $\sqrt{1e9^2 + 1} = 1e9 + 5e$ 
     $-10$ .
22c921 const double EPS = 1e-12;
d41d8c
d41d8c // Point struct implementation. Some methods are useful only when
    using this to represent vectors.
4fce64 template <class T>
4befb0 struct point
f95b70 {
5dcf91     typedef point<T> P;
645c5d     T x, y;
d41d8c
571f13     explicit point(T x = 0, T y = 0) : x(x), y(y) {}
0d0d56     bool operator<(P p) const { return tie(x, y) < tie(p.x, p.y); }
ec7475     bool operator==(P p) const { return tie(x, y) == tie(p.x, p.y); }
2798c7     P operator+(P p) const { return P(x + p.x, y + p.y); }
40d57e     P operator-(P p) const { return P(x - p.x, y - p.y); }
e03fa4     P operator*(T d) const { return P(x * d, y * d); }
0b99e8     P operator/(T d) const { return P(x / d, y / d); }
57bee4     T dot(P p) const { return x * p.x + y * p.y; }
460881     T cross(P p) const { return x * p.y - y * p.x; }
b3fab9     T cross(P a, P b) const { return (a - *this).cross(b - *this); } //
    product sign: right hand rule from a to b.
f681d2     T dist2() const { return x * x + y * y; } // Distance
    squared to origin.
18b7a8     double dist() const { return sqrt((double)dist2()); } //
    Vector norm (distance to origin).
9073ff     double angle() const { return atan2(y, x); } // angle to
    x-axis in interval [-pi, pi]
6f5d42     point<double> unit() const { return *this / dist(); } // makes
    dist()==1 (unit vector).
200c8f     P perp() const { return P(-y, x); } // rotates +90
    degrees around origin.
567be8     point<double> normal() const { return perp().unit(); } //
    perpendicular unit vector.
82fcdd     point<double> rotate(double a) const // returns point
    rotated 'a' radians ccw around the origin.
f95b70     {

```

```

80d6a0     return P(x * cos(a) - y * sin(a), x * sin(a) + y * cos(a));
cbb184     }
8ad6e5     double angle(P p) const { return p.rotate(-angle()).angle(); } //
    Angle between the vectors in interval [-pi, pi]. Positive if p is ccw from
    this.
2145c1     };
d41d8c
d41d8c     // Solves the linear system {a * x + b * y = e
d41d8c     //                               {c * x + d * y = f
d41d8c     // Returns {1, {x, y}} if solution is unique, {0, {0,0}} if no
    solution and {-1, {0,0}} if infinite solutions.
d41d8c     // If using integer function type, this will give wrong answer if
    answer is not integer.
d41d8c     // TODO: test me with integer and non-integer.
4fce64     template <class T>
562c39     pair<int, point<T>> linear_solve2(T a, T b, T c, T d, T e, T f)
f95b70     {
468cb9         point<T> retv;
256940         T det = a * d - b * c;
d41d8c
57f40d         if (det == 0) // Maybe do EPS compare if using floating point.
f95b70         {
cdd981             if (b * f == d * e && a * f == c * e)
3d7337                 return {-1, point<T>()};
37dde3             return {0, point<T>()};
cbb184         }
d41d8c
d41d8c         // In case solution needs to be integer, use something like the
    line below.
d41d8c         // assert((e * d - f * b) % det == 0 && (a * f - c * e) % det == 0)
    ;
d41d8c
848480     return {1, point<T>((e * d - f * b) / det, (a * f - c * e) / det)};
cbb184     }
d41d8c
d41d8c     // Represents line segments defined by two points.
4fce64     template <class T>
4b2ec6     struct segment
f95b70     {
5dcf91         typedef point<T> P;
efb78f         P pi, pf; // Initial and final points.
d41d8c
a76c62         explicit segment(P a = P(), P b = P()) : pi(a), pf(b) {}
d41d8c
d41d8c         // Distance from this segment to a given point. TODO: test me.
325177         double dist(P p)
f95b70         {
58fd41             if (pi == pf)
adefd2                 return (p - pi).dist();
96a4f0             auto d = (pf - pi).dist2();

```

```

486c32     auto t = min(d, max(.0, (p - pi).dot(pf - pi)));
5dab06     return ((p - pi) * d - (pf - pi) * t).dist() / d;
cbb184 }
d41d8c
d41d8c     // Checks if given point belongs to segment. Use dist(p) <= EPS
instead when using point<double>.
0e3dba     bool on_segment(P p)
f95b70     {
50f719         return p.cross(pi, pf) == 0 && (pi - p).dot(pf - p) <= 0;
cbb184     }
d41d8c
d41d8c     // If a unique intersection point between the line segments exists
then it is returned.
d41d8c     // If no intersection point exists an empty vector is returned.
d41d8c     // If infinitely many exist a vector with 2 elements is returned,
containing the endpoints of the common line segment.
d41d8c     // The wrong position will be returned if P is point<ll> and the
intersection point does not have integer coordinates.
d41d8c     // However, no problem in using it to check if intersects or not in
this case (size of vector will be correct).
d41d8c     // Products of three coordinates are used in intermediate steps
so watch out for overflow if using int or long long.
f3f800     vector<P> intersect(segment rhs)
f95b70     {
9b1730         auto oa = rhs.pi.cross(rhs.pf, pi), ob = rhs.pi.cross(rhs.pf, pf)
,
1d46ec         oc = pi.cross(pf, rhs.pi), od = pi.cross(pf, rhs.pf);
d41d8c
d41d8c         // Checks if intersection is single non-endpoint point.
288e4c         if (sign(oa) * sign(ob) < 0 && sign(oc) * sign(od) < 0)
655339             return {(pi * ob - pf * oa) / (ob - oa)};
d41d8c
4c122f         set<P> s;
0373dd         if (rhs.on_segment(pi))
f07e25             s.insert(pi);
6725fe         if (rhs.on_segment(pf))
3c93ab             s.insert(pf);
3ad8fc         if (on_segment(rhs.pi))
522b2f             s.insert(rhs.pi);
f425cd         if (on_segment(rhs.pf))
d1c5a5             s.insert(rhs.pf);
d2dd66         return vector<P>(s.begin(), s.end());
cbb184     }
2145c1 };
d41d8c
d41d8c     // Represents a line by its equation in the form  $a * x + b * y = c$ .
d41d8c     // Can be created from two points or directly from constants.
4fce64     template <class T>
3fe318     struct line
f95b70     {

```



```

5dcf91     typedef point<T> P;
52d831     T a, b, c; // line a * x + b * y = c
d41d8c
f4f0fd     explicit line(P p1, P p2) // TODO: test me.
f95b70     {
4c2f1e         assert(!(p1 == p2));
6a88e5         a = p2.y - p1.y;
82330e         b = p1.x - p2.x;
cfae8e         c = a * p1.x + b * p1.y;
d41d8c
d41d8c         // In case of int, it is useful to scale down by gcd (e.g to use
in a set).
d41d8c         // Might be useful to normalize here.
cbb184     }
d41d8c
510551     explicit line(T _a, T _b, T _c) : a(_a), b(_b), c(_c) {}
d41d8c
d41d8c         // Distance from this line to a given point. TODO: test me.
325177     double dist(P p)
f95b70     {
d37216         return abs(a * p.x + b * p.y - c) / sqrt((double)(a * a + b * b))
;
cbb184     }
d41d8c
d41d8c         // Intersects this line with another given line. See linear_solve2
for usage. TODO: test me.
4a5d8e     pair<int, P> intersect(line rhs)
f95b70     {
6c76dc         return linear_solve2(a, b, rhs.a, rhs.b, c, rhs.c);
cbb184     }
d41d8c
d41d8c         // Normalize line to c >= 0, a*a + b*b == 1. Only use with double.
050345     line normalize()
f95b70     {
22b5e2         double d = P(a, b).dist() * (c < 0 ? -1 : 1);
7c9abe         return line(a / d, b / d, c / d);
cbb184     }
2145c1 };
d41d8c
d41d8c         // Represents a circle by its center and radius. Mostly only works
with double.
4fce64     template <class T>
0b1113     struct circle
f95b70     {
5dcf91         typedef point<T> P;
1ab228         P center;
c3df30         T r;
d41d8c
d41d8c         // Intersects circle with a given line. This does not work with
integer types.

```

```

d41d8c    // If there is no intersection, returns 0 and retv is whatever.
d41d8c    // If intersection is a single point, returns 1 and retv is a pair
    of equal points.
d41d8c    // If intersection is two points, return 2 and retv is the two
    intersection points.
d41d8c    // Assume points are given in no particular order. If you really
    need it, should be leftmost first when looking from center of the circle.
ec2c6b    int intersect(line<T> l, pair<P, P> &retv)
f95b70    {
800175        l = l.normalize();
f543ca        l.c -= l.a * center.x + l.b * center.y; // Recenter so that we
    can consider circle center in origin.
18b956        P v(l.a, l.b);
cf8231        P p0 = v * l.c; // p0 is the point in the line closest to origin.
d41d8c
2d9566        if (p0.dist() > r + EPS) // No intersection.
bb30ba            return 0;
40b0e2        else if (p0.dist() > r - EPS) // dist in [r - EPS, r + EPS] ->
    single point intersection at p0.
f95b70            {
de0c90                retv = {p0, p0};
6a5530                return 1;
cbb184            }
d41d8c
85b09c        double d = sqrt(r * r - l.c * l.c); // d is distance from p0 to
    the intersection points.
c4bf3f        retv = {center + p0 + v.normal() * d, center + p0 - v.normal() *
    d};
18b932        return 2;
cbb184    }
d41d8c
d41d8c    // Intersects circle with another circle. This does not work with
    integer types.
d41d8c    // This assumes the circles do not have the same center. Check this
    case if needed, can have 0 or infinite intersection points.
d41d8c    // If there is no intersection, returns 0 and retv is whatever.
d41d8c    // If intersection is a single point, returns 1 and retv is a pair
    of equal points.
d41d8c    // If intersection is two points, return 2 and retv is the two
    intersection points.
d41d8c    // Assume points are given in no particular order. If you really
    need it, should be leftmost first when looking from center of the rhs
    circle.
f2bab0    int intersect(circle rhs, pair<P, P> &retv)
f95b70    {
db42cd        rhs.center = rhs.center - center;
2adf3a        int num = rhs.intersect(line<T>(2 * rhs.center.x, 2 * rhs.center.
    y, rhs.center.x * rhs.center.x + rhs.center.y * rhs.center.y + r * r - rhs
    .r * rhs.r), retv);
2a6a69        retv.first = retv.first + center;

```

```

e34010      retv.second = retv.second + center;
fcc01b      return num;
cbb184    }
d41d8c
d41d8c      // Returns a pair of the two points on the circle whose tangent
lines intersect p.
d41d8c      // If p lies within the circle NaN-points are returned. P is
intended to be Point<double>.
d41d8c      // The first point is the one to the right as seen from the point p
towards the circle.
163627    pair<P, P> tangents(P p)
f95b70    {
75ad6b      p = p - center;
28b73b      double k1 = r * r / p.dist2();
f84c08      double k2 = sqrt(k1 - k1 * k1);
a64b03      return {center + p * k1 + p.perp() * k2, center + p * k1 - p.perp
() * k2};
cbb184    }
d41d8c
d41d8c      // TODO: find pair of tangent lines passing two circles.
2145c1  };
d41d8c
d41d8c      // The circumcircle of a triangle is the circle intersecting all
three vertices.
d41d8c      // Returns the unique circle going through points A, B and C (given
in no particular order).
d41d8c      // This assumes that the triangle has non-zero area.
d41d8c      // TODO: test specifically.
11308f    circle<double> circumcircle(const point<double> &A, const point<
double> &B, const point<double> &C)
f95b70    {
b10dc9      circle<double> retv;
6d2418      point<double> a = C - B, b = C - A, c = B - A;
1d9440      retv.r = a.dist() * b.dist() * c.dist() / abs(c.cross(b)) / 2;
0d1695      retv.center = A + (b * c.dist2() - c * b.dist2()).perp() / b.cross(
c) / 2;
6272cf      return retv;
cbb184    }
Full file hash: af0eba

```

### 3.2 Graham Scan (convex hull)

```

d41d8c  /*
6ccc59    Solution for convex hull problem (minimum polygon covering a set of
points) based on ordering points by angle.
3248ac    * Finds the subset of points in the convex hull in O(Nlog(N)).
687c39    * This version works if you either want intermediary points in
segments or not (see comments delimited by //)
01b744    * This version works when all points are collinear
246d86    * This version works for repeated points if you add a label to

```

```

    struct, and use this label in overloaded +, - and =.
d41d8c
1d1558    Source: https://github.com/kth-competitive-programming/kactl/blob/
    master/content/contest/template.cpp
c4c9bd    */
d41d8c
2b74fa    #include<bits/stdc++.h>
d41d8c
ad1153    typedef long long ll;
d41d8c
ca417d    using namespace std;
d41d8c
67a100    template<typename T>
4befb0    struct point
f95b70    {
5dcf91        typedef point<T> P;
645c5d        T x, y;
d41d8c
571f13        explicit point(T x = 0, T y = 0) : x(x), y(y) {}
d41d8c        //Double version: bool operator<(P p) const { return fabs(x - p.x)
    < EPS ? y < p.y : x < p.x; }
0d0d56        bool operator<(P p) const { return tie(x, y) < tie(p.x, p.y); }
d41d8c        //Double version: bool operator==(P p) const { return fabs(x - p.x)
    < EPS && fabs(y - p.y) < EPS; }
ec7475        bool operator==(P p) const { return tie(x, y) == tie(p.x, p.y); }
2798c7        P operator+(P p) const { return P(x + p.x, y + p.y); }
f681d2        T dist2() const { return x*x + y*y; }
40d57e        P operator-(P p) const { return P(x - p.x, y - p.y); }
57bee4        T dot(P p) const { return x * p.x + y * p.y; }
460881        T cross(P p) const { return x * p.y - y * p.x; }
b3fab9        T cross(P a, P b) const { return (a - *this).cross(b - *this); }
5b4ebf        long double dist() const { return sqrt((long double)dist2()); }
2145c1    };
d41d8c
d41d8c    /*Compara primeiro por angulo em relacao a origem e depois por
    distancia para a origem*/
67a100    template<typename T>
d74eff    bool cmp(point<T> a, point<T> b){
a9b570        if(a.cross(b) != 0)
c33606            return a.cross(b) > 0;
ba7b3a        return a.dist2() < b.dist2();
cbb184    }
d41d8c
67a100    template<typename T>
3c7876    vector<point<T> > CH(vector<point<T> > points){
d41d8c        /*Encontra pivo (ponto extremos que com ctz faz parte do CH)*/
95b799        point<T> pivot = points[0];
e409fb        for(auto p : points)
e01c07            pivot = min(pivot, p);
d41d8c

```

```

d41d8c    /*Desloca conjunto para pivo ficar na origem e ordena pontos pelo
          angulo e distancia do pivo*/
9ac126    for(int i = 0; i < (int) points.size(); i++)
3010bd        points[i] = points[i] - pivot;
d41d8c
e2c4e0    sort(points.begin(), points.end(), cmp<ll>);
d41d8c
9ac126    for(int i = 0; i < (int) points.size(); i++)
eda5a9        points[i] = points[i] + pivot;
d41d8c
d41d8c    /*Ponto extra para fechar o poligono*/
36b3da    points.push_back(points[0]);
d41d8c
620533    vector<point<T> > ch;
d41d8c
b7f960    for(auto p : points){
d41d8c        /*Enquanto o proximo ponto gera uma curva para a direita, retira
          ultimo ponto atual*/
d41d8c        /*Segunda comparaçãõ serve para caso especial de pontos
          colineares quando se quer eliminar os intermediarios*/
d41d8c        //Trocar terceira comparacao pra <= para descartar pontos do meio
          de arestas no ch
d41d8c        //Double: trocar terceira comparaçãõ por < EPS (descarta pontos
          em arestas) ou < -EPS (mantem ponto em aresta
29fcb4        while(ch.size() > 1 && !(p == ch[ch.size() - 2]) && ch[ch.size()
          - 2].cross(ch[ch.size() - 1] , p) < 0)
9d9654            ch.pop_back();
d2ebaf            ch.push_back(p);
cbb184    }
d41d8c
d41d8c    /*Elimina ponto extra*/
9d9654    ch.pop_back();
d41d8c
66cc3c    return ch;
cbb184 }
d41d8c
e8d76f int main(){
1a88fd     int n;
f4c120     scanf("%d", &n);
76374e     vector<point<ll> > p(n);
d41d8c
d41d8c     /*Le poligono*/
83008c     for(int i = 0; i < n; i++)
3daa4c         scanf("%lld %lld", &p[i].x, &p[i].y);
d41d8c
d41d8c     /*Encontra CH*/
680587     vector<point<ll> > ch = CH(p);
d41d8c
d41d8c     /*Imprime resultado*/
3b846e     printf("%d\n", (int)ch.size());

```

```

c857e7     for(int i = 0; i < (int)ch.size(); i++)
fc3047         printf("%d% %d\n", ch[i].x, ch[i].y);
d41d8c
cbb184     }
Full file hash: b6794f

```

### 3.3 Min Enclosing Circle (randomized)

```

ad578e     #include "../2d/2d.cpp"
d41d8c
d41d8c     /*
744027         Minimum Enclosing Circle:
2d38cc         Given a list of points, returns a circle of minimum radius such
            that all given
2602de         points are within the circle.
0c4a3b         Runs in O(n) expected time (in practice 200 ms for 10^5 points).
d41d8c
ca2095         Constraints:
99b71a         Non-empty list of points.
d41d8c
3db72f         Author: Arthur Pratti Dadalto
c4c9bd     */
d41d8c
e89126     #define point point<double>
0f3aa0     #define circle circle<double>
d41d8c
41ee07     circle min_enclosing_circle(vector<point> p)
f95b70     {
b4da45         shuffle(p.begin(), p.end(), mt19937(time(0)));
2e09de         point o = p[0];
76160f         double r = 0, eps = 1 + 1e-8;
fe16ed         for (int i = 0; i < sz(p); i++)
197ee7             if ((o - p[i]).dist() > r * eps)
f95b70             {
ba37a5                 o = p[i], r = 0;
c791cd                 for (int j = 0; j < i; j++)
f5972f                     if ((o - p[j]).dist() > r * eps)
f95b70                     {
d2b545                         o = (p[i] + p[j]) / 2;
0657ce                         r = (o - p[i]).dist();
674051                         for (int k = 0; k < j; k++)
355d4d                             if ((o - p[k]).dist() > r * eps)
f95b70                             {
7fb807                                 o = circumcircle(p[i], p[j], p[k]).center;
0657ce                                 r = (o - p[i]).dist();
cbb184                             }
cbb184                         }
cbb184                     }
d41d8c             }
645c1d         return {o, r};

```

```
cbb184  }
d41d8c
Full file hash: 5d3836
```

### 3.4 Min Enclosing Circle (ternary search)

```
ad578e  #include "../2d/2d.cpp"
2729c3  #include "../misc/ternary_search/ternary_search_continuous.cpp"
d41d8c
d41d8c  /*
744027   Minimum Enclosing Circle:
2d38cc   Given a list of points, returns a circle of minimum radius such
        that all given
2602de   points are within the circle.
a1c921   Runs in  $O(n * \log^2((top - bot) / eps))$  (in practice 2.5s at best
        for  $10^5$  points).
d41d8c
ca2095   Constraints:
99b71a   Non-empty list of points.
d41d8c
b95cae   Usage:
ca9f24   The coordinates of the circle's center must be in the range [bot,
        top].
a09e29   eps specifies the precision of the result, but set it to a higher
        value
53006e   than necessary since the error in x affects the y value.
d41d8c
3db72f   Author: Arthur Pratti Dadalto
c4c9bd  */
d41d8c
e89126  #define point point<double>
0f3aa0  #define circle circle<double>
d41d8c
e1710f  circle min_enclosing_circle(const vector<point> &p, double bot = -1e9
        , double top = 1e9, double eps = 1e-9)
f95b70  {
1841a9   circle retv;
d41d8c
0a37af   auto f1 = [&](double x) {
d9991d       auto f2 = [&](double y)
f95b70       {
996834           double r = 0;
fe16ed           for (int i = 0; i < sz(p); i++)
62adf4               r = max(r, (p[i].x - x)*(p[i].x - x) + (p[i].y - y)*(p[i].y -
        y));
4c1f3c           return r;
2145c1       };
410f57   retv.center.y = ternary_search(f2, bot, top, eps);
50ac50   return f2(retv.center.y);
2145c1   };
```

```
d41d8c
596ad7    retv.center.x = ternary_search(f1, bot, top, eps);
3b2a60    retv.r = sqrt(f1(retv.center.x));
d41d8c
6272cf    return retv;
cbb184    }
Full file hash: 2acede
```



## 4 Graph

### 4.1 Biconnected Components

```

5d1131  #include "../contest/header.hpp"
d41d8c
d41d8c  /*
399f8b    Finding bridges, articulation points and biconnected components in
        0(V + E):
5d1e77    A bridge is an edge whose removal splits the graph in two
        connected components.
8dd98b    An articulation point is a vertex whose removal splits the graph
        in two connected components.
d41d8c
8254c5    A biconnected component (or 2VCC) is a maximal subgraph where the
        removal of any vertex doesn't
44c916    make the subgraph disconnected. In other words, it is a
        maximal 2-vertex-connected (2VC) subgraph.
d41d8c
deb2a1    A 2-connected graph is a 2VC one, except that a---b is
        considered 2VC but not 2-connected.
d41d8c
3a585b    Useful theorems:
d41d8c
45e89c    A 2-edge connected (2EC) graph is a graph without bridges.
        Any 2-connected graph is also 2EC.
d41d8c
20dbb2    Let G be a graph on at least 2 vertices. The following
        propositions are equivalent:
81e567    * (i) G is 2-connected;
959a93    * (ii) any two vertices are in a cycle; (a cycle can't
        repeat vertices)
3f2e8a    * (iii) any two edges are in a cycle and  $\hat{\tau}(G) \leq 2$ ;
57a10e    * (iv) for any three vertices x,y et z, there is a (x,z)-
        path containing y.
561b74    Let G be a graph on at least 3 vertices. The following
        propositions are equivalent:
346e5e    * (i) G is 2-edge-connected;
4883af    * (ii) any edge is in a cycle;
368e3a    * (iii) any two edges are in a tour and  $\hat{\tau} \leq 1$ ;
d6bb5f    * (iv) any two vertices are in a tour (a tour can repeat
        vertices)
d41d8c
8c2af3    If G is 2-connected and not bipartite, all vertices belong to
        some odd cycle. And any two vertices are in a odd cycle (not really proven
        ).
d41d8c
bcc5c9    If G is 2-edge-connected (proof by AC):
33bf43    For any two vertices x, y and one edge e, there is a (x,
        y)-walk containing e without repeating edges.

```

```

d41d8c
ab50d8      A graph admits a strongly connected orientation if and only
            if it is 2EC.
d3ea79      A strong orientation of a given bridgeless undirected graph
            may be found in linear time by performing
0ec987      a depth first search of the graph, orienting all edges in the
            depth first search tree away from the
d494fc      tree root, and orienting all the remaining edges (which must
            necessarily connect an ancestor and a
030955      descendant in the depth first search tree) from the
            descendant to the ancestor.
d41d8c
ca2095      Constraints:
b9aa54      ***undirected*** graph.
80b2d0      Vertices are labeled from 0 to n (inclusive).
568d46      Graph is connected (but for unconnected just replace single dfs
            call with a loop).
d41d8c
b95cae      Usage:
4436dc      Create the struct setting the starting vertex (a), the maximum
            vertex label (n),
e1993f      the graph adjacency list (graph) and a callback f to apply on
            the biconnected components.
f8f25e      Afterwards, art[i] == true if i is an articulation point.
e9a79e      If the pair {a, i} is on the bridges list, then the edge {a,
            graph[a][i]} is a bridge.
ccfd29      The callback must receive a vector of edges {a, b} that are
            in the same biconnected component.
a32c3f      Remember that for a single vertex, the biconnected callback will
            not be called.
d41d8c
e152b4      Sample Usage:
0ec6ee      auto rdm = apb(1, n, graph, [&](vector<pii> v){
f4ecd5      set<int> s;
9ad08e      for (int i = 0; i < sz(v); i++)
f95b70      {
f19ef4      s.insert(v[i].first);
0858fa      s.insert(v[i].second);
cbb184      }
d41d8c
0fe299      ans = max(ans, sz(s));
c0c97e      });
c4c9bd      */
d41d8c
f117a6      struct apb
f95b70      {
9cf2b9      vector<int> *graph;
9cf143      vector<bool> art;
c9001f      vector<int> num /* dfs order of vertices starting at 1 */, low;
c83796      vector<pii> bridges;

```

```

91936b vector<pii> st;
53e65f int id;
d41d8c
044d82 template<class F>
09caad apb(int a, int n, vector<int> graph[], const F &f) : graph(graph),
      art(n + 1, false), num(n + 1), low(n + 1)
f95b70 {
0f6720     id = 1;
ccac4e     dfs(a, a, f);
cbb184 }
d41d8c
044d82 template<class F>
dc584b void dfs(int a, int p, const F &f)
f95b70 {
7be506     low[a] = num[a] = id++;
34863b     int comp = 0;
d41d8c
1429ef     for (int i = 0; i < sz(graph[a]); i++)
f95b70     {
b7a810         if (num[graph[a][i]] == 0)
f95b70         {
d40410             int si = sz(st);
f309f5             comp++;
8ece2e             st.push_back({a, graph[a][i]}); // Tree edge.
d41d8c
fc5941             dfs(graph[a][i], a, f);
085d64             low[a] = min(low[a], low[graph[a][i]]);
d41d8c
bb63a0             if (low[graph[a][i]] >= num[a])
f95b70             {
558f81                 if (a != 1)
016392                 art[a] = true;
d41d8c
b91456                 f(vector<pii>(st.begin() + si, st.end()));
901921                 st.resize(si);
cbb184             }
d41d8c
0e9ddb             if (low[graph[a][i]] > num[a])
b3cacb                 bridges.push_back({a, i});
cbb184         }
624580     else if (graph[a][i] != p && num[graph[a][i]] < num[a]) // Back
edge.
f95b70     {
066898         low[a] = min(low[a], num[graph[a][i]]);
8ece2e         st.push_back({a, graph[a][i]});
cbb184     }
cbb184 }
d41d8c
85e3a2 if (a == p && comp > 1)
016392     art[a] = true;

```

```
cbb184  }
2145c1  };
d41d8c
Full file hash: 5cb0b8
```

## 4.2 Bipartite Matching (Hopcroft Karp)

```
2b74fa  #include <bits/stdc++.h>
ca417d  using namespace std;
d41d8c
d41d8c  /*
ec23c9      Hopcroft-Karp:
eaeddf          Bipartite Matching O(sqrt(V)E)
d41d8c
ca2095      Constraints:
998cc9          Vertices are labeled from 1 to l + r (inclusive).
682ff0          DO NOT use vertex 0.
968b86          Vertices 1 to l belong to left partition.
a6a4c4          Vertices l + 1 to l + r belong to right partition.
d41d8c
b95cae      Usage:
d86132          Set MAXV if necessary.
70636b          Call init passing l and r.
0f3b71          Add edges to the graph from left side to right side.
5263f1          Call hopcroft to get the matching size.
a0da8e          Then, each vertex v has its pair indicated in p[v] (or 0 for not
paired).
c4c9bd  */
d41d8c
dde07b  namespace hopcroft
f95b70  {
998014  const int inf = 0x3f3f3f3f;
ed5ed2  const int MAXV = 112345;
d41d8c
3098d4  vector<vector<int>> graph;
0a3d29  int d[MAXV], q[MAXV], p[MAXV], l, r;
d41d8c
4025e1  void init(int _l, int _r)
f95b70  {
0ebd66      l = _l, r = _r;
2213c3      graph = vector<vector<int>>(l + r + 1);
cbb184  }
d41d8c
6a1cf9  bool bfs()
f95b70  {
18753f      int qb = 0, qe = 0;
4f2bde      memset(d, 0x3f, sizeof(int) * (l + 1));
a89ba9      for (int i = 1; i <= l; i++)
8b3877          if (p[i] == 0)
248d2f              d[i] = 0, q[qe++] = i;
```

```

d41d8c
2caa87     while (qb < qe)
f95b70     {
e8e8a0         int a = q[qb++];
0087d7         if (a == 0)
8a6c14             return true;
c4fff3         for (int i = 0; i < graph[a].size(); i++)
68367c             if (d[p[graph[a][i]]] == inf)
a8cd28                 d[q[qe++]] = p[graph[a][i]] = d[a] + 1;
cbb184     }
d41d8c
d1fe4d     return false;
cbb184 }
d41d8c
0752c9 bool dfs(int a)
f95b70 {
0087d7     if (a == 0)
8a6c14         return true;
c4fff3     for (int i = 0; i < graph[a].size(); i++)
7d85df         if (d[a] + 1 == d[p[graph[a][i]]])
a2f815             if (dfs(p[graph[a][i]]))
f95b70                 {
460f0a                     p[a] = graph[a][i];
51e040                     p[graph[a][i]] = a;
8a6c14                     return true;
cbb184                 }
d41d8c
343737     d[a] = inf;
d1fe4d     return false;
cbb184 }
d41d8c
68fd9d int hopcroft()
f95b70 {
9e3790     memset(p, 0, sizeof(int) * (l + r + 1));
fc833c     int matching = 0;
d594a7     while (bfs())
f95b70     {
a89ba9         for (int i = 1; i <= l; i++)
8b3877             if (p[i] == 0)
57e7a2                 if (dfs(i))
730cbb                     matching++;
cbb184     }
d41d8c
2afcbe     return matching;
cbb184 }
cbb184 } // namespace hopcroft
Full file hash: 976bec

```

## 4.3 Bridges/Articulation Points

```

5d1131 #include "../contest/header.hpp"
d41d8c
d41d8c /*
62784d     Finding bridges and articulation points in  $O(V + E)$ :
5d1e77     A bridge is an edge whose removal splits the graph in two
        connected components.
8dd98b     An articulation point is a vertex whose removal splits the graph
        in two connected components.
8b8ace     This can also be adapted to generate the biconnected components
        of a graph, since the
14a784     articulation points split components.
d41d8c
d41d8c
ca2095     Constraints:
b9aa54     ***undirected*** graph.
80b2d0     Vertices are labeled from 0 to n (inclusive).
1e6120     Graph is connected (otherwise it doesn't make sense).
d41d8c
b95cae     Usage:
668bef     Create the struct setting the starting vertex (a), the maximum
        vertex label (n)
8acebe     and the graph adjacency list (graph).
6ffc91     Afterwards, art[i] == true if i is an articulation point.
e9a79e     If the pair {a, i} is on the bridges list, then the edge {a,
        graph[a][i]} is a bridge.
c4c9bd */
d41d8c
f117a6 struct apb
f95b70 {
9cf2b9 vector<int> *graph;
9cf143 vector<bool> art;
c9001f vector<int> num /* dfs order of vertices starting at 1 */, low;
c83796 vector<pii> bridges;
53e65f int id;
d41d8c
4dc736 apb(int a, int n, vector<int> graph[]) : graph(graph), art(n + 1,
        false), num(n + 1), low(n + 1)
f95b70 {
0f6720     id = 1;
bb407e     dfs(a, a);
cbb184 }
d41d8c
69c421 void dfs(int a, int p)
f95b70 {
7be506     low[a] = num[a] = id++;
34863b     int comp = 0;
d41d8c
c4fff3     for (int i = 0; i < graph[a].size(); i++)
f95b70     {
b7a810         if (num[graph[a][i]] == 0)

```

```

f95b70      {
f309f5      comp++;
783129      dfs(graph[a][i], a);
085d64      low[a] = min(low[a], low[graph[a][i]]);
d41d8c
b28b5f      if (a != 1 && low[graph[a][i]] >= num[a])
016392      art[a] = true;
d41d8c
0e9ddb      if (low[graph[a][i]] > num[a])
b3cacb      bridges.push_back({a, i});
cbb184      }
2cae7a      else if (graph[a][i] != p && num[graph[a][i]] < low[a])
ed0d8a      low[a] = num[graph[a][i]];
cbb184      }
d41d8c
85e3a2      if (a == p && comp > 1)
016392      art[a] = true;
cbb184      }
2145c1      };
d41d8c
Full file hash: 780b6d

```

#### 4.4 Max Flow (Dinic)

```

2b74fa      #include <bits/stdc++.h>
ca417d      using namespace std;
d41d8c
d41d8c      /*
908d2f      Dinic:
67cbe4      Max-flow  $O(V^2E)$ 
eaeddf      Bipartite Matching  $O(\sqrt{V}E)$ 
d41d8c
ca2095      Constraints:
80b2d0      Vertices are labeled from 0 to n (inclusive).
8f4ce8      Edge capacities must fit int (flow returned is long long).
d41d8c
b95cae      Usage:
d86132      Set MAXV if necessary.
148d9c      Call init passing n, the source and the sink.
2d6398      Add edges to the graph by calling put_edge(_undirected).
bb3825      Call max_flow to get the total flow. Then, individual edge flows
can be retrieved in the graph.
22c3c2      Note that flow will be negative in return edges.
c4c9bd      */
d41d8c
82657b      namespace dinic
f95b70      {
729806      struct edge
f95b70      {
bf6256      int dest, cap, re, flow;

```

```

2145c1 };
d41d8c
998014 const int inf = 0x3f3f3f3f;
8550b5 const int MAXV = 312345;
d41d8c
8a367b int n, s, t, d[MAXV], q[MAXV], next[MAXV];
d8f9f2 vector<vector<edge>> graph;
d41d8c
bc6f23 void init(int _n, int _s, int _t)
f95b70 {
c992e9     n = _n, s = _s, t = _t;
b72d19     graph = vector<vector<edge>>(n + 1);
cbb184 }
d41d8c
7c85eb void put_edge(int u, int v, int cap)
f95b70 {
506964     graph[u].push_back({v, cap, (int)graph[v].size(), 0});
68ec95     graph[v].push_back({u, 0, (int)graph[u].size() - 1, 0});
cbb184 }
d41d8c
d6a592 void put_edge_undirected(int u, int v, int cap)
f95b70 {
506964     graph[u].push_back({v, cap, (int)graph[v].size(), 0});
fce495     graph[v].push_back({u, cap, (int)graph[u].size() - 1, 0});
cbb184 }
d41d8c
6a1cf9 bool bfs()
f95b70 {
18753f     int qb = 0, qe = 0;
3c6658     q[qe++] = s;
98fde3     memset(d, 0x3f, sizeof(int) * (n + 1));
d66185     d[s] = 0;
2caa87     while (qb < qe)
f95b70     {
e8e8a0         int a = q[qb++];
c9a55a         if (a == t)
8a6c14             return true;
3352c6         for (int i = 0; i < (int)graph[a].size(); i++)
f95b70         {
10e42b             edge &e = graph[a][i];
d948dd             if (e.cap - e.flow > 0 && d[e.dest] == inf)
f4063b                 d[q[qe++] = e.dest] = d[a] + 1;
cbb184         }
cbb184     }
d41d8c
d1fe4d     return false;
cbb184 }
d41d8c
1a19d4 int dfs(int a, int flow)
f95b70 {

```



```

c9a55a     if (a == t)
99d2e8         return flow;
10647a     for (int &i = next[a]; i < (int)graph[a].size(); i++)
f95b70     {
10e42b         edge &e = graph[a][i];
c6fb85         if (d[a] + 1 == d[e.dest] && e.cap - e.flow > 0)
f95b70         {
5f308a             int x = dfs(e.dest, min(flow, e.cap - e.flow));
5f75db             if (x == 0)
5e2bd7                 continue;
7f9751             e.flow += x;
4c55a5             graph[e.dest][e.re].flow -= x;
ea5659             return x;
cbb184         }
cbb184     }
d41d8c
343737     d[a] = inf;
bb30ba     return 0;
cbb184 }
d41d8c
afa2f7 long long max_flow()
f95b70 {
f013d3     long long total_flow = 0;
d594a7     while (bfs())
f95b70     {
ba90c2         memset(next, 0, sizeof(int) * (n + 1));
60616b         while (int path_flow = dfs(s, inf))
a0d8d9             total_flow += path_flow;
cbb184     }
d41d8c
793f63     return total_flow;
cbb184 }
cbb184 } // namespace dinic
Full file hash: 73c17c

```

## 4.5 Max Flow (Dinic w/ Scaling)

```

2b74fa #include <bits/stdc++.h>
ca417d using namespace std;
d41d8c
d41d8c /*
3678e9     Dinic with Scaling:
476fd1         Max-flow  $O(VE * \log(\text{MAX\_CAP}))$ , but usually slower than regular
        Dinic.
d41d8c
ca2095     Constraints:
80b2d0         Vertices are labeled from 0 to n (inclusive).
8f4ce8         Edge capacities must fit int (flow returned is long long).
d41d8c
b95cae     Usage:

```

```

d86132      Set MAXV if necessary.
148d9c      Call init passing n, the source and the sink.
2d6398      Add edges to the graph by calling put_edge(_undirected).
bb3825      Call max_flow to get the total flow. Then, individual edge flows
            can be retrieved in the graph.
22c3c2      Note that flow will be negative in return edges.
c4c9bd  */
d41d8c
82657b  namespace dinic
f95b70  {
729806  struct edge
f95b70  {
bf6256      int dest, cap, re, flow;
2145c1  };
d41d8c
998014  const int inf = 0x3f3f3f3f;
8550b5  const int MAXV = 312345;
d41d8c
19c361  int n, s, t, lim, d[MAXV], q[MAXV], next[MAXV];
d8f9f2  vector<vector<edge>> graph;
d41d8c
bc6f23  void init(int _n, int _s, int _t)
f95b70  {
c992e9      n = _n, s = _s, t = _t;
b72d19      graph = vector<vector<edge>>(n + 1);
cbb184  }
d41d8c
7c85eb  void put_edge(int u, int v, int cap)
f95b70  {
506964      graph[u].push_back({v, cap, (int)graph[v].size(), 0});
68ec95      graph[v].push_back({u, 0, (int)graph[u].size() - 1, 0});
cbb184  }
d41d8c
d6a592  void put_edge_undirected(int u, int v, int cap)
f95b70  {
506964      graph[u].push_back({v, cap, (int)graph[v].size(), 0});
fce495      graph[v].push_back({u, cap, (int)graph[u].size() - 1, 0});
cbb184  }
d41d8c
6a1cf9  bool bfs()
f95b70  {
18753f      int qb = 0, qe = 0;
3c6658      q[qe++] = s;
98fde3      memset(d, 0x3f, sizeof(int) * (n + 1));
d66185      d[s] = 0;
2caa87      while (qb < qe)
f95b70      {
e8e8a0          int a = q[qb++];
c9a55a          if (a == t)
8a6c14              return true;

```

```

3352c6     for (int i = 0; i < (int)graph[a].size(); i++)
f95b70     {
10e42b         edge &e = graph[a][i];
21aca9         if (e.cap - e.flow >= lim && d[e.dest] == inf)
f4063b             d[q[qe++]] = e.dest; d[a] = 1;
cbb184     }
cbb184     }
d41d8c
d1fe4d     return false;
cbb184 }
d41d8c
1a19d4 int dfs(int a, int flow)
f95b70 {
c9a55a     if (a == t)
99d2e8         return flow;
10647a     for (int &i = next[a]; i < (int)graph[a].size(); i++)
f95b70     {
10e42b         edge &e = graph[a][i];
cbf046         if (d[a] + 1 == d[e.dest] && e.cap - e.flow >= lim /* >= 1 ? */)
f95b70         {
5f308a             int x = dfs(e.dest, min(flow, e.cap - e.flow));
5f75db             if (x == 0)
5e2bd7                 continue;
7f9751             e.flow += x;
4c55a5             graph[e.dest][e.re].flow -= x;
ea5659             return x;
cbb184         }
cbb184     }
d41d8c
343737     d[a] = inf;
bb30ba     return 0;
cbb184 }
d41d8c
afa2f7 long long max_flow()
f95b70 {
f013d3     long long total_flow = 0;
aab413     for (lim = (1 << 30); lim >= 1; lim >>= 1)
d594a7         while (bfs())
f95b70         {
ba90c2             memset(next, 0, sizeof(int) * (n + 1));
60616b             while (int path_flow = dfs(s, inf))
a0d8d9                 total_flow += path_flow;
cbb184         }
d41d8c
793f63     return total_flow;
cbb184 }
cbb184 } // namespace dinic
Full file hash: c74595

```

## 4.6 Min Cost Max Flow

```

2b74fa #include <bits/stdc++.h>
ca417d using namespace std;
d41d8c
d41d8c /*
dfc480     Min-Cost Max-Flow:  $O(V^2E^2)$ 
078f0d     Finds the maximum flow of minimum cost.
d41d8c
ca2095     Constraints:
80b2d0     Vertices are labeled from 0 to n (inclusive).
247575     Edge cost and capacities must fit int (flow and cost returned are
long long).
75ef18     Edge Cost must be non-negative.
d41d8c
b95cae     Usage:
d86132     Set MAXV if necessary.
148d9c     Call init passing n, the source and the sink.
909583     Add edges to the graph by calling put_edge.
553d3e     Call mincost_maxflow to get the total flow and its cost (in this
order).
0e374d     Individual edge flows can be retrieved in the graph. Note that
flow will be negative in return edges.
c4c9bd */
d41d8c
ad1153 typedef long long ll;
d29b14 typedef pair<long long, long long> pll;
d41d8c
e3de19 namespace mcmf
f95b70 {
729806 struct edge
f95b70 {
60f183     int dest, cap, re, cost, flow;
2145c1 };
d41d8c
ed5ed2 const int MAXV = 112345;
6a4c6c const ll infll = 0x3f3f3f3f3f3f3f3fLL;
998014 const int inf = 0x3f3f3f3f;
d41d8c
128c92 int n, s, t, p[MAXV], e_used[MAXV];
97e5bb bool in_queue[MAXV];
c97378 ll d[MAXV];
d41d8c
d8f9f2 vector<vector<edge>> graph;
d41d8c
bc6f23 void init(int _n, int _s, int _t)
f95b70 {
c992e9     n = _n, s = _s, t = _t;
b72d19     graph = vector<vector<edge>>(n + 1);
cbb184 }
d41d8c
a4abfa void put_edge(int u, int v, int cap, int cost)

```

```

f95b70 {
bd3e8b     graph[u].push_back({v, cap, (int)graph[v].size(), cost, 0});
2b8b8d     graph[v].push_back({u, 0, (int)graph[u].size() - 1, -cost, 0});
cbb184 }
d41d8c
b34984 bool spfa()
f95b70 {
664c61     memset(in_queue, 0, sizeof(bool) * (n + 1));
9eef50     memset(d, 0x3f, sizeof(ll) * (n + 1));
26a528     queue<int> q;
d66185     d[s] = 0;
e2828b     p[s] = s;
08bec3     q.push(s);
ee6bdd     while (!q.empty())
f95b70     {
0930a5         int a = q.front();
833270         q.pop();
e7249b         in_queue[a] = false;
d41d8c
c4fff3         for (int i = 0; i < graph[a].size(); i++)
f95b70         {
10e42b             edge &e = graph[a][i];
6fa321             if (e.cap - e.flow > 0 && d[e.dest] > d[a] + e.cost)
f95b70             {
3bf598                 d[e.dest] = d[a] + e.cost;
6d6530                 p[e.dest] = a;
183d83                 e_used[e.dest] = i;
27788c                 if (!in_queue[e.dest])
b34293                     q.push(e.dest);
04f0f7                 in_queue[e.dest] = true;
cbb184             }
cbb184         }
cbb184     }
d41d8c
d1cd45     return d[t] < infll;
cbb184 }
d41d8c
99658d pll mincost_maxflow()
f95b70 {
f04b2a     pll retv = pll(0, 0);
d9383f     while (spfa())
f95b70     {
e98031         int x = inf;
c9b315         for (int i = t; p[i] != i; i = p[i])
d4a316             x = min(x, graph[p[i]][e_used[i]].cap - graph[p[i]][e_used[i]].
            flow);
c9b315         for (int i = t; p[i] != i; i = p[i])
dc731d             graph[p[i]][e_used[i]].flow += x, graph[i][graph[p[i]][e_used[i]
            ]].re].flow -= x;
d41d8c

```

```

75907a      retv.first += x;
1be465      retv.second += x * d[t];
cbb184    }
d41d8c
6272cf      return retv;
cbb184    }
cbb184    } // namespace mcmf
Full file hash: 5bd8df

```

## 4.7 Heavy-Light Decomposition

```

2b74fa      #include<bits/stdc++.h>
d41d8c
ca417d      using namespace std;
d41d8c
eed838      #define ll long long
efe13e      #define pb push_back
d41d8c
3a6c63      typedef vector<ll> vll;
990dd5      typedef vector<int> vi;
d41d8c
e06cc0      #define MAXN 100010
d41d8c
d41d8c      //Vetor que guarda a arvore
698e25      vector<vi> adj;
d41d8c
9e6e6d      int subsize[MAXN], parent[MAXN];
d41d8c      //Inciar chainHead com -1; e chainSize e chainNo com 0.
080553      int chainNo = 0, chainHead[MAXN], chainPos[MAXN], chainInd[MAXN],
          chainSize[MAXN];
42a605      void hld(int cur){
cb42fb          if(chainHead[chainNo] == -1)
6591fe              chainHead[chainNo] = cur;
d41d8c
3a4605          chainInd[cur] = chainNo;
220e91          chainPos[cur] = chainSize[chainNo];
6f00fb          chainSize[chainNo]++;
d41d8c
89108d          int ind = -1, mai = -1;
9d9afd          for(int i = 0; i < (int)adj[cur].size(); i++){
9ff2fa              if(adj[cur][i] != parent[cur] && subsize[adj[cur][i]] > mai){
31fcc6                  mai = subsize[adj[cur][i]];
b9b7e9                  ind = i;
cbb184              }
cbb184          }
d41d8c
27d206          if(ind >= 0)
f23581              hld(adj[cur][ind]);
d41d8c
e506c6          for(int i = 0; i < (int)adj[cur].size(); i++)

```

```

6f7286         if(adj[cur][i] != parent[cur] && i != ind){
959ef6             chainNo++;
270563             hld(adj[cur][i]);
cbb184         }
cbb184     }
d41d8c
d41d8c //usar LCA para garantir que v eh pai de u!!
f179f7 ll query_up(int u, int v){
c20c7b     int uchain = chainInd[u], vchain = chainInd[v];
bdd5ea     ll ans = 0LL;
d41d8c
31e3cd     while(1){
f523c5         if(uchain == vchain){
d41d8c             //Query deve ir de chainPos[i] ate chainPos[v]
7d2150             ll cur = /*sum(chainPos[u], uchain) - (chainPos[u] == 0? 0LL :
                sum(chainPos[v] - 1, vchain))*/;
d133d8             ans += cur;
c2bef1             break;
cbb184         }
d41d8c
d41d8c         //Query deve ir de chainPos[i] ate o fim da estrutura
d41d8c         //ll cur = sum(chainPos[u], uchain);
d133d8         ans += cur;
a258cd         u = chainHead[uchain];
8039e1         u = parent[u];
cab24         uchain = chainInd[u];
cbb184     }
ba75d2     return ans;
cbb184 }
d41d8c
b7aa64 int dfs0(int pos, int prev = -1){
c92501     int res = 1;
97817b     for(int i = 0; i < (int)adj[pos].size(); i++){
ec49a3         int nx = adj[pos][i];
773904         if(nx != prev){
3f20e3             res += dfs0(nx, pos);
522845             parent[nx] = pos;
cbb184         }
cbb184     }
a1881c     return subsize[pos] = res;
cbb184 }
d41d8c
0b8977 int main()
f95b70 {
d41d8c     //Salvar arvore em adj
d41d8c
d41d8c     //Inicializa estrutura de dados
b75143     memset(chainHead, -1, sizeof(chainHead));
d41d8c
d41d8c     //Ou 0, se for o no raiz

```

## 4.8 Strongly Connected Components

```

5d1131 #include "../..//contest/header.hpp"
d41d8c
d41d8c /*
eaba86     Strongly connected components in  $O(V + E)$ :
970b0b     Finds all strongly connected components of a graph.
3f48a9     A strongly connected component is a maximal set of vertices
such that
de0185     every vertex can reach every other vertex in the component.
0a2f52     The graph where the SCCs are considered vertices is a DAG.
d41d8c
ca2095     Constraints:
000269     Vertices are labeled from 1 to n (inclusive).
d41d8c
b95cae     Usage:
ee0d06     Create the struct setting the maximum vertex label (n) and the
graph adjacency list (graph).
862124     Afterwards, ncomp has the number of SCCs in the graph and scc[i]
indicates the SCC i
67d138     belongs to ( $1 \leq \text{scc}[i] \leq \text{ncomp}$ ).
d41d8c
38224f     sorted is a topological ordering of the graph, byproduct of
the algorithm.
484a00     if edge  $a \rightarrow b$  exists, a appears before b in the sorted list.
c4c9bd */
d41d8c
d41d8c
73e60a struct scc_decomp
f95b70 {
9cf2b9     vector<int> *graph;
00b6a0     vector<vector<int>> tgraph;
1b013f     vector<int> scc;
1ee615     vector<bool> been;
8d35d1     int ncomp;
2035f3     list<int> sorted;
d41d8c
4875fb     scc_decomp(int n, vector<int> graph[]) : graph(graph), tgraph(n +
1), scc(n + 1, 0), been(n + 1, false), ncomp(0)
f95b70     {
5359f3         for (int i = 1; i <= n; i++)
6376e8             for (int j = 0; j < graph[i].size(); j++)
14234d                 tgraph[graph[i][j]].push_back(i);

```



```

d41d8c
5359f3     for (int i = 1; i <= n; i++)
018df6     if(!been[i])
1e5da3         dfs(i);
d41d8c
16ef86     for(int a : sorted)
f49735         if(scc[a] == 0)
f95b70     {
a8f1f2         ncomp++;
4dd966         dfst(a);
cbb184     }
cbb184 }
d41d8c
0cbab3 void dfs(int a)
f95b70 {
1689c6     been[a] = true;
c4fff3     for(int i = 0; i < graph[a].size(); i++)
b0c443         if(!been[graph[a][i]])
fded7f             dfs(graph[a][i]);
ddb66     sorted.push_front(a);
cbb184 }
d41d8c
9b760a void dfst(int a)
f95b70 {
1689c6     been[a] = true;
d28fcc     scc[a] = ncomp;
9c28b7     for(int i = 0; i < tgraph[a].size(); i++)
c480c5         if(scc[tgraph[a][i]] == 0)
caa482             dfst(tgraph[a][i]);
cbb184     }
2145c1 };
Full file hash: 20fe5c

```

## 5 Misc

### 5.1 DP Optimization - Binary Search

```

d41d8c // https://codeforces.com/contest/321/problem/E
d41d8c
5d1131 #include "../contest/header.hpp"
d41d8c
d41d8c /*
4e8f0c     Binary Search Optimization for DP:
6aaf9d     Optimizes dp of the form (or similar)  $dp[i][j] = \min_{k < i}(dp[k][j-1] + c(k + 1, i))$ .
78d4c1     The classical case is a partitioning dp, where k determines the
            break point for the next partition.
0467c8     In this case, i is the number of elements to partition and j is
            the number of partitions allowed.
d41d8c
537168     Let  $opt[i][j]$  be the values of k which minimize the function. (in
            case of tie, choose the smallest)
765123     To apply this optimization, you need  $opt[i][j] \leq opt[i+1][j]$ .
6c4b0a     That means the when you add an extra element (i + 1), your
            partitioning choice will not be to include more elements
efc594     than before (e.g. will no go from choosing [k, i] to [k-1, i+1]).
242554     This is usually intuitive by the problem details.
d41d8c
4d883e     Time goes from  $O(n^2m)$  to  $O(nm \log n)$ .
d41d8c
895144     To apply try to write the dp in the format above and verify if
            the property holds.
d41d8c
3db72f     Author: Arthur Pratti Dadalto
c4c9bd */
d41d8c
3494e9 #define MAXN 4123
fc36c1 #define MAXM 812
d41d8c
14e0a7 int n, m;
1590eb int u[MAXN][MAXN];
2bbe6d int tab[MAXN][MAXM];
d41d8c
65a7b7 inline int c(int i, int j)
f95b70 {
229880     return (u[j][j] - u[j][i - 1] - u[i - 1][j] + u[i - 1][i - 1]) / 2;
cbb184 }
d41d8c
d41d8c // This is responsible for computing  $tab[l...r][j]$ , knowing that  $opt[$ 
             $l...r][j]$  is in range  $[low\_opt...high\_opt]$ 
30d71a void compute(int j, int l, int r, int low_opt, int high_opt)
f95b70 {
c30a4b     int mid = (l + r) / 2, opt = -1; // mid is equivalent to i in the

```

original dp.

```

d41d8c
7222d3     tab[mid][j] = inf;
0e2f2c     for (int k = low_opt; k <= high_opt && k < mid; k++)
6f6e42         if (tab[k][j - 1] + c(k + 1, mid) < tab[mid][j])
f95b70         {
451068             tab[mid][j] = tab[k][j - 1] + c(k + 1, mid);
613f3c             opt = k;
cbb184         }
d41d8c
d41d8c     // New bounds on opt for other pending computation.
42c8a1     if (l <= mid - 1)
c7dd31         compute(j, l, mid - 1, low_opt, opt);
8b4e40     if (mid + 1 <= r)
8aa379         compute(j, mid + 1, r, opt, high_opt);
cbb184 }
d41d8c
13a4b1 int main(void)
f95b70 {
d69917     scanf("%d %d", &n, &m);
5359f3     for (int i = 1; i <= n; i++)
947790         for (int j = 1; j <= n; j++)
f95b70         {
433ab9             getchar();
512e3d             u[i][j] = getchar() - '0';
cbb184         }
d41d8c
5359f3     for (int i = 1; i <= n; i++)
947790         for (int j = 1; j <= n; j++)
a10370             u[i][j] += u[i - 1][j] + u[i][j - 1] - u[i - 1][j - 1];
d41d8c
5359f3     for (int i = 1; i <= n; i++)
5c5410         tab[i][0] = inf;
d41d8c
d41d8c     // Original dp
d41d8c     // for (int i = 1; i <= n; i++)
d41d8c     //   for (int j = 1; j <= m; j++)
d41d8c     //   {
d41d8c     //       tab[i][j] = inf;
d41d8c     //       for (int k = 0; k < i; k++)
d41d8c     //           tab[i][j] = min(tab[i][j], tab[k][j-1] + c(k + 1,i);
d41d8c     //   }
d41d8c
2e2a5d     for (int j = 1; j <= m; j++)
fdaa69         compute(j, 1, n, 0, n - 1);
d41d8c
721eeb     cout << tab[n][m] << endl;
cbb184 }
Full file hash: f2bb43

```

## 5.2 DP Optimization - CHT

```

d41d8c // https://codeforces.com/contest/319/problem/C
d41d8c
ad67d1 #include "../data_structures/line_container/line_container.cpp"
d41d8c
d41d8c /*
082e10     Convex Hull Trick for DP:
366d02     Transforms dp of the form (or similar)  $dp[i] = \min_{\{j < i\}}(dp[j]$ 
      +  $b[j] * a[i])$ .
ab5453     Time goes from  $O(n^2)$  to  $O(n \log n)$ , if using online line
      container, or  $O(n)$  if
56c021     lines are inserted in order of slope and queried in order of  $x$ .
d41d8c
3ffb56     To apply try to find a way to write the factor inside
      minimization as a linear function
ac2c47     of a value related to  $i$ . Everything else related to  $j$  will become
      constant.
c4c9bd */
d41d8c
69abfb #define MAXN 112345
d41d8c
a58cd5 int a[MAXN];
c4b25f int b[MAXN];
d41d8c
f80900 ll tab[MAXN];
d41d8c
13a4b1 int main(void)
f95b70 {
1a88fd     int n;
f4c120     scanf("%d", &n);
83008c     for (int i = 0; i < n; i++)
9376f3         scanf("%d", &a[i]);
83008c     for (int i = 0; i < n; i++)
264aeb         scanf("%d", &b[i]);
d41d8c
a447b8     tab[0] = 0;
79ab5f     line_container l;
c01116     l.add(-b[0], -tab[0]);
d41d8c
aa4866     for (int i = 1; i < n; i++)
f95b70     {
23bd61         tab[i] = -l.query(a[i]);
8fd447         l.add(-b[i], -tab[i]);
cbb184     }
d41d8c
d41d8c     // Original DP  $O(n^2)$ .
d41d8c     // for (int i = 1; i < n; i++)
d41d8c     // {
d41d8c     //     tab[i] = inf;

```

```

d41d8c // for (int j = 0; j < i; j++)
d41d8c //     tab[i] = min(tab[i], tab[j] + a[i] * b[j]);
d41d8c // }
d41d8c
cf6e15 cout << tab[n - 1] << endl;
cbb184 }
Full file hash: 722d84

```

### 5.3 DP Optimization - Knuth

```

d41d8c // https://www.spoj.com/problems/BRKSTRNG/
d41d8c
5d1131 #include "../contest/header.hpp"
d41d8c
d41d8c /*
c97188     Knuth Optimization for DP:
a41b1d     Optimizes dp of the form (or similar)  $dp[i][j] = \min_{i \leq k \leq j} \{dp[i][k-1] + dp[k+1][j] + c(i, j)\}$ .
f833b2     The classical case is building a optimal binary tree, where k
            determines the root.
d41d8c
c8aa2c     Let  $opt[i][j]$  be the value of k which minimizes the function. (in
            case of tie, choose the smallest)
c472ed     To apply this optimization, you need  $opt[i][j - 1] \leq opt[i][j] \leq opt[i+1][j]$ .
0eeb73     That means the when you remove an element form the left ( $i + 1$ ),
            you won't choose a breaking point more to the left than before.
b38eb7     Also, when you remove an element from the right ( $j - 1$ ), you won'
            t choose a breking point more to the right than before.
242554     This is usually intuitive by the problem details.
d41d8c
cbb42a     Time goes from  $O(n^3)$  to  $O(n^2)$ .
d41d8c
895144     To apply try to write the dp in the format above and verify if
            the property holds.
f76c09     Be careful with edge cases for opt.
d41d8c
3db72f     Author: Arthur Pratti Dadalto
c4c9bd */
d41d8c
dbf7e4 #define MAXN 1123
d41d8c
c4b25f int b[MAXN];
1ee552 ll tab[MAXN][MAXN];
38ab0d int opt[MAXN][MAXN];
ef864b int l, n;
d41d8c
5a7750 int c(int i, int j)
f95b70 {
33e24b     return b[j + 1] - b[i - 1];

```

```

cbb184 }
d41d8c
13a4b1 int main(void)
f95b70 {
57a598     while (scanf("%d %d", &l, &n) != EOF)
f95b70     {
5359f3         for (int i = 1; i <= n; i++)
264aeb             scanf("%d", &b[i]);
665bd2         b[n + 1] = l;
00d08b         b[0] = 0;
d41d8c
da41df         for (int i = 1; i <= n + 1; i++)
d6bc61             tab[i][i - 1] = 0, opt[i][i - 1] = i;
d41d8c
586d50         for (int i = n; i > 0; i--)
5d4199             for (int j = i; j <= n; j++)
f95b70             {
639af9                 tab[i][j] = infll;
823124                 for (int k = max(i, opt[i][j - 1]); k <= j && k <= opt[i +
1][j]; k++)
9e9168                     if (tab[i][k - 1] + tab[k + 1][j] + c(i, j) < tab[i][j])
f95b70                     {
680c31                         tab[i][j] = tab[i][k - 1] + tab[k + 1][j] + c(i, j);
14da03                         opt[i][j] = k;
cbb184                     }
cbb184             }
d41d8c
ea7bd9     printf("%lld\n", tab[1][n]);
cbb184 }
cbb184 }
Full file hash: 17b7c8

```

## 5.4 Ternary Search (continuous)

```

d41d8c  /*
0a2f9f     Ternary Search:
28ea2d     Finds x such that f(x) is minimum in range [bot, top] in O(lg((
top - bot) / eps)).
6f01ba     Value is correct within the specified precision eps.
d41d8c
ca2095     Constraints:
7474cd     f(x) is strictly decreasing for some interval [bot, x1], constant
in an interval [x1, x2]
60dab3     and strictly increasing in a interval [x2, top]. x1 <= x2 are
arbitrary values where [x1, x2]
7523fb     is a plateau of optimal solutions.
d41d8c
b95cae     Usage:
5b60e3     Call the function passing a lambda expression or function f.
8bc9f4     If there are multiple possible solutions, assume that an

```

arbitrary one in the plateau is returned.

```
d41d8c
3db72f    Author: Arthur Pratti Dadalto
c4c9bd    */
d41d8c
398727    template <typename F>
ca64a1    double ternary_search(const F &f, double bot = -1e9, double top = 1e9
    , double eps = 1e-9)
f95b70    {
14d91b        while (top - bot > eps)
f95b70        {
8e37d7            double x1 = (0.55*bot + 0.45*top); // (2*bot + top) / 3 is more
    stable, but slower.
3f8318            double x2 = (0.45*bot + 0.55*top);
9482ba            if (f(x1) > f(x2))
443914                bot = x1;
2954e9            else
16b1b8                top = x2;
cbb184        }
d41d8c
05eb81        return (bot + top) / 2;
cbb184    }
Full file hash: 082811
```

## 6 Number Theory

### 6.1 Euclid

```

5d1131 #include "../contest/header.hpp"
d41d8c
d41d8c /*
4b5b70     Extended Euclidean Algorithm:
71fa74     Returns the gcd of a and b.
2c19ff     Also finds numbers x and y for which  $a * x + b * y = \text{gcd}(a, b)$  (
           not unique).
a0c250     All pairs can be represented in the form  $(x + k * b / \text{gcd}, y - k$ 
            $* a / \text{gcd})$  for k an arbitrary integer.
57ad55     If there are several such x and y, the function returns the pair
           for which  $|x| + |y|$  is minimal.
2d446b     If there are several x and y satisfying the minimal criteria, it
           outputs the pair for which  $X \leq Y$ .
d41d8c
3997db     Source: modified from https://cp-algorithms.com/algebra/extended-
           euclid-algorithm.html
d41d8c
b95cae     Usage:
6475da     For non-extendend version, c++ has __gcd and __lcm.
d41d8c
ca2095     Constraints:
30a9e9     Produces correct results for negative integers as well.
c4c9bd */
d41d8c
4fce64     template<class T>
94606e     T gcd(T a, T b, T &x, T &y)
f95b70     {
fcbb63         if (b == 0)
f95b70         {
483406             x = 1;
01dbf4             y = 0;
3f5343             return a;
cbb184         }
d41d8c
32895f         T x1, y1;
254183         T d = gcd(b, a % b, x1, y1);
711e33         x = y1;
a2a46d         y = x1 - y1 * (a / b);
be245b         return d;
cbb184     }
Full file hash: 0c35ae

```

### 6.2 Pollard rho

```

d41d8c /*
baee35     Description: Pollard-rho randomized factorization algorithm.
           Returns prime

```



9849b1 factors of a number, in arbitrary order (e.g. 2299 -> {11, 19, 11})

c21e7b Time:  $O(n^{1/4})$  gcd calls, less for numbers with small factors.

d41d8c

1d1558 Source: <https://github.com/kth-competitive-programming/kactl/blob/master/content/number-theory/Factor.h>

c4c9bd \*/

f4cf5b typedef unsigned long long ull;

088cf4 typedef long double ld;

d41d8c

ae330e ull mod\_mul(ull a, ull b, ull M) {

053258 ll ret = a \* b - M \* ull(ld(a) \* ld(b) / ld(M));

964402 return ret + M \* (ret < 0) - M \* (ret >= (ll)M);

cbb184 }

97f234 ull mod\_pow(ull b, ull e, ull mod) {

c1a4a1 ull ans = 1;

4d1884 for (; e; b = mod\_mul(b, b, mod), e /= 2)

0c1ecc if (e & 1) ans = mod\_mul(ans, b, mod);

ba75d2 return ans;

cbb184 }

d41d8c

da49ed bool isPrime(ull n) {

32f8ec if (n < 2 || n % 6 % 4 != 1) return n - 2 < 2;

43a246 ull A[] = {2, 325, 9375, 28178, 450775, 9780504, 1795265022},

c17dd6 s = \_\_builtin\_ctzll(n-1), d = n >> s;

d236d6 for(auto &a : A) { // ^ count trailing zeroes

8a86e5 ull p = mod\_pow(a, d, n), i = s;

274cbc while (p != 1 && p != n - 1 && a % n && i--)

2cbb80 p = mod\_mul(p, p, n);

e2871b if (p != n-1 && i != s) return 0;

cbb184 }

6a5530 return 1;

cbb184 }

d41d8c

7eb30f ull pollard(ull n) {

4de5da auto f = [n](ull x) { return (mod\_mul(x, x, n) + 1) % n; };

68eade if (!(n & 1)) return 2;

7b6fa7 for (ull i = 2;; i++) {

e17462 ull x = i, y = f(x), p;

332fe8 while ((p = \_\_gcd(n + y - x, n)) == 1)

b789c2 x = f(x), y = f(f(y));

c493bb if (p != n) return p;

cbb184 }

cbb184 }

d41d8c

bc43a4 vector<ull> factorize(ull n) {

1b90e8 if (n == 1) return {};

6b5b32 if (isPrime(n)) return {n};

bc6125 ull x = pollard(n);

b3b29a auto l = factorize(x), r = factorize(n / x);

```

7af87c    l.insert(l.end(), all(r));
792fd4    return l;
cbb184 }
Full file hash: 93549d

```

## 6.3 Modular Inverse

```

771bdd    #include "../euclid/euclid.cpp"
d41d8c
d41d8c    /*
18a91e        Modular Inverse:
76e032            Returns an integer x such that (a * x) % m == 1.
4e4745            The modular inverse exists if and only if a and m are relatively
                prime.
ff1c03            Modular inverse is also equal to a^(phi(m) - 1) % m.
261f2c            In particular, if m is prime a^(-1) == a^(m-2), which might be
                faster to code.
d41d8c
3997db        Source: modified from https://cp-algorithms.com/algebra/module-
                inverse.html
c4c9bd    */
d41d8c
4fce64    template<class T>
b267c1    T mod_inverse(T a, T m)
f95b70    {
645c5d        T x, y;
6553c1        assert(gcd(a, m, x, y) == 1); // Or return something, if gcd is not
                1 the inverse doesn't exist.
08ffd4        return (x % m + m) % m;
cbb184    }
Full file hash: 7efa11

```

## 6.4 Phi

```

5d1131    #include "../contest/header.hpp"
d41d8c
d41d8c    /*
bf26c1        Euler's totient function (PHI):
fc5093            Euler's totient function, also known as phi-function PHI(n),
                counts the number of integers
d7ef61            between 1 and n inclusive, which are coprime to n. Two numbers
                are coprime if their greatest
bf3431            common divisor equals 1 (1 is considered to be coprime to any
                number).
d41d8c
d41d8c
3997db        Source: modified from https://cp-algorithms.com/algebra/phi-
                function.html
e1fde7            and https://github.com/kth-competitive-programming/kactl/blob/
                master/content/number-theory/phiFunction.h
d41d8c

```

```

b95cae    Usage:
a1b248    Some useful properties:
9f5d84    - If p is a prime number, PHI(p)=p-1.
d4d311    - If a and b are relatively prime, PHI(ab)=PHI(a)*PHI(b).
65a02c    - In general, for not coprime a and b, PHI(ab)=PHI(a)*PHI(b)*d/
    PHI(d), with d=gcd(a,b) holds.
417c3d    - PHI(PHI(m)) <= m / 2
037fb5    - Euler's theorem: a^PHI(m) === 1 (mod m), for a and m coprime.
bffb1c    - For a and m coprime: a^n === a^(n % PHI(m)) (mod m)
986ce1    - For arbitrary x,m and n >= log_2(m): x^n === x^(PHI(m)+[n % PHI
    (m)]) (mod m)
8d568b    The one above allows computing modular exponentiation for really
    large exponents.
ec5d4e    - If d is a divisor of n, then there are phi(n/d) numbers i <= n
    for which gcd(i,n)=d
137411    - sum_{d|n} phi(d) = n
c228b3    - sum_{1 <= k <= n, gcd(k,n)=1} k = n * phi(n) / 2, for n > 1
c4c9bd    */
d41d8c
d41d8c    // Use this one for few values of phi.
b5f6f9    int phi(int n)
f95b70    {
efa47a        int result = n;
83f497        for (int i = 2; i * i <= n; i++)
f95b70        {
775f6d            if (n % i == 0)
f95b70            {
49edb8                while (n % i == 0)
1358bf                    n /= i;
21cd49                    result -= result / i;
cbb184            }
cbb184        }
f3d362        if (n > 1)
e48781            result -= result / n;
dc8384        return result;
cbb184    }
d41d8c
4fee4d    namespace totient
f95b70    {
2d637a        const int MAXV = 1000001; // Takes ~0.03 s for 10^6.
6e559b        int phi[MAXV];
d41d8c
b2a56e        void init()
f95b70        {
9484bb            for (int i = 0; i < MAXV; i++)
ed1f90                phi[i] = i & 1 ? i : i / 2;
9be7d6            for (int i = 3; i < MAXV; i += 2)
a2252f                if (phi[i] == i)
8a4437                    for (int j = i; j < MAXV; j += i)
a9ba36                        phi[j] -= phi[j] / i;

```

```
cbb184 }
cbb184 } // namespace totient
Full file hash: e79764
```

## 6.5 Sieve

```
5d1131 #include "../contest/header.hpp"
d41d8c
d41d8c /*
d5cfbe     Sieve of Eratosthenes:
0a5343     Finds all primes in interval [2, MAXP] in O(MAXP) time.
7e51df     Also finds lp[i] for every i in [2, MAXP], such that lp[i] is the
           minimum prime factor of i.
6dbf8e     Particularly useful for factorization.
d41d8c
3997db     Source: modified from https://cp-algorithms.com/algebra/prime-sieve-linear.html
d41d8c
b95cae     Usage:
9290fb     Set MAXP and call init.
85265b     Sieve for 10^7 should run in about 0.2 s.
c4c9bd */
d41d8c
2ca8b0 namespace sieve
f95b70 {
bace98     const int MAXP = 10000000; // Will find primes in interval [2, MAXP].
39b809     int lp[MAXP + 1]; // lp[i] is the minimum prime factor of i.
632f82     vector<int> p; // Ordered list of primes up to MAXP.
d41d8c
b2a56e     void init()
f95b70     {
008cd3         for (int i = 2; i <= MAXP; i++)
f95b70         {
d4a1cc             if (lp[i] == 0)
b6fab7                 p.push_back(lp[i] = i);
d41d8c
9d854a             for (int j = 0; j < (int)p.size() && p[j] <= lp[i] && i * p[j] <=
           MAXP; j++)
fb7d48                 lp[i * p[j]] = p[j];
cbb184         }
cbb184     }
cbb184 } // namespace sieve
Full file hash: e9076d
```

## 7 Numerical

### 7.1 Big Int

```

5d1131  #include "../contest/header.hpp"
d41d8c
d41d8c  // This code is not meant to be written in icpc contests. This is
      just here to fill a void for now.
d41d8c  // Source: someone on CF
d41d8c
d41d8c  // NOTE:
d41d8c  // This code contains various bug fixes compared to the original
      version from
d41d8c  // indy256 (github.com/indy256/codelibrary/blob/master/cpp/
      numbertheory/bigint-full.cpp),
d41d8c  // including:
d41d8c  // - Fix overflow bug in mul_karatsuba.
d41d8c  // - Fix overflow bug in fft.
d41d8c  // - Fix bug in initialization from long long.
d41d8c  // - Optimized operators + - *.
d41d8c  //
d41d8c  // Tested:
d41d8c  // - https://www.e-olymp.com/en/problems/266: Comparison
d41d8c  // - https://www.e-olymp.com/en/problems/267: Subtraction
d41d8c  // - https://www.e-olymp.com/en/problems/271: Multiplication
d41d8c  // - https://www.e-olymp.com/en/problems/272: Multiplication
d41d8c  // - https://www.e-olymp.com/en/problems/313: Addition
d41d8c  // - https://www.e-olymp.com/en/problems/314: Addition/Subtraction
d41d8c  // - https://www.e-olymp.com/en/problems/317: Multiplication (simple
      / karatsuba / fft)
d41d8c  // - https://www.e-olymp.com/en/problems/1327: Multiplication
d41d8c  // - https://www.e-olymp.com/en/problems/1328
d41d8c  // - VOJ BIGNUM: Addition, Subtraction, Multiplication.
d41d8c  // - SGU 111: sqrt
d41d8c  // - SGU 193
d41d8c  // - SPOJ MUL, VFMUL: Multiplication.
d41d8c  // - SPOJ FDIV, VFIV: Division.
d41d8c
d73a77  const int BASE_DIGITS = 9;
82e97b  const int BASE = 1000000000;
d41d8c
6acb6c  struct BigInt {
d65d12      int sign;
a9d078      vector<int> a;
d41d8c
d41d8c      // ----- Constructors -----
d41d8c      // Default constructor.
1acfca      BigInt() : sign(1) {}
d41d8c
d41d8c      // Constructor from long long.

```

```

ccf902     BigInt(long long v) {
324222         *this = v;
cbb184     }
235125     BigInt& operator = (long long v) {
ce6fc2         sign = 1;
ea2149         if (v < 0) {
6a74a9             sign = -1;
6fab41             v = -v;
cbb184         }
22838a         a.clear();
fefe2d         for (; v > 0; v = v / BASE)
c237f1             a.push_back(v % BASE);
357a55         return *this;
cbb184     }

d41d8c     // Initialize from string.
d41d8c     BigInt(const string& s) {
c710ec         read(s);
cbb184     }

d41d8c     // ----- Input / Output -----
d41d8c     void read(const string& s) {
6c30c4         sign = 1;
ce6fc2         a.clear();
22838a         int pos = 0;
bec7a6         while (pos < (int) s.size() && (s[pos] == '-' || s[pos] == '+'
a68fdf             ')) {
dbe226             if (s[pos] == '-')
2b8bd1                 sign = -sign;
17dad0             ++pos;
cbb184         }
7959ef         for (int i = s.size() - 1; i >= pos; i -= BASE_DIGITS) {
c67d6f             int x = 0;
d343c4             for (int j = max(pos, i - BASE_DIGITS + 1); j <= i; j++)
cfc7e4                 x = x * 10 + s[j] - '0';
7c6978             a.push_back(x);
cbb184         }
0ebb65         trim();
cbb184     }

bd2995     friend istream& operator>>(istream &stream, BigInt &v) {
ac0066         string s;
e0c759         stream >> s;
c4002a         v.read(s);
a87cf7         return stream;
cbb184     }

d41d8c     friend ostream& operator<<(ostream &stream, const BigInt &v) {
44647f         if (v.sign == -1 && !v.isZero())
b5c525             stream << '-';
27bc2a         stream << (v.a.empty() ? 0 : v.a.back());
4fda68

```

```

fce618         for (int i = (int) v.a.size() - 2; i >= 0; --i)
018b85             stream << setw(BASE_DIGITS) << setfill('0') << v.a[i];
a87cf7         return stream;
cbb184     }
d41d8c
d41d8c     // ----- Comparison -----
7014c0     bool operator<(const BigInt &v) const {
eb909f         if (sign != v.sign)
603965             return sign < v.sign;
a2765e         if (a.size() != v.a.size())
f7d303             return a.size() * sign < v.a.size() * v.sign;
305fef         for (int i = ((int) a.size()) - 1; i >= 0; i--)
00d0de             if (a[i] != v.a[i])
2441c5                 return a[i] * sign < v.a[i] * sign;
d1fe4d         return false;
cbb184     }
d41d8c
426053     bool operator>(const BigInt &v) const {
54bd3a         return v < *this;
cbb184     }
65677c     bool operator<=(const BigInt &v) const {
0fe7a0         return !(v < *this);
cbb184     }
605209     bool operator>=(const BigInt &v) const {
d9c542         return !(*this < v);
cbb184     }
880606     bool operator==(const BigInt &v) const {
7f44a6         return !(*this < v) && !(v < *this);
cbb184     }
062171     bool operator!=(const BigInt &v) const {
6c55aa         return *this < v || v < *this;
cbb184     }
d41d8c
d41d8c     // Returns:
d41d8c     // 0 if |x| == |y|
d41d8c     // -1 if |x| < |y|
d41d8c     // 1 if |x| > |y|
ce6386     friend int __compare_abs(const BigInt& x, const BigInt& y) {
e78df5         if (x.a.size() != y.a.size()) {
c86c62             return x.a.size() < y.a.size() ? -1 : 1;
cbb184         }
d41d8c
a552ab         for (int i = ((int) x.a.size()) - 1; i >= 0; --i) {
a5b2df             if (x.a[i] != y.a[i]) {
b1ec3d                 return x.a[i] < y.a[i] ? -1 : 1;
cbb184             }
cbb184         }
bb30ba         return 0;
cbb184     }
d41d8c

```

```

d41d8c      // ----- Unary operator - and operators +-
-----
1e3c00      BigInt operator-() const {
18bf1f          BigInt res = *this;
b9607c          if (isZero()) return res;
d41d8c
290faa          res.sign = -sign;
b5053e          return res;
cbb184      }
d41d8c
d41d8c      // Note: sign ignored.
d60e6f      void __internal_add(const BigInt& v) {
f7247c          if (a.size() < v.a.size()) {
2ce41c              a.resize(v.a.size(), 0);
cbb184          }
1addcf          for (int i = 0, carry = 0; i < (int) max(a.size(), v.a.size()
) || carry; ++i) {
df4512              if (i == (int) a.size()) a.push_back(0);
d41d8c
85e77e              a[i] += carry + (i < (int) v.a.size() ? v.a[i] : 0);
49bff0              carry = a[i] >= BASE;
1791a8              if (carry) a[i] -= BASE;
cbb184          }
cbb184      }
d41d8c
d41d8c      // Note: sign ignored.
8b47dc      void __internal_sub(const BigInt& v) {
65cb2e          for (int i = 0, carry = 0; i < (int) v.a.size() || carry; ++i
) {
a1437d              a[i] -= carry + (i < (int) v.a.size() ? v.a[i] : 0);
e0b1f1              carry = a[i] < 0;
da53a6              if (carry) a[i] += BASE;
cbb184          }
0e329b          this->trim();
cbb184      }
d41d8c
89fb6b      BigInt operator += (const BigInt& v) {
8ea459          if (sign == v.sign) {
570069              __internal_add(v);
9d9745          } else {
ae3659              if (__compare_abs(*this, v) >= 0) {
e9815a                  __internal_sub(v);
9d9745              } else {
dcc3fe                  BigInt vv = v;
3c5f43                  swap(*this, vv);
fe0d8d                  __internal_sub(vv);
cbb184              }
cbb184          }
357a55          return *this;
cbb184      }

```



```

d41d8c
6b1a22     BigInt operator -= (const BigInt& v) {
8ea459         if (sign == v.sign) {
ae3659             if (__compare_abs(*this, v) >= 0) {
e9815a                 __internal_sub(v);
9d9745             } else {
dcc3fe                 BigInt vv = v;
3c5f43                 swap(*this, vv);
fe0d8d                 __internal_sub(vv);
0db96d                 this->sign = -this->sign;
cbb184             }
9d9745         } else {
570069             __internal_add(v);
cbb184         }
357a55         return *this;
cbb184     }
d41d8c
d41d8c     // Optimize operators + and - according to
d41d8c     // https://stackoverflow.com/questions/13166079/move-semantics-
    and-pass-by-rvalue-reference-in-overloaded-arithmetic
f1e02d     template< typename L, typename R >
81c687         typename std::enable_if<
4eceb0             std::is_convertible<L, BigInt>::value &&
c0db24             std::is_convertible<R, BigInt>::value &&
061102             std::is_lvalue_reference<R&&>::value,
6b2030             BigInt>::type friend operator + (L&& l, R&& r) {
46b960     BigInt result(std::forward<L>(l));
fbef75     result += r;
dc8384     return result;
cbb184 }
f1e02d     template< typename L, typename R >
81c687         typename std::enable_if<
4eceb0             std::is_convertible<L, BigInt>::value &&
c0db24             std::is_convertible<R, BigInt>::value &&
bcc2f             std::is_rvalue_reference<R&&>::value,
6b2030             BigInt>::type friend operator + (L&& l, R&& r) {
5f09ae     BigInt result(std::move(r));
a5a040     result += l;
dc8384     return result;
cbb184 }
d41d8c
f1e02d     template< typename L, typename R >
81c687         typename std::enable_if<
4eceb0             std::is_convertible<L, BigInt>::value &&
6ca6cc             std::is_convertible<R, BigInt>::value,
1612ea             BigInt>::type friend operator - (L&& l, R&& r) {
46b960     BigInt result(std::forward<L>(l));
1d15a0     result -= r;
dc8384     return result;
cbb184 }

```

```

d41d8c
d41d8c // ----- Operators * / % -----
a179f4 friend pair<BigInt, BigInt> divmod(const BigInt& a1, const BigInt
& b1) {
872d46     assert(b1 > 0); // divmod not well-defined for b < 0.
d41d8c
25f4e9     long long norm = BASE / (b1.a.back() + 1);
7c41dc     BigInt a = a1.abs() * norm;
ecd4f4     BigInt b = b1.abs() * norm;
da5ddc     BigInt q = 0, r = 0;
90ee93     q.a.resize(a.a.size());
d41d8c
72b5b8     for (int i = a.a.size() - 1; i >= 0; i--) {
79aca3         r *= BASE;
0caac0         r += a.a[i];
0eeb4e         long long s1 = r.a.size() <= b.a.size() ? 0 : r.a[b.a.
size()];
bcl1a99         long long s2 = r.a.size() <= b.a.size() - 1 ? 0 : r.a[b.a
.size() - 1];
0ebba0         long long d = ((long long) BASE * s1 + s2) / b.a.back();
5d4f85         r -= b * d;
612239         while (r < 0) {
bd3902             r += b, --d;
cbb184         }
5898c8         q.a[i] = d;
cbb184     }
d41d8c
535024     q.sign = a1.sign * b1.sign;
a29af3     r.sign = a1.sign;
36a918     q.trim();
9a35fd     r.trim();
38a539     auto res = make_pair(q, r / norm);
458098     if (res.second < 0) res.second += b1;
b5053e     return res;
cbb184 }
547e4b BigInt operator/(const BigInt &v) const {
ce8f7c     return divmod(*this, v).first;
cbb184 }
d41d8c
ee46c3 BigInt operator%(const BigInt &v) const {
7a671a     return divmod(*this, v).second;
cbb184 }
d41d8c
c2998e void operator/=(int v) {
d1ee66     assert(v > 0); // operator / not well-defined for v <= 0.
dd9f94     if (llabs(v) >= BASE) {
85cc00         *this /= BigInt(v);
505b97         return ;
cbb184     }
8e679f     if (v < 0)

```

```

20198f         sign = -sign, v = -v;
8e5533     for (int i = (int) a.size() - 1, rem = 0; i >= 0; --i) {
cbe153         long long cur = a[i] + rem * (long long) BASE;
8d1e71         a[i] = (int) (cur / v);
cb35e0         rem = (int) (cur % v);
cbb184     }
0ebb65     trim();
cbb184 }

d41d8c
49658a     BigInt operator/(int v) const {
d1ee66         assert(v > 0); // operator / not well-defined for v <= 0.
d41d8c
dd9f94         if (llabs(v) >= BASE) {
ed0225             return *this / BigInt(v);
cbb184         }
18bf1f         BigInt res = *this;
37184f         res /= v;
b5053e         return res;
cbb184     }
3b4fa6     void operator/=(const BigInt &v) {
e51f70         *this = *this / v;
cbb184     }
d41d8c

54c35d     long long operator%(long long v) const {
d1ee66         assert(v > 0); // operator / not well-defined for v <= 0.
a1e888         assert(v < BASE);
cbcd95         int m = 0;
947442         for (int i = a.size() - 1; i >= 0; --i)
95269a             m = (a[i] + m * (long long) BASE) % v;
9af577         return m * sign;
cbb184     }
d41d8c

a0b62a     void operator*=(int v) {
dd9f94         if (llabs(v) >= BASE) {
014cdd             *this *= BigInt(v);
505b97             return ;
cbb184         }
8e679f         if (v < 0)
20198f             sign = -sign, v = -v;
c6279c         for (int i = 0, carry = 0; i < (int) a.size() || carry; ++i)
        {
74ab7d             if (i == (int) a.size())
ddfb75                 a.push_back(0);
d09f08             long long cur = a[i] * (long long) v + carry;
98cd39             carry = (int) (cur / BASE);
861843             a[i] = (int) (cur % BASE);
d41d8c             //asm("divl %%ecx" : "=a"(carry), "=d"(a[i]) : "A"(cur),
        "c"(base));
d41d8c             /*
97f03f             int val;

```

```

ab8362         __asm {
bab6b5         lea esi, cur
6cd1f3         mov eax, [esi]
d5ad3f         mov edx, [esi+4]
378c50         mov ecx, base
d88250         div ecx
e3e615         mov carry, eax
6f8726         mov val, edx;
cbb184         }
26a9ce         a[i] = val;
c4c9bd         */
cbb184     }
0ebb65     trim();
cbb184 }
d41d8c
d1d185     BigInt operator*(int v) const {
dd9f94         if (llabs(v) >= BASE) {
42696e             return *this * BigInt(v);
cbb184         }
18bf1f         BigInt res = *this;
6b38f1         res *= v;
b5053e         return res;
cbb184     }
d41d8c
d41d8c     // Convert BASE 10^old --> 10^new.
ead252     static vector<int> convert_base(const vector<int> &a, int
    old_digits, int new_digits) {
943071         vector<long long> p(max(old_digits, new_digits) + 1);
c4bbd4         p[0] = 1;
85cf8d         for (int i = 1; i < (int) p.size(); i++)
7cc6c9             p[i] = p[i - 1] * 10;
02fb60         vector<int> res;
c6278d         long long cur = 0;
6427c9         int cur_digits = 0;
c0e004         for (int i = 0; i < (int) a.size(); i++) {
b28c31             cur += a[i] * p[cur_digits];
e4696c             cur_digits += old_digits;
5ebda5             while (cur_digits >= new_digits) {
6f203f                 res.push_back((long long)(cur % p[new_digits]));
1cec8a                 cur /= p[new_digits];
318982                 cur_digits -= new_digits;
cbb184             }
cbb184         }
a5eaaa         res.push_back((int) cur);
c5a021         while (!res.empty() && !res.back())
efcb65             res.pop_back();
b5053e         return res;
cbb184     }
d41d8c
009dfc     void fft(vector<complex<double> > & a, bool invert) const {

```

```

8ec808         int n = (int) a.size();
d41d8c
677a94         for (int i = 1, j = 0; i < n; ++i) {
4af5d7             int bit = n >> 1;
425aec             for (; j >= bit; bit >>= 1)
b39a0f                 j -= bit;
297413             j += bit;
9dcc5c             if (i < j)
33275d                 swap(a[i], a[j]);
cbb184         }
d41d8c
eb733a         for (int len = 2; len <= n; len <= 1) {
2f82ea             double ang = 2 * 3.14159265358979323846 / len * (invert ?
-1 : 1);
a0b444             complex<double> wlen(cos(ang), sin(ang));
6c8781             for (int i = 0; i < n; i += len) {
c2eaad                 complex<double> w(1);
876230                 for (int j = 0; j < len / 2; ++j) {
371eda                     complex<double> u = a[i + j];
0c0391                     complex<double> v = a[i + j + len / 2] * w;
6c3014                     a[i + j] = u + v;
273255                     a[i + j + len / 2] = u - v;
3e4104                     w *= wlen;
cbb184                 }
cbb184             }
cbb184         }
2111a0         if (invert)
6cb8cc             for (int i = 0; i < n; ++i)
b098a6                 a[i] /= n;
cbb184     }
d41d8c
0d5969         void multiply_fft(const vector<int> &a, const vector<int> &b,
vector<int> &res) const {
58dd64             vector<complex<double> > fa(a.begin(), a.end());
249aaa             vector<complex<double> > fb(b.begin(), b.end());
43ec81             int n = 1;
727e5e             while (n < (int) max(a.size(), b.size()))
c149a4                 n <= 1;
c149a4             n <= 1;
37aa6c             fa.resize(n);
870070             fb.resize(n);
d41d8c
3a13f2             fft(fa, false);
c76760             fft(fb, false);
6cb8cc             for (int i = 0; i < n; ++i)
940eb7                 fa[i] *= fb[i];
959d01             fft(fa, true);
d41d8c
f38aa2             res.resize(n);
6e20af             long long carry = 0;

```

```

baeb9e         for (int i = 0; i < n; ++i) {
6e6901             long long t = (long long) (fa[i].real() + 0.5) + carry;
9e18f0             carry = t / 1000;
bb5b3b             res[i] = t % 1000;
cbb184         }
cbb184     }
d41d8c
d64466     BigInt mul_simple(const BigInt &v) const {
02a624         BigInt res;
325cfe         res.sign = sign * v.sign;
4bc9af         res.a.resize(a.size() + v.a.size());
7a7093         for (int i = 0; i < (int) a.size(); ++i)
b40a68             if (a[i])
761845                 for (int j = 0, carry = 0; j < (int) v.a.size() ||
carry; ++j) {
df3e98                     long long cur = res.a[i + j] + (long long) a[i] *
(j < (int) v.a.size() ? v.a[j] : 0) + carry;
98cd39                     carry = (int) (cur / BASE);
ff01d5                     res.a[i + j] = (int) (cur % BASE);
cbb184                 }
d7ee6d             res.trim();
b5053e             return res;
cbb184     }
d41d8c
ad1556     typedef vector<long long> vll;
d41d8c
4d42f9     static vll karatsubaMultiply(const vll &a, const vll &b) {
94d5f8         int n = a.size();
1fb0e0         vll res(n + n);
44d3ec         if (n <= 32) {
83008c             for (int i = 0; i < n; i++)
f90a6b                 for (int j = 0; j < n; j++)
8dd9af                     res[i + j] += a[i] * b[j];
b5053e             return res;
cbb184         }
d41d8c
af0b16         int k = n >> 1;
f9fca2         vll a1(a.begin(), a.begin() + k);
72c0c7         vll a2(a.begin() + k, a.end());
48ebf6         vll b1(b.begin(), b.begin() + k);
88c9a6         vll b2(b.begin() + k, b.end());
d41d8c
03c868         vll a1b1 = karatsubaMultiply(a1, b1);
e56678         vll a2b2 = karatsubaMultiply(a2, b2);
d41d8c
40d6ad         for (int i = 0; i < k; i++)
c20ed7             a2[i] += a1[i];
40d6ad         for (int i = 0; i < k; i++)
b009cc             b2[i] += b1[i];
d41d8c

```

```

6a2f29         vll r = karatsubaMultiply(a2, b2);
be9bd2         for (int i = 0; i < (int) a1b1.size(); i++)
47fef2             r[i] -= a1b1[i];
cf04ec         for (int i = 0; i < (int) a2b2.size(); i++)
00a00c             r[i] -= a2b2[i];
d41d8c
5951a9         for (int i = 0; i < (int) r.size(); i++)
1bf61e             res[i + k] += r[i];
be9bd2         for (int i = 0; i < (int) a1b1.size(); i++)
d6cf88             res[i] += a1b1[i];
cf04ec         for (int i = 0; i < (int) a2b2.size(); i++)
ab9916             res[i + n] += a2b2[i];
b5053e         return res;
cbb184     }
d41d8c
287510     BigInt mul_karatsuba(const BigInt &v) const {
48c647         vector<int> a6 = convert_base(this->a, BASE_DIGITS, 6);
f64a05         vector<int> b6 = convert_base(v.a, BASE_DIGITS, 6);
e1cb30         vll a(a6.begin(), a6.end());
5ed74f         vll b(b6.begin(), b6.end());
1a813e         while (a.size() < b.size())
ddf7b5             a.push_back(0);
0d118e         while (b.size() < a.size())
c40831             b.push_back(0);
634b60         while (a.size() & (a.size() - 1))
eed3fb             a.push_back(0), b.push_back(0);
16bf35         vll c = karatsubaMultiply(a, b);
02a624         BigInt res;
325cfe         res.sign = sign * v.sign;
6e20af         long long carry = 0;
7dbc9f         for (int i = 0; i < (int) c.size(); i++) {
dc97b8             long long cur = c[i] + carry;
cdf472             res.a.push_back((int) (cur % 1000000));
735fb2             carry = cur / 1000000;
cbb184         }
7b10c4         res.a = convert_base(res.a, 6, BASE_DIGITS);
d7ee6d         res.trim();
b5053e         return res;
cbb184     }
d41d8c
933d02     void operator*=(const BigInt &v) {
fa4bc1         *this = *this * v;
cbb184     }
24478f     BigInt operator*(const BigInt &v) const {
de6792         if (a.size() * v.a.size() <= 1000111) return mul_simple(v);
fec548         if (a.size() > 500111 || v.a.size() > 500111) return mul_fft(
    v);
a67c32         return mul_karatsuba(v);
cbb184     }
d41d8c

```

```

0f0ce5     BigInt mul_fft(const BigInt& v) const {
02a624         BigInt res;
325cfe         res.sign = sign * v.sign;
d1a018         multiply_fft(convert_base(a, BASE_DIGITS, 3), convert_base(v.
    a, BASE_DIGITS, 3), res.a);
74be5c         res.a = convert_base(res.a, 3, BASE_DIGITS);
d7ee6d         res.trim();
b5053e         return res;
cbb184     }
d41d8c
d41d8c     // ----- Misc -----
9f0aff     BigInt abs() const {
18bf1f         BigInt res = *this;
3ccc69         res.sign *= res.sign;
b5053e         return res;
cbb184     }
a0fac1     void trim() {
b03a9b         while (!a.empty() && !a.back())
4685a5             a.pop_back();
e28510         if (a.empty())
ce6fc2             sign = 1;
cbb184     }
d41d8c
88d324     bool isZero() const {
5c0518         return a.empty() || (a.size() == 1 && !a[0]);
cbb184     }
d41d8c
e7ccd6     friend BigInt gcd(const BigInt &a, const BigInt &b) {
183a15         return b.isZero() ? a : gcd(b, a % b);
cbb184     }
7977e6     friend BigInt lcm(const BigInt &a, const BigInt &b) {
8b81ac         return a / gcd(a, b) * b;
cbb184     }
d41d8c
2f7166     friend BigInt sqrt(const BigInt &a1) {
b25149         BigInt a = a1;
53b77e         while (a.a.empty() || a.a.size() % 2 == 1)
8a6b34             a.a.push_back(0);
d41d8c
0c5896         int n = a.a.size();
d41d8c
f9194d         int firstDigit = (int) sqrt((double) a.a[n - 1] * BASE + a.a[
    n - 2]);
3c7b49         int norm = BASE / (firstDigit + 1);
b65c20         a *= norm;
b65c20         a *= norm;
53b77e         while (a.a.empty() || a.a.size() % 2 == 1)
8a6b34             a.a.push_back(0);
d41d8c
8a28a4         BigInt r = (long long) a.a[n - 1] * BASE + a.a[n - 2];

```



```

4e5685         firstDigit = (int) sqrt((double) a.a[n - 1] * BASE + a.a[n -
    2]);
97c0e8         int q = firstDigit;
02a624         BigInt res;
d41d8c
a1054f         for(int j = n / 2 - 1; j >= 0; j--) {
e63f29             for(; ; --q) {
592185                 BigInt r1 = (r - (res * 2 * BigInt(BASE) + q) * q) *
    BigInt(BASE) * BigInt(BASE) + (j > 0 ? (long long) a.a[2 * j - 1] * BASE +
    a.a[2 * j - 2] : 0);
60f563                 if (r1 >= 0) {
01144f                     r = r1;
c2bef1                     break;
cbb184                 }
cbb184             }
d2c0d8         res *= BASE;
f2637e         res += q;
d41d8c
e79d0e         if (j > 0) {
febb34             int d1 = res.a.size() + 2 < r.a.size() ? r.a[res.a.
    size() + 2] : 0;
baacce             int d2 = res.a.size() + 1 < r.a.size() ? r.a[res.a.
    size() + 1] : 0;
78b193             int d3 = res.a.size() < r.a.size() ? r.a[res.a.size()
    ] : 0;
7d925d             q = ((long long) d1 * BASE * BASE + (long long) d2 *
    BASE + d3) / (firstDigit * 2);
cbb184         }
cbb184     }
d41d8c
d7ee6d         res.trim();
28ae5c         return res / norm;
cbb184     }
2145c1 };
Full file hash: f1f35b

```

## 7.2 FFT

```

5d1131 #include "../contest/header.hpp"
d41d8c
d41d8c /*
18a91e     Modular Inverse:
f7b296     FFT allows multiplication of two polynomials in  $O(n \log n)$ .
420c7a     This can also be used to multiply two long numbers faster.
c00ff6     Other applications:
c35b73     - All possible sums of two arrays.
1da5a4     - Dot product of vector a with every cyclic shift of vector b.
49afe3     - Attaching two boolean stripes without two 1s next to each other
.
52f6a3     - String matching.

```

```

d41d8c
b95cae      Usage:
afb0f5      long double is a lot slower. 3s with ld and 0.7 with double for
10^6 size vectors.
d41d8c
1d1558      Source: https://cp-algorithms.com/algebra/fft.html
c4c9bd      */
d41d8c
f4f8e6      using cd = complex<long double>;
c4f8de      const ld PI = acos(-1.0L);
d41d8c
9b5b94      void fft(vector<cd> &a, bool invert)
f95b70      {
94d5f8          int n = a.size();
d41d8c
d94885          for (int i = 1, j = 0; i < n; i++)
f95b70          {
4af5d7              int bit = n >> 1;
474fac              for (; j & bit; bit >>= 1)
53c7ca                  j ^= bit;
53c7ca                  j ^= bit;
d41d8c
9dcc5c              if (i < j)
33275d                  swap(a[i], a[j]);
cbb184          }
d41d8c
2fe9ad          for (int len = 2; len <= n; len <<= 1)
f95b70          {
c19c97              ld ang = 2 * PI / len * (invert ? -1 : 1);
808a0b              cd wlen(cos(ang), sin(ang));
3dd9d3              for (int i = 0; i < n; i += len)
f95b70              {
8c3c80                  cd w(1);
5594fb                  for (int j = 0; j < len / 2; j++)
f95b70                  {
cf0824                      cd u = a[i + j], v = a[i + j + len / 2] * w;
6c3014                      a[i + j] = u + v;
273255                      a[i + j + len / 2] = u - v;
3e4104                      w *= wlen;
cbb184                  }
cbb184              }
cbb184          }
d41d8c
2111a0          if (invert)
f95b70          {
0b5665              for (cd &x : a)
b6d31b                  x /= n;
cbb184          }
cbb184      }
d41d8c

```

```

d41d8c // Input a[0] + a[1]x + a[2]x^2 ...
d41d8c // Returns polynomial of size equal to the smallest power of two at
least
d41d8c // as large as a.size() + b.size(). This can have some extra zeros.
d41d8c // Use long double if using long long.
4fce64 template <class T>
a3a2ed vector<T> multiply(vector<T> const &a, vector<T> const &b)
f95b70 {
6fa6b9     vector<cd> fa(a.begin(), a.end()), fb(b.begin(), b.end());
43ec81     int n = 1;
cd5a64     while (n < a.size() + b.size())
c149a4         n <= 1;
37aa6c     fa.resize(n);
870070     fb.resize(n);
d41d8c
3a13f2     fft(fa, false);
c76760     fft(fb, false);
83008c     for (int i = 0; i < n; i++)
940eb7         fa[i] *= fb[i];
959d01     fft(fa, true);
d41d8c
ebf3b6     vector<T> result(n);
83008c     for (int i = 0; i < n; i++)
4c9bb2         result[i] = round(fa[i].real()); // Remember to remove rounding
if working with floats.
dc8384     return result;
cbb184 }
Full file hash: 0402dc

```

## 7.3 Fraction

```

5d1131 #include "../contest/header.hpp"
d41d8c
d41d8c /*
390211     Fraction representation:
4d1181         All operations run in  $O(\gcd) = O(\log)$ .
d41d8c
b95cae     Usage:
70e7d7         Don't modify internal values, use constructor.
408ef0         Some nice things about the constructor: frac() = 0/1, frac(5) =
5/1.
d41d8c
b8d28c         Be careful that the numerator and denominator might overflow if
lcm is too big.
20fa30         In those cases, you can always do frac<big_int>, but that will be
painful to code.
d41d8c
3db72f     Author: Arthur Pratti Dadalto
c4c9bd */
d41d8c

```

```

4fce64 template <class T>
4cf1ca struct frac
f95b70 {
e75828     T a, b; // b can't be negative, very important.
d41d8c
191fc6     explicit frac(T a = 0, T b = 1) : a(a), b(b) { simpl(); }
d41d8c
7d70f7     void simpl()
f95b70     {
8eb5bb         T g = __gcd(abs(a), abs(b)) * sign(b); // Make b positive.
fe7245         a /= g;
ee2d42         b /= g;
cbb184     }
d41d8c
d59b8a     bool operator<(const frac &rhs) const
f95b70     {
5c6427         return a * rhs.b < rhs.a * b;
cbb184     }
d41d8c
7ebf19     bool operator>(const frac &rhs) const
f95b70     {
2ab79c         return rhs < *this;
cbb184     }
d41d8c
d60bf3     bool operator==(const frac &rhs) const // TODO: untested.
f95b70     {
77c0b8         return !(*this < rhs) && !(rhs < *this);
cbb184     }
d41d8c
473b74     frac operator*(const frac &rhs) const
f95b70     {
f0117d         return frac(a * rhs.a, b * rhs.b);
cbb184     }
d41d8c
04b5a1     frac operator+(const frac &rhs) const
f95b70     {
3ff11f         T m = (b * rhs.b) / __gcd(b, rhs.b);
24edd6         return frac(a * (m / b) + rhs.a * (m / rhs.b), m);
cbb184     }
d41d8c
c8ca1d     frac operator-(void) const
f95b70     {
132fb3         return frac(-a, b);
cbb184     }
d41d8c
de243f     frac operator-(const frac &rhs) const
f95b70     {
111760         return (*this) + (-rhs);
cbb184     }
d41d8c

```

```

d63a85     frac operator/(const frac &rhs) const
f95b70     {
f5299b         return (*this) * frac(rhs.b, rhs.a);
cbb184     }
d41d8c
9e018a     friend ostream &operator<<(ostream &os, const frac &f)
f95b70     {
891d94         return os << f.a << "/" << f.b;
cbb184     }
2145c1 };
d41d8c
Full file hash: c8862e

```

## 7.4 Integration

```

d41d8c  /*
f64ead    Numerical Integration:
49d5e8    Given a function f and an interval [a, b] estimates integral of f
          (x) dx from a to b.
bfe460    Error is in theory inversely proportional to n^4.
d41d8c
b95cae    Usage:
be1ead    n, the number of intervals must be even.
d41d8c
3db72f    Author: Arthur Pratti Dadalto
c4c9bd  */
d41d8c
044d82  template <class F>
7d9945  double simpsons(const F &f, int n /* even */, double a, double b)
f95b70  {
46af34      double retv = f(a) + f(b);
d025af      double h = (b - a) / n;
acfc81      for (int i = 1; i < n; i += 2)
900086          retv += 4 * f(a + i * h);
1c3900      for (int i = 2; i < n; i += 2)
6c1313          retv += 2 * f(a + i * h);
d41d8c
055fe5      retv *= h / 3;
6272cf      return retv;
cbb184  }
d41d8c
d41d8c  // Sample usage:
d41d8c  // int main(void)
d41d8c  // {
d41d8c  //     printf("%.20lf\n", simpsons([](double x) { return pow(sin(M_PI *
          x / 2.0), 3.2);}, 2000, 0, 2));
d41d8c  // }
Full file hash: caa0e5

```

## 7.5 linalg

```

5d1131 #include "../contest/header.hpp"
d41d8c
d41d8c /*
f92339     Vector and matrix operations:
687bbc     Details are given in each function.
3ab55f     vec inherits from vector<T>, so there is a lot you can do with it
        .
5ae524     Also, mat inherits from vector<vec<T>>.
d41d8c
3db72f     Author: Arthur Pratti Dadalto
d41d8c
1ef4c8     Source: some of it from https://github.com/kth-competitive-
        programming/kactl/blob/master/content/numerical/MatrixInverse.h
c4c9bd */
d41d8c
4fce64 template <class T>
fe4002 struct vec : vector<T>
f95b70 {
469362     vec(int n) : vector<T>(n) {}
d41d8c
d41d8c     // c = a*x + b*y
e918cb     static void linear_comb(const vec &a, T x, const vec &b, T y, vec &
        c)
f95b70     {
8fe753         for (int i = 0; i < sz(a); i++)
75e753             c[i] = a[i] * x + b[i] * y;
cbb184     }
d41d8c
d41d8c     // return a*x + b*y
250f88     static vec linear_comb(vec a, T x, const vec &b, T y)
f95b70     {
4fec85         linear_comb(a, x, b, y, a);
3f5343         return a;
cbb184     }
2145c1 };
d41d8c
4fce64 template <class T>
dade1f struct mat : vector<vec<T>>
f95b70 {
d41d8c     // Creates a zero-filled matrix of n rows and m columns.
2d2b5d     mat(int n, int m) : vector<vec<T>>(n, vec<T>(m)) {}
d41d8c
d41d8c     // c = a * x + b * y
762fbc     static void linear_comb(const mat &a, T x, const mat &b, T y, mat &
        c)
f95b70     {
8fe753         for (int i = 0; i < sz(a); i++)
f47ed7             for (int j = 0; j < sz(a[i]); j++)
4f844b             c[i][j] = a[i][j] * x + b[i][j] * y;
cbb184     }

```

```

d41d8c
d41d8c    // return a * x + b * y
08e6ea    static mat linear_comb(mat a, T x, const mat &b, T y)
f95b70    {
4fec85        linear_comb(a, x, b, y, a);
3f5343        return a;
cbb184    }
d41d8c
13fd2a    mat operator-(const mat &b) const { return linear_comb(*this, T(1),
    b, T(-1)); }
d41d8c
0138fa    mat operator+(const mat &b) const { return linear_comb(*this, T(1),
    b, T(1)); }
d41d8c
93d3e8    mat operator*(const T &x) { return linear_comb(*this, x, *this, T
    (0)); }
d41d8c
d41d8c    // Absolutely does not work for int.
72c1fd    mat operator/(const T &x) const { return linear_comb(*this, T(1) /
    x, *this, T(0)); }
d41d8c
d41d8c    // Returns inverse of matrix (assuming it is square and non-
    singular). Runs in O(n^3).
d41d8c    // Absolutely does not work for int.
14566d    mat inverse() // TODO: test singular.
f95b70    {
d23a72        int n = sz(*this);
bca455        mat a(n, 2 * n); // A is Nx2N: X|I.
f7f2d1        vector<int> col(n); // Will be using column pivoting, so need to
    remember original columns.
83008c        for (int i = 0; i < n; i++)
f95b70        {
f90a6b            for (int j = 0; j < n; j++)
c1c7c0                a[i][j] = (*this)[i][j];
34ac5b            a[i][i + n] = T(1);
6dcd38            col[i] = i;
cbb184        }
d41d8c
83008c        for (int i = 0; i < n; i++)
f95b70        {
903ccf            int r = i, c = i;
775cab            for (int j = i; j < n; j++)
90f1d8                for (int k = i; k < n; k++)
f78c7f                    if (abs(a[j][k]) > abs(a[r][c]))
d4c894                        r = j, c = k;
d41d8c
d41d8c            // assert(abs(a[r][c]) > EPS); Uncomment to check singular
    matrix
a2fa24            swap(a[i], a[r]);
d41d8c

```

```

f90a6b         for (int j = 0; j < n; j++)
c8cc8f             swap(a[j][i], a[j][c]), swap(a[j][i + n], a[j][c + n]);
c1d48e         swap(col[i], col[c]);
d41d8c
b70d15         vec<T>::linear_comb(a[i], T(1) / a[i][i], a[i], T(0), a[i]);
67830d         a[i][i] = T(1);
d41d8c
197ab1         for (int j = i + 1; j < n; j++)
3704dc             vec<T>::linear_comb(a[j], T(1), a[i], -a[j][i], a[j]);
cbb184     }
d41d8c
d41d8c     // Right now A is:
d41d8c     //
d41d8c     //  1 * *
d41d8c     //  0 1 *
d41d8c     //  0 0 1
d41d8c     //
d41d8c     // Next we remove non-1s from right to left.
d41d8c
917d8b         for (int i = n - 1; i > 0; i--)
c791cd             for (int j = 0; j < i; j++)
3704dc                 vec<T>::linear_comb(a[j], T(1), a[i], -a[j][i], a[j]);
d41d8c
c70ad2         mat retv(n, n);
83008c         for (int i = 0; i < n; i++)
f90a6b             for (int j = 0; j < n; j++)
4eb40a                 retv[col[i]][col[j]] = a[i][j + n];
6272cf         return retv;
cbb184     }
2145c1 };
Full file hash: 2457f9

```

## 7.6 Simplex

```

5d1131 #include "../contest/header.hpp"
d41d8c
d41d8c /*
458b90     Simplex:
6956ec         Optimizes a linear program of the form:
15b127             maximize c*x, s.t. a*x <ops> b, x >= 0.
7b88d6             Each constraint can use a different operator from {<= >= ==}.
8aa76d             Not polynomial, but got AC 150 ms with 4000 constraints and 200
                variables.
d41d8c
b95cae     Usage:
e8b3b7         Call run_simplex, with the number of constraints and variables, a
                , b, ops and c (as specified above).
34036d         Return value is ok if solution was found, unbounded if objective
                value can be infinitely large
eb42f2         or infeasible if there is no solution given the constraints.

```



```

d41d8c
2baa60      The value of each variable is returned in vector res. Objective
            function optimal value is also returned.
060dc4      Sample usage is commented below.
d41d8c
3db72f      Author: Arthur Pratti Dadalto
c4c9bd      */
d41d8c
4fce64      template <class T>
fe4002      struct vec : vector<T>
f95b70      {
469362          vec(int n) : vector<T>(n) {}
d41d8c
d41d8c          // c = a*x + b*y
e918cb      static void linear_comb(const vec &a, T x, const vec &b, T y, vec &
            c)
f95b70      {
8fe753          for (int i = 0; i < sz(a); i++)
75e753              c[i] = a[i] * x + b[i] * y;
cbb184      }
2145c1      };
d41d8c
4fce64      template <class T>
dade1f      struct mat : vector<vec<T>>
f95b70      {
d41d8c          // Creates a zero-filled matrix of n rows and m columns.
2d2b5d          mat(int n, int m) : vector<vec<T>>(n, vec<T>(m)) {}
d41d8c
d41d8c          // Erase row 0(n^2).
82436c          void erase_row(int i)
f95b70          {
7c9f9f              this->erase(this->begin() + i);
cbb184          }
d41d8c
d41d8c          // Erase column 0(n^2).
1b22c6          void erase_col(int j)
f95b70          {
798fc8              for (int i = 0; i < sz(*this); i++)
a7796a                  (*this)[i].erase((*this)[i].begin() + j);
cbb184          }
2145c1      };
d41d8c
d3ff82      namespace simplex
f95b70      {
d41d8c          // Any value within [-EPS, +EPS] will be considered equal to 0.
05667a          const double EPS = 1e-6;
d41d8c
5e6f5b          enum op { ge, le, eq };
d41d8c
242dbb          enum optimization_status { ok, unbouded, infeasible };

```

```

d41d8c
4d9580 int get_entering_var(mat<double> &tab)
f95b70 {
d41d8c     // Get first non-artificial variable with negative objective
        coefficient. If none, return -1. (could instead return most negative, but
        that could cycle)
682f62     for (int i = 0; i < sz(tab[0]) - 1; i++)
72e0d2         if (tab[0][i] < -EPS)
d9a594             return i;
daa4d1     return -1;
cbb184 }
d41d8c
201003 int get_exiting_var_row(mat<double> &tab, int entering_var)
f95b70 {
d41d8c     // Get smallest value of val and first in case of tie. If none,
        return -1.
fcb2fc     int retv = -1;
6213b9     double val = -1.0;
a07064     for (int i = 1; i < sz(tab); i++)
f95b70     {
d41d8c         // If strictly positive, it bounds the entering var.
dcda72         if (tab[i][entering_var] > EPS)
f95b70         {
d41d8c             // Entering var will be bounded by tab[i][tab.size().second -
                1] / tab[i][entering_var].
d41d8c             // val could be slightly negative if tab[i][tab.size().second -
                1] = -0.
393d3f             if (val == -1.0 || tab[i][sz(tab[i]) - 1] / tab[i][entering_var
                ] < val)
f95b70                 {
78d87c                     val = tab[i][sz(tab[i]) - 1] / tab[i][entering_var];
52cece                     retv = i;
cbb184                 }
cbb184             }
cbb184         }
d41d8c
6272cf     return retv;
cbb184 }
d41d8c
ed25d2 optimization_status solve_tab(mat<double> &tab, vector<int> &
        basic_var)
f95b70 {
d41d8c     // artificial_count is the number of variables at the end we should
        ignore.
a17ec7     int entering_var;
6b7846     while ((entering_var = get_entering_var(tab)) != -1)
f95b70     {
6c0a23         int exiting_var_row = get_exiting_var_row(tab, entering_var);
d41d8c
d41d8c         // If no exiting variable bounds the entering variable, the

```

```

    objective is unbounded.
813335     if (exiting_var_row == -1)
914a2e         return optimization_status::unbounded;
d41d8c
d41d8c         // Set new basic var coefficient to 1.
89c7a2         vec<double>::linear_comb(tab[exiting_var_row], (1.0 / tab[
    exiting_var_row][entering_var]), tab[exiting_var_row], 0.0, tab[
    exiting_var_row]);
d41d8c
d41d8c         // Gaussian elimination of the other rows.
c7a773         for (int i = 0; i < sz(tab); i++)
81c379             if (i != exiting_var_row)
ed2730                 if (abs(tab[i][entering_var]) > EPS)
7ad878                     vec<double>::linear_comb(tab[i], 1.0, tab[exiting_var_row],
    -tab[i][entering_var], tab[i]);
d41d8c
64dd6a         basic_var[exiting_var_row] = entering_var;
cbb184     }
d41d8c
c52f1c     return optimization_status::ok;
cbb184 }
d41d8c
d41d8c // maximize c*x, s.t. a*x <ops> b. x >= 0.
f1a105 optimization_status run_simplex(int num_constraints, int num_vars,
    mat<double> a, vec<op> ops, vec<double> b, vec<double> c, vec<double> &res
    , double &obj_val)
f95b70 {
334f46     for (int i = 0; i < num_constraints; i++)
5f946c         if (ops[i] == op::ge)
f95b70             {
d41d8c                 // Beyond this point "ge" constraints won't exist.
44438f                 vec<double>::linear_comb(a[i], -1, a[i], 0, a[i]); // a[i] *=
    -1;
250b4d                 b[i] *= -1;
1c38d4                 ops[i] = op::le;
cbb184             }
d41d8c
0264da     int num_artificial_variables = 0;
371f2b     int num_slack_variables = 0;
334f46     for (int i = 0; i < num_constraints; i++)
f95b70     {
0ec40f         if (ops[i] == op::le)
f95b70             {
37acf9             num_slack_variables++;
cbb184         }
d41d8c
359aa4         if ((ops[i] == op::le && b[i] < -EPS) || ops[i] == op::eq)
f95b70         {
d41d8c             // If we have rhs strictly negative in a inequality or an
    equality constraint, we need an artificial val.

```

```

fc36e6      num_artificial_variables++;
cbb184      }
cbb184      }
d41d8c
854c33      mat<double> tab(num_constraints + 1, num_vars + num_slack_variables
+ num_artificial_variables + 1);
9a9a70      vector<int> basic_var(num_constraints + 1);
775265      vector<int> slack_cols, artificial_cols;
7f63aa      for (int i = num_vars; i < num_vars + num_slack_variables; i++)
10c71f          slack_cols.push_back(i);
e0b615      for (int i = num_vars + num_slack_variables; i < num_vars +
num_slack_variables + num_artificial_variables; i++)
eafbfb          artificial_cols.push_back(i);
c70a50      int rhs_col = num_vars + num_slack_variables +
num_artificial_variables;
d41d8c
d41d8c      // First objective will be to have artificial variables equal to 0.
017565      for (int i : artificial_cols)
b98201          tab[0][i] = 1;
d41d8c
9c49f5      for (int i = 0, k = 0, l = 0; i < num_constraints; i++)
f95b70      {
861a15          for (int j = 0; j < num_vars; j++)
e3832e              tab[i + 1][j] = a[i][j];
d41d8c
0ec40f          if (ops[i] == op::le)
141495              tab[i + 1][slack_cols[l++]] = 1;
d41d8c
142f37          tab[i + 1][rhs_col] = b[i];
d41d8c
359aa4          if ((ops[i] == op::le && b[i] < -EPS) || ops[i] == op::eq) //
Basic var will be artificial
f95b70          {
2a6978              if (b[i] < -EPS)
009fda                  vec<double>::linear_comb(tab[i + 1], -1, tab[i + 1], 0, tab[i
+ 1]); // a[i] *= -1;
d41d8c
86fab4              tab[i + 1][artificial_cols[k++]] = 1;
116454              basic_var[i + 1] = artificial_cols[k - 1];
d41d8c
06db08              vec<double>::linear_comb(tab[0], 1.0, tab[i + 1], -1.0, tab[0])
;
cbb184          }
2954e9          else // Basic var will be slack var.
f95b70          {
ae77b6              basic_var[i + 1] = slack_cols[l - 1];
cbb184          }
cbb184      }
d41d8c
df8d17      assert(solve_tab(tab, basic_var) == optimization_status::ok);

```

```

d41d8c
d41d8c    // Best solution could not bring artificial variables to 0 (
objective max Z = sum(-xa)).
fe0d64    if (tab[0][sz(tab[0]) - 1] < -EPS)
94b8a3        return optimization_status::infeasible;
d41d8c
d41d8c    // If we have an artificial variable on the base with xb = 0, we
need to remove it.
e6411b    for (int i = 1; i < sz(basic_var); i++)
0778cb        if (basic_var[i] >= num_vars + num_slack_variables)
f95b70        {
d41d8c            // Find non-artificial replacement.
e2f213            for (int j = 0; j < sz(tab[i]) - 1 - num_artificial_variables;
j++)
f95b70            {
d41d8c                // If non-zero value in row, we can replace.
a8880b                if (j != basic_var[i] && abs(tab[i][j]) > EPS)
f95b70                {
d41d8c                    // Remove from the other rows.
b5fa44                    vec<double>::linear_comb(tab[i], 1.0 / tab[i][j], tab[i],
0, tab[i]);
d41d8c
443db5                    for (int k = 0; k < sz(tab); k++)
635b4c                        if (k != i)
f95b70                            {
e76184                                if (abs(tab[k][j]) > EPS)
4b6b27                                    vec<double>::linear_comb(tab[k], 1, tab[i], -tab[k][j
], tab[k]);
cbb184                            }
d41d8c
d41d8c                // Basic variable replacement done, so proceed to next
basic_var.
7e0f27                basic_var[i] = j;
c2bef1                break;
cbb184            }
cbb184        }
cbb184    }
d41d8c
ca2210    for (int i = sz(tab) - 1; i > 0; i--)
0778cb        if (basic_var[i] >= num_vars + num_slack_variables)
f95b70        {
d41d8c            // Could not replace basic var, so constraint is redundant.
2cd1fb            tab.erase_row(i);
fe14c7            basic_var.erase(basic_var.begin() + i);
cbb184        }
d41d8c
d41d8c    // Remove artificial variable columns.
5c3178    for (int i = sz(artificial_cols) - 1; i >= 0; i--)
9a226e        tab.erase_col(artificial_cols[i]);
d41d8c

```

```

1311b7     for (int i = 0; i < sz(tab[0]); i++)
d2677f         tab[0][i] = 0;
f17293     for (int i = 0; i < num_vars; i++)
94256d         tab[0][i] = -c[i];
d41d8c
a07064     for (int i = 1; i < sz(tab); i++)
b39526         vec<double>::linear_comb(tab[0], 1, tab[i], -tab[0][basic_var[i
    ]], tab[0]);
d41d8c
54ad02     optimization_status status = solve_tab(tab, basic_var);
d41d8c
b68670     res = vec<double>(num_vars);
e6411b     for (int i = 1; i < sz(basic_var); i++)
047b20         if (basic_var[i] < num_vars)
81f54e             res[basic_var[i]] = tab[i][sz(tab[i]) - 1];
d41d8c
a3473e     obj_val = tab[0][sz(tab[0]) - 1];
d41d8c
62d3d5     return status;
cbb184 }
cbb184 } // namespace simplex
d41d8c
d41d8c /*
13a4b1 int main(void)
f95b70 {
14e0a7     int n, m;
aa3380     cin >> n >> m;
d41d8c
37ce14     int num_constraints = m, num_vars = n;
d41d8c
d41d8c     // maximize c*x, s.t. a*x <ops> b. x >= 0.
2626bb     mat<double> a(num_constraints, num_vars);
84d434     vec<double> b(num_constraints);
01b2af     vec<simplex::op> ops(num_constraints);
dabb12     vec<double> c(num_vars);
40ca17     vec<double> res(num_vars);
d41d8c
83008c     for (int i = 0; i < n; i++)
a733f7         cin >> c[i];
d41d8c
94f72b     for (int i = 0; i < m; i++)
f95b70     {
7ba74c         int l, r, x;
15994b         cin >> l >> r >> x;
0dfebd         for (int j = l - 1; j <= r - 1; j++)
a21125             a[i][j] = 1;
df0b9d         b[i] = x;
80367f         ops[i] = simplex::op::le;
cbb184     }
d41d8c

```

```
1afc12    double ans;  
dd6c28    simplex::run_simplex(num_constraints, num_vars, a, ops, b, c, res,  
          ans);  
d41d8c  
530b75    cout << ((long long)(ans + 0.5)) << endl;  
cbb184    }  
c4c9bd    */  
Full file hash: 46f321
```

## 8 String

### 8.1 KMP

```

5d1131 #include "../contest/header.hpp"
d41d8c
d41d8c /*
8dec4f     Prefix Function and KMP:
e45403     Computes prefix function for a given string in O(n).
16bb22     String matching in O(n + m).
37f784     No need to be strings, you can use vector<int> since the
be2fe6     algorithms don't depend on the alphabet size, they only perform
equality comparisons.
b5efd9     Usage is explained in each function.
d41d8c
3db72f     Author: Arthur Pratti Dadalto
c4c9bd */
d41d8c
d41d8c // Returns the prefix function for the given string.
d41d8c // pi[i] for 0 <= i <= s.size() (s.size() + 1 elements).
d41d8c // pi[i] considers the prefix of string s having size i.
d41d8c // pi[i] is the size of its (the prefix's) largest proper prefix
which is also a suffix.
d41d8c // For "aabaaab", pi is is {0,0,1,0,1,2,2,3}
4fce64 template <class T>
8fa849 vector<int> prefix_function(T s)
f95b70 {
d2c5d5     vector<int> pi(s.size() + 1, 0);
a94e4a     for (int i = 2; i <= s.size(); i++)
f95b70     {
3f878c         int j = pi[i - 1];           // j is the size of the candidate
prefix to expand.
4b3f35         while (j > 0 && s[j] != s[i - 1]) // While we still have a
candidate prefix and it can't be expanded.
187475             j = pi[j];                 // Go to the next candidate prefix.
d41d8c
d41d8c         // If candidate prefix can be expanded, do it. Otherwise, there
is no prefix that is also a suffix.
f986f8         pi[i] = s[j] == s[i - 1] ? j + 1 : 0;
cbb184     }
d41d8c
81d1a2     return pi;
cbb184 }
d41d8c
d41d8c // Returns a sorted list of all positions in the text string where
begins an occurrence of the key string.
d41d8c // e.g. kmp("aabaaab", "aab") returns {0, 4}.
4fce64 template <class T>
15b377 vector<int> kmp(T text, T key)
f95b70 {

```



```

aeb888     vector<int> retv;
7fa638     vector<int> pi = prefix_function(key);
5d936d     for (int i = 0, match = 0; i < text.size(); i++) // There is no
            need to have the entire text in memory, you could do this char by char.
f95b70     {
d41d8c         // match stores the size of the prefix of the key which is a
            suffix of the current processed text.
9d984d         while (match > 0 && text[i] != key[match])
7eb4cc             match = pi[match];
db8319         if (text[i] == key[match])
24b638             match++;
d41d8c
dd8c14         if (match == key.size())
f95b70             {
7b8421                 retv.push_back(i - match + 1);
7eb4cc                 match = pi[match]; // To avoid access to key[key.size()] in
            next iteration.
cbb184             }
cbb184         }
d41d8c
6272cf     return retv;
cbb184 }
Full file hash: 415801

```

## 8.2 Aho Corasick

```

5d1131     #include "../contest/header.hpp"
d41d8c
d41d8c     /*
30562e     Aho-Corasick: O(alpha_size * string_sum)
4e9057     In general, multiple pattern string matching tree/automaton.
d41d8c
fbc6b5     Keep in mind that find_all can be O(N*sqrt(N)) if no duplicate
            patterns. (N is total string length)
d41d8c
ca2095     Constraints:
00d37a     chars in the string are all in the interval [first, first +
            alpha_size - 1].
3da079     This will not free some memory on object destruction.
390590     Duplicate patterns are allowed, empty patterns are not.
d41d8c
b95cae     Usage:
df3a72     Set alpha_size and the first char in the alphabet.
e98cb2     Call constructor passing the list of pattern strings.
0a657b     Use one of find, find_all ... to process a text or do your own
            thing.
acdd39     To find the longest words that start at each position, reverse
            all input.
3439d1     Bottleneck in this code is memory allocation.
91a84c     For 10^6 total string size, memory usage can be up to 300 Mb.

```

```

d41d8c
b34145     You can save time:
93df09     list_node, match_list, match_list_last are only needed to list
all matches.
57e4db     atm automaton table can be cut to reduce memory usage.
018d49     The text processing stuff is also optional.
02e3ad     Node memory can be one big array instead of vector.
d41d8c
3db72f     Author: Arthur Pratti Dadalto
c4c9bd    */
d41d8c
e7f92a    struct aho_corasick
f95b70    {
da45ec        enum
f95b70        {
033315            alpha_size = 26, // Number of chars in the alphabet.
b3d02f            first = 'a'      // First char.
2145c1        };
d41d8c
fc487b    struct list_node // Simple linked list node struct.
f95b70    {
53e65f        int id;
6ec94b        list_node *next;
ff56a7        explicit list_node(int id, list_node *next) : id(id), next(next)
        {}
2145c1    };
d41d8c
e4accb    struct node
f95b70    {
ca8b7e        int fail = -1; // node failure link (aka suffix link).
2eb620        int nmatches = 0; // Number of matches ending in this node.
9005b9        int next[alpha_size]; // Next node in trie for each letter.
Replace with unordered_map or list if memory is tight.
c0f747        int atm[alpha_size]; // Optional: Automaton state transition
table. Simpler text processing.
d41d8c
44edb6        list_node *match_list = nullptr; // Pointer to first node in
linked list of matches. List ends with null pointer.
01009c        list_node *match_list_last = nullptr; // Internal: pointer to
last node in list of matches (before bfs), or null if empty list.
d41d8c
e6fb82        node() { memset(next, -1, sizeof(next)); } // Start with all
invalid transitions.
2145c1    };
d41d8c
b9ea22        vector<node> nodes;
d41d8c
9b61f6        aho_corasick(const vector<string> &pats)
f95b70        {
225eb3            nodes.emplace_back(); // Make root node 0.

```

```

b5bf96     for (int i = 0; i < sz(pats); i++)
f95b70     {
b3da3c         int cur = 0;  // Start from root.
9f5c69         for (int j = 0; j < sz(pats[i]); j++)
f95b70         {
ec0388             int k = pats[i][j] - first;
d41d8c
10937b             if (nodes[cur].next[k] <= 0)  // Make new node if needed.
f95b70             {
976fa3                 nodes[cur].next[k] = sz(nodes);
225eb3                 nodes.emplace_back();
cbb184             }
d41d8c
47b49f             cur = nodes[cur].next[k];
cbb184         }
d41d8c
d41d8c         // Add logic here if additional data is needed on matched
        strings.
4daeea         nodes[cur].nmatches++;
45f177         nodes[cur].match_list = new list_node(i, nodes[cur].match_list)
        ; // Add string to node list of matches.
fe38fe         if (nodes[cur].nmatches == 1)
947da5             nodes[cur].match_list_last = nodes[cur].match_list;
cbb184     }
d41d8c
26a528     queue<int> q;
6733a6     for (int i = 0; i < alpha_size; i++) // Define fail for first
        level.
f95b70     {
e8dc83         if (nodes[0].next[i] == -1) // Invalid transitions from 0 now
        become valid self transitions.
fb628f             nodes[0].next[i] = 0;
d41d8c
7d3171         nodes[0].atm[i] = nodes[0].next[i]; // Automaton state
        transition table.
d41d8c
bc34bf         if (nodes[0].next[i] > 0) // Single letter nodes have fail = 0
        and go in the queue.
f95b70         {
eded92             q.push(nodes[0].next[i]);
9b22e6             nodes[nodes[0].next[i]].fail = 0;
cbb184         }
cbb184     }
d41d8c
ee6bdd     while (!q.empty()) // Use bfs to compute fail for next level.
f95b70     {
69faa7         int cur = q.front();
833270         q.pop();
d41d8c
6733a6         for (int i = 0; i < alpha_size; i++)

```

```

af4a6e         if (nodes[cur].next[i] > 0) // Don't use -1 and don't use
               transition to root.
f95b70         {
3ecdd3         nodes[cur].atm[i] = nodes[cur].next[i]; // Unrelated to
               code below, filling automaton.
d41d8c
d41d8c         // Computing fail for next node and putting it in the queue
               .
3ae7da         int prox = nodes[cur].next[i];
53ef92         q.push(prox);
d41d8c
f252cb         int state = nodes[cur].fail;
c66324         while (nodes[state].next[i] == -1)
d712e2             state = nodes[state].fail;
d41d8c
7836db         nodes[prox].fail = nodes[state].next[i];
d41d8c
d41d8c         // Add logic here if additional data is needed on matched
               strings.
2940ed         nodes[prox].nmatches += nodes[nodes[prox].fail].nmatches;
d41d8c
d41d8c         // Add in O(1) list from fail link to next node's list.
               Operation: a->b->null c->null to a->b->c->null.
59ed4d         (nodes[prox].match_list_last ? nodes[prox].match_list_last
               ->next : nodes[prox].match_list) = nodes[nodes[prox].fail].match_list;
cbb184         }
2954e9         else
f95b70         {
a04598             nodes[cur].atm[i] = nodes[nodes[cur].fail].atm[i];
cbb184         }
cbb184     }
cbb184 }
d41d8c
d41d8c // Optional
d41d8c // Returns a vector retv such that, for each text position i:
d41d8c // retv[i] is the index of the largest pattern ending at position i
               in the text.
d41d8c // If retv[i] == -1, no pattern ends at position i.
32246d vector<int> find(const string &text)
f95b70 {
107323     vector<int> retv(sz(text));
b3da3c     int cur = 0;
d41d8c
77447e     for (int i = 0; i < sz(text); i++)
f95b70     {
13dae2         cur = nodes[cur].atm[text[i] - first];
29e58f         retv[i] = (nodes[cur].match_list ? nodes[cur].match_list->id :
               -1);
cbb184     }
d41d8c

```

```

6272cf         return retv;
cbb184     }
d41d8c
d41d8c     // Optional
d41d8c     // Returns a vector retv such that, for each text position i:
d41d8c     // retv[i] is the number of pattern matches ending at position i in
the text.
48d0f2     vector<int> count(const string &text)
f95b70     {
107323         vector<int> retv(sz(text));
b3da3c         int cur = 0;
d41d8c
77447e         for (int i = 0; i < sz(text); i++)
f95b70         {
13dae2             cur = nodes[cur].atm[text[i] - first];
1a43d3             retv[i] = nodes[cur].nmatches;
cbb184         }
d41d8c
6272cf         return retv;
cbb184     }
d41d8c
d41d8c     // Optional
d41d8c     // Returns a vector retv such that, for each text position i:
d41d8c     // retv[i] is a list of indexes to the patterns ending at position
i in the text.
d41d8c     // These lists will be sorted from largest to smallest pattern
length.
d41d8c     // Keep in mind that find_all can be O(N*sqrt(N)) if no duplicate
patterns. (N is total string length)
4e5a4c     vector<vector<int>> find_all(const string &text)
f95b70     {
77b54a         vector<vector<int>> retv(sz(text));
b3da3c         int cur = 0;
d41d8c
77447e         for (int i = 0; i < sz(text); i++)
f95b70         {
13dae2             cur = nodes[cur].atm[text[i] - first];
d82b0e             for (auto n = nodes[cur].match_list; n != nullptr; n = n->next)
4c4784                 retv[i].push_back(n->id);
cbb184         }
d41d8c
6272cf         return retv;
cbb184     }
d41d8c
d41d8c     // Optional
d41d8c     // Returns a vector retv such that:
d41d8c     // retv is a list of indexes to the patterns ending at position pos
in the text.
d41d8c     // This list will be sorted from largest to smallest pattern length
.

```

```

251c66     vector<int> find_all_at_pos(const string &text, int pos)
f95b70     {
aeb888         vector<int> retv;
b3da3c         int cur = 0;
d41d8c
77447e         for (int i = 0; i < sz(text); i++)
f95b70         {
13dae2             cur = nodes[cur].atm[text[i] - first];
d41d8c
c57c6f             if (i == pos)
d82b0e                 for (auto n = nodes[cur].match_list; n != nullptr; n = n->
next)
1ad617                     retv.push_back(n->id);
cbb184         }
d41d8c
6272cf         return retv;
cbb184     }
2145c1 };
Full file hash: 2ec64b

```

### 8.3 Suffix Array

```

d41d8c
5d1131 #include "../contest/header.hpp"
d41d8c
d41d8c /*
1f77e9     Suffix array:
be00ca         Build suffix array and LCP array in  $O((n + \text{lim}) \log n)$  using  $O(n + \text{lim})$  memory, where lim is the alphabet size.
d41d8c
643e15         sa[i] is the starting index of the suffix which is i-th in the
sorted suffix array.
0d5e62         The returned vector is of size s.size()+1, and sa[0] == s.size().
The '\0' char at the end is considered
29ea73         part of the string, so sa[0] = "\0", the prefix starting at index
s.size().
d41d8c
ee7035         The lcp array contains longest common prefixes for neighbouring
strings
317e94         in the suffix array: lcp[i] = lcp(sa[i], sa[i-1]), lcp[0] = 0.
d41d8c
81eeab     Example:
981b73         Computing the LCP and the SA of "GATAGACA"
d33e7b         i sa[i] lcp[i] suffix
fd774f         0 8 0 ""
cac682         1 7 0 "A"
430b3a         2 5 1 "ACA"
d30cc0         3 3 1 "AGACA"
c895f5         4 1 1 "ATAGACA"
1a04b3         5 6 0 "CA"

```

```

b1b780      6 4 0 "GACA"
2999cd      7 0 2 "GATAGACA"
08e6dc      8 2 0 "TAGACA"
d41d8c
b95cae      Usage:
1c63e0      Important: the input string must not contain any zero values.
            Must use C++11 or above.
b847d4      You can use this for strings of integers, just change the
            alphabet size.
d41d8c
1d1558      Source: https://github.com/kth-competitive-programming/kactl/blob/
            master/content/strings/SuffixTree.h
c4c9bd      */
d41d8c
15a9b6      struct suffix_array
f95b70      {
71675a      vector<int> sa, lcp;
092958      suffix_array(const string &s, int lim = 256) // or basic_string<int
            > for integer strings.
f95b70      {
e72340      int n = sz(s) + 1, k = 0, a, b;
f6a0db      vector<int> x(s.begin(), s.end() + 1), y(n), ws(max(n, lim)),
            rank(n);
85469f      sa = lcp = y;
eb75f9      iota(sa.begin(), sa.end(), 0);
7707f7      for (int j = 0, p = 0; p < n; j = max(1, j * 2), lim = p)
f95b70      {
8dff9b      p = j;
00aec0      iota(y.begin(), y.end(), n - j);
83008c      for (int i = 0; i < n; i++)
e9b19c      if (sa[i] >= j)
d0873d      y[p++] = sa[i] - j;
450a8a      fill(ws.begin(), ws.end(), 0);
83008c      for (int i = 0; i < n; i++)
799bb0      ws[x[i]]++;
7d6bd3      for (int i = 1; i < lim; i++)
f256af      ws[i] += ws[i - 1];
5df399      for (int i = n; i--;)
d01b67      sa[--ws[x[y[i]]]] = y[i];
9dd20c      swap(x, y);
017be6      p = 1;
16ab1b      x[sa[0]] = 0;
d41d8c
aa4866      for (int i = 1; i < n; i++)
f95b70      {
fcb940      a = sa[i - 1];
2d820b      b = sa[i];
0cc036      x[b] = (y[a] == y[b] && y[a + j] == y[b + j]) ? p - 1 : p++;
cbb184      }
cbb184      }

```

```
d41d8c
aa4866     for (int i = 1; i < n; i++)
2f33c5         rank[sa[i]] = i;
d41d8c
05cb2b     for (int i = 0, j; i < n - 1; lcp[rank[i++]] = k)
487069         for (k && k--, j = sa[rank[i] - 1]; s[i + k] == s[j + k]; k++)
9eecb7             ;
cbb184     }
2145c1 };
Full file hash: 87092f
```