

Question 1

Suppose we have a set of blocks encrypted with the RSA algorithm and we don't have the private key. The public key is $\{n, e\}$. You are an attacker and are trying to break the private key. Suppose someone tells you they know one of the plaintext blocks has a common factor with n . Does this help you in any way to determine the private key? If yes, describe how.

SOLUTION:

In RSA, the form of the public and the private key is as follows:

Private key - $\{n, d\}$

Public key - $\{n, e\}$

We know the public key value i.e. (n, e) .

Thus we can get one component of the private key since we know that ' n ' is common in both the keys.

We just need to calculate the value ' d ' to find the private key.

We know that, in RSA:

$$d = (e^{-1}) \bmod (\text{PHI})$$

$$\text{PHI} = (p-1) \times (q-1)$$

[where p and q are prime numbers].

Yes! If someone tells me that one of the plaintext box has a common factor with ' n ', it does help me to determine the private key.

The steps I would use to determine the private key are:

1. We are told that one of the plain text box has a common factor with n .
2. This will help us in factoring n .
3. We can factor n to find ' p ' and ' q '.
4. We can use ' p ' and ' q ' to calculate ' PHI ' as $\text{PHI} = (p-1) \times (q-1)$.
5. Now using Euclid's algorithm, we can compute $d = (e^{-1}) \bmod (\text{PHI})$.
6. Thus, we can find ' d ' and determine the private key.

It is important to note that 'step 3' which is the factoring of ' n ' is a non trivial step. However, If a plaintext block has a common factor with n modulo n then the encoded block will also have a common factor with n modulo n . Because we encode blocks, which are smaller than ' p & q ', the factor must be p or q and the plaintext block must be a multiple of p or q . We can test each block for primality. If prime, it is p or q . In this case we divide into n to find the other factor. If not prime, we factor it and try the factors as divisors of n .

Knowing that the plaintext box has a common factor with ' n ', it helps in factorization and thus reduces the effort required to perform cryptanalysis to determine the private key.

NOTE:

Reference to PPT - Computer Security 7 : Cryptography — Public Key Version: 2012/02/15 16:15:24

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Link - <http://www.cs.arizona.edu/~collberg/Teaching/466-566/2012/Handouts/Handout-7.pdf>

Reference - Cryptography and Network Security: Principles and Practice - William Stallings

Chapter 9 - Public Key Cryptography and RSA