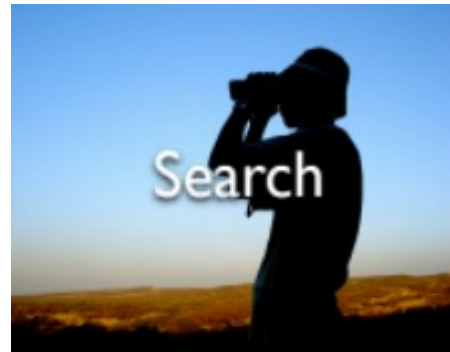


Topics for this Lecture

- Manually Test generation
- Automatic Test Generation Approaches
 - ~~Random Testing~~ (Guess)
 - ~~Search Based Software Testing~~ (Search)
 - Constraint-Based (Symbolic execution) Testing \forall (Deduce)



1 2 3 4 5 6 7 8 9 10



Constraint-Based Testing (CBT)

- **Constraint-Based Testing (CBT)** is the process of generating test cases against a testing objective by using constraint solving techniques.
- **Test objective:** generating a test case that covers a given testing criterion (such as statement, branch, batch, etc.)
- Constraints on inputs: if inputs satisfy constraints, then testing objective will be satisfied
- **CBT** improves significantly code-coverage (as constraints capture hard-to-reach test objectives).
- **CBT** is fully automated test data generation methods.
-

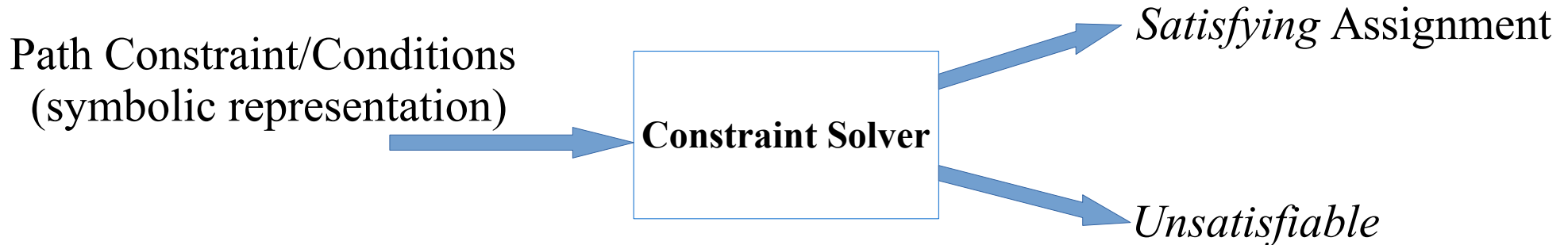


Symbolic Execution

- **Symbolic execution** is a program analysis technique that analyzes a program's code to automatically generate test data for the program.
- **Symbolic execution** analysis program's code with unspecified inputs.
- **Symbolic execution** uses symbolic values, instead of concrete values, as program inputs, and represents the values of program variables as symbolic expressions of those inputs.
- For each path, Symbolic Execution builds a path constraint/path condition.
- The **path constraint/path condition** (PC) is a Boolean formula over the symbolic inputs, which is an accumulation of the constraints that the inputs must satisfy for an execution to follow that path.



Constraint Solvers

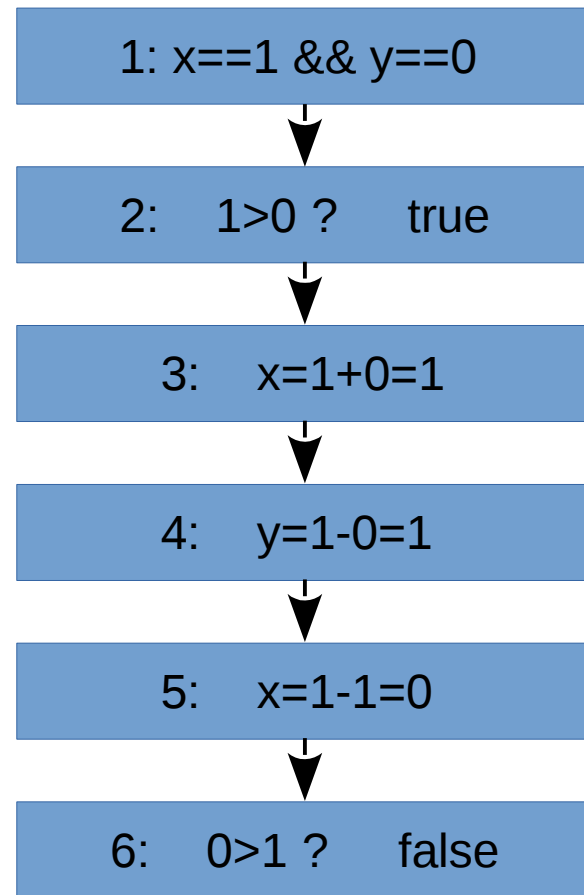


- A path condition is *satisfiable* (i.e., **feasible**) if there is a way to assign values to variables and make the **path condition** (i.e., boolean formula) true.
e.g., $(x > 0 \ \&\& \ x < 20 \ \&\& \ x == y + y)$ is satisfied by $(x:10, y:5)$
- A path condition is *unsatisfiable* (i.e., **infeasible**) if every assignment of values to variables makes the formula false, which means that there is no program input for which the program will take that (infeasible) path.
e.g., $(x > 0 \ \&\& \ x < 20 \ \&\& \ x == y + y \ \&\& \ x \% 2 == 1)$ is unsatisfiable.

Standard Execution

```
int x,y;  
1:  if (x>y) {  
2:      x = x+y;  
3:      y = x-y;  
4:      x=x-y;  
5:      if (x>y)  
6:          assert false;  
7:  }  
8:  print x,y
```

Code that swaps two integers
variables x and y



Symbolic Execution

```

int x, y;
1 if(x > y){
2   x = x+y;
3   y = x-y;
4   x = x-y;
5   if(x - y > 0)
6     assert false;
7 }
8 print(x, y)
    
```

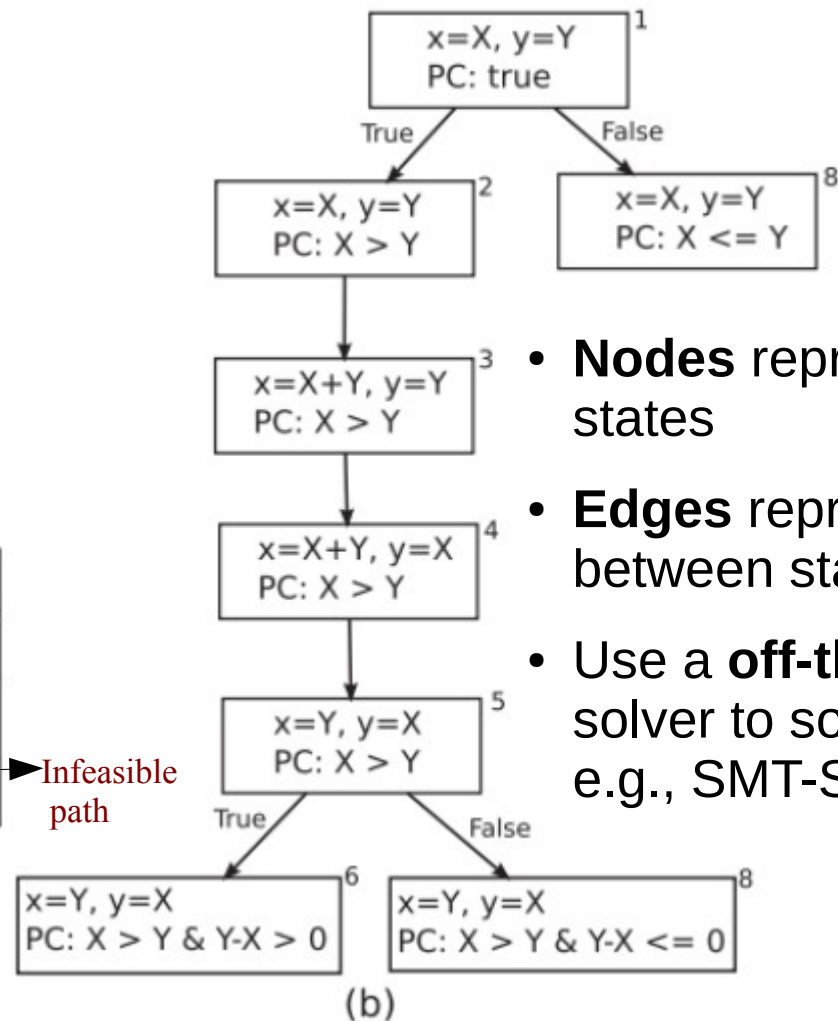
(a)

Code that swaps two integers

Path	PC	Program Input
1,8	$X \leq Y$	$X=1 \ Y=1$
1,2,3,4,5,8	$X > Y \ \& \ Y - X \leq 0$	$X=2 \ Y=1$
1,2,3,4,5,6	$X > Y \ \& \ Y - X > 0$	none

(c)

The PC's and their solutions (if they exist)



The symbolic execution tree for the code

- **Nodes** represent program states
- **Edges** represent transitions between states
- Use a **off-the-shelf** constraint solver to solve path conditions, e.g., SMT-Solver / SAT-Solver



Symbolic Execution Limitations

- The technique suffers from some problems that limit its effectiveness on real world software:
 1. **Path explosion**: most real world software have an extremely large number of paths, and it is difficult to analyze all program paths.
 2. **Complex constraints**: it is difficult to analyze constraints involving mathematical functions such as `sin()` and `log()`

The third branch contains a problematic non-linear constraint, and if the **Math** library is not available in source code or bytecode, then deriving constraints can be difficult in the first place.

```
1 double example(int x, int y, double z) {  
2   boolean flag = y > 1000;  
3   // ...  
4   if(x + y == 1024)  
5     if(flag)  
6       if(Math.cos(z) - 0.95 < Math.exp(z))  
7         // target branch  
8   // ...  
9 }
```

Native code: no access to source code

Tools

- Symbolic Execution - Java PathFinder (NASA)
<http://javapathfinder.sourceforge.net/extensions/symbc/doc/>
- SAGE (Microsoft)
https://patricegodefroid.github.io/public_psfiles/ndss2008.pdf
- PEX (C#) <https://pexforfun.com/default.aspx>
- CUTE(C)
<http://osl.cs.illinois.edu/publications/conf/sigsoft/SenMA05.html>
- Z3 is one of the most powerful SMT solvers currently available.
<http://channel9.msdn.com/posts/Peli/The-Z3-Constraint-Solver/>



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