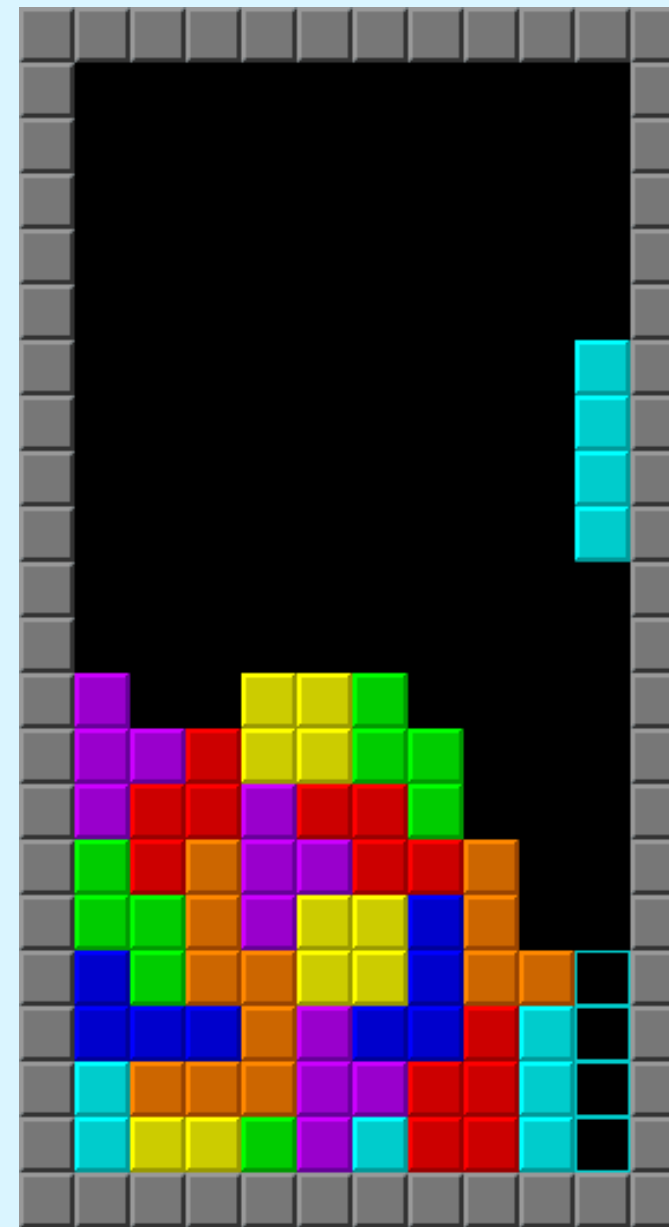
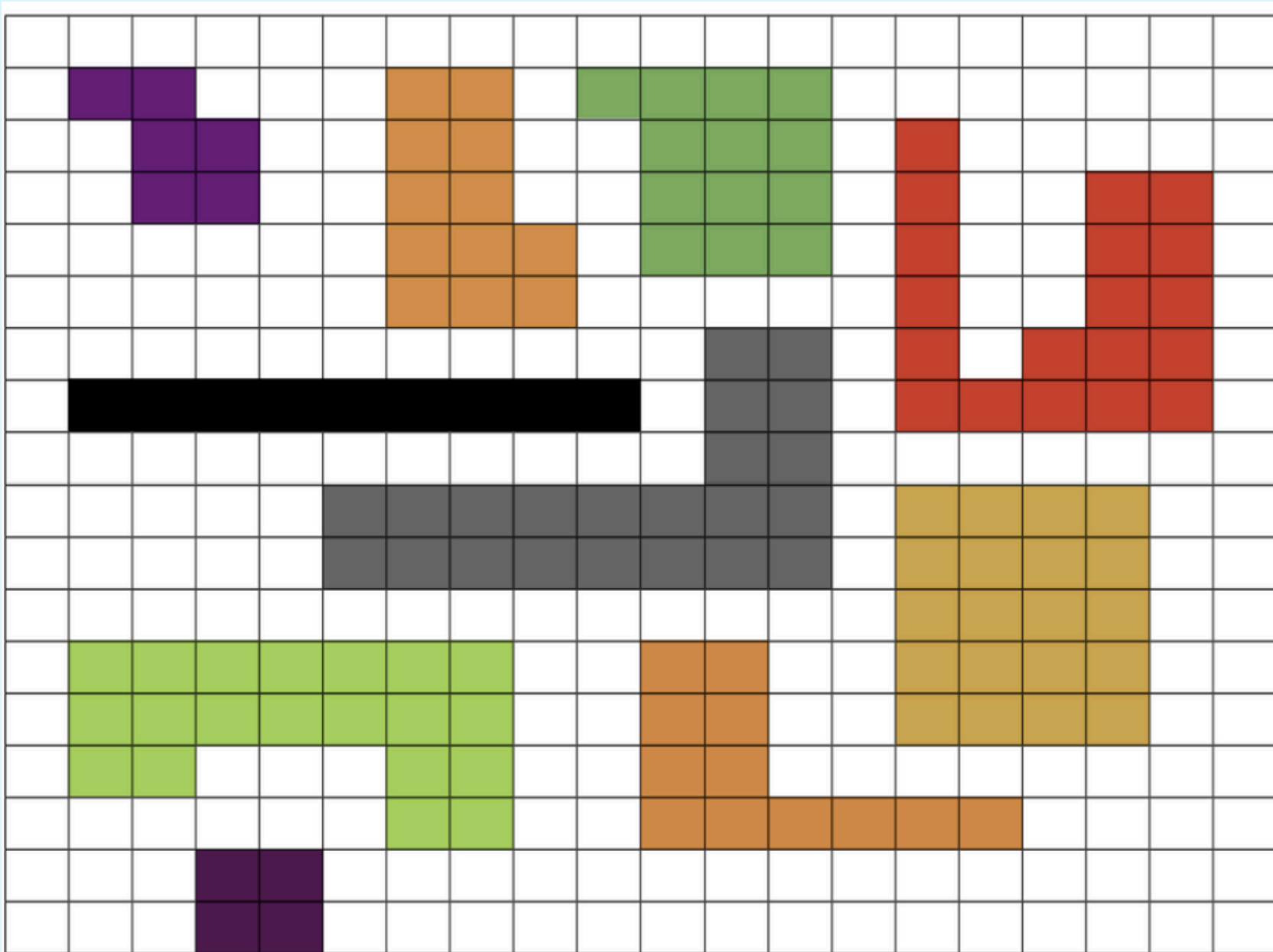


The background is a pixel art scene. At the top, there are three clusters of pixelated clouds in shades of light blue and white. On the left, a small patch of ground with green grass and brown soil has a single yellow sunflower with a brown center. On the right, a similar patch of ground has a flagpole with a blue flag. The bottom of the image is a wide, horizontal strip of pixelated ground, consisting of a thin layer of green grass on top of a thicker layer of brown soil.

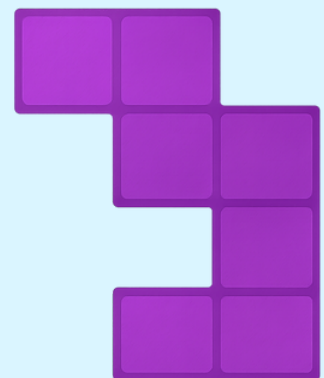
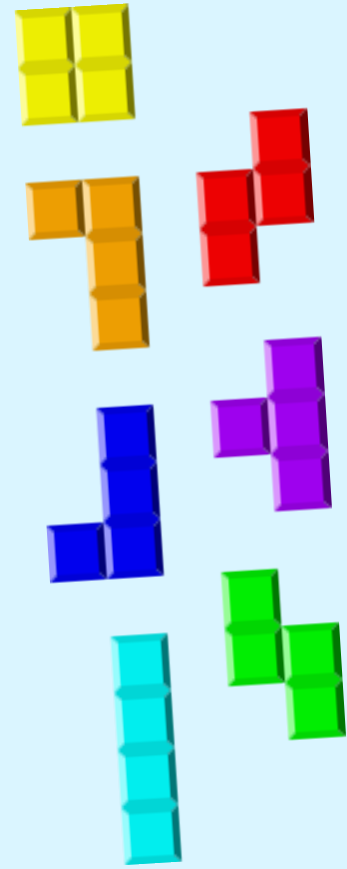
# SQUARE PACKING

ARTUR SAHAKYAN

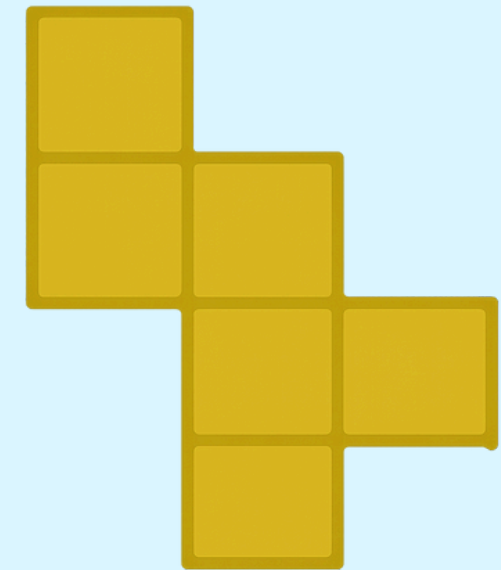
# PROBLEM



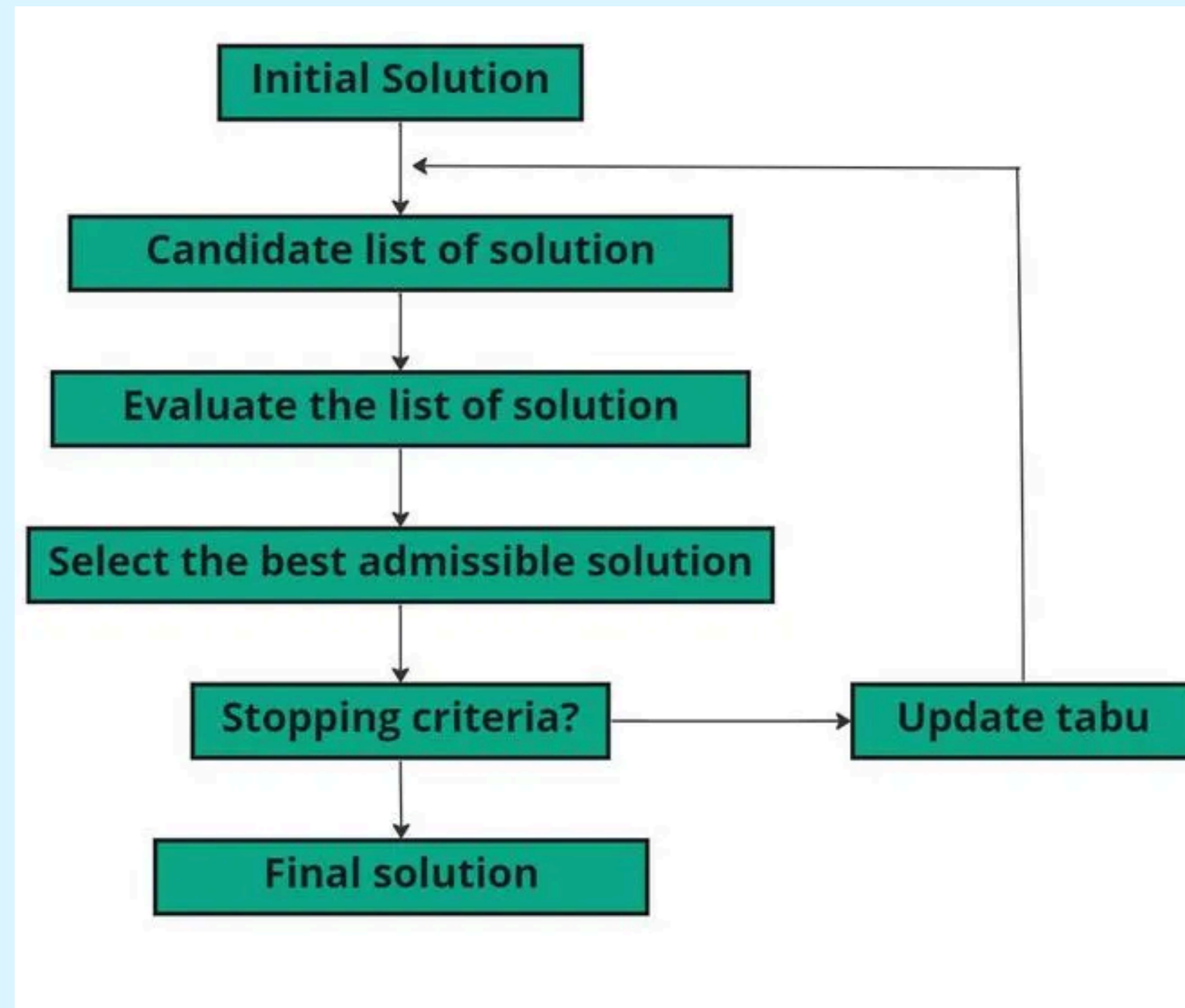
# LOGICAL TO REPRESENT AS "LABELED" MATRIX



```
matrix = [  
  [0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0],  
  [0,1,1,0,0,0,2,2,0,3,3,3,3,0,0,0,0,0],  
  [0,0,1,1,0,0,2,2,0,0,3,3,3,0,4,0,0,0],  
  [0,0,1,1,0,0,2,2,0,0,3,3,3,0,4,0,0,4],  
  [0,0,0,0,0,0,2,2,2,0,3,3,3,0,4,0,0,4],  
  [0,0,0,0,0,0,2,2,2,0,0,0,0,0,4,0,0,4],  
  [0,0,0,0,0,0,0,0,0,0,0,0,0,5,5,0,4,4],  
  [0,6,6,6,6,6,6,6,6,6,0,5,5,0,4,4,4,4],  
  [0,0,0,0,0,0,0,0,0,0,0,5,5,0,0,0,0,0],  
  [0,0,0,0,0,5,5,5,5,5,5,5,5,0,7,7,7,7],  
  [0,0,0,0,0,5,5,5,5,5,5,5,5,0,7,7,7,7],  
  [0,0,0,0,0,0,0,0,0,0,0,0,0,0,7,7,7,7],  
  [0,9,9,9,9,9,9,9,0,0,8,8,0,0,7,7,7,7],  
  [0,9,9,9,9,9,9,9,0,0,8,8,0,0,7,7,7,7],  
  [0,9,9,0,0,0,9,9,0,0,8,8,0,0,0,0,0,0],  
  [0,0,0,0,0,0,9,9,0,0,8,8,8,8,8,8,0,0],  
  [0,0,0,10,10,0,0,0,0,0,0,0,0,0,0,0,0],  
  [0,0,0,10,10,0,0,0,0,0,0,0,0,0,0,0,0],  
]
```



# METAHEURISTIC APPROACH WITH TABU SEARCH



# TABU SEARCH APPLIED ON SQUARE PACKING

## Step 1: Random Valid Placement

We randomly place each shape in the grid within the board bounds, even if it overlaps.

## Step 2: Define the Fitness Function

We want to minimize the number of illegal cells, i.e., cells that are: Out of bounds, or overlapping with other shapes

$$\text{fitness} = 1000 + \text{number\_of\_conflict\_cells}$$

(The 1000 ensures all imperfect packings are worse than any perfect one.)

## Step 3: Explore the Neighborhood

We generate neighbors by:

Choosing one shape at random, moving it to a different random legal position on the board. This gives a new potential solution.

#### Step 4: Tabu List

To avoid cycling back to recent placements:

We store recent moves as (shape\_index, old\_position)

These moves are tabu (forbidden) for TABU\_TENURE steps

Aspiration rule: we allow a tabu move if it improves the best solution ever found

This helps the algorithm:

- Escape local optima
- Avoid infinite loops
- Explore the search space better
- 

#### Step 5: Update Best Solution

If a neighbor is better (lower fitness), it becomes the new current solution.

If it's the best so far (even if TABU), we remember it.

#### Step 6: Repeat

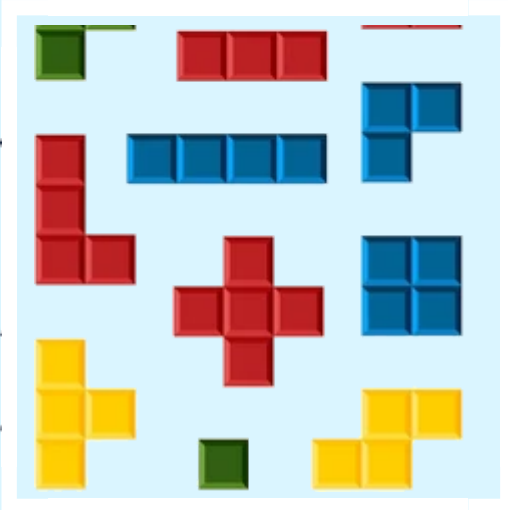
We do this for a fixed number of iterations (say 500), and restart the process if needed.

# OVERLAPS, VIOLATIONS

**WE CHECK FOR COLLISIONS, THERE ARE MANY METHODS TO DO SO**

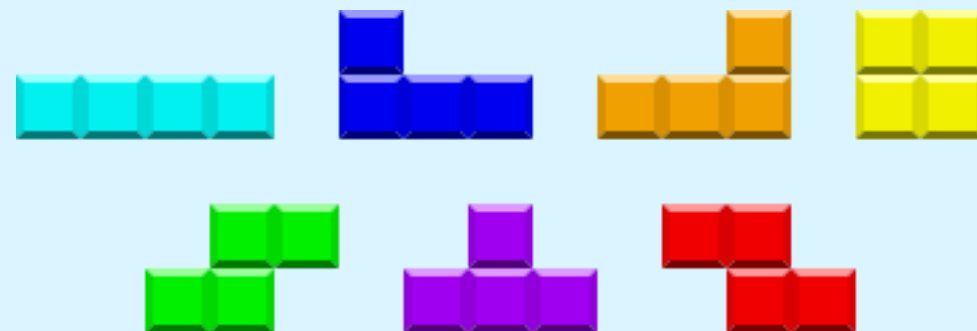
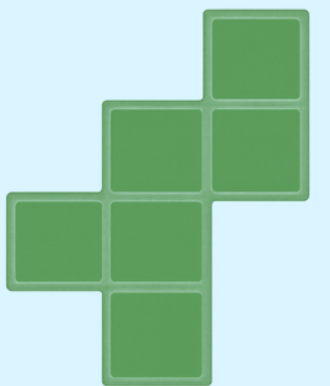
Search space has lower bound, [length of side of smallest square for packing, we can increment it in case of failure to pack]

$$S(A) = \sqrt{A}$$

$$S(A) = \sqrt{A}$$


```
def build_board(anchors, shape_list, S):  
    ## blank board  
    board = np.zeros((S, S), int)  
    bad = 0  
    for (r0, c0), (k, sh) in zip(anchors, shape_list):  
        for dr, dc in sh["cells"]:  
            r, c = r0 + dr, c0 + dc  
  
            ## cell out of bounds ??  
            if r < 0 or r >= S or c < 0 or c >= S:  
                bad += 1  
  
            ## cell already occupied ??  
            elif board[r, c]:  
                bad += 1  
            else:  
                board[r, c] = k  
    return board, bad
```

```
def fitness(anchors, shape_list, S):  
    _, bad = build_board(anchors, shape_list, S)  
    if bad == 0:  
        return 0.0  
    return 1000 + bad
```



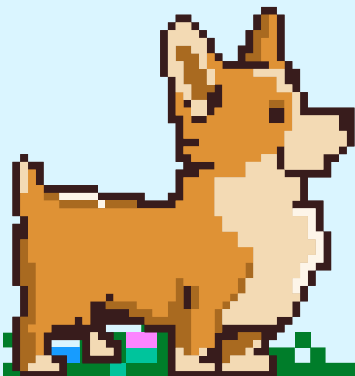
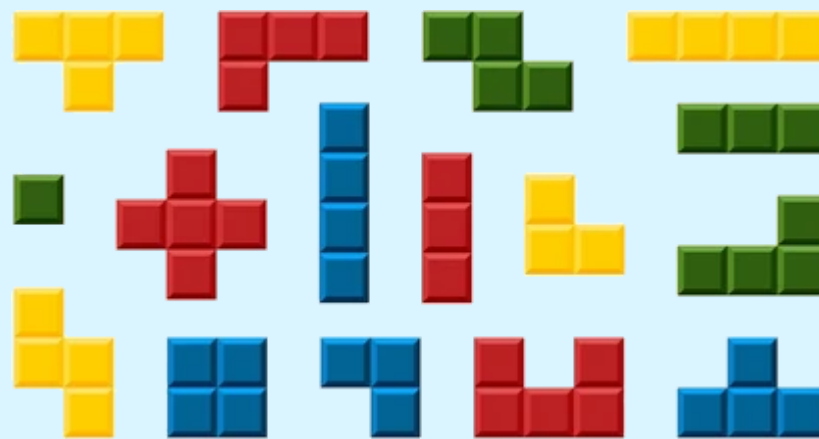
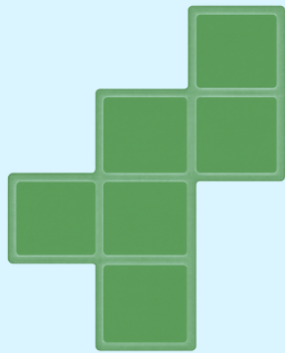
## RESULT WITH LOW-COST PARAMETERS FOR TABU SEARCH





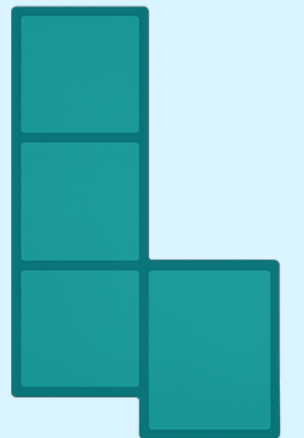
⚠ THIS IS HEURISTIC – WE FIND A LIKELY MINIMAL PACKING, NOT A PROVABLY MINIMAL ONE.

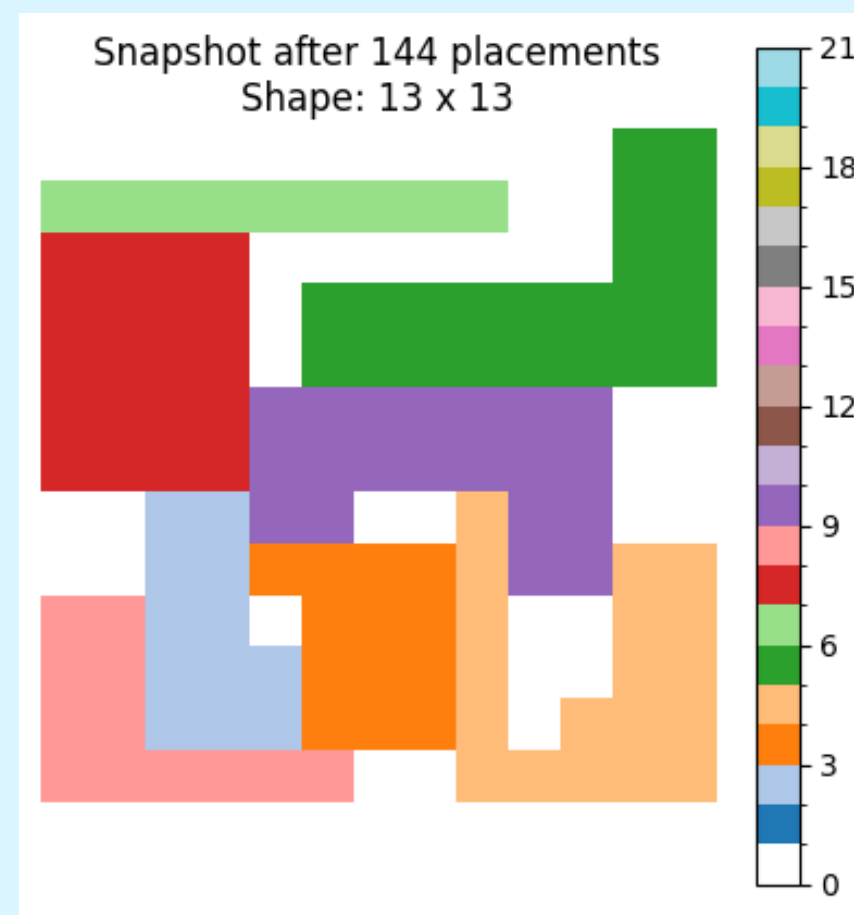
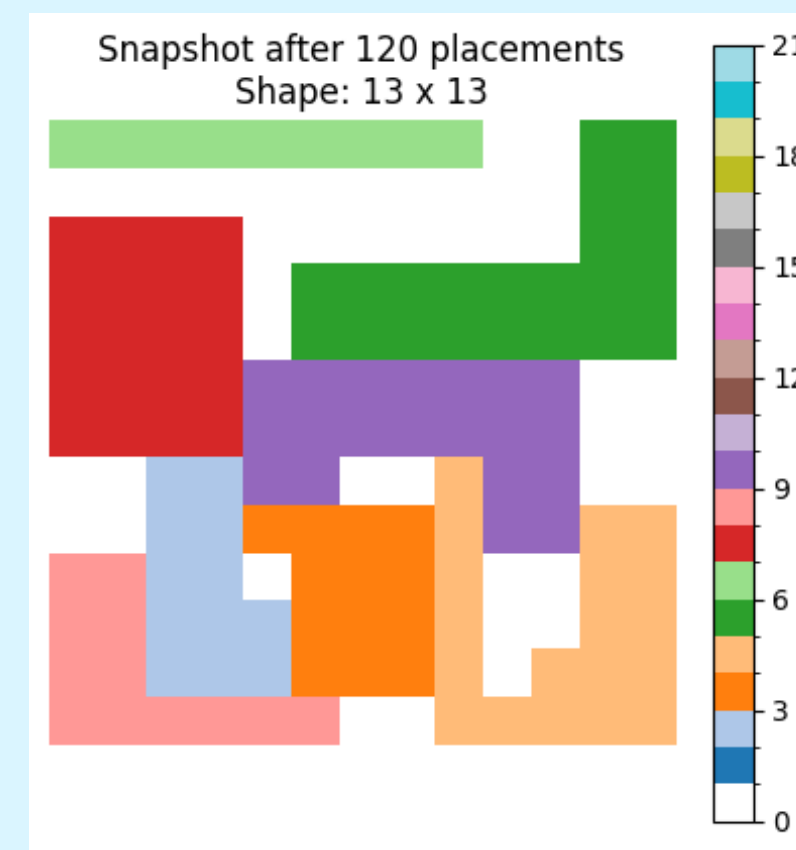
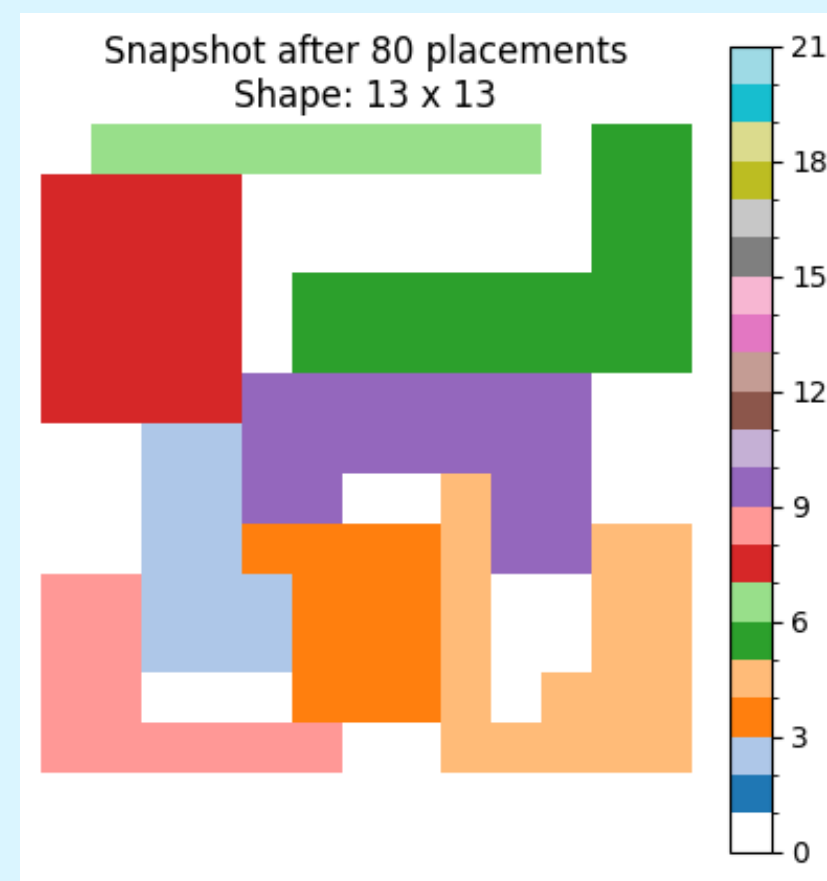
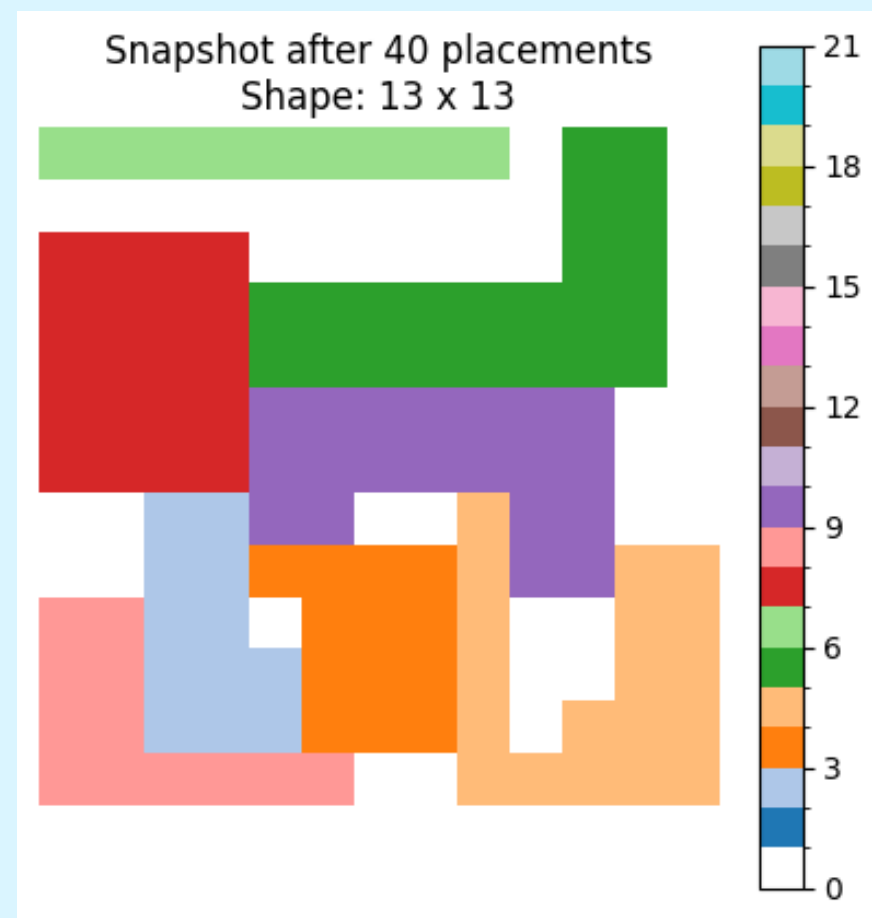
TABU SEARCH DOES USE THE SAME LOWER-BOUND STRATEGY TO START ITS SEARCH, BUT UNLIKE BACKTRACKING, IT CAN'T GUARANTEE OPTIMALITY – IT JUST TRIES TO FIND THE SMALLEST  $S$  WHERE A FEASIBLE PACKING CAN BE FOUND HEURISTICALLY.



# DETERMINISTIC SOLUTION

- SORT SHAPES
- START WITH A LOWER BOUND OF  $S = \text{SQRT}(A)$  AND INCREMENT IF NEEDED
- DO A DEPTH-FIRST SEARCH, PLACE LARGER SHAPES FIRST
- BACKTRACK IN CASE OF FAILING TO PACK THE NEXT SHAPE
- AFTER SOME N PLACEMENTS WE SEE NO IMPROVEMENT (NO REDUCTION IN SMALLEST SQUARE AREA)





# COMPARISON



- 
- ⌚ DETERMINISTIC SOLUTION COMPLETED IN 7.021 SECONDS
  - ⌚ TABU SEARCH COMPLETED IN 4.631 SECONDS



THANK  
YOU

