



**POLITECNICO**  
MILANO 1863

# Three-Dimensional Bin Packing with Vertical Support

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# Case study

- Large warehouses
- Mixed-case palletization
- No control over items' shape (strongly heterogeneous)
- Pallets wrapped during loading procedure



Figure: Example of pit palletization  
(Schäfer Case Picking — SSI SCHÄFER)

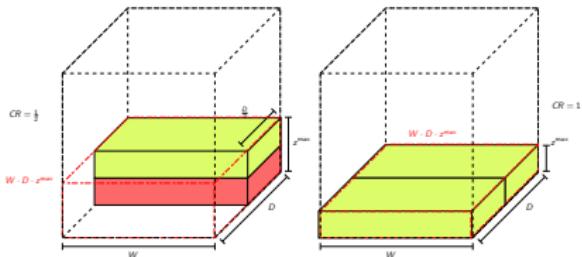


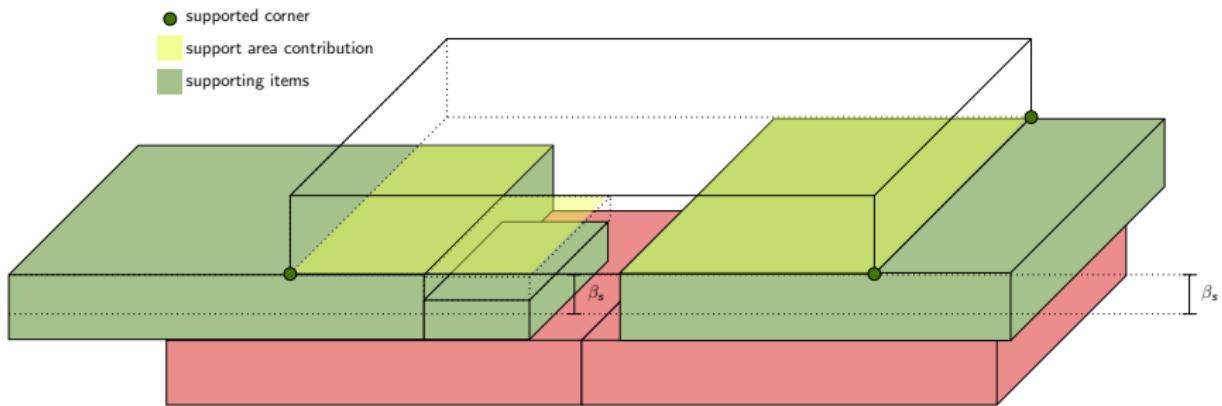
Figure: Cage ratio of two different bin configurations

### Definition

An item has vertical support if one of the following conditions hold:

- **Condition 1:** at least a percentage  $\alpha_s$  of its base area is resting on other items
- **Condition 2:** at least 3 of its vertices are resting over other items and **Condition 1** holds with a lower percentage

# Vertical Support



**Figure:** Representation of an item with conditions 1 and 2 of vertical support given  
 $\alpha_s = 0.5, \beta_s$

- The problem is NP-Hard
- Exact methods only for small instances
- Existing 3D-BPP heuristics don't consider practical constraints
- Solutions for container loading and pallet loading problems are layer based

**minimize** number of used bins

**then, maximize** average cage ratio of the used bins

**subject to** all items are assigned to one and only one bin

all items are inside the bin's bounds

no overlaps between items in the same bin

all items have vertical support

## MILP Proxy Model - Objective Function

$$\begin{array}{ll}
 \min & \sum_{b \in B} (Hv_b + z_b^{\max}) \\
 \text{s.t.} & \sum_{b \in B} u_{ib} + \sum_{b \in B} u_{jb} = 1 \quad \forall (i, j) \in I^{OR} \\
 & u_{ib} \leq v_b \quad \forall i \in I, \forall b \in B \\
 & v_b \geq v_c \quad \forall (b, c) \in B : b < c \\
 & x_i + w_i \leq W \quad \forall i \in I \\
 & y_i + d_i \leq D \quad \forall i \in I \\
 & z_i + h_i \leq H \quad \forall i \in I \\
 & (x_i + w_i) - x_j \leq W(1 - x_{ij}^P) \quad \forall i, j \in I \\
 & x_j - (x_i + w_i) + 1 \leq Wx_{ij}^P \quad \forall i, j \in I \\
 & (y_i + d_i) - y_j \leq D(1 - y_{ij}^P) \quad \forall i, j \in I \\
 & y_j - (y_i + d_i) + 1 \leq Dy_{ij}^P \quad \forall i, j \in I \\
 & (z_i + h_i) - z_j \leq H(1 - z_{ij}^P) \quad \forall i, j \in I \\
 & z_j - (z_i + h_i) + 1 \leq Hz_{ij}^P \quad \forall i, j \in I \\
 & x_{ij}^P + x_{ji}^P + y_{ij}^P + y_{ji}^P + z_{ij}^P + z_{ji}^P \geq u_{ib} + u_{jb} - 1 \quad \forall i, j \in I, \forall b \in B \\
 & z_b^{\max} \geq (z_i + h_i) - H(1 - u_{ib}) \quad \forall i \in I, \forall b \in B
 \end{array}$$

**minimize**      number of used bins  
**then, maximize**    average cage ratio of the used bins  
**subject to**    all items are assigned to one and only one bin  
                      all items are inside the bin's bounds  
                      no overlaps between items in the same bin  
                      all items have vertical support

## MILP Proxy Model - Geometric Constraints 1

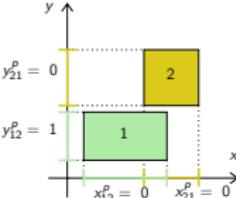
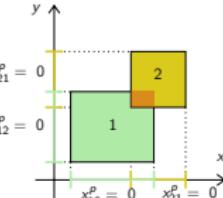
$$\begin{array}{ll}
 \min & \sum_{b \in B} (Hv_b + z_b^{max}) \\
 \text{s.t.} & \sum_{b \in B} u_{ib} + \sum_{b \in B} u_{jb} = 1 \quad \forall (i, j) \in I^{OR} \\
 & u_{ib} \leq v_b \quad \forall i \in I, \forall b \in B \\
 & v_b \geq v_c \quad \forall (b, c) \in B : b < c \\
 & x_i + w_i \leq W \quad \forall i \in I \\
 & y_i + d_i \leq D \quad \forall i \in I \\
 & z_i + h_i \leq H \quad \forall i \in I \\
 & (x_i + w_i) - x_j \leq W(1 - x_{ij}^P) \quad \forall i, j \in I \\
 & x_j - (x_i + w_i) + 1 \leq Wx_{ij}^P \quad \forall i, j \in I \\
 & (y_i + d_i) - y_j \leq D(1 - y_{ij}^P) \quad \forall i, j \in I \\
 & y_j - (y_i + d_i) + 1 \leq Dy_{ij}^P \quad \forall i, j \in I \\
 & (z_i + h_i) - z_j \leq H(1 - z_{ij}^P) \quad \forall i, j \in I \\
 & z_j - (z_i + h_i) + 1 \leq Hz_{ij}^P \quad \forall i, j \in I \\
 & x_{ij}^P + x_{ji}^P + y_{ij}^P + y_{ji}^P + z_{ij}^P + z_{ji}^P \geq u_{ib} + u_{jb} - 1 \quad \forall i, j \in I, \forall b \in B \\
 & z_b^{max} \geq (z_i + h_i) - H(1 - u_{ib}) \quad \forall i \in I, \forall b \in B
 \end{array}$$

minimize                    number of used bins  
 then, maximize            average cage ratio of the used bins  
**subject to**    all items are assigned to one and only one bin  
                   all items are inside the bin's bounds  
                   no overlaps between items in the same bin  
                   all items have vertical support

## MILP Proxy Model - Geometric Constraints 2

$$\begin{array}{ll}
 \min & \sum_{b \in B} (Hv_b + z_b^{\max}) \\
 \text{s.t.} & \sum_{b \in B} u_{ib} + \sum_{b \in B} u_{jb} = 1 \quad \forall (i, j) \in I^{OR} \\
 & u_{ib} \leq v_b \quad \forall i \in I, \forall b \in B \\
 & v_b \geq v_c \quad \forall (b, c) \in B : b < c \\
 & x_i + w_i \leq W \quad \forall i \in I \\
 & y_i + d_i \leq D \quad \forall i \in I \\
 & z_i + h_i \leq H \quad \forall i \in I \\
 & (x_i + w_i) - x_j \leq W(1 - x_{ij}^P) \quad \forall i, j \in I \\
 & x_j - (x_i + w_i) + 1 \leq Wx_{ij}^P \quad \forall i, j \in I \\
 & (y_i + d_i) - y_j \leq D(1 - y_{ij}^P) \quad \forall i, j \in I \\
 & y_j - (y_i + d_i) + 1 \leq Dy_{ij}^P \quad \forall i, j \in I \\
 & (z_i + h_i) - z_j \leq H(1 - z_{ij}^P) \quad \forall i, j \in I \\
 & z_j - (z_i + h_i) + 1 \leq Hz_{ij}^P \quad \forall i, j \in I \\
 & x_{ij}^P + x_{ji}^P + y_{ij}^P + y_{ji}^P + z_{ij}^P + z_{ji}^P \geq u_{ib} + u_{jb} - 1 \quad \forall i, j \in I, \forall b \in B \\
 & z_b^{\max} \geq (z_i + h_i) - H(1 - u_{ib}) \quad \forall i \in I, \forall b \in B
 \end{array}$$

minimize                    number of used bins  
 then, maximize            average cage ratio of the used bins  
**subject to**            all items are assigned to one and only one bin  
                           all items are inside the bin's bounds  
                           no overlaps between items in the same bin  
                           all items have vertical support

		$\min \quad \sum_{b \in B} (Hv_b + z_b^{max})$	
		$\text{s.t.} \quad \sum_{b \in B} u_{ib} + \sum_{b \in B} u_{jb} = 1 \quad \forall (i, j) \in I^{OR}$	
minimize then, maximize	number of used bins average cage ratio of the used bins	$u_{ib} \leq v_b \quad \forall i \in I, \forall b \in B$	
<b>subject to</b>	all items are assigned to one and only one bin all items are inside the bin's bounds no overlaps between items in the same bin all items have vertical support	$v_b \geq v_c \quad \forall (b, c) \in B : b < c$ $x_i + w_i \leq W \quad \forall i \in I$ $y_i + d_i \leq D \quad \forall i \in I$ $z_i + h_i \leq H \quad \forall i \in I$ $(x_i + w_i) - x_j \leq W(1 - x_{ij}^P) \quad \forall i, j \in I$ $x_j - (x_i + w_i) + 1 \leq Wx_{ij}^P \quad \forall i, j \in I$ $(y_i + d_i) - y_j \leq D(y_{ij}^P) \quad \forall i, j \in I$ $y_j - (y_i + d_i) + 1 \leq Dy_{ij}^P \quad \forall i, j \in I$ $(z_i + h_i) - z_j \leq H(z_{ij}^P) \quad \forall i, j \in I$ $z_j - (z_i + h_i) + 1 \leq Hz_{ij}^P \quad \forall i, j \in I$ $x_{ij}^P + x_{ji}^P + y_{ij}^P + y_{ji}^P + z_{ij}^P + z_{ji}^P \geq u_{ib} + u_{jb} - 1 \quad \forall i, j \in I, \forall b \in B$ $z_b^{max} \geq (z_i + h_i) - H(1 - u_{ib}) \quad \forall i \in I, \forall b \in B$	
			

**Figure:** Precedences variables (2D case)

**minimize** number of used bins

then, maximize average cage ratio of the used bins

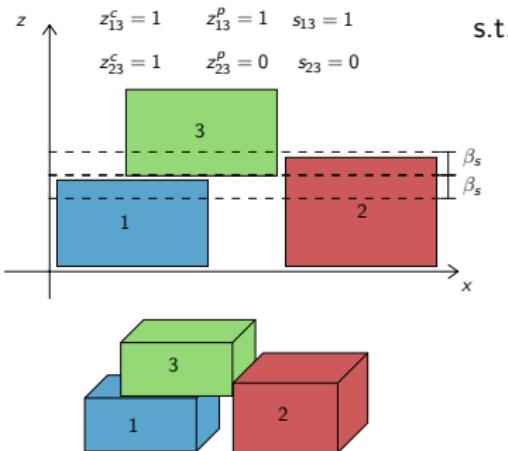
**subject to** all items are assigned to one and only one bin

all items are inside the bin's bounds

no overlaps between items in the same bin

all items have vertical support

## MILP Proxy Model - Closeness



$$\begin{aligned}
 & z \\
 & z_{13}^c = 1 \quad z_{13}^p = 1 \quad s_{13} = 1 \\
 & z_{23}^c = 1 \quad z_{23}^p = 0 \quad s_{23} = 0 \\
 & \text{s.t.} \quad z_j - (z_i + h_i) \leq \beta_s + H(1 - z_{ij}^c) \quad \forall (i, j) \in I : i \neq j \\
 & z_j - (z_i + h_i) \geq -\beta_s - H(1 - z_{ij}^c) \quad \forall (i, j) \in I : i \neq j \\
 & s_{ij} \leq z_{ij}^p \quad \forall (i, j) \in I \\
 & s_{ij} \leq z_{ij}^c \quad \forall (i, j) \in I \\
 & s_{ij} \geq z_{ij}^p + z_{ij}^c - 2 \quad \forall (i, j) \in I : i \neq j \\
 & \sum_{j \in I} s_{ij} \leq \sum_{b \in B} u_{ib} \quad \forall i \in I
 \end{aligned}$$

**Figure:** Closeness variables example

## MILP Proxy Model - Discretized Vertical Support

Pre-Computed Parameter

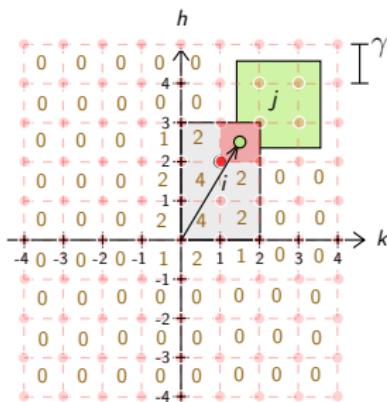
$$O(i, j, k, h)$$

$$O(i, j, 1, 2) = 1$$

Variables

$$s_{ijb}^{12} = 1$$

$$(x_j - x_i, y_j - y_i)$$



s.t.

$z_i \leq H(1 - g_i)$	$\forall i \in I$
$\sum s_{ijb}^{kh} \leq s_{ij}$	$\forall (i, j) \in I$
$\sum_{(k,h) \in \Delta: O(i,j,k,h) \neq 0} s_{ijb}^{kh} \leq u_{ib}$	$\forall (i, j, b) \in I^B$
$\sum_{(k,h) \in \Delta: O(i,j,k,h) \neq 0} s_{ijb}^{kh} \leq u_{jb}$	$\forall (i, j, b) \in I^B$
$x_j - x_i \geq \gamma k - 2W(1 - s_{ijb}^{kh})$	$\forall (k, h) \in \Delta, \forall (i, j, b) \in I^B : O(i, j, k, h) \neq 0$
$x_j - x_i \leq \gamma(k+1) + 2W(1 - s_{ijb}^{kh})$	$\forall (k, h) \in \Delta, \forall (i, j, b) \in I^B : O(i, j, k, h) \neq 0$
$y_j - y_i \geq \gamma h - 2D(1 - s_{ijb}^{kh})$	$\forall (k, h) \in \Delta, \forall (i, j, b) \in I^B : O(i, j, k, h) \neq 0$
$y_j - y_i \leq \gamma(h+1) + 2D(1 - s_{ijb}^{kh})$	$\forall (k, h) \in \Delta, \forall (i, j, b) \in I^B : O(i, j, k, h) \neq 0$
$\sum_{(k,h) \in \Delta, b \in B, j \in I: i \neq j \wedge O(i,j,k,h) \neq 0} O(i, j, k, h) s_{ijb}^{kh} \geq \alpha_s w_i d_i - w_i d_i g_i$	$\forall i \in I$

Figure: Space discretization

# Proposed Heuristic Overview

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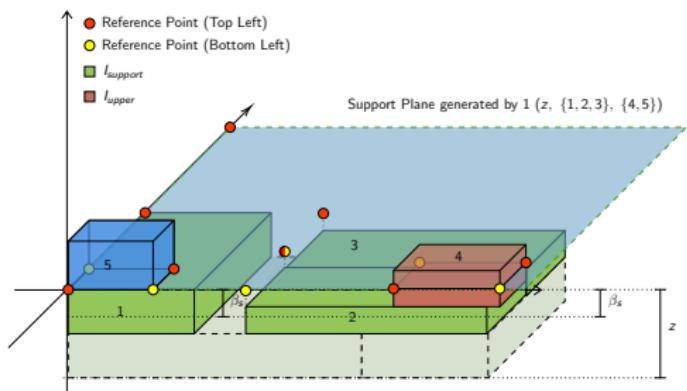
Composed of:

- Constructive heuristic (Support Planes)
- Beam-Search

Each node is a partial solution, starts from the empty solution

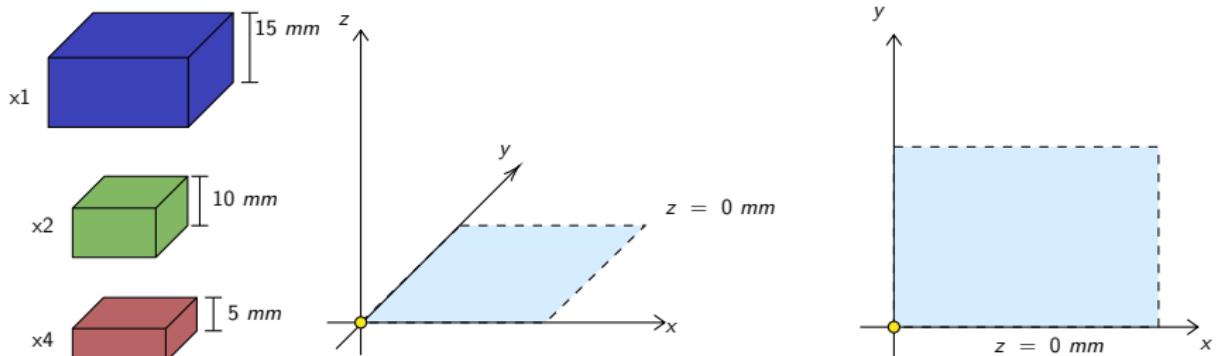
# Support Planes

- Operates on a single bin
- Exploits a modified 2D-BPP heuristic
- No explicit layers
- Guarantees vertical support

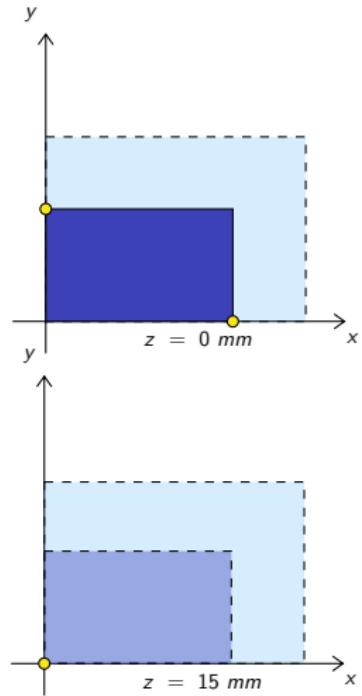
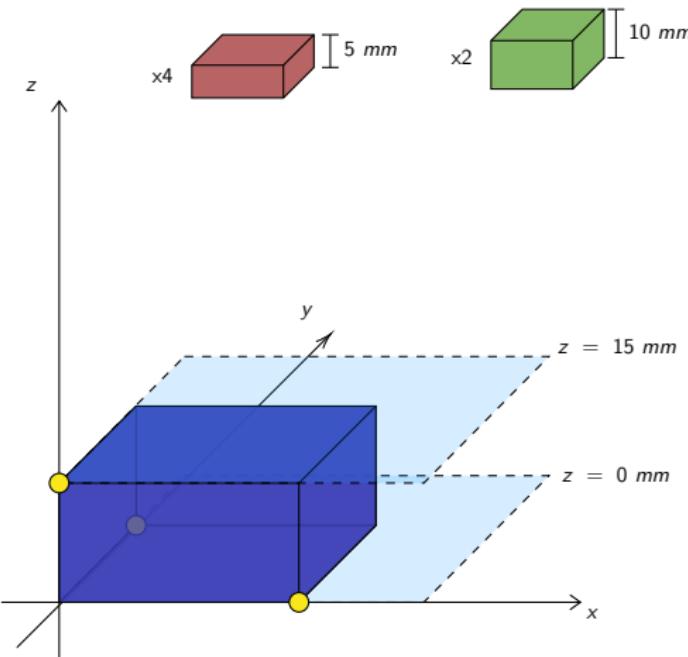


**Figure:** An example of a support plane generated by item 1

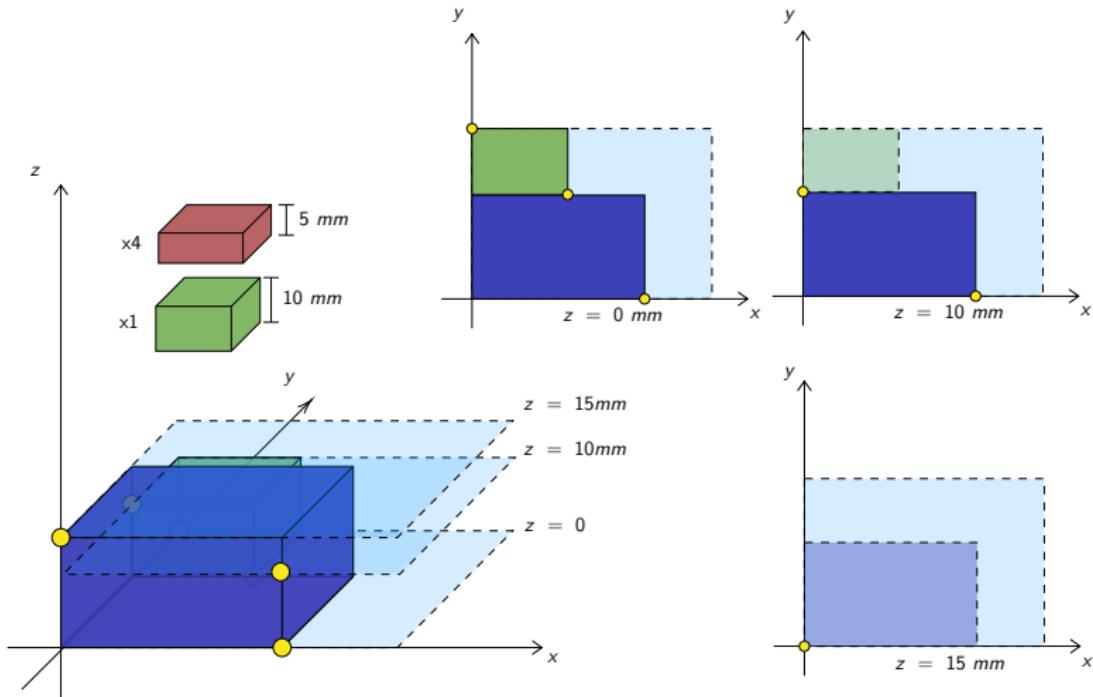
# Support Planes



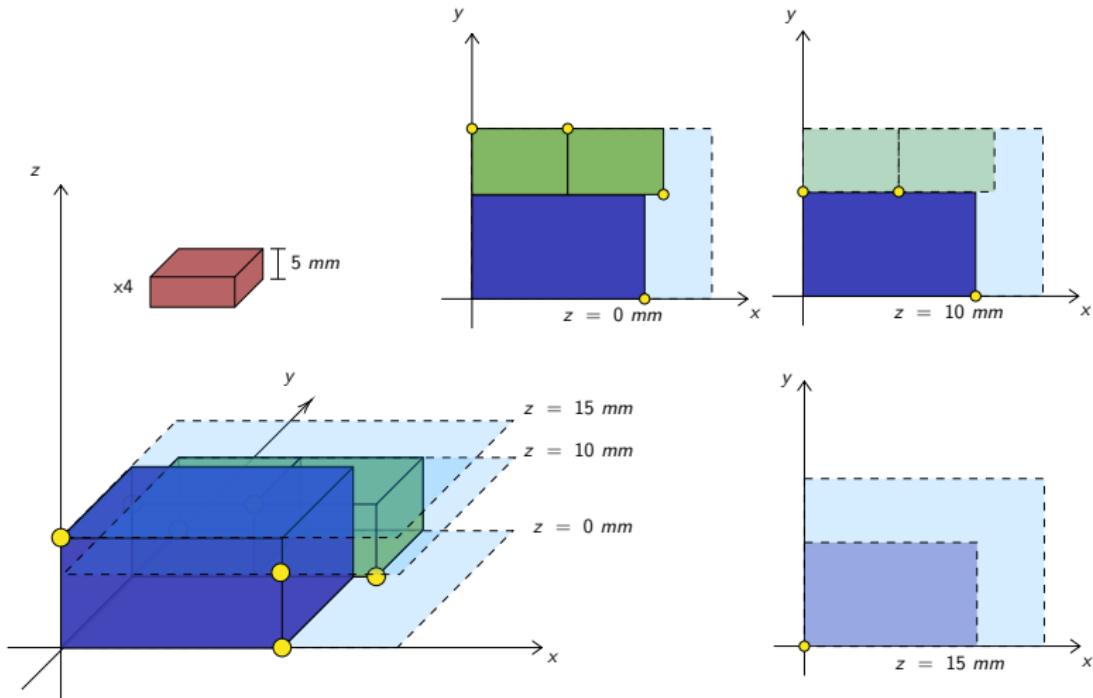
# Support Planes



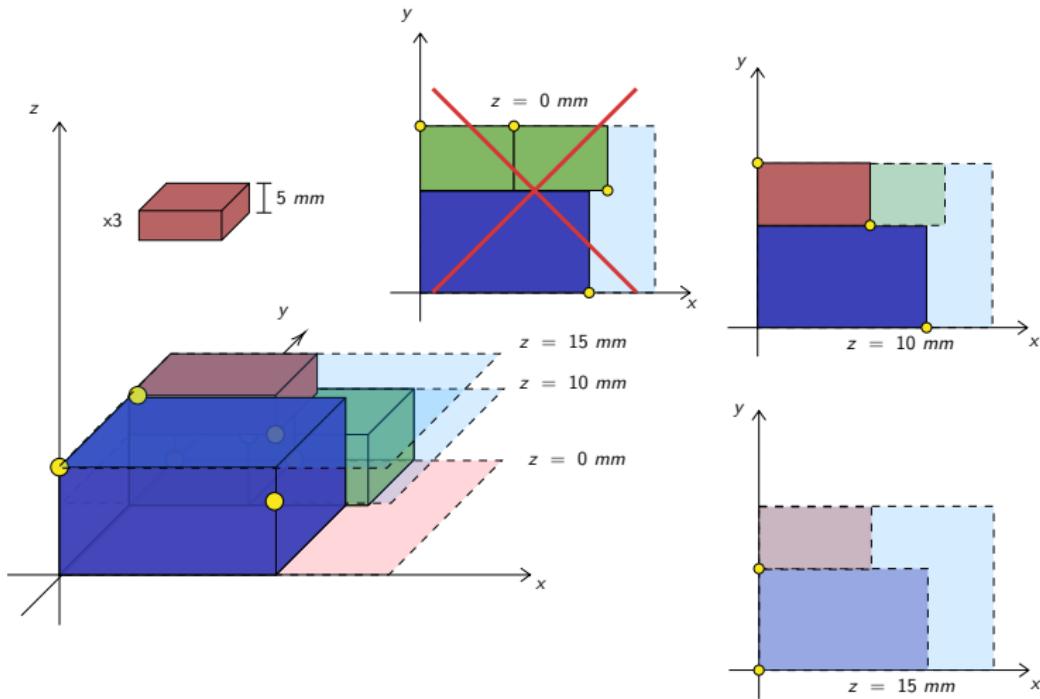
# Support Planes



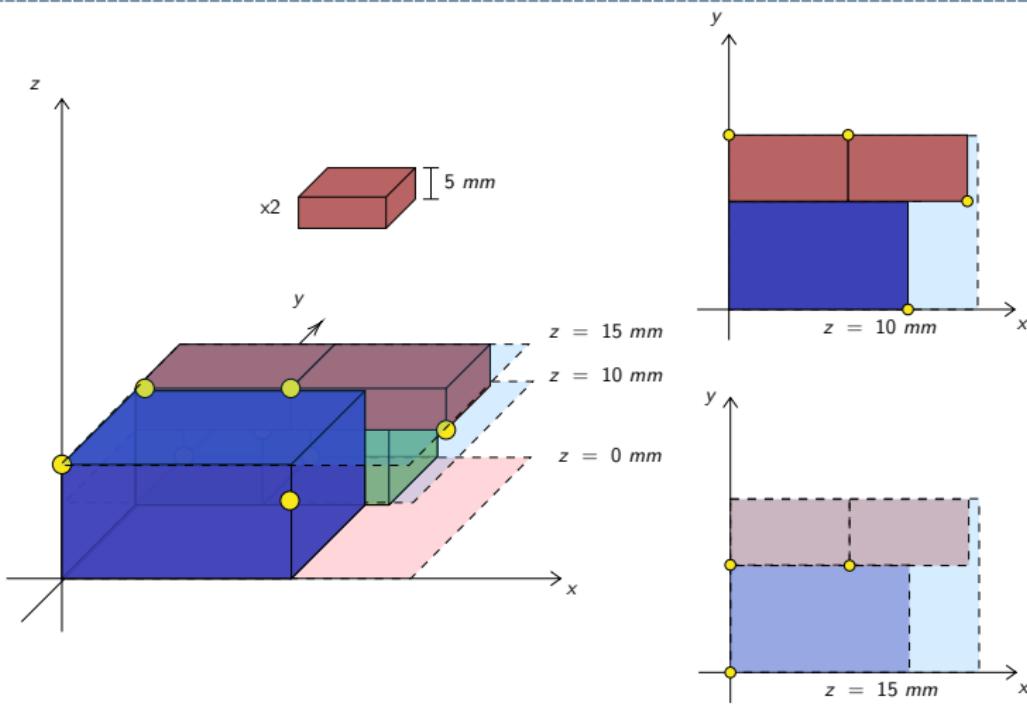
# Support Planes



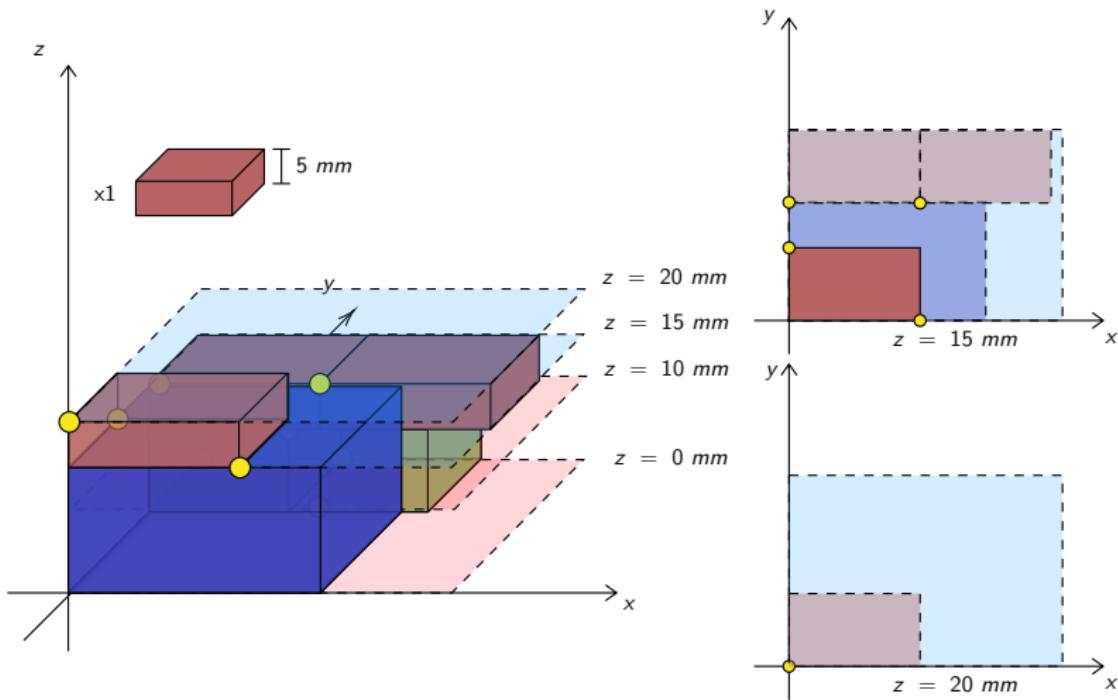
# Support Planes



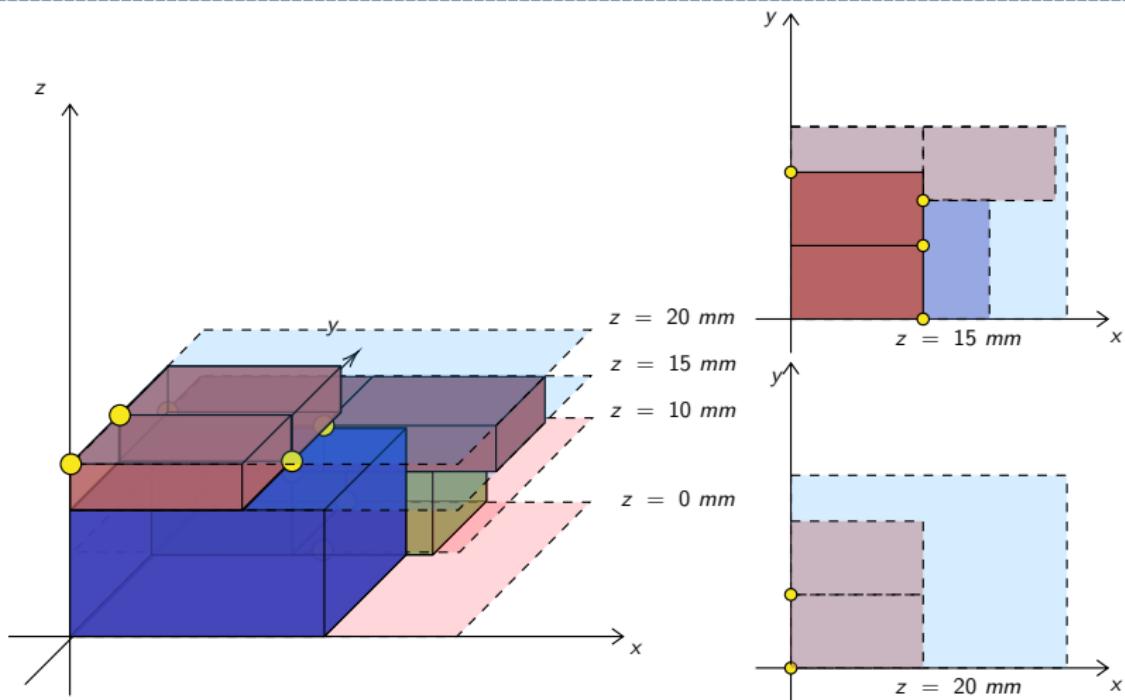
# Support Planes



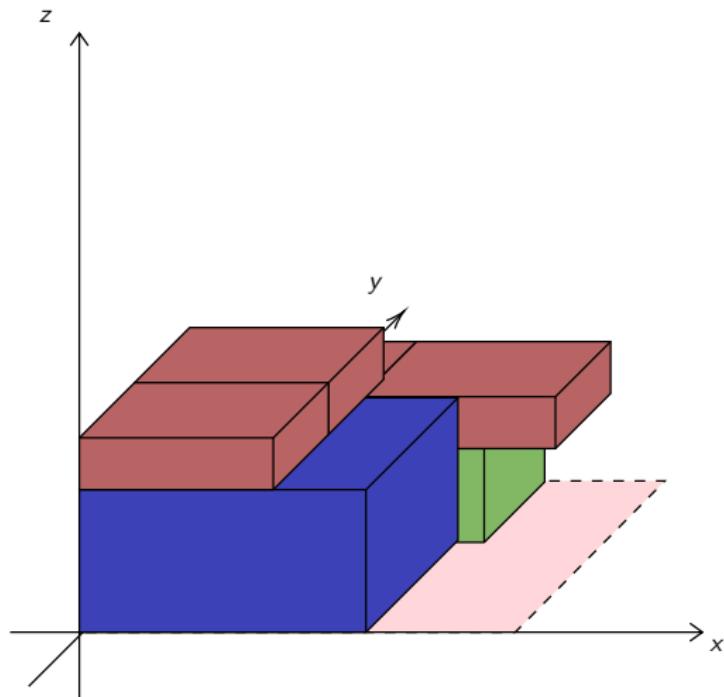
# Support Planes



## Support Planes

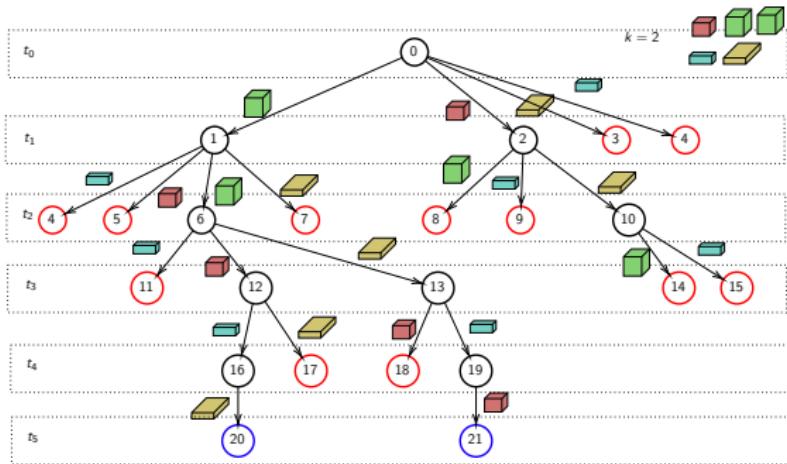


## Support Planes



## Beam Search

- Exploits support planes
- Explores different sequences of placements
- Limited to  $k$  best placements



## Beam Search

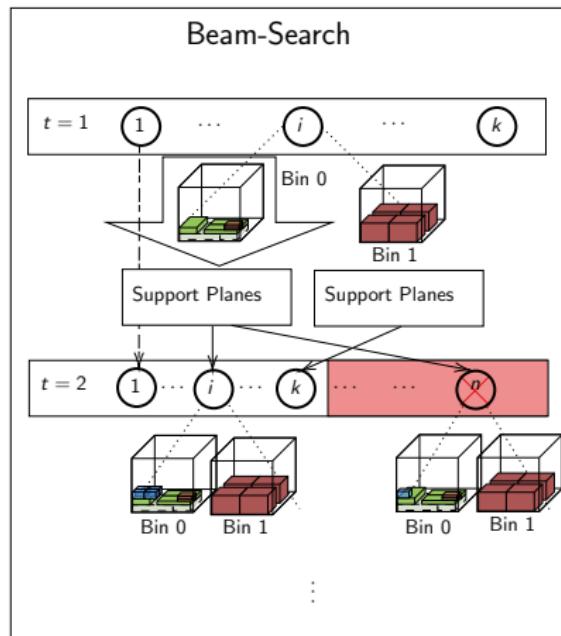


Figure: Conceptual representation of the proposed heuristic

## Beam Search

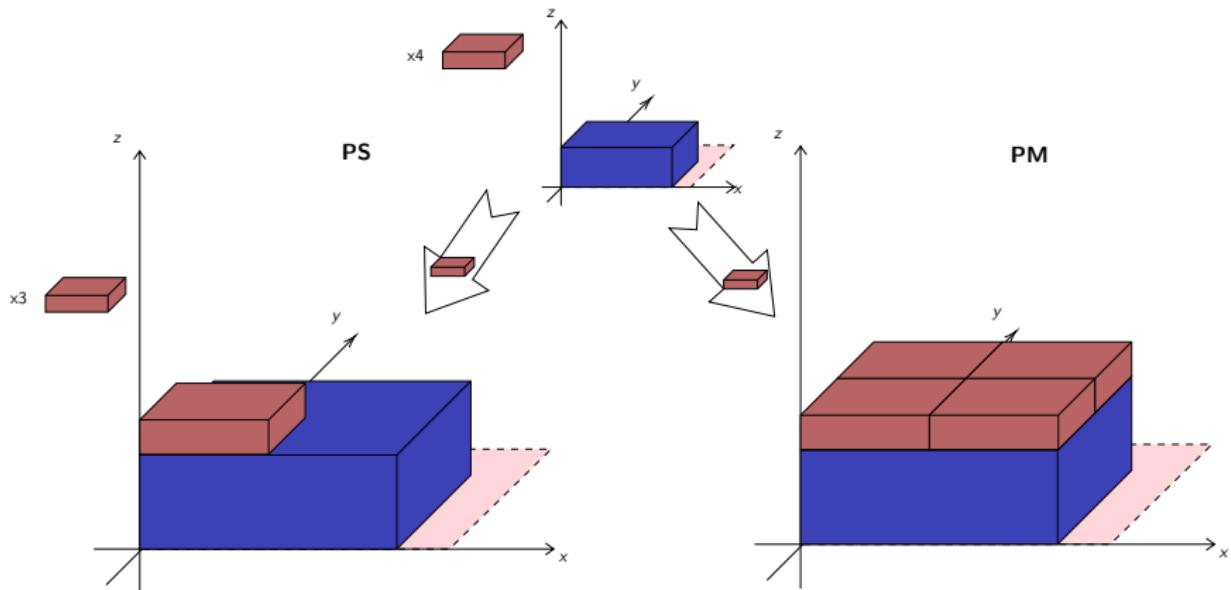


Figure: Proposed placement modes

# Proposed Heuristic Optimizations

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- Duplicates removal
- Fast overlap checks
- Memory management

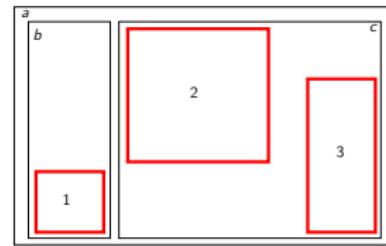
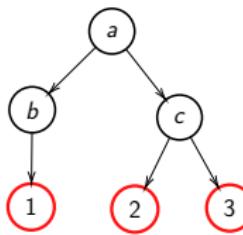


Figure: AABB Tree

# Computational Experiments

**Table:** Comparison with MILP model on limited set of boxes

<i>n</i>	MILP Model			PM		PS	
	Max Z	TT(s)	Gap(%)	Max Z	TT(s)	Max Z	TT(s)
1	85	0.01	0.00	85	0.00	85	0.00
2	85	0.07	0.00	85	0.00	85	0.00
3	85	0.13	0.00	85	0.00	85	0.00
4	85	0.20	0.00	85	0.01	85	0.01
5	85	2.02	0.00	85	0.02	85	0.02
6	158	90.58	0.00	158	0.06	158	0.05
7	158	1,369.24	0.00	158	0.07	158	0.08
8	161*	3,600.00	1.86	160	0.10	160	0.08
9	-	-	-	169	0.09	161	0.10
10	-	-	-	218	0.12	218	0.13
11	-	-	-	240	0.12	240	0.12
12	-	-	-	310	0.13	316	0.16
13	-	-	-	310	0.15	333	0.18
14	-	-	-	310	0.20	333	0.22
15	-	-	-	406	0.21	397	0.27
16	-	-	-	435	0.23	452	0.36
17	-	-	-	429	0.27	515	0.41
18	-	-	-	432	0.32	522	0.47
19	-	-	-	458	0.35	522	0.55
20	-	-	-	539	0.37	564	0.62

\* Some boxes had lower support than expected due to discretization errors.

# Computational Experiments

Table: Average execution time of literature results with bin gap

Heuristic		Execution Time (s)				Bin Gap (%)
		$n = 50$	$n = 100$	$n = 150$	$n = 200$	
<b>PM</b>	$k = 1$	0.05	0.11	0.28	0.55	4.57
	$k = 5$	0.08	0.39	1.02	2.16	4.32
	$k = 10$	0.15	0.74	1.98	4.12	4.29
	$k = 20$	0.29	1.45	3.89	8.07	4.05
	$k = 50$	0.72	3.63	9.72	20.47	3.95
<b>PS</b>	$k = 1$	0.04	0.18	0.51	1.08	4.35
	$k = 5$	0.12	0.74	2.19	4.79	4.01
	$k = 10$	0.23	1.43	4.19	9.39	3.94
	$k = 20$	0.47	2.81	8.48	18.93	3.74
	$k = 50$	1.15	6.74	21.03	45.78	3.52
<b>BRKGA-VD</b>		17.13	80.63	190.50	369.75	0.00

# Computational Experiments

**Table:** Summary of case study tests

k	PS			PM		
	TT (ms)	B	CR (%)	TT (ms)	B	CR (%)
1	423.87	1.37	65.87	65.18	1.31	<b>70.70</b>
5	1,597.54	1.34	69.19	185.22	1.29	<b>73.08</b>
10	2,627.52	1.32	70.35	344.90	1.27	<b>73.56</b>
20	5,373.79	1.34	70.78	620.95	1.27	<b>74.57</b>
50	14,203.10	1.31	72.11	1,279.96	1.29	<b>74.61</b>
100	26,934.21	1.31	73.23	2,340.37	1.26	<b>75.36</b>
200	48,944.90	1.30	73.89	4,465.78	1.25	<b>76.39</b>

# Computational Experiments

**Table:** Case study experiments trade off between average execution times and average cage ratio

k	PS		PM	
	CR* – CR (%)	TT – TT* (ms)	CR* – CR (%)	TT – TT* (ms)
1	10.56	358.69	5.73	0.00
5	7.24	1,532.36	3.35	120.04
10	6.08	2,562.34	2.87	279.72
<b>20</b>	5.65	5,308.61	<b>1.85</b>	<b>555.77</b>
<b>50</b>	4.32	14,137.92	<b>1.82</b>	<b>1,214.78</b>
100	3.20	26,869.03	1.07	2,275.19
200	2.54	48,879.72	0.04	4,400.60

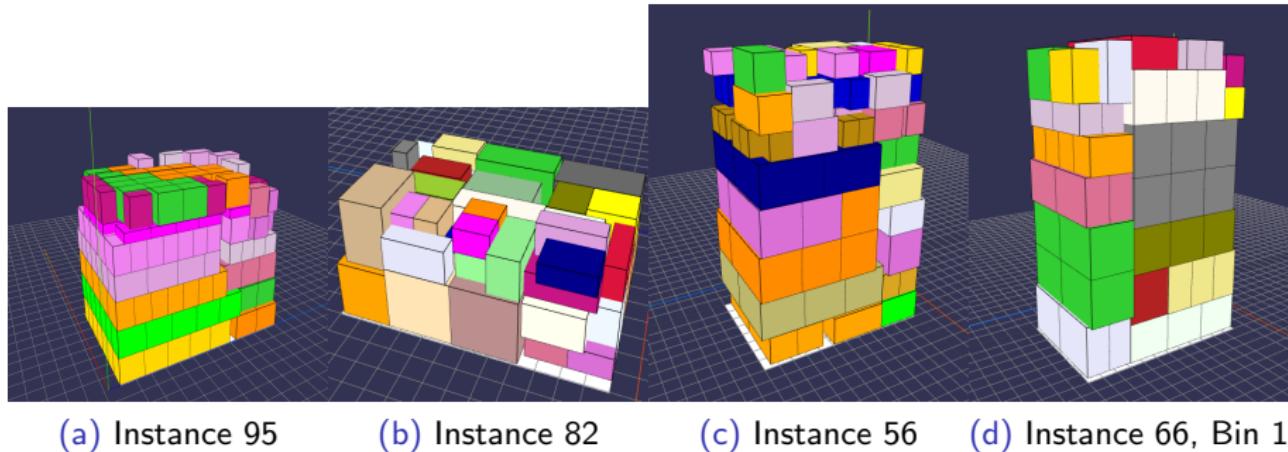


Figure: Solutions of case study tests with the "PM" placement and  $k = 200$

## Future Developments

- More practical constraints
- Solution optimization heuristics
- Improvements to the underlying 2D-BPP heuristic

Thank you for your attention!  
Questions?