CSP FAST TRACK

ARTURO

1. INTRODUCTION

Let R be a set of relation symbols. Let $\mathcal{A}$ be a relational structure over R. Let X be a countable set of variables. By the **constraint satisfaction problem** $\mathrm{CSP}(\mathcal{A})$[1] we mean the following decision problem: given a finite set Σ of atomic formulas over R, decide whether there is an assignment $(-)^{\mathcal{A}} : X \to A$ such that $\mathcal{A} \models \Sigma$; i.e. for all $r \in R_n$ and for all $x, y, x_1, \ldots, x_n \in X$

$$(1) \qquad r(x_1, \ldots, x_n) \in \Sigma \implies (x_1^{\mathcal{A}}, \ldots, x_n^{\mathcal{A}}) \in r^{\mathcal{A}}$$

$$(2) \qquad x \equiv y \in \Sigma \implies x^{\mathcal{A}} = y^{\mathcal{A}}$$

Clearly, it is enough to find an assignment only for those variables that appear in Σ.

Remark 1.1. Usually we deal with *indexed* relational structures, that is we fix a set of relation symbols R and we consider a set A with a set P of relations on A indexed by the elements of R. Sometimes is useful to deal with *non-indexed* relational structures (A, P). In this case P will serve as the index set as well. The same applies to algebraic structures (i.e. algebras) and function symbols.

Definition 1.2. Let $\rho \in A^k$, and $\varphi : A^n \to A$. We say that φ **preserves** ρ if given a matrix
$$\begin{bmatrix} a_1^1 & \cdots & a_n^1 \\ \vdots & \ddots & \vdots \\ a_1^k & \cdots & a_n^k \end{bmatrix}$$
with $(a_1^1, \ldots, a_1^k) \in \rho, \ldots, (a_n^1, \ldots, a_n^k) \in \rho$,
$$(\varphi(a_1^1, \ldots, a_n^1), \ldots, \varphi(a_1^k, \ldots, a_n^k)) \in \rho$$

Definition 1.3. If Γ is a set of relations on A and Φ is a set of operations on A we denote by

(1) $\mathrm{Inv}(\Phi)$ the set of relations on A that are preserved by all the elements of Φ;

(2) $\mathrm{Pol}(\Gamma)$ the set of operations on A that preserve all the elements of Γ.

Moreover, if $\mathbf{A}$ is an algebra and $\mathcal{A}$ a relational structure on the same set A:

(1) $\mathrm{Inv}(\mathbf{A})$ is the set of relations that are preserved by all the $f^{\mathbf{A}}$;

(2) $\mathrm{Pol}(\mathcal{A})$ is the set of operations that preserve all the $r^{\mathcal{A}}$.

---

[1]More often denoted by $\mathrm{CSP}(P)$.

Let $\mathbf{A}$ be an algebra. Let $\mathcal{A} := (A, \mathrm{Inv}(\mathbf{A}))$. By $\mathrm{CSP}(\mathbf{A})$ we mean the decision problem $\mathrm{CSP}(\mathcal{A})$.

Definition 1.4. Let F be a set of function symbols and $\mathbf{A}$ be an algebra over F. We denote by $\mathrm{Clo}(\mathbf{A})$ the smallest set containing

$$\{f^{\mathbf{A}} : f \in F\} \quad \text{and} \quad \{\pi_i^n : A^n \to A, 1 \leq i \leq n, n \in \omega\}$$

and closed under composition. The elements of $\mathrm{Clo}(\mathbf{A})$ are called **term** operations. We say that two algebras $\mathbf{A}$ and $\mathbf{B}$ on the same carrier are **term equivalent** if $\mathrm{Clo}(\mathbf{A}) = \mathrm{Clo}(\mathbf{B})$.

Recall that a variety is a class of algebras closed under homomorphic images, subalgebras and products. If V is a variety, we define $\mathrm{Clo}(\mathsf{V})$ to be $\mathrm{Clo}(\mathbf{F}_{\mathsf{V}}(\omega))$, where $\mathbf{F}_{\mathsf{V}}(\omega)$ is the V-free algebra generated by a countable number of generators. When V is Set, the variety of sets, $\mathrm{Clo}(\mathsf{Set})$ is the clone of projections $\{\pi_i^n : 1 \leq i \leq n, n \in \omega\}$, that we denote by $\mathbf{N}$.

Goal: prove

Theorem 1.5. *Let $\mathbf{A}$ be a finite idempotent algebra. Then the following are equivalent:*

(1) $\mathrm{CSP}(\mathbf{A})$ *is polynomial-time decidable;*

(2) $\mathrm{Clo}(\mathbf{A})$ *contains a weak near-unanimity operation;*

(3) *for every $\mathbf{B} \in HS(\mathbf{A})$, $\mathrm{Clo}(\mathbf{B}) \neq \mathbf{N}$.*

Otherwise, $\mathrm{CSP}(\mathbf{A})$ is NP-complete.

2. Kinds of Operations

Definition 2.1. An operation $\varphi : A^n \to A$ is called

(1) **essentially unary** if there is an index i and a non-constant function $\psi : A \to A$ such that
$$\varphi(a_1, \ldots, a_n) = \psi(a_i)$$
for all $a_1, \ldots, a_n \in A$.

(2) **idempotent** if $\varphi(a, \ldots, a) = a$ for all $a \in A$.

(3) a **near unanimity** operation if for all $a, b \in A$
$$\varphi(b, a, \ldots, a) = \varphi(a, b, \ldots, a) = \cdots = \varphi(a, \ldots, a, b) = a$$

Example 2.2. A ternary near unanimity operation is called a **majority** operation. For instance the ternary function defined as

$$\delta(a, b, c) = \begin{cases} b & \text{if } b = c \\ a & \text{otherwise} \end{cases}$$

is a majority operation called **dual discriminator**.

3. Relational Clones

Definition 3.1. Let R be a set of relation symbols and $\mathcal{A}$ be a relational structure over R. We denote by $\mathrm{Clo}(\mathcal{A})$ the smallest set containing

$$\{r^{\mathcal{A}} : r \in R\} \quad \text{and} \quad \{\Delta^{(n)} : n \in \omega\}$$

and closed under

(1) **permutation**: if $\rho \in \mathrm{Clo}(\mathcal{A})$, then also

$$\{(a_{\sigma(1)}, \ldots, a_{\sigma(n)}) : \sigma \in S_n, (a_1, \ldots, a_n) \in \rho\} \in \mathrm{Clo}\,\mathcal{A}$$

(2) **extension**: if $\rho \in \mathrm{Clo}(\mathcal{A})$, then also

$$\{(a_1, \ldots, a_n, a_{n+1}) : (a_1, \ldots, a_n) \in \rho, a_{n+1} \in A\} \in \mathrm{Clo}\,\mathcal{A}$$

(3) **truncation**: if $\rho \in \mathrm{Clo}(\mathcal{A})$, then also

$$\{(a_1, \ldots, a_{n-1}) : (a_1, \ldots, a_{n-1}, a_n) \in \rho, \text{ for some } a_n \in A\} \in \mathrm{Clo}\,\mathcal{A}$$

(4) intersection.

Remark 3.2. Observe that $\mathrm{Clo}(\mathcal{A})$ is given by all the relations ρ of A definable by a first-order primitive positive formula (that is, involving only conjunctions and existential quantifications). Recall that $\rho \subseteq A^n$ is definable if there is a formula $\varphi(x_1, \ldots, x_n)$ such that

$$\mathcal{A} \models \varphi(a_1, \ldots, a_n) \iff (a_1, \ldots, a_n) \in \rho$$

Theorem 3.3 ([2]). *For any pair of relational structures $\mathcal{A} = (A, \Gamma)$ and $\mathcal{B} = (A, H)$ such that H is finite and $H \subseteq \mathrm{Clo}(\mathcal{A})$, $\mathrm{CSP}(\mathcal{B})$ is polynomial-time reducible to $\mathrm{CSP}(\mathcal{A})$.*

Proof. Let Σ be a set of atomic formulas over H. Let $\eta(x_1, \ldots, x_n) \in \Sigma$. Then there is $\varphi(x_1, \ldots, x_n)$ of the form

$$\exists y_1, \ldots, y_m \, (\gamma_1(z_1^1, \ldots, z_{n_1}^1) \wedge \cdots \wedge \gamma_k(z_1^k, \ldots, z_{n_k}^k))$$

(where $\gamma_1, \ldots, \gamma_k \in \Gamma$ and $z_j^i \in \{x_1, \ldots, x_n, y_1, \ldots, y_m\}$) such that for $a_1, \ldots, a_n \in A$

$$(3) \qquad \mathcal{B} \models \eta(a_1, \ldots, a_n) \iff \mathcal{A} \models \varphi(a_1, \ldots, a_n)$$

We can assume (up to renaming of variables) that $y_1, \ldots, y_m$ do not appear in any formula of Σ.

Now, for each $\eta(x_1, \ldots, x_n) \in \Sigma$ perform the following steps:

- add $\{\gamma_1(z_1^1, \ldots, z_{n_1}^1), \ldots, \gamma_k(z_1^k, \ldots, z_{n_k}^k)\}$ to Σ;

- remove $\eta(x_1, \ldots, x_n)$ from Σ.

At the the end we obtain a set of equations T over G. This is a polynomial-time reduction. By (3) it is clear that we can find an assignment $X \to A$ such that $\mathcal{B} \models \Sigma$ iff we can find an assignment such that $\mathcal{A} \models T$. $\qquad \square$

Remark 3.4. We observe that in the above result it is not necessary that A is a finite set.

Corollary 3.5. *Let $\mathcal{A}$ be a relational structure and $\mathcal{B} = (A, \mathrm{Clo}(\mathcal{A}))$. Then*

(1) $\mathrm{CSP}(\mathcal{A})$ *is polynomial-time decidable iff* $\mathrm{CSP}(\mathcal{B})$ *is.*

(2) $\mathrm{CSP}(\mathcal{A})$ *is NP-complete iff* $\mathrm{CSP}(\mathcal{B})$ *is.*

4. Surjective and Idempotent Algebras

Definition 4.1. An algebra $\mathbf{A}$ is **surjective** if all the element of $\mathrm{Clo}(\mathbf{A})$ are surjective.

Remark 4.2. Let $\mathbf{A}$ be a finite algebra over F. Every element of $\mathrm{Clo}(\mathbf{A})$ is surjective iff every element of $\mathrm{Clo}_1(\mathbf{A})$ is. In this case $\mathrm{Clo}_1(\mathbf{A})$ is a group. It is enough to show that for every $f \in F$, there is $\varphi \in \mathrm{Clo}_1(\mathbf{A})$ such that $f^{\mathbf{A}}\varphi = 1_{\mathbf{A}}$. Let $m := |A|$. Then $(f^{\mathbf{A}})^{m!} = 1_{\mathbf{A}}$. Let n_f be the least n such that $(f^{\mathbf{A}})^{n_f} = 1_{\mathbf{A}}$. Let $\varphi := (f^{\mathbf{A}})^{n_f - 1}$.

Let $\mathbf{A}$ be a finite algebra and let B be a subset of A. Recall that if $\varphi : A^n \to A$, $\varphi|B$ is the function $B^n \to A$, $(b_1, \ldots, b_n) \mapsto \varphi(b_1, \ldots, b_n)$; Let

$$\mathrm{Clo}(\mathbf{A})|B := \{\varphi|B : \varphi \in \mathrm{Clo}(\mathbf{A}) \text{ and } \varphi[B^n] \subseteq B\}$$

We denote by $\mathbf{A}|B$ the algebra $(B, \mathrm{Clo}(\mathbf{A})|B)$.

Lemma 4.3 ([2])**.** *Let $\mathcal{A} = (A, P)$ be a relational structure. Let $\varphi \in \mathrm{Pol}_1(\mathcal{A})$. For $\rho \in P_n$ let*

$$\varphi(\rho) := \{(\varphi(a_1), \ldots, \varphi(a_n)) : (a_1, \ldots, a_n) \in \rho\}$$

Let $\mathcal{B} := (A, \varphi(P))$ where $\varphi(P) := \{\varphi(\rho) : \rho \in P\}$. Then

(1) $\mathrm{CSP}(\mathcal{A})$ *is polynomial-time decidable iff* $\mathrm{CSP}(\mathcal{B})$ *is;*

(2) $\mathrm{CSP}(\mathcal{A})$ *is NP-complete iff* $\mathrm{CSP}(\mathcal{B})$ *is.*

Proof. We show that $\mathrm{CSP}(\mathcal{A})$ is polynomial-time equivalent to $\mathrm{CSP}(\mathcal{B})$. Let Σ be a finite set of atomic formulas over P. Replace every occurrence of $\rho(x_1, \ldots, x_n)$ with $\varphi(\rho)(x_1, \ldots, x_n)$, obtaining a set of formulas T over $\varphi(P)$. Given any assignment $(-)^{\mathcal{B}} : X \to A$ such that $\mathcal{B} \models T$, define $(x)^{\mathcal{A}} := (x)^{\mathcal{B}}$; then $\mathcal{A} \models \Sigma$. Indeed, given $\rho(x_1, \ldots, x_n) \in \Sigma$, since $\varphi \in \mathrm{Pol}_1(\mathcal{A})$:

$$\begin{aligned}
\mathcal{B} \models \varphi(\rho)(x_1, \ldots, x_n) &\implies (x_1^{\mathcal{B}}, \ldots, x_n^{\mathcal{B}}) \in \varphi(\rho) \\
&\implies (x_1^{\mathcal{B}}, \ldots, x_n^{\mathcal{B}}) \in \rho \\
&\implies (x_1^{\mathcal{A}}, \ldots, x_n^{\mathcal{A}}) \in \rho \\
&\implies \mathcal{A} \models \rho(x_1, \ldots, x_n)
\end{aligned}$$

Conversely, given any assignment $(-)^{\mathcal{A}} : X \to A$ such that $\mathcal{A} \models \Sigma$, defining $(x)^{\mathcal{B}} := \varphi(x^{\mathcal{A}})$ is enough to have $\mathcal{B} \models T$. Indeed, given $\varphi(\rho)(x_1, \ldots, x_n) \in T$, by

definition of $\varphi(P)$:

$$
\begin{aligned}
\mathcal{A} \models \rho(x_1, \ldots, x_n) &\implies (x_1^{\mathcal{A}}, \ldots, x_n^{\mathcal{A}}) \in \rho \\
&\implies (\varphi(x_1^{\mathcal{A}}), \ldots, \varphi(x_n^{\mathcal{A}})) \in \varphi(\rho) \\
&\implies (x_1^{\mathcal{B}}, \ldots, x_n^{\mathcal{B}}) \in \varphi(\rho) \\
&\implies \mathcal{B} \models \varphi(\rho)(x_1, \ldots, x_n) \qquad \qquad \square
\end{aligned}
$$

Theorem 4.4. *Let* $\mathbf{A}$ *be a finite algebra. Then there is* $B \subseteq A$ *such that* $\mathbf{A}|B$ *is surjective and*

(1) $\mathrm{CSP}(\mathbf{A})$ *is polynomial-time decidable iff* $\mathrm{CSP}(\mathbf{A}|B)$ *is;*

(2) $\mathrm{CSP}(\mathbf{A})$ *is NP-complete iff* $\mathrm{CSP}(\mathbf{A}|B)$.

Proof. Assume that $\mathbf{A}$ is not surjective. Then by Remark 4.2 there is $\psi \in \mathrm{Clo}_1(\mathbf{A})$ not surjective. Let $\varphi \in \mathrm{Clo}_1(\mathbf{A})$ such that φ is not surjective and $\varphi[A]$ has minimal cardinality. Define $B := \varphi[A]$. We show that $\mathbf{A}|B$ is surjective. Let $\psi \in \mathrm{Clo}_1(\mathbf{A}|B) = \mathrm{Clo}_1(\mathbf{A})|B$; if, towards a contradiction, $\psi[B] \subset B$, then $\psi\varphi[A] \subset \varphi[A] \subset A$ contradicting the minimality. The last part of the statement follows immediately from Lemma 4.3. $\qquad \square$

Definition 4.5. Let $\mathbf{A}$ be an algebra. Let $\mathrm{id}(A)$ be the set of idempotent operations on A. We define $\mathrm{Clo}_{\mathrm{id}}(\mathbf{A}) := \mathrm{Clo}(\mathbf{A}) \cap \mathrm{id}(A)$. We say that $\mathbf{A}$ is **idempotent** of all the elements of $\mathrm{Clo}(\mathbf{A})$ are.

Remark 4.6. Let $\mathbf{A}$ be an algebra. Observe that an operation $\varphi \in \mathrm{Clo}(\mathbf{A})$ is idempotent iff it preserves the relations in $\Delta^{(1)} = \{\{a\} : a \in A\}$. Hence $\mathrm{Inv}(\mathrm{Clo}_{\mathrm{id}}(\mathbf{A})) = \mathrm{Clo}(\mathcal{A})$ where $\mathcal{A} = (A, \mathrm{Inv}(\mathbf{A}) \cup \Delta^{(1)})$.

Lemma 4.7. *Let* $A = \{a_1, \ldots, a_k\}$ *be a finite set. Let* $\mathbf{A}$ *be an algebra over* F. *Then the relation*

$$
(4) \qquad \qquad \sigma := \{(\psi(a_1), \ldots, \psi(a_k)) : \psi \in \mathrm{Clo}_1(\mathbf{A})\}
$$

belongs to $\mathrm{Inv}(\mathbf{A})$.

Proof. We show that for every $f \in F_n$ and for every matrix M

$$
\begin{bmatrix}
a_1^1 & \cdots & a_n^1 \\
\vdots & \ddots & \vdots \\
a_1^k & \cdots & a_n^k
\end{bmatrix}
$$

such that $(a_1^1, \ldots, a_1^k) \in \sigma, \ldots, (a_n^1, \ldots, a_n^k) \in \sigma$ we have

$$
(f^{\mathbf{A}}(a_1^1, \ldots, a_n^1), \ldots, f^{\mathbf{A}}(a_1^k, \ldots, a_n^k)) \in \sigma
$$

By hypothesis we can write M as

$$
\begin{bmatrix}
\psi_1(a_1) & \cdots & \psi_n(a_1) \\
\vdots & \ddots & \vdots \\
\psi_1(a_k) & \cdots & \psi_n(a_k)
\end{bmatrix}
$$

but then

$$(f^{\mathbf{A}}(a_1^1, \ldots, a_n^1), \ldots, f^{\mathbf{A}}(a_1^k, \ldots, a_n^k))$$
$$= (f^{\mathbf{A}}(\psi_1(a_1), \ldots, \psi_n(a_1)), \ldots, f^{\mathbf{A}}(\psi_1(a_k), \ldots, \psi_n(a_k)))$$
$$= (f^{\mathbf{A}}[\psi_1, \ldots, \psi_n](a_1), \ldots, f^{\mathbf{A}}[\psi_1, \ldots, \psi_n](a_k))$$

and we conclude since $f^{\mathbf{A}}[\psi_1, \ldots, \psi_n] \in \mathrm{Clo}_1(\mathbf{A})$. $\qquad\square$

Theorem 4.8. *Let $\mathbf{A} = A$ be a finite surjective algebra. Let $\mathbf{B} := (A, \mathrm{Clo}_{\mathrm{id}}(\mathbf{A}))$. Then*

(1) $\mathrm{CSP}(\mathbf{A})$ *is polynomial-time decidable iff* $\mathrm{CSP}(\mathbf{B})$ *is.*

(2) $\mathrm{CSP}(\mathbf{A})$ *is NP-complete iff* $\mathrm{CSP}(\mathbf{B})$ *is.*

Proof. Let $A = \{a_1, \ldots, a_k\}$ and let $\Gamma := \Delta^{(1)}$. Let $\mathcal{A} := (A, \mathrm{Inv}(\mathbf{A}))$ and $\mathcal{B} := (A, \mathrm{Inv}(\mathbf{A}) \cup \Gamma)$.

By definition and by Remark 4.6 $\mathrm{CSP}(\mathbf{B})$ is polynomial-time equivalent to $\mathrm{CSP}(\mathbf{A})$ iff $\mathrm{CSP}(\mathcal{B})$ is polynomial-time equivalent to $\mathrm{CSP}(\mathcal{A})$.

That $\mathrm{CSP}(\mathcal{A})$ is polynomial-time reducible to $\mathrm{CSP}(\mathcal{B})$ is obvious. Let Σ be a finite set of atomic formulas over $\mathrm{Inv}(\mathbf{A}) \cup \Gamma$ and let $\{x_1, \ldots, x_k\}$ be variables that do not appear in Σ. By Remark 4.2, since $\mathbf{A}$ is surjective, $\mathrm{Clo}_1(\mathbf{A})$ forms a group. Moreover, the relation σ of Lemma 4.7 belongs to $\mathrm{Inv}(\mathbf{A})$. Now, perform the following steps:

- replace every formula $\gamma_i(x)$ with $x \equiv x_i$;

- add the formula $\sigma(x_1, \ldots, x_k)$.

At the the end we obtain a set of equations T over $\mathrm{Inv}(\mathbf{A})$. This is a polynomial-time reduction. We finally show that we can find an assignment such that $\mathcal{A} \models T$ iff we can find an assignment such that $\mathcal{B} \models \Sigma$. Let $(-)^{\mathcal{B}} : X \to A$ be an assignment such that $\mathcal{B} \models \Sigma$. Consider the assignment

$$x^{\mathcal{A}} = \begin{cases} x^{\mathcal{B}} & \text{if } x \neq x_i \\ a_i & \text{if } x = x_i \end{cases}$$

Then $(-)^{\mathcal{A}}$ is such that $\mathcal{A} \models T$. Conversely, assume that there is an assignment $(-)^{\mathcal{A}}$ such that $\mathcal{A} \models T$. By definition of σ, there is $\psi \in \mathrm{Clo}_1(\mathbf{A})$ such that $x_i^{\mathcal{A}} = \psi(a_i)$ for all i. Consider $(-)'^{\mathcal{A}} := \psi^{-1} \circ (-)^{\mathcal{A}}$. Define $(x)^{\mathcal{B}} := (x)'^{\mathcal{A}}$. Let $\rho(y_1, \ldots, y_n), \gamma_i(x) \in \Sigma$. Every relation in $\mathrm{Inv}(\mathbf{A})$ is invariant under ψ^{-1}, hence

$$\mathcal{A} \models \rho(y_1, \ldots, y_n) \implies (y_1^{\mathcal{A}}, \ldots, y_n^{\mathcal{A}}) \in \rho \qquad\qquad \mathcal{A} \models x \equiv x_i \implies x^{\mathcal{A}} = x_i^{\mathcal{A}}$$
$$\implies (\psi^{-1}(y_1^{\mathcal{A}}), \ldots, \psi^{-1}(y_n^{\mathcal{A}})) \in \rho \qquad\qquad \implies x^{\mathcal{A}} = \psi(a_i)$$
$$\implies (y_1^{\mathcal{B}}, \ldots, y_n^{\mathcal{B}}) \in \rho \qquad\qquad\qquad\qquad \implies \psi^{-1}(x^{\mathcal{A}}) = a_i$$
$$\implies \mathcal{B} \models \rho(y_1, \ldots, y_n) \qquad\qquad\qquad\qquad \implies x^{\mathcal{B}} = a_i$$

Thus $\mathcal{B} \models \Sigma$. $\qquad\square$

Remark 4.9. By the above result, given a finite algebra $\mathbf{A}$, to the purpose of the study of $\mathrm{CSP}(\mathbf{A})$, we can assume without loss of generality that $\mathbf{A}$ is idempotent.

5. Subalgebras and Images

Theorem 5.1. *Let* **A** *be a finite algebra.*

(1) *if* CSP(**A**) *is polynomial-time decidable, so is* CSP(**B**) *for every* **B** $\leq$ **A***;*

(2) *if there is* **B** $\leq$ **A** *such that* CSP(**B**) *is NP-complete, so is* CSP(**A**)*.*

Proof. Let $\mathcal{A} = (A, \text{Inv}(\mathbf{A}))$ and $\mathcal{B} = (A, \text{Inv}(\mathbf{B}))$. By definition CSP(**B**) is polynomial-time reducible to CSP(**A**) iff CSP($\mathcal{B}$) is polynomial-time reducible to CSP($\mathcal{A}$). But that CSP($\mathcal{B}$) is polynomial-time reducible to CSP($\mathcal{A}$) is obvious since $\text{Inv}(\mathbf{B}) \subseteq \text{Inv}(\mathbf{A})$. $\square$

Theorem 5.2. *Let* **A**, **B** *be two finite algebras of the same type.*

(1) *if* CSP(**A**) *is polynomial-time decidable, so is* CSP(**B**) *for every surjective homomorphism* $\alpha : \mathbf{A} \to \mathbf{B}$*;*

(2) *if there is a surjective homomorphism* $\alpha : \mathbf{A} \to \mathbf{B}$ *such that* CSP(**B**) *is NP-complete, so is* CSP(**A**)*.*

Proof. Let $\mathcal{B} = (B, \text{Inv}(\mathbf{B}))$. We show that there is $\Gamma \subseteq \text{Inv}(\mathbf{B})$ such that, defining $\mathcal{A} = (A, \Gamma)$, CSP($\mathcal{B}$) is polynomial-time reducible to CSP($\mathcal{A}$). For every $\rho \in \text{Inv}(\mathbf{B})_n$ let

$$\alpha^{-1}(\rho) := \{(a_1, \ldots, a_n) \in A^n : (\alpha(a_1), \ldots, \alpha(a_n)) \in \rho\}$$

Clearly, $\alpha^{-1}(\rho) \in \text{Inv}(\mathbf{A})_n$ and therefore, letting $\Gamma := \{\alpha^{-1}(\rho) : \rho \in \Gamma\}$, $\Gamma \subseteq \text{Inv}(\mathbf{A})$. Let Σ be a set of atomic formulas over $\text{Inv}(\mathbf{B})$. Replace every formula $\rho(x_1, \ldots, x_n)$ with $\alpha^{-1}(\rho)(x_1, \ldots, x_n)$. We obtain a set of equations T over Γ by a polynomial-time reduction. Let β be a section of α.
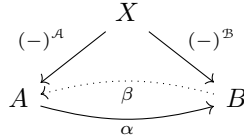
Referring to the assignments in the picture, each defined in terms of the other so that the diagram commute, it is clear that $\mathcal{A} \models T$ iff $\mathcal{B} \models \Sigma$. $\square$

Lemma 5.3 ([2])**.** *Let* $\mathcal{A}$ *be a relational structure. If* $\text{Pol}(\mathcal{A})$ *contains essentially unary operations only,* CSP($\mathcal{A}$) *is NP-complete.*

Corollary 5.4. *Let* **A** *be a finite algebra. If there is* **B** $\in HS(\mathbf{A})$ *such that* $\text{Clo}(\mathbf{B}) = \mathbf{N}$*, then* CSP(**A**) *is NP-complete.*

Proof. If **B** is such that $\text{Clo}(\mathbf{B}) = \mathbf{N}$, then CSP(**B**) is NP-complete by Lemma 5.3. We conclude using the second cluases of Theorems 5.1 and 5.2. $\square$

6. Omitting Types and Complexity

Refer to the Appendices.

Theorem 6.1. *Let* $\mathbf{A}$ *be a finite idempotent algebra. If* $\mathrm{CSP}(\mathbf{A})$ *is decidable in polynomial-time, then* $\mathbf{1} \notin \mathrm{typ}\{\mathbf{A}\}$.

Proof. Assume for the sake of contradiction that $\mathbf{1} \in \mathrm{typ}\{\mathbf{A}\}$. We show that $\mathrm{CSP}(\mathbf{A})$ is NP-complete. If $\mathbf{1} \in \mathrm{typ}\{\mathbf{A}\}$, then there are $\alpha, \beta \in \mathrm{Con}(\mathbf{A})$ such that (α, β) is tame and $\mathrm{typ}(\alpha, \beta) = \mathbf{1}$. Let N be a (α, β)-trace. Then $\mathbf{M} := (\mathbf{A}||N)/(\alpha|N)$ is minimal and $\mathrm{typ}(\mathbf{M}) = \mathbf{1}$. Thus $\mathbf{M}$ is polynomial equivalent to a $\mathbf{G}$-set for some finite group $\mathbf{G}$. By construction $\mathbf{M}$ is term equivalent to a $\mathbf{G}$-set, and therefore, by Theorem 5.3, NP-complete.

By definition there is $U \in \mathrm{M}_{\mathbf{A}}(\alpha, \beta)$ with $N \subseteq U$. By Lemma **??**, there is $\varepsilon \in \mathrm{E}(\mathbf{A})$ such that $\varepsilon[A] = U$. By Lemma **??**, $\mathbf{M} = ((\mathbf{A}||U)||N)/((\alpha|U)|N)$. Consider the algebra $\mathbf{A}' = (A, \mathrm{Clo}(\mathbf{A}) \cup \{\varepsilon\})$. Let $\mathbf{M}' := (\mathbf{A}'||N)/(\alpha|N)$.

$$
\begin{aligned}
\mathrm{CSP}(\mathbf{M}) \text{ NP-complete} \implies & \ \mathrm{CSP}(\mathbf{M}') \text{ NP-complete} && \text{obvious} \\
\implies & \ \mathrm{CSP}((\mathbf{A}'||U)||N) \text{ NP-complete} && \text{by Thm. 5.2} \\
\implies & \ \mathrm{CSP}((\mathbf{A}'||U)|N) \text{ NP-complete} && \text{same clone} \\
\implies & \ \mathrm{CSP}(\mathbf{A}'||U) \text{ NP-complete} && \\
\implies & \ \mathrm{CSP}(\mathbf{A}'|U) \text{ NP-complete} && \text{same clone} \\
\implies & \ \mathrm{CSP}(\mathbf{A}') \text{ NP-complete} && \text{by proof of Thm. 4.4} \\
\implies & \ \mathrm{CSP}(\mathbf{A}) \text{ NP-complete} && \mathrm{Clo}(\mathbf{A}) \subseteq \mathrm{Clo}(\mathbf{A}')
\end{aligned}
$$

The third implication is justified as follows: since N is an (α, β)-trace is closed under the operations of $\mathbf{A}'$,

$$
\begin{aligned}
\mathrm{Clo}((\mathbf{A}'||U)|N) &= \mathrm{Clo}(\mathrm{Pol}(\mathbf{A}'|U)|N) \\
&= \mathrm{Clo}(\mathrm{Pol}(\mathbf{A}'|U)) \\
&= \mathrm{Pol}(\mathbf{A}'|U) \\
&= \mathrm{Clo}(\mathbf{A}'||U). \qquad\qquad \square
\end{aligned}
$$

References

[1] Bulatov, A. A., Jeavons, P., Krokhin, A. A. (2005). Classifying the complexity of constraints using finite algebras. *SIAM Journal of Computation* 34(3), 720-742.

[2] Jeavons, P. (1998). On the algebraic structure of combinatorial problems, *Theoretical Computer Science* 200, 185-204.

[3]