# 6.01: Introduction to EECS I

# Search Algorithms

Week 12

November 24, 2009

# Looking Ahead

- Reaction: Use a rule to determine the 'action' to take, as a direct function of the state
  - wall-following
  - proportional controller
- **Planning:** Choose action based on 'looking ahead': exploring alternative sequences of actions

Methods for planning require us to specify an explicit model of the effects of our actions in the world.

# Using models to choose actions

Assume states and actions are discrete.

### Given

- · A state-machine model of the world
- A start state
- A goal test

Find a sequence of actions (inputs to the state machine) to reach a goal state from the start state.

# **Application: Navigation**

What are good definitions of states, actions?

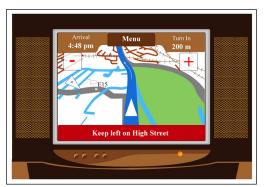
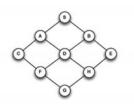


Figure by MIT OpenCourseWare.

What makes a path good?

### Abstraction: Labeled graph

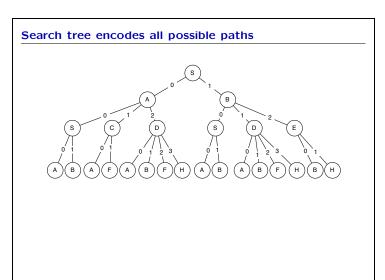


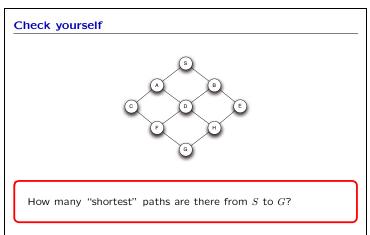
Lots of possible paths! Could enumerate them and evaluate each one. Too hard...

Assume  $additive\ cost.$  For now, each segment has a cost of one. We want to find the shortest path.

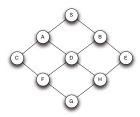
### Formal model

- States: {'S', 'A', 'B', 'C', 'D', 'E', 'F', 'G', 'H'}.
- Starting state is 'S'.
- Goal test: lambda x: x == 'H'
- Legal actions and successors:





# Check yourself



How many "shortest" paths are there from S to G?

6

# A numeric example

- States: integers
- Start state: 1
- Legal actions (and successors) in state n:  $\{2n, n+1, n-1, n^2, -n\}$
- Goal test: x = 10

How long is the longest path in this domain?

# Search trees in Python

```
Node in search tree, not the same as a state!
```

class SearchNode:

```
def __init__(self, action, state, parent):
    self.state = state
    self.action = action
```

self.parent = parent

Search node encodes a whole path

def path(self):

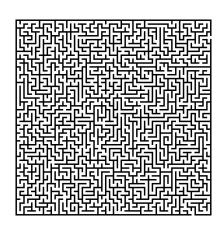
```
if self.parent == None:
    return [(self.action, self.state)]
```

else:

return self.parent.path() + \

[(self.action, self.state)]

# Finding your way



# Search algorithm

Until we find the goal or the agenda is empty:

- Extract a node from the agenda
- Expand it (find its children)
- Add its children to the agenda (visit its children)

**Goal:** Search as few nodes as possible while guaranteeing that we still find the shortest path.

```
Return a path:
```

```
[(None,S_0),(A_1,S_1),(A_2,S_2),...,(A_n,S_n)] where S_n satisfies the goal test.
```

# Search in Python

```
def search(initialState, goalTest, actions, successor):
   if goalTest(initialState):
      return [(None, initialState)]
```

### Search in Python

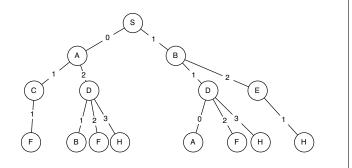
```
def search(initialState, goalTest, actions, successor):
   if goalTest(initialState):
       return [(None, initialState)]
   agenda = [SearchNode(None, initialState, None)]
   while not empty(agenda):
```

### Search in Python

```
def search(initialState, goalTest, actions, successor):
   if goalTest(initialState):
        return [(None, initialState)]
   agenda = [SearchNode(None, initialState, None)]
   while not empty(agenda):
        parent = getElement(agenda)
        for a in actions:
            newS = successor(parent.state, a)
            newN = SearchNode(a, newS, parent)
        if goalTest(newS):
            return newN.path()
        else:
            add(newN, agenda)
        return None
```

# Don't be totally stupid!

Pruning Rule 1. Don't consider any path that visits the same state twice.

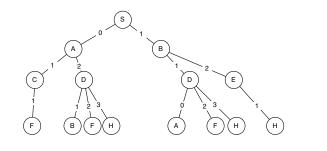


# Not being totally stupid, in Python

```
def search(initialState, goalTest, actions, successor):
    \verb|if goalTest(initialState)|:\\
        return [(None, initialState)]
    agenda = [SearchNode(None, initialState, None)]
    while not empty(agenda):
        parent = getElement(agenda)
        for a in actions:
            newS = successor(parent.state, a)
            newN = SearchNode(a, newS, parent)
            if goalTest(newS):
                return newN.path()
            elif parent.inPath(newS):
                pass
            else:
               add(newN, agenda)
    return None
```

# Another pruning rule

Pruning Rule 2. If there are multiple actions that lead from a state r to a state s, consider only one of them.



### Stack and Queue using Lists class Stack: def \_\_init\_\_(self): self.data = [] def push(self, item): self.data.append(item) def pop(self): return self.data.pop() def isEmpty(self): return self.data is [] class Queue: def \_\_init\_\_(self): self.data = [] def push(self, item): self.data.append(item) def pop(self): return self.data.pop(0) def isEmpty(self): return self.data is []

## Stack data structure

```
Last in, first out

>>> s = Stack()
>>> s.push(1)
>>> s.push(9)
>>> s.push(3)
>>> s.pop()
3
>>> s.pop()
9
>>> s.push(-2)
>>> s.pop()
-2
```

# Queue data structure

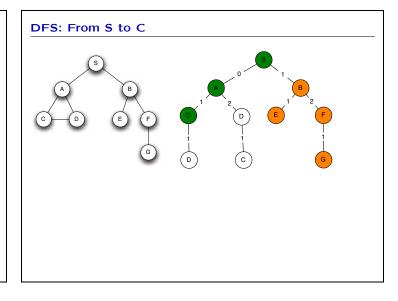
```
First in, first out

>>> q = Queue()
>>> q.push(1)
>>> q.push(9)
>>> q.push(3)
>>> q.pop()
1

>>> q.pop()
9
>>> q.push(-2)
>>> q.pop()
3
```

### Depth-First search

```
{\tt def \ depthFirstSearch(initialState, \ goalTest, \ actions, \ successor):}
    agenda = Stack()
    if goalTest(initialState):
       return [(None, initialState)]
    agenda.push(SearchNode(None, initialState, None))
    while not agenda.isEmpty()
        parent = agenda.pop()
        newChildStates = []
        for a in actions:
            newS = successor(parent.state, a)
            newN = SearchNode(a, newS, parent)
            if goalTest(newS):
                return newN.path()
            elif newS in newChildStates: # pruning rule 2
                pass
            elif parent.inPath(newS):
                                           # pruning rule 1
                pass
            else:
                {\tt newChildStates.append(newS)}
                agenda.push(newN)
    return None
```



### **DFS** properties

- May run forever if we don't apply pruning rule 1.
- May run forever in an infinite domain.
- Doesn't necessarily find the shortest path.
- Efficient in the amount of space it requires to store the agenda.

### Breadth-First search $\tt def\ breadthFirstSearch(initialState,\ goalTest,\ actions,\ successor):$ agenda = Queue() $\hbox{if goalTest(initialState):}\\$ return [(None, initialState)] agenda.push(SearchNode(None, initialState, None)) while not agenda.isEmpty(): parent = agenda.pop() newChildStates = [] for a in actions: newS = successor(parent.state, a) newN = SearchNode(a, newS, parent) if goalTest(newS): return newN.path() elif newS in newChildStates: pass elif parent.inPath(newS): pass

# BFS: From S to G

## **BFS** properties

return None

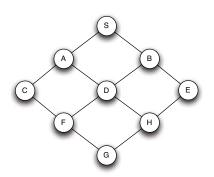
else:

- Always returns a shortest path to a goal state, if a goal state exists in the set of states reachable from the start state.
- May run forever in an infinite domain if there is no solution.
- Requires more space than depth-first search.

newChildStates.append(newS)
agenda.push(newN)

### **Dynamic Programming**

When happened when we did BFS in this city with goal G?



Visits 16 nodes, but there are only 9 states!!

# **Dynamic Programming Principle**

The  $\mathit{shortest}$  path from X to Z that goes through Y is made up of

- ullet the *shortest* path from X to Y and
- the *shortest* path from Y to Z.

So, we only need to remember the *shortest* path from the start state to each other state.

### DP in breadth-first search

The  $\mathit{first}$  path that BFS finds from start to X is the  $\mathit{shortest}$  path from start to X.

So, we only need to remember the *first* path we find from the start state to each other state.

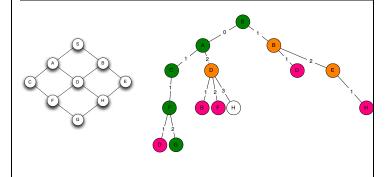
# DP as a pruning technique

Pruning Rule 3. Don't consider any path that visits a state that you have already visited via some other path.

### BFS with DP

```
{\tt def\ breadthFirstDP(initialState,\ goalTest,\ actions,\ successor):}
    agenda.push(SearchNode(None, initialState, None))
    visited = {initialState: True}
    while not agenda.isEmpty():
        parent = agenda.pop()
        newChildStates = []
        for a in actions:
           newS = successor(parent.state, a)
            newN = SearchNode(a, newS, parent)
            if goalTest(newS):
                return newN.path()
            elif visited.has_key(newS): # rules 1, 2, 3
                pass
            else:
                visited[newS] = True:
                newChildStates.append(newS)
                agenda.push(newN)
    return None
```

### BFS-DP: From S to G



Visits 9 states.

Can never expand more nodes than there are states.

Can be used with DFS as well.

### State machines as world models

Add two new features:

- done(self, state): returns True if the machine has terminated;
   we can use this as a goal test
- legalInputs: list of possible legal inputs to the machine; we can
  use this as the set of possible actions

For now, we will ignore the output of the state machine. Later we will put it to good use.

### Planning in a state machine

Question: Given a state machine in its initial state, what sequence of inputs can we feed to it, in order to cause it to enter a done state?

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### A numeric example

- States: integers
- Start state: 1
- Legal actions (and successors) in state n:  $\{2n, n+1, n-1, n^2, -n\}$
- Goal test: x = 10

# A numeric example – state machine

```
class NumberTestSM(sm.SM):
    startState = 1
    legalInputs = ['x*2', 'x+1', 'x-1', 'x**2', '-x']
    def __init__(self, goal):
       self.goal = goal
    def nextState(self, state, action):
       if action == 'x*2':
           return state*2
        elif action == 'x+1':
           return state+1
        elif action == 'x-1':
           return state-1
        elif action == 'x**2':
           return state**2
        elif action == '-x':
           return -state
   def getNextValues(self. state. action):
       nextState = self.nextState(state, action)
        return (nextState, nextState)
    def done(self, state):
       return state == self.goal
```

### Numeric - Breadth First

```
>>> smSearch(NumberTestSM(10), initialState = 1,
            depthFirst = False, DP = False)
   expanding: 1
   expanding: 1-x*2->2
   expanding: 1-x-1->0
   expanding: 1--x->-1
  expanding: 1-x*2->2-x*2->4
   expanding: 1-x*2->2-x+1->3
   expanding: 1-x*2->2--x->-2
   expanding: 1-x-1->0-x-1->-1
   expanding: 1--x->-1-x*2->-2
   expanding: 1--x->-1-x+1->0
   expanding: 1-x*2->2-x*2->4-x*2->8
   expanding: 1-x*2->2-x*2->4-x+1->5
33 states visited
[(None, 1), ('x*2', 2), ('x*2', 4), ('x+1', 5), ('x*2', 10)]
```

### Numeric - Breadth First with DP

# Numeric – Depth First

# Computational complexity

### Let

- *b* be the *branching factor* of the graph; that is, the number of successors a node can have.
- *d* be the *maximum depth* of the graph; that is, the length of the longest path in the graph.
- l be the solution depth of the problem; that is, the length of the shortest path from the start state to the shallowest goal state.
- n be the state space size of the graph; that is the total number of states in the domain.

There are  $b^d$  paths at depth d.

The number of nodes in the tree of depth d is about  $b^{d+1}$ .

# Without dynamic programming

- Depth first:
  - may have to search every path  $(b^{d+1} \ \mathsf{nodes})$ , but
  - agenda is small (bd)
- Breadth first:
  - may have to search to depth l ( $b^{l+1}$  nodes),
  - $-\,\,$  agenda may be as large as  $b^l$

# With dynamic programming

Visit at most n states!

Sometimes  $n<< b^l$  (in a road network, for example), sometimes not (small problem in large space).

DP is almost always an improvement in running time.

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