



Memory Layout

Linux Userspace Process Memory Layout

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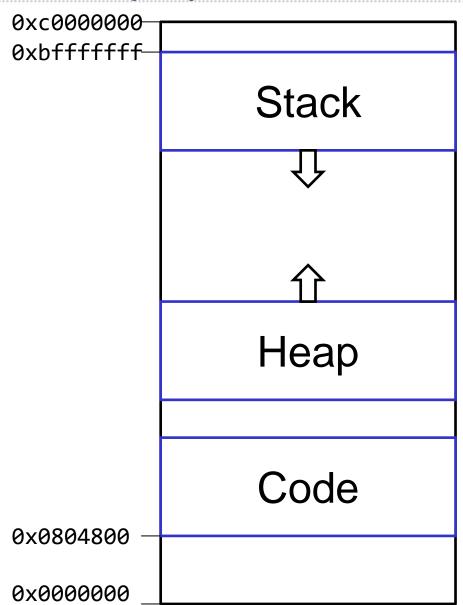
Userspace Memory Layout

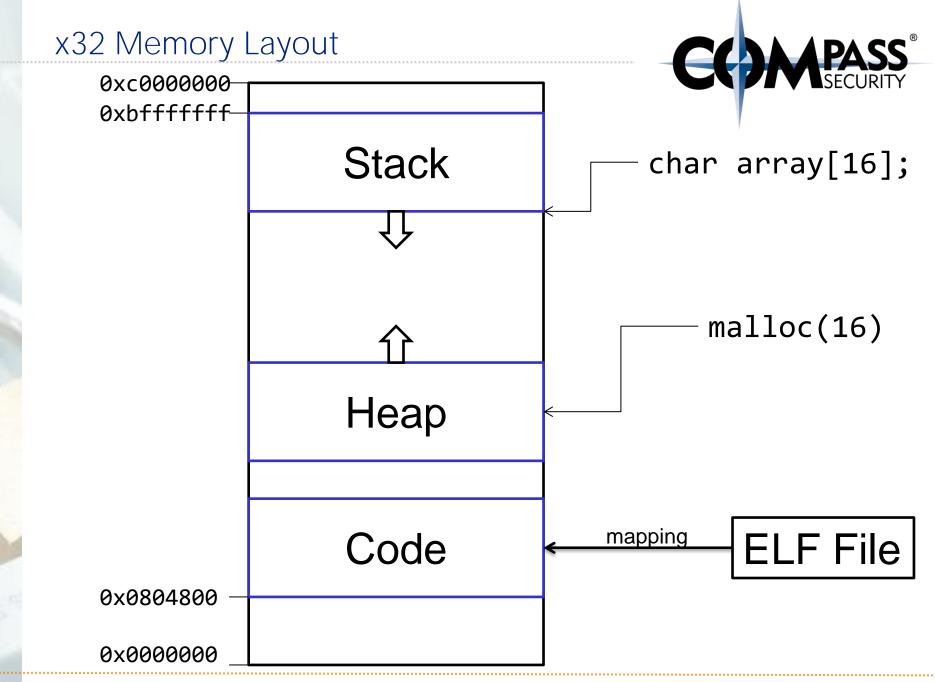
In x32

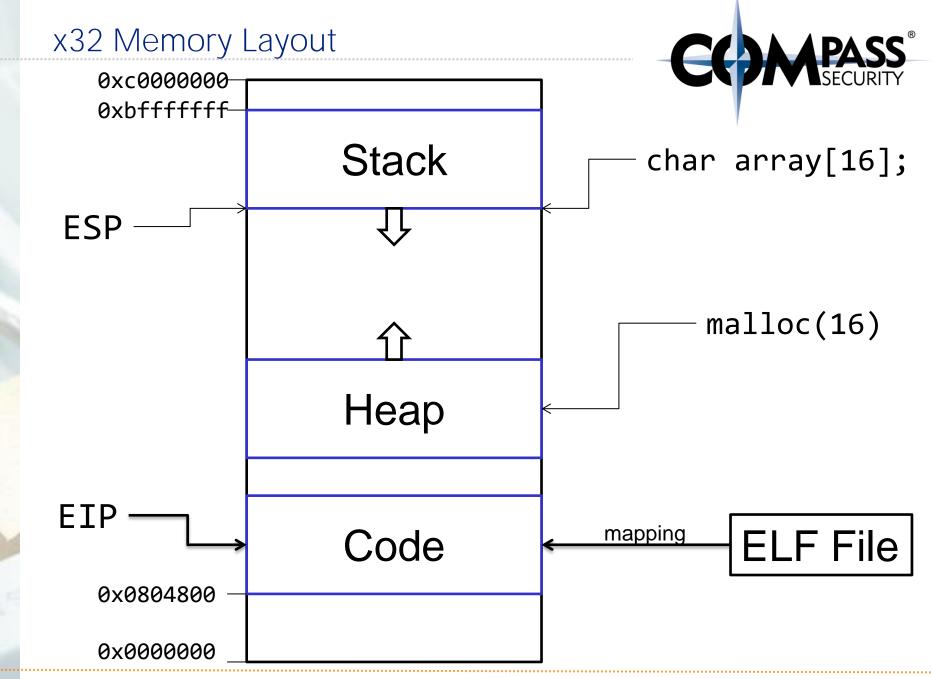
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x32 Memory Layout









x32 Memory Layout



Memory regions:

Stack

- There's one contiguous memory region containing the stack for the process
- ◆ LIFO Last in, First Out
- Contains function local variables
- Also contains: Saved Instruction Pointer (SIP)
- Current function adds data to the top (bottom) of the stack

Heap

- → There's one contiguous memory region containing the heap
- Memory allocator returns specific pieces of the memory region
- ✦ For malloc()
- Also contains: heap management data

x32 Memory Layout



Memory regions:

Code

→ Compiled program code





How do programs on disk look like



Programs are stored in ELF files

ELF: Executable and Linkable Format

- Previously: "a.out" (Linux 1.2)
- → Like COFF, PE (EXE), COM, ...

ELF types:

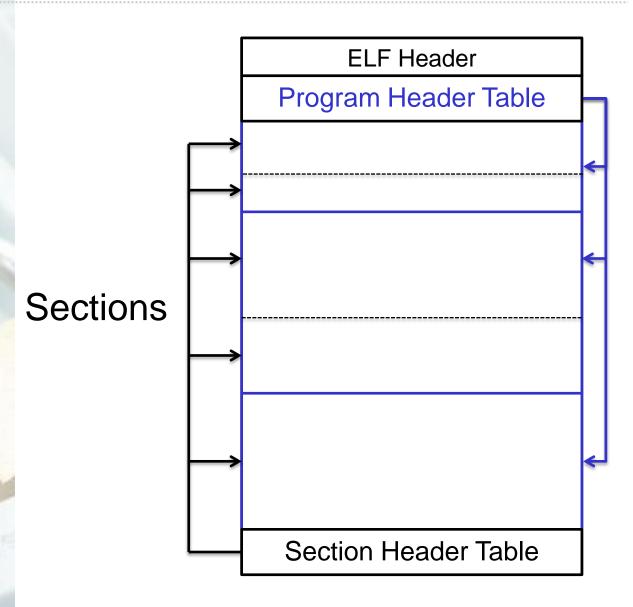
- ★ ET_EXEC: Executable File
- ★ ET_REL: Relocatable File
- → ET_DYN: Shared Object File

ELF "views":

- ★ Sections
- → Segments

\$ readelf -I < binary>





Segments



Program Headers:

Offset	VirtAddr	PhysAddr	
FileSiz	MemSiz	Flags	Align
0x0000040	0x0000400040	0x00000	000000400040
0x00001c0	0x0000001c0	R E	8
0x00000200	0x0000400200	0x000000000400200	
0x000001c	0x00000001c	R	1
0x0000000	0x0000400000	0x00000	00000400000
0x00000b24	0x000000b24	R E	200000
0x00000b28	0x0000600b28	0x00000	00000600b28
0x00000270	0x000000278	RW	200000
0x00000b40	0x0000600b40	0x0000000000600b40	
0x000001e0	0x0000001e0	RW	8
0x0000021c	0x000040021c	0x000000000040021c	
0x0000044	0x000000044	R	4
0x000009ac	0x00004009ac	0x00000000004009ac	
0x0000044	0x000000044	R	4
0x0000000	0x000000000	0x00000000000000000	
0x0000000	0x000000000	RW	10
	FileSiz 0x0000001c0 0x000001c0 0x0000001c 0x00000000 0x00000000 0x00000024 0x000000270 0x00000040 0x0000001e0 0x0000001c 0x0000001c	FileSiz MemSiz 0x00000040 0x00000400040 0x0000001c0 0x00000001c0 0x000000200 0x00000400200 0x00000001c 0x000000001c 0x00000000 0x00000400000 0x000000b24 0x0000000b24 0x000000b28 0x000000000278 0x000000b40 0x0000000001e0 0x0000001e0 0x000000001e0 0x00000044 0x0000000044 0x00000044 0x0000000044 0x00000044 0x0000000044 0x000000044 0x0000000000	FileSiz MemSiz Flags 0x00000040 0x0000400040 0x000000 0x000001c0 0x00000001c0 R E 0x0000001c 0x000000001c R 0x00000000 0x000000000 0x0000000 0x00000000 0x000000000 0x000000 0x000000024 0x0000000024 R E 0x000000270 0x00000000278 RW 0x00000040 0x000000001e0 RW 0x0000001e0 0x000000001e0 RW 0x00000044 0x0000000044 R 0x00000044 0x000000044 R 0x00000044 0x0000000044 R 0x00000004 0x000000000 0x0000000000000



\$ readelf -1 challenge0

Section to Segment mapping:

```
Segment Sections...
 00
 01
        .interp
        .interp .note.ABI-tag .note.gnu.build-id .gnu.hash
 02
        .dynsym .dynstr .gnu.version .gnu.version r
        .rela.dyn .rela.plt .init .plt .text .fini .rodata
        .eh frame hdr .eh frame
 03
        .init array .fini array .jcr .dynamic .got .got.plt
        .data .bss
 0.4
        .dynamic
 05
        .note.ABI-tag .note.gnu.build-id
 06
        .eh frame hdr
 07
```



Sections:

- → .text: Executable instructions
- → .bss: Unitialized data (usually the heap)
- → .data: initialized data
- → .rodata: Read-Only data
- → .got: Global Offset Table
- → .plt: Procedure Linkage Table
- → .init/.fini: Initialization instructions ("glibc")



Program Headers:

	Type	Offset	PhysAddr	
		FileSiz	Flags	Align
(02)	LOAD	0x00000000000000	0x0000	000000400000
		0x000000000000b24	R E	200000
(03)	LOAD	0x000000000000b28	0x0000	000000600ь28
		0x000000000000270	RW	200000
(07)	GNU_STACK	0x00000000000000	0x0000	000000000000
		0x00000000000000	RW	10

```
.init .plt .text .fini .rodata
```

.got .got.plt .data .bss

07





ELF Loader

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ELF Header
Program Header Table
.plt
.text
.init
.got
.data
.bss
Section Header Table

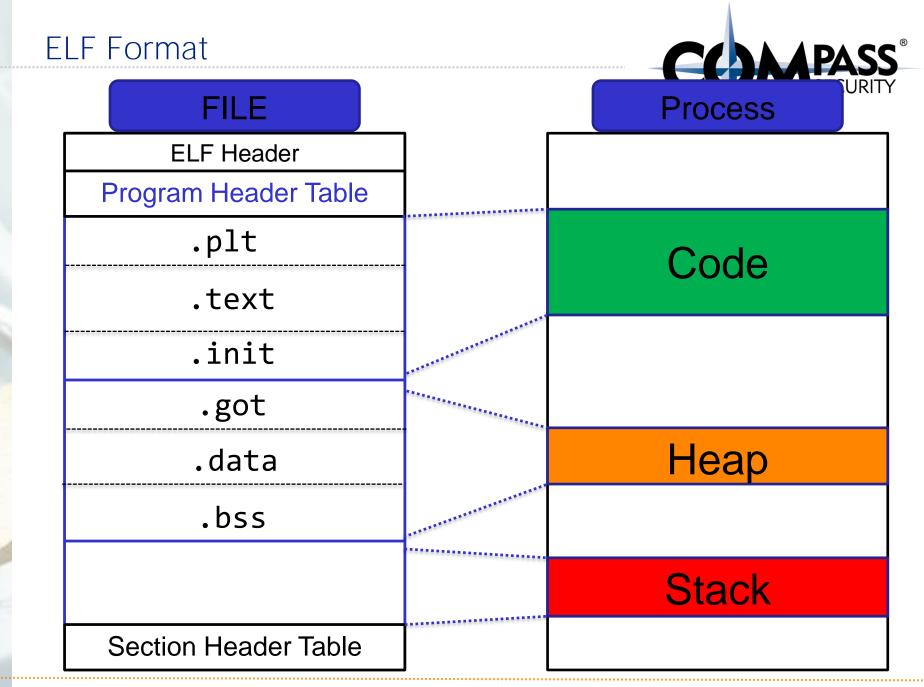
02 Executable Segment r-x

03 Data Segment

rw-

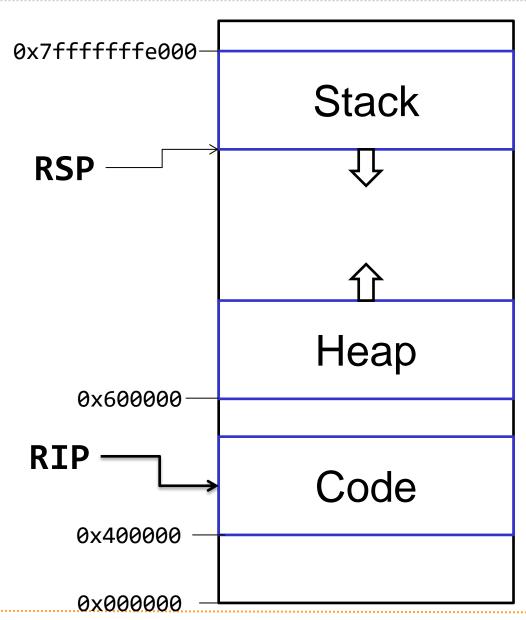
07 Stack

rw-



x64 Memory Layout







Lets do an example some static and dynamic binary analysis

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```
char *globalVar = "Global";
void main(void) {
        char stackVar[16];
        char *heapVar = (char *) malloc(4);
        printf("Global var: %p\n", globalVar);
        printf("Heap var: %p\n", heapVar);
        printf("Stack var: %p\n", stackVar);
```



Global var: 0x400654

Heap var: 0x601010

Stack var: 0x7ffffffe990

(2) LOAD $0 \times 0000000000400000$

R E 200000

(3) LOAD 0x000000000000000028

RW 200000

(7) **GNU_STACK** 0x00000000000000

RW 10



See it at runtime

```
# cat /proc/self/maps
```

```
      00400000-0040c000
      r-xp
      000000000
      08:01
      391694
      /bin/cat

      0060b000-0060c000
      r--p
      0000b000
      08:01
      391694
      /bin/cat

      0060c000-0060d000
      rw-p
      0000c000
      08:01
      391694
      /bin/cat
```

•••

7ffffffde000-7ffffffff000 rw-p 00000000 00:00 0 [stack]



Show Code section, and disassemble:

```
$ objdump -d ./challenge1
./challenge1: file format elf64-x86-64
Disassembly of section .init:
0000000000400588 < init>:
000000000040077f <handleData>:
  40077f: 55
                                 %rbp
                          push
  400780: 48 89 e5
                                 %rsp,%rbp
                          mov
  400783: 48 83 ec 30
                          sub
                                 $0x30,%rsp
  400787: 48 89 7d d8
                                 %rdi, -0x28(%rbp)
                          mov
  40078b: 48 89 75 d0
                                 %rsi,-0x30(%rbp)
                          mov
```



The process of creating a process from an ELF file is called:

"Linking and Loading"

Sections:

★ Are for compiler (gcc), to link several object files together (.o)

Segments:

- ★ Are for the loader, to create the process
- Consists of one ore more sections



Recap:

- → Program Code is stored in ELF Files
- ★ ELF Files contain segments
- → Segments are copied 1:1 in the memory

Challenges



Challenges:

https://exploit.courses

- → Challenge 0: Introduction to memory layout basic
- → Challenge 1: Introduction to memory layout advanced