## CS3231 Tutorial 11–12

- 1. Suppose  $L_1, L_2$  are recursively enumerable sets.
  - (a) Is  $L_1 \cap L_2$  recursively enumerable?
  - (b) Is  $L_1 L_2$  recursively enumerable?
- 2. The state entry problem for Turing Machines can be described as follows. Suppose a Turing machine  $M = (Q, \Sigma, \Gamma, \delta, q_0, B, F)$ , a state  $q \in Q$ , and input  $w \in \Sigma^*$  is given. Does M ever enter the state q when input is string w?

Show that above problem is undecidable using appropriate reduction from the Halting problem.

- 3. Recall that  $W_i$  denotes the language accepted by the Turing Machine with code i. That is,  $W_i = L(M_i)$ .
  - (a) Show that  $L_{inf} = \{1^i \mid W_i \text{ is infinite }\}$  is not RE.
  - (b) Show that  $L_{fin} = \{1^i \mid W_i \text{ is finite }\}$  is not RE.
  - (c) Show that  $\{w_i \mid a \in W_i \text{ or } b \notin W_i\}$  is not RE, where  $a \neq b$ .
  - (d) Show that  $L_a = \{w_i \mid a \notin W_i\}$  is not RE.

Note that Rice's theorem cannot be used to show a language over machines is not RE. So you would need to use some of the techniques used in class.

- 4. Which of the following problems are decidable and why?
  - (a) Given two DFA's  $M_1$  and  $M_2$ , is  $L(M_1) \cap L(M_2) = \emptyset$ ?
  - (b) Given a CFG G, and a DFA M, is  $L(G) \cap L(M) = \emptyset$ ?
  - (c) Given a TM M', and a DFA M, is  $L(M') \cap L(M) = \emptyset$ ?
- 5. Show that the following languages are undecidable, using Rice's theorem.
  - (a)  $\{M \mid \text{number of elements in } L(M) \leq 20\}.$
  - (b)  $\{M \mid a \in L(M)\}.$
  - (c)  $\{M \mid a \in L(M) \text{ and } b \notin L(M)\}.$
  - (d)  $\{M \mid a \in L(M) \text{ or } b \notin L(M)\}.$
- 6. Prove or disprove:

If  $L^*$  is recursive. Then L is recursive.

- 7. Show that the following problems are undecidable:
  - (a) Given two CFG,  $G_1$  and  $G_2$ , is  $L(G_1) = L(G_2)$ ?
  - (b) Given two CFG,  $G_1$  and  $G_2$ , is  $L(G_1) \subseteq L(G_2)$ ?

- 8. Give unrestricted grammar for
  - (a)  $\{a^ib^jc^k \mid i \le j \le k\}.$
  - (b)  $\{w \mid w \in \{a,b,c\}^* \text{ and } w \text{ contains equal number of } a\text{'s, } b\text{'s and } c\text{'s}\}.$
  - (c)  $\{ww \mid w \in \{a, b\}^*\}.$
- 9. Show that every RE language can be generated by an unrestricted grammar.