

Binary Cyclic codes

- Binary cyclic codes form a sub class of Linear Block codes
- They have two distinct **advantages** over others
 1. Encoding and syndrome calculating circuits can be easily implemented with simple shift register with feedback connections and some basic gates
 2. Cyclic codes have a mathematical structure(algebraic structure)that makes it easy to design

Algebraic structure of Cyclic codes

- Definition
- A (n,k) Linear block code is said to be cyclic if it exhibits the 2 properties
 1. Cyclic property - every cyclic shifts of the code is also a code vector of C
 2. Linearity – The sum of any 2 code words is also a code word

For example : Let $C_1 = 0111001$ be a code-vector of C . If $C_2 = 1011100$ [the last '1' of C_1 has moved into the first position] is also a code-vector of C , then it is called as "*Cyclic Code*".

Similarly $C_3 = 0101110$, $C_4 = 0010111$ etc. will also be code vectors of C

In general, let the n -tuple be represented by

$$V = (v_0 \ v_1 \ v_2 \ \dots \ v_{n-1}) \quad \dots (5.60)$$

If v belongs to a cyclic code, then

$$\text{and } \left. \begin{aligned} V^{(1)} &= (v_{n-1} \ v_0 \ v_1 \ v_2 \ \dots \ v_{n-2}) \\ V^{(2)} &= (v_{n-2} \ v_{n-1} \ v_0 \ v_1 \ v_2 \ \dots \ v_{n-3}) \\ &\vdots \\ V^{(i)} &= (v_{n-i} \ v_{n-i+1} \ \dots \ v_0 \ v_1 \ v_2 \ \dots \ v_{n-i-1}) \end{aligned} \right\} \quad \dots (5.61)$$

obtained by shifting V cyclically successively, are also code-vectors of C .

This property of cyclic codes allows us to treat the elements of each code-vector as the coefficients of a polynomial of degree $(n - 1)$.

Equation (5.60) can now be represented as a polynomial given by

$$V(x) = v_0 + v_1x + v_2x^2 + \dots + v_{n-1}x^{n-1} \quad \dots\dots (5.62)$$

Similarly, set of equations (5.61) are also represented as polynomials given by

$$\left. \begin{aligned} V^{(1)}(x) &= v_{n-1} + v_0x + v_1x^2 + \dots\dots\dots + v_{n-2}x^{n-1} \\ V^{(2)}(x) &= v_{n-2} + v_{n-1}x + v_0x^2 + \dots\dots\dots + v_{n-3}x^{n-1} \\ &\vdots \\ V^{(i)}(x) &= v_{n-i} + v_{n-i+1}x + v_{n-i+2}x^2 + \dots\dots\dots + v_{n-i-1}x^{n-1} \end{aligned} \right\} \quad \dots\dots (5.63)$$

Field with the following rules of addition

Modulo 2 Algebra

0 + 0 = 0	0 . 0 = 0
0 + 1 = 1	0 . 1 = 0
1 + 0 = 0	1 . 0 = 0
1 + 1 = 0	1 . 1 = 1
Modulo-2 addition	Modulo-2 Multiplication



Modulo-2 Algebra:

Let us discuss addition, subtraction (same as addition in modulo-2 arithmetic), multiplication and division of polynomials with suitable examples.

Addition : The quantity $x + x$ can be written as

$$x + x = x(1 + 1) = x.0 = 0$$

..... (5.64)

since $1 + 1 = 0$ in modulo-2 addition

Similarly $x^2 + x^2 = x^2(1+1) = x^2.0 = 0$

and $x^3 + x^3 = x^3(1+1) = x^3.0 = 0$ and so on.

Subtraction of polynomials is same as addition in modulo-2 algebra.

Multiplication : The quantity $x.x = x^2$, $x^2.x = x^3$ and so on. Let us now consider multiplication of two polynomials.

Example 5.19 : Find the product of polynomials $f_1(x) = (x + 1)$ and $f_2(x) = x^3 + x + 1$ using modulo-2 algebra.

Solution

$$\begin{aligned} f_1(x).f_2(x) &= (x + 1)(x^3 + x + 1) \\ &= x^4 + x^2 + x + x^3 + x + 1 \\ &= x^4 + x^3 + x^2 + x + x + 1 \\ &= x^4 + x^3 + x^2 + x(1 + 1) + 1 \\ &= x^4 + x^3 + x^2 + 1 \quad \text{using equation (5.64).} \end{aligned}$$

Example 5.20 : Multiply $f_1(x) = 1 + x + x^3$ and $f_2(x) = (1 + x + x^2 + x^4)$ using modulo-2 algebra.

Solution

$$\begin{aligned} f_1(x).f_2(x) &= (1 + x + x^3)(1 + x + x^2 + x^4) \\ &= 1 + x + x^2 + x^4 \\ &= \quad + x + x^2 + x^3 + x^5 \\ &= \quad \quad + x^3 + x^4 + x^5 + x^7 \\ &= 1 + x^7 \text{ using equation (5.64).} \end{aligned}$$

Example 5.21 : Divide $f_2(x) = x^6 + x^5 + x^2$ by $f_1(x) = x^3 + x + 1$ using modulo-2 algebra.

Solution

$$x^3 + x + 1 \mid x^6 + x^5 + x^2 \quad (x^3 + x^2 + x \leftarrow Q(x))$$

$$\underline{x^6 + x^4 + x^3}$$

$$x^5 + x^4 + x^3 + x^2 \quad \text{since subtraction is same as addition}$$

$$\underline{x^5 \quad \quad + x^3 + x^2}$$

$$x^4$$

$$\underline{x^4 + x^2 + x}$$

$$x^2 + x \leftarrow R(x)$$