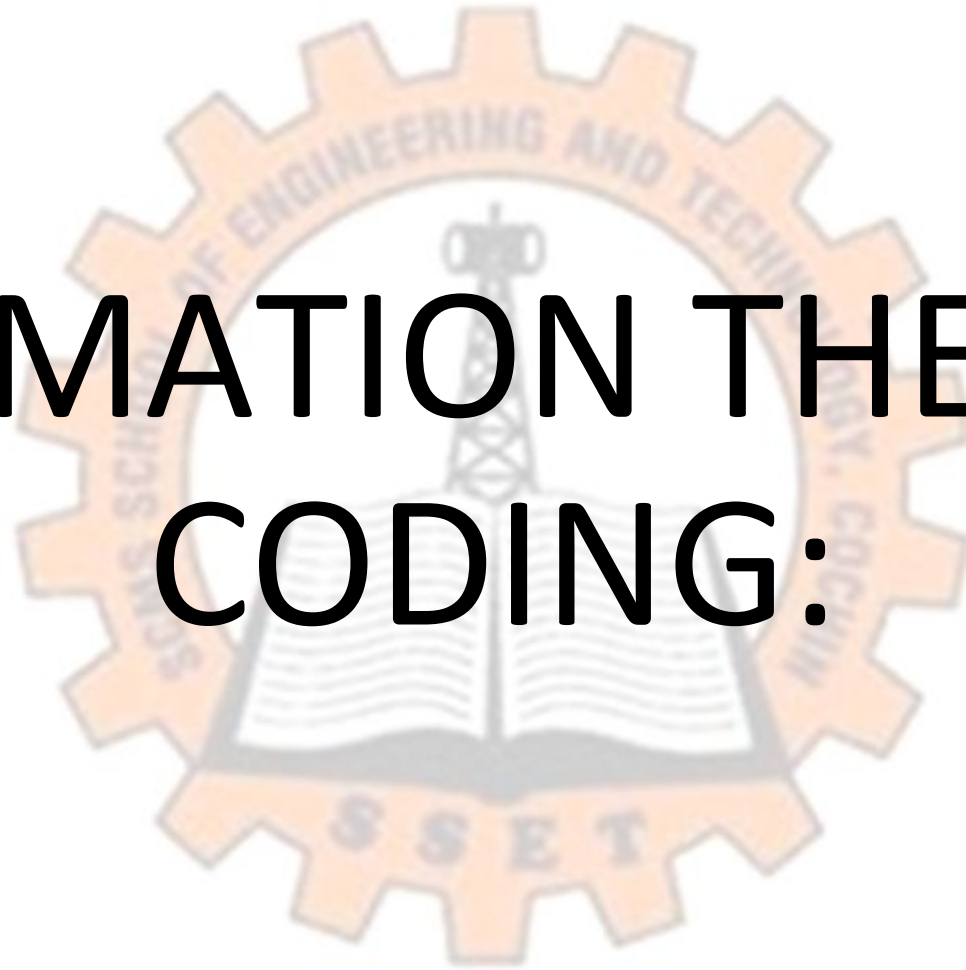


# INFORMATION THEORY & CODING:



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- Quick recap
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# Shannon's Hartley law / Shannon's Third Theorem

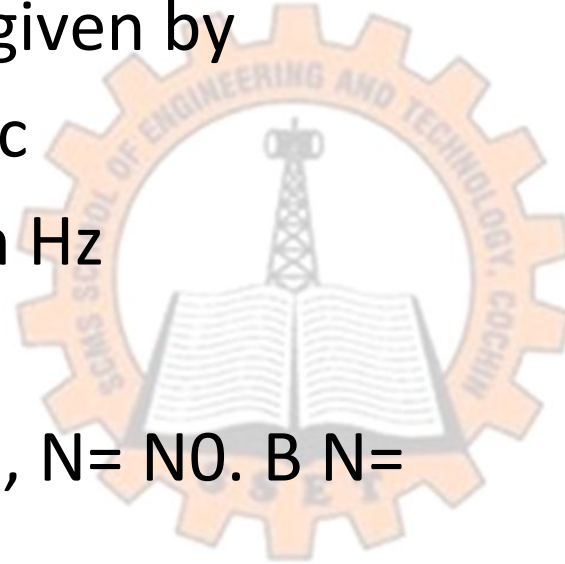
- The Theorem states that the capacity of the bandlimited Gaussian Channel with AWGN is given by

- $C = B \log(1 + S/N)$  bits/sec

$B$  = Channel bandwidth in Hz

$S$  = signal power in watts

$N$  = noise power in watts ,  $N = N_0 \cdot B$   $N =$



# Shannon – Hartley Law Implications

$$C = B \log_2 \left( 1 + \frac{S}{N} \right) \quad \text{bits/sec} \quad \text{--- ①}$$

indicates that a noiseless Gaussian channel  $\frac{S}{N} = \infty$  has an infinite capacity

On the other hand, while the channel capacity does increase, it does not become infinite because with an increase in BW, noise power also increases ( $N = N_0 B$ ).

# Upper Limit for Channel Capacity

Thus for a fixed signal power and in presence of  $wgn$ ;  
the channel capacity approaches an upper limit with  
increasing BW.

The limit

$$N = N_0 B \quad \text{or} \quad N = \eta B \quad \text{in eqn (1) above} \quad \text{PSD} = \frac{N_0}{2}$$

$$C = B \left\{ \log_2 \left( 1 + \frac{S}{\eta B} \right) \right\}$$

$$= \frac{S}{\eta} \frac{1}{B} \cdot B \left\{ \log_2 \left( 1 + \frac{S}{\eta B} \right) \right\}$$



$$= \frac{S}{\eta} \cdot \log_2 \left[ 1 + \frac{S}{\eta B} \right]^{\frac{\eta B}{S}}$$

**Log  $m^n$**

Substitute  $x = \frac{S}{\eta B}$  as  $B \rightarrow \infty$   $x \rightarrow 0$ .

$$C = \frac{S}{\eta} \log_2 (1+x)^{1/x}$$

when  $B \rightarrow \infty$ ,  $x \rightarrow 0$ .

x)

$$C = \lim_{\substack{B \rightarrow \infty \\ x \rightarrow 0}} \frac{S}{\eta} \log_2 (1+x)^{1/x}$$

$$C_{\infty} = \frac{S}{\eta} \log_2 \lim_{x \rightarrow 0} [(1+x)^{1/x}]$$

$$= \frac{S}{\eta} \log_2 e$$

$$C_{\infty} = 1.44 \frac{S}{\eta}$$

This is the upper limit for channel capacity

## 2. Bandwidth S/N trade off

Let us consider the trade off b/w BW and S/N.

$$\text{Let } \frac{S}{N} = 7 \quad B = 4 \text{ KHz}$$

$$C = 4 \text{ KHz} \cdot \log_2 (1 + 7) = \underline{\underline{12 \times 10^3 \text{ bits/sec}}}$$

$$\begin{aligned} \therefore \text{Channel capacity } C_1 &= B_1 \log \left( 1 + \frac{S_1}{N_1} \right) \\ &= 4 \times 10^3 \log(1 + 7) \\ &= 12 \times 10^3 \text{ bits/sec.} \end{aligned}$$

Keeping the channel capacity  $C_2$  same as  $C_1$  and if signal-to-noise ratio is increased to 15, then

$$\begin{aligned} C_2 &= C_1 = 12 \times 10^3 = B_2 \log \left( 1 + \frac{S_2}{N_2} \right) \\ &= B_2 \log(1 + 15) \end{aligned}$$

$$\therefore B_2 = 3 \text{ KHz}$$



Since the noise power  $N = \eta B$ , as the bandwidth gets reduced from 4 to 3 KHz, the noise also decreases indicating an increase in signal power as shown below.

$$\text{We have } N_1 = \eta B_1 = (\eta) (4 \text{ KHz})$$

$$\text{and } N_2 = \eta B_2 = (\eta) (3 \text{ KHz})$$

$$\text{Consider } \frac{(S_2/N_2)}{(S_1/N_1)} = \frac{15}{7}$$

$$\therefore \frac{S_2}{S_1} = \frac{15N_2}{7N_1} = \frac{(15)(\eta)(3\text{KHz})}{(7)(\eta)(4\text{KHz})} = \frac{45}{28} = 1.6\text{---}$$

# Inference

→ With 3 KHz BW, noise power will be  $\frac{3}{4}$  as large as with 4 KHz  
→ Thus S/N power will have to be increased by the factor

$$\frac{\frac{3}{4} \times 15}{7} = 1.60.$$

∴ 25% reduction in bandwidth requires 60% increase in signal power

Thus to decrease the BW, the signal power has to be increased;  
If to decrease signal power, bandwidth must be increased.

# Bandwidth S/N trade off curve

BW, S/N Trade off Curve.

To draw the trade off curve from  
Shannon Hartley law;  $C = B \log_2 \left(1 + \frac{S}{N}\right)$ .

$C \Rightarrow \text{bits/sec}.$

$$\frac{B}{C} = \frac{1}{\log_2 \left(1 + \frac{S}{N}\right)}.$$

$$\log_2(1 + 50) = 5.64$$



# Table for B/C for different values of S/N

$\frac{S}{N}$	0.5	1	2	5	10	20	50
B/C	1.71	1	0.63	0.37	0.289	0.23	0.176

$\log_2(1+50) = 5.64$

$\frac{S}{N} = 15$      $B = 5 \text{ KHz}$   
 $C = 20 \text{ kb/s}$

$\frac{S}{N} = 31$      $C = 20 \text{ kb/s}$   
 $B = 4 \text{ KHz}$

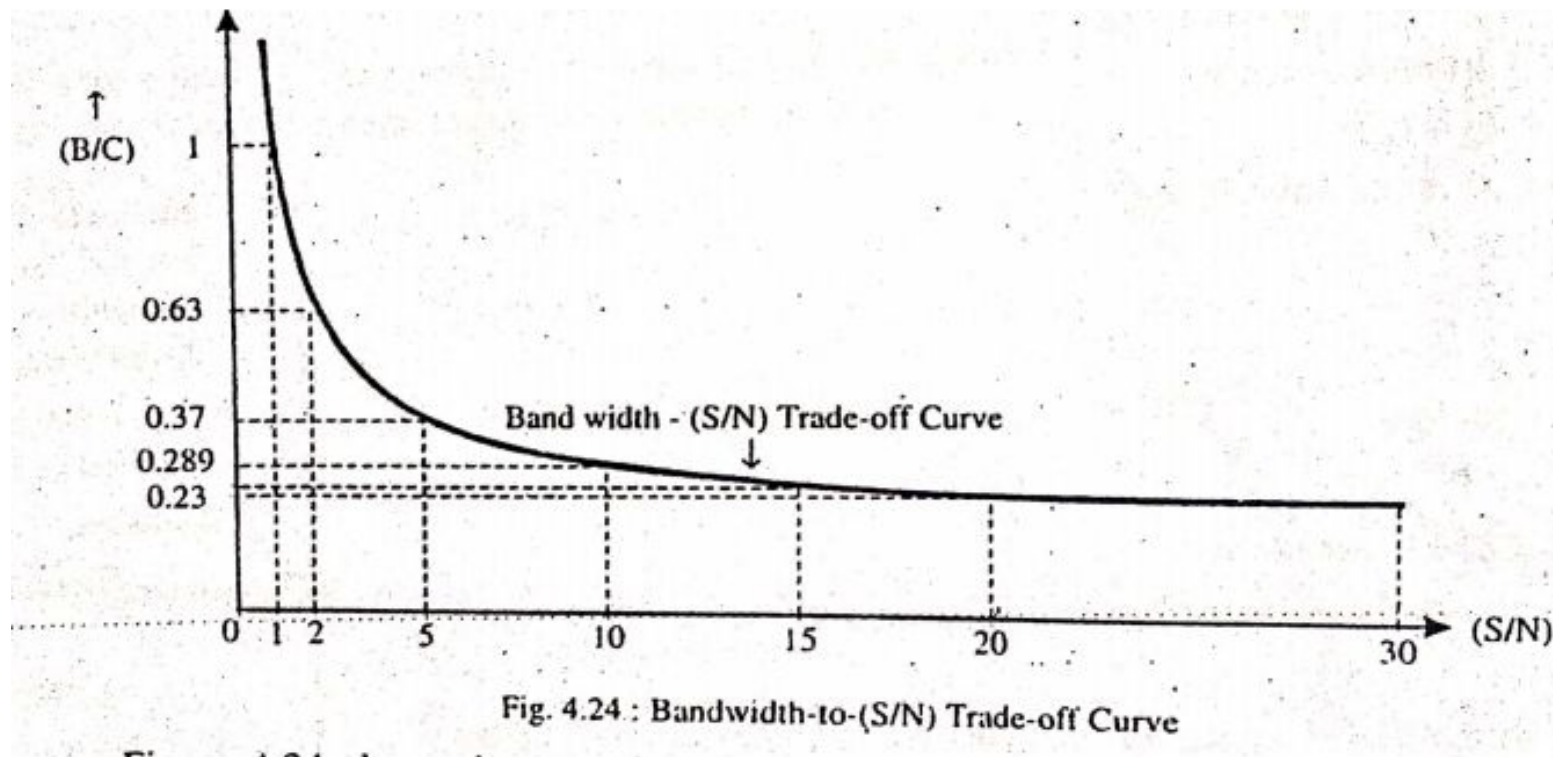
did req bit rate, in dB, answer.

↑ ↑



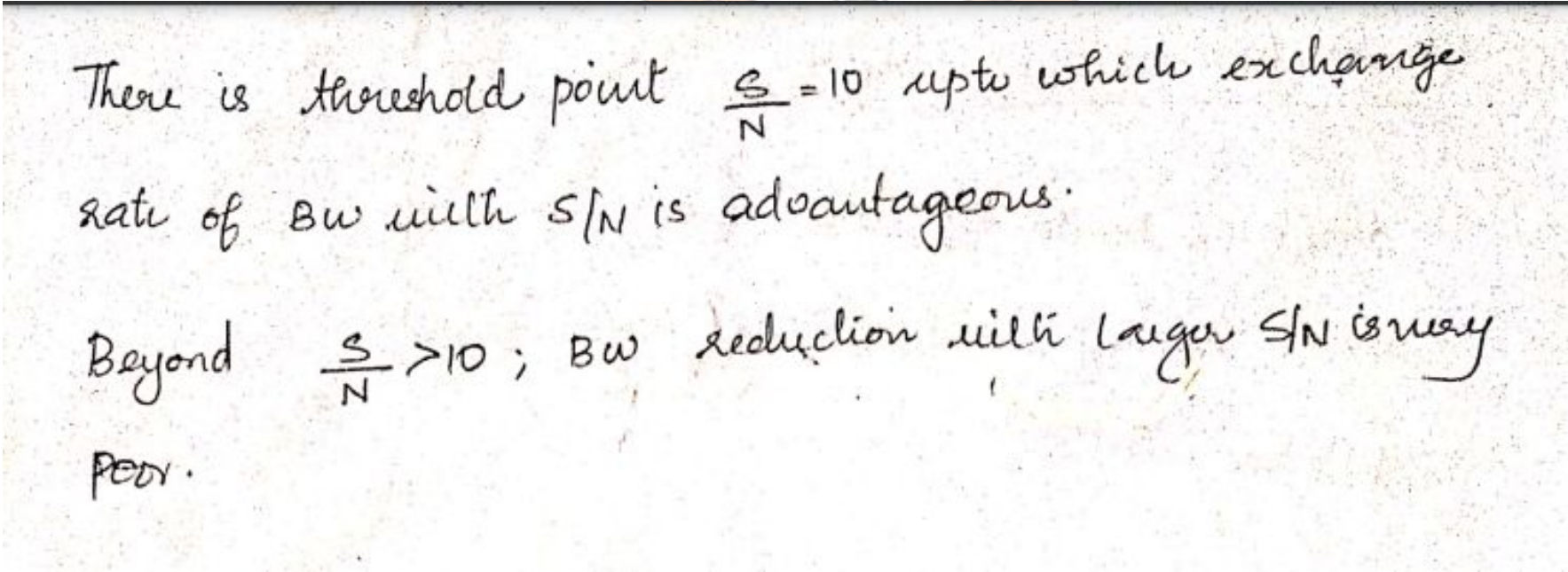
# B/c vs S/N trade off curve

The plot shows B/C as a function of S/C



# Inference of the curve

1. Using the curve same channel capacity may be obtained by increasing bandwidth if  $S/N$  is small



There is threshold point  $\frac{S}{N} = 10$  upto which exchange rate of BW with  $S/N$  is advantageous.

Beyond  $\frac{S}{N} > 10$ ; BW reduction with larger  $S/N$  is very poor.

# Conclusion

- Upper limit for Shannon's capacity
- Trade off curve



# THANK YOU

