

①

module 1 :

→ 2 classes are ~~there~~ present

i) Attr

ii) AssignModule 1.

↳ This contains the main function & Scanner class.

→ class "Attr" contains all the stuff.

→ This class contains.

i) 15 variables. (excluding method specific variables).

ii) 22 methods.

Algorithm;

②

I have divided the work into 5 parts.

- i) computing FD closure
- ii) computing minimal keys.
- iii) computing Normal Form of each FD.
- iv) Decomposing.
- v) Testing for Dependency preserving.

1) Computing FD closure:

we are using " iR_6 " to compute FD closure because it includes " iR_3 within itself." and " iR_6 " is ~~more~~ having more information when compared with iR_3 .

iR_1	$A \rightarrow A$
iR_2	$x \rightarrow y \Rightarrow xz \rightarrow yz$
iR_3	$x \rightarrow y \ \& \ y \rightarrow z \Rightarrow x \rightarrow z$
iR_4	$x \rightarrow yz \Rightarrow x \rightarrow y \ \& \ x \rightarrow z$
iR_5	$x \rightarrow A, x \rightarrow B \Rightarrow x \rightarrow AB$
iR_6	$x \rightarrow y \ \& \ wy \rightarrow z \text{ then } wx \rightarrow z$

→ iR_1, iR_2, iR_4, iR_5 are anyways useless.

→ In order to understand how it works, PTO.

write all the FDs in 2 columns like this...

(3)

(i)

$AB \rightarrow C$

$C \rightarrow D$

$CE \rightarrow FD$

$B \rightarrow H$

(j)

$AB \rightarrow C$

$C \rightarrow D$

$CE \rightarrow FD$

$B \rightarrow H$

use ^{nested} 2 for loops - one for 'i' & another for 'j'

& check whether iR_k is applicable.

ie.) $R_{hs}[i]$ should ~~contain~~ be present in $lhs[j]$

then we can apply iR_k .

Eg:

$AB \rightarrow C$ & $CE \rightarrow FD$

$\Rightarrow ABE \rightarrow FD$ \rightarrow new FD created.

append it at the bottom.

check other way round also.

ie.)

if $R_{hs}[j]$ ~~should~~ is present in $lhs[i]$
then also we can apply iR_k .

used methods;

(4)

o) Fill FDs.

i) public void fd closure

ii) sort

iii) remove Duplicates

iv) remove Redundant FDs (overloaded once)

v) remove null spaces.

2) computing minimal keys;

i) To compute this, FD closure should be ready.

ii) All $2^n - 2$ combinations of attributes should also be ready.

→ we store we first store all these $(2^n - 2)$ combinations in a string[] called as "attr-combns".

→ To compute all the $(2^n - 2)$ combinations, I used Recursion concept.

→ methods used are

i) fillAttrCombs.

ii) nChooseP

iii) store

(5)

→ computing keys:

→ ~~Take every combination~~

→ Take appropriate combination from
"attr-combns" & check whether it is a
key.

→ Don't check for all combinations.

ie, if AB is already a key

then checking for ABC, ABD,

ABCDE is
waste of time

store the found keys in "keys" array..

used methods:

i) fillKeyList

ii) getCombinations()

iii) Computing Normal Form of Each FD; (6)

Basically Every FD belong to any of the following 9 cases- (Assuming only one attribute on RHS).

i) $\text{key} \rightarrow \text{key}$ (BCNF)

eg: $AB \rightarrow C$

(let us say both are keys)

ii) $\text{key} \rightarrow \text{Another Key's Attribute}$ (BCNF)

eg: \exists 2 keys
 AB, DE

$\& AB \rightarrow D$

iii) $\text{key} \rightarrow \text{non key attribute}$ (BCNF)

iv) $\text{key attribute} \rightarrow \text{key}$

eg: \exists 2 keys AB, C

$B \rightarrow C$

Impossible case

because if key attribute
 \downarrow
key
then that key attr
becomes key

v) keyattribute \rightarrow keyattribute (3NF)

(7)

Eg: \exists 2 keys i) ABC

ii) DE

$E \rightarrow C$

vi) keyattribute \rightarrow nonkeyattribute (1NF)

Eg: \exists a key AB.

$A \rightarrow C$

vii) non-keyattribute \rightarrow key (impossible case).

viii) non-keyattribute \rightarrow key attribute (3NF).

Eg: \exists a key ABC

$D \rightarrow B$

ix) non-keyattribute \rightarrow non key attribute (2NF).

~~Q~~

In order to calculate NF, take a FD, ⑧
scan it for the above 9 cases & find out
the Normal Form. & put it in nf-list[]

Note:

out of 9,

— 2 impossible

— 1 2NF case → can be decomposed to BCNF

→ 1 1NF case → " " "

→ 2 3NF cases →

v) can't be decomposed with
dependency preserving
ie, this is final

viii) can be brought to BCNF
by decomposition.

→ 3 BCNF.

used methods:

i) fillNFList()

ii) isAKey()

iii) isAKeyAttribute()

v) Decomposing:

- from the above NF,
- decompose...

used methods:

- decomposeRel()
- completeFormalities()

↳ To compute keys & FDs
for newly formed
relations

v) Testing for Dependency preserving:

- Take FD closures from all the decomposed rels
and store ~~it~~ in a temporary string [].
all of them
- calculate FD closure.
- count the number of FDs from initial &
final cases.
- compare every FD.

used method:

- calculate FD closure from Decomposed And Compare with
Original().

Total methods used explained
till now!

~~17~~ + 1
17 + 1

(10)

Remaining methods:

- i) public void printrell()
 - ii) public void printNF()
 - iii) printKeys
 - iv) printFDS.
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