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CS 30 - HW #8

- ① a)  $\{S, F\}$  is the sample space, where  $S$  refers to success on a given play, and  $F$  refers to failure.
- b) Event  $E_S = \{S\}$  and event  $E_F = \{F\}$
- c)  $P(S) = \frac{1}{38}$ ,  $P(F) = \frac{37}{38}$
- d) Random variable  $X_i$  is a random variable whose value is 1 if the  $i$ th trial is  $S$  and 0 if the  $i$ th trial is a failure. As such,  $X_i$  is an indicator variable for  $E_S$ .

e) The probability distribution  $X_i$  is:  $P(X_i = 1) = \frac{1}{38}$  and  $P(X_i = 0) = \frac{37}{38}$

f) Let us define a random variable  $Y = \sum_{i=1}^{105} X_i$ . As such  $Y$  is the

number of successful trials of the 105 total trials. We know the reward we receive is  $36Y$ . We also know that we lose a total of 105 dollars over the course of 105 trials. So, the amount of money we have at the end is  $105 - 105 + 36Y$ , which must be greater than 105.

$$105 - 105 + 36Y > 105$$

$$36Y > 105; Y > 2.92$$

Since  $Y$  may only be a whole number,  $Y \geq 3$ . So we are calculating  $P(Y \geq 3)$ .

Since  $Y$  has a binomial since it is a sum of 105 bernoulli trials:

$$P(Y=0) = \left(\frac{37}{38}\right)^{105}; P(Y=1) = 105 \left(\frac{37}{38}\right)^{104} \left(\frac{1}{38}\right)$$

$$P(Y=2) = \binom{105}{2} \left(\frac{37}{38}\right)^{103} \left(\frac{1}{38}\right)^2$$

$$\text{So } P(Y \geq 3) = 1 - \left(\frac{37}{38}\right)^{105} - 105 \left(\frac{37}{38}\right)^{104} \left(\frac{1}{38}\right) - \binom{105}{2} \left(\frac{37}{38}\right)^{103} \left(\frac{1}{38}\right)^2$$
$$= 0.5242$$

- (2) Sample space:  $\{R, L, M, D\} = N$ . The sample space is  
 $S = \{s \in N^* \mid \text{every element of } N \text{ appears at least once and the last element of the string appears exactly once}\}$ .

$l_i(s)$  = the length of the shortest prefix of the string  $s$  such that the  $i$ th distinct element of  $N$  first appears.  
 $l_0(s) = 0$ .

$$P(R) = P(L) = P(M) = P(D) = \frac{1}{4}.$$

- a) Random Variable  $X$  denotes the number of times you have eaten a happy meal before you first encounter Leonardo. That is, it is the number of Happy Meals consumed until  $E_L$  occurs.

Since  $P(L) = \frac{1}{4}$  and so  $P(\bar{L}) = \frac{3}{4}$ , this may be viewed as a binomially distributed variable, with  $P(X=1) = \frac{1}{4}$ ,  $P(X=2) = \left(\frac{3}{4}\right)\left(\frac{1}{4}\right)$  and  $P(X=n) = \left(\frac{3}{4}\right)^{n-1}\left(\frac{1}{4}\right)$ , with  $x$  assuming all integer values such that  $x \geq 1$ .

$$\begin{aligned} E[X] &= \sum_{x=1}^{\infty} x \left(\frac{3}{4}\right)^{x-1} \left(\frac{1}{4}\right) = \frac{1}{4} \sum_{x=1}^{\infty} (x) \left(\frac{3}{4}\right)^{x-1} \\ &= \frac{1}{4} (1 + 2\left(\frac{3}{4}\right) + 3\left(\frac{3}{4}\right)^2 + \dots) = S \\ &\quad 1 + 2\left(\frac{3}{4}\right) + 3\left(\frac{3}{4}\right)^2 + \dots = 4S \\ &\quad \frac{3}{4} + 2\left(\frac{3}{4}\right)^2 + 3\left(\frac{3}{4}\right)^3 = (4S)\left(\frac{3}{4}\right) = 3S \\ &\quad 4S - 3S = 1 + (2-1)\frac{3}{4} + (3-2)\left(\frac{3}{4}\right)^2 + \dots \\ &\quad = 1 + \frac{3}{4} + \left(\frac{3}{4}\right)^2 + \left(\frac{3}{4}\right)^3 + \dots \\ &\quad = 1 + \sum_{i=1}^{\infty} \left(\frac{3}{4}\right)^i \\ &\quad = 1 + \frac{3/4}{1-(3/4)} = 1 + \frac{3/4}{1/4} = 1 + 3 \\ &\quad = 4 = S \end{aligned}$$



Since  $S = E[X]$ , the expected number of happy meals before encountering Leonardo is 4.

- b) Let Random Variable  $X_i$  be the number of happy meals until the first instance of the  $i^{\text{th}}$  toy is encountered, where  $i=1$  if the toy is the first toy of all 4 to be in the series,  $i=2$  if it is the second, and so on. As such  $X_i(S) = L_i(S) - L_{i-1}(S)$

$X_1$ :  $P(X_1=1) = 1$ , and  $P(X_1=x) = 0$  for all  $x \neq 1$ . So  $E[X_1] = 1$ .

Once the first toy is encountered, the probability of encountering the second toy in a given happy meal is  $\frac{3}{4}$ . Since there are 3 undiscovered toys. So  $E[X_2] = 1/p$ , where  $p = \frac{3}{4}$ , so  $E[X_2] = 4/3$ .

By the same logic, after the second toy has been found, the expected number of happy meals to be consumed before the third one is found is 2, because the probability of discovering the third toy in a given happy meal is  $\frac{1}{2}$ . So  $E[X_3] = 2$ .

Also, after the third toy has been found, since the probability of finding the fourth toy in a given happy meal is  $\frac{1}{4}$ ,  $E[X_4] = 4$  by the same logic.

Let us define the random variable  $Y = X_1 + X_2 + X_3 + X_4$  as the number of happy meals to be eaten before finding all 4 toys.

$E[Y] = E[X_1 + X_2 + X_3 + X_4] = E[X_1] + E[X_2] + E[X_3] + E[X_4]$ , by linearity of expectation.

$$\text{So } E[Y] = 1 + \frac{4}{3} + 2 + 4 = 8\frac{1}{3} \text{ or } \frac{25}{3}$$

c) For  $n$  toys, let us assume that  $Y = \sum_{i=1}^n X_i$ , where  $X_i$

is a random variable that represents the number of happy meals that must be eaten to discover the  $i^{\text{th}}$  toy after the  $(i-1)^{\text{th}}$  toy has been discovered.  $X_i = l_i(S) - l_{i-1}(S)$

Prob (finding the  $i^{\text{th}}$  in a given happy meal) =  $\frac{n-i+1}{n}$ , because

we know Prob (finding the 1<sup>st</sup> toy) is  $\frac{n}{n}$ .

By the logic discussed in part b)  $E[X_i] = \frac{n}{n-i+1}$

$$\begin{aligned} E[Y] &= E\left[\sum_{i=1}^n X_i\right] = \sum_{i=1}^n E[X_i] = \sum_{i=1}^n \frac{n}{n-i+1} \\ &= \left[ n \sum_{i=1}^n \frac{1}{n-i+1} \right], \text{ since } n \text{ is constant} \end{aligned}$$



3) Proof by counterexample:

$$P(X=1 \wedge Y=1) = \text{Prob}(\text{having one of the coins be heads and the other be tails})$$

Let  $G_1$  and  $G_2$  be the two different coins; the sample space is therefore  $\{(H,H), (H,T), (T,H), (T,T)\}$ , where the first number in the pair denotes whether  $G_1$  had heads or tails, and the second number denotes whether  $G_2$  turned up heads or tails.

$$\text{Let event } E_H = \{(H,H), (H,T)\}, E_T = \{(T,H), (T,T)\}, \\ F_H = \{(H,H), (T,H)\}, F_T = \{(H,T), (T,T)\}.$$

$E_H$  is the event that  $G_1$  turns up heads,  $E_T$  is the event that  $G_1$  turns up tails.  $F_H$  is the event that  $G_2$  turns up heads and  $F_T$  is the event that  $G_2$  turns up tails.

$$\begin{aligned} P(X=1 \wedge Y=1) &= P(E_H \cap F_T) + P(E_T \cap F_H) \\ &= P(E_H|F_T) \cdot P(F_T) + P(E_T|F_H) \cdot P(F_H) \\ &= \left(\frac{1}{2}\right) \cdot \left(\frac{1}{2}\right) + \left(\frac{1}{2}\right) \cdot \left(\frac{1}{2}\right) \\ &= \frac{1}{2} \end{aligned}$$

$P(X=1) = \text{Prob}(\text{having exactly 1 heads, and therefore 1 tails}) = \frac{1}{2}$ , as calculated above.  $P(X=0) = \frac{1}{4}$  and  $P(X=2) = \frac{1}{4}$

$P(Y=1) = \text{Prob}(\text{having exactly 1 tails and therefore 1 heads}) = \frac{1}{2}$ , as calculated above.  $P(Y=0) = \frac{1}{4}$  and  $P(Y=2) = \frac{1}{4}$

$$P(X=1) \cdot P(Y=1) = \frac{1}{2} \cdot \frac{1}{2} = \frac{1}{4} \neq \frac{1}{2} = P(X=1 \wedge Y=1).$$

So, the events are not independent.

- 4) A dodecahedral pair of dice have 12 faces each, with numbers ranging from 1 to 12.  
As such, the sample space is the set  $[1 \dots 12] \times [1 \dots 12] = S$

$X_1$  is a random variable that represents the number that comes up on the first die, and  $X_2$  is a random variable that represents the number that comes up on the second die.

$X_1$  and  $X_2$  may assume all values  $a$  such that  $a \in \mathbb{N}$ , and  $1 \leq a \leq 12$ .  $P(X_1 = a) = \frac{1}{12}$  for all  $a$  and  $P(X_2 = a) = \frac{1}{12}$  for all  $a$ .

- a) Let  $X$  be the random variable that represents the sum of the die rolls. As such  $X = X_1 + X_2$ .

$$E[X] = E[X_1 + X_2] = E[X_1] + E[X_2] \text{ by linearity of expectation}$$

$$E[X_1] = \sum_{i=1}^{12} \left(\frac{1}{12}\right) i = \frac{1}{12} \sum_{i=1}^{12} i = \frac{1}{12} \left(\frac{12+1}{2}\right)(12) = \frac{12+1}{2} = \left\lfloor \frac{13}{2} \right\rfloor$$

$$E[X_2] = \sum_{i=1}^{12} \frac{1}{12} i = \frac{13}{2} \text{ by the logic above for } E[X_1]$$

$$E[X_1 + X_2] = E[X] = 13$$

- b)  $X = X_1 + X_2$  as in part a)

$$V[X] = V[X_1 + X_2] = V(X_1) + V(X_2) \text{ since } X_1 \text{ and } X_2 \text{ are independent}$$

$$\begin{aligned} V[X_1] &= E[X_1^2] - E^2[X_1] \\ &= \sum_{i=1}^{12} \left(i^2 \left(\frac{1}{12}\right)\right) - \left(\frac{13}{2}\right)^2 = \frac{1}{12} \sum_{i=1}^{12} i^2 = \frac{1}{12} \left(\frac{n(n+1)(2n+1)}{6}\right), n=12 \\ &= \frac{1}{12} \left(\frac{12(13)(25)}{6}\right) = \frac{(13)(25)}{6} = \frac{325}{6} \end{aligned}$$

$$V[X] = \frac{325}{6} - \left(\frac{13}{2}\right)^2 \text{ from part a) } =$$



$$= \frac{325}{6} - \frac{169}{4} = \frac{143}{12}$$

$$V[X_2] = \sum_{i=1}^{12} (i)(\frac{1}{2}) - \left(\frac{13}{2}\right)^2, \text{ since } E[X_2^2] - E[X_2]^2 = V[X_2]$$

$$= \frac{143}{12} \text{ from the same logic as above for } V[X_1]$$

$$V[X_1 + X_2] = V[X_1] + V[X_2] = \frac{143}{12} + \frac{143}{12} = \boxed{\frac{143}{6}}$$

5)  $V[X] = E[X^2] - E^2[X]$

Sample space is  $\{S, F\}$  and  $P(S) = p$  and  $P(F) = q = 1-p$

$E_S$  is the event that represents success or  $\{S\}$ , and  $E_F$  represents  $\{F\}$

$X$  is a random variable such that  $X=1$  when event  $E_S$  occurs and  $X=0$  when  $E_F$  occurs.

$$P(X=0) = q = 1-p \text{ and } P(X=1) = p$$

$$E[X] = (1)(p) + 0(1-p) = p, \text{ so } E^2[X] = p^2.$$

$$E[X^2] = (1)^2(p) + 0^2(1-p) = p.$$

$$E[X^2] - E^2[X] = p - p^2$$

From calculus, we know that the max of  $p - p^2$  can be calculated as follows:

$$\frac{d}{dp} [p - p^2] = 1 - 2p = 0 \text{ when } p = \frac{1}{2}$$

$$\text{Since } \frac{d}{dp} [1 - 2p] = -2 < 0, p = \frac{1}{2} \text{ is a max for } p - p^2.$$

$$\text{So } p - p^2 \text{ is at a max when } p = \frac{1}{2}, \text{ and } p - p^2 = \frac{1}{2} - \frac{1}{4} = \boxed{\frac{1}{4}}$$

$$\text{So } p - p^2 = E[X^2] - E^2[X] = V[X] \leq \frac{1}{4}.$$



not

(b) Sample space:  $M = [1 \dots m]$ ,  $S = \{s \in M^k \mid \text{every element of } M \text{ appears at least once in } s \text{ and the last element of } s \text{ appears exactly once}\}$ .

Let  $Y$  be a random variable that represents the elements needed to fill all slots of a hash table. Let  $l_i(s)$  be the length of the shortest prefix of the sequence such that the  $i^{\text{th}}$  distinct slot is filled.

We now define  $X_i(s) = l_i(s) - l_{i-1}(s)$ , which is equal to the length of the string that lies between the  $(i-1)^{\text{th}}$  distinct hashing and the  $i^{\text{th}}$  distinct hashing.  $l_0(s) = 0$ .  $\therefore$  therefore is  $\sum_{i=1}^m X_i(s)$

$$E[Y] = E\left[\sum_{i=1}^m X_i\right] = \sum_{i=1}^m E[X_i], \text{ by linearity of expectation}$$

We know that, from problem (2), since the probability of finding the  $i^{\text{th}}$  empty slot is  $\frac{k-i+1}{k}$  (since the distribution is uniform),

$$E[X_i] = \frac{k}{k-i+1}, \text{ as shown in problem (2)}$$

$$\sum_{i=1}^k E[X_i] = \sum_{i=1}^k \left(\frac{k}{k-i+1}\right) = \left[ k \sum_{i=1}^k \frac{1}{k-i+1} \right], \text{ since } k \text{ is a constant}$$

$$\boxed{\text{So } E[Y] = k \sum_{i=1}^k \frac{1}{k-i+1}}$$

(7) Assume that  $X_1, \dots, X_n$  are pairwise independent. That is,  $E[X_i X_j] = E[X_i] E[X_j]$  for all  $i, j \in \{1, \dots, n\}$ ,  $i \neq j$ .

$$\text{As such } V\left(\sum_{i=1}^n X_i\right) = E\left(\left(\sum_{i=1}^n X_i\right)^2\right) - E^2\left(\sum_{i=1}^n X_i\right)$$

$$= E((X_1 + X_2 + \dots + X_n)^2) - E^2(X_1 + X_2 + \dots + X_n)$$

$$= E((X_1 + X_2 + \dots + X_n)^2) - (E(X_1) + E(X_2) + \dots + E(X_n))^2 \text{ by}$$

$$= E\left(\sum_{i=1}^n \sum_{j=1}^n X_i X_j\right) - \sum_{i=1}^n \sum_{j=1}^n E(X_i) E(X_j), \text{ because } (a_1 + a_2 + \dots + a_n)^2 = \sum_{i=1}^n \sum_{j=1}^n a_i a_j$$

$$= E\left(\sum_{i=1}^n X_i^2\right) + E\left(\sum_{i=1}^n \sum_{j=1, j \neq i}^n X_i X_j\right) - \left(\sum_{i=1}^n E^2(X_i) + \sum_{i=1}^n \sum_{j=1, j \neq i}^n E(X_i) E(X_j)\right)$$

This is done by splitting the square cases, and using linearity of expectation

$$= \sum_{i=1}^n E(X_i^2) + \sum_{i=1}^n \sum_{j=1, j \neq i}^n E(X_i X_j) - \sum_{i=1}^n E^2(X_i) - \sum_{i=1}^n \sum_{j=1, j \neq i}^n E(X_i) E(X_j)$$

using linearity of expectation.

$$= \sum_{i=1}^n E(X_i^2) + \sum_{i=1}^n \sum_{j=1, j \neq i}^n E(X_i) E(X_j) - \sum_{i=1}^n E^2(X_i) - \sum_{i=1}^n \sum_{j=1, j \neq i}^n E(X_i) E(X_j)$$

Since we know that  $X_1, \dots, X_n$  are pairwise independent, and  $E(X_i X_j) = E(X_i) E(X_j)$  for independent random variables  $X_i$  and  $X_j$

$$= \sum_{i=1}^n E(X_i^2) - \sum_{i=1}^n E^2(X_i)$$

$$= (E(X_1^2) - E^2(X_1)) + (E(X_2^2) - E^2(X_2)) + \dots + (E(X_n^2) - E^2(X_n))$$

$$= \text{Var}(X_1) + \text{Var}(X_2) + \dots + \text{Var}(X_n)$$

$$= \sum_{i=1}^n \text{Var}(X_i)$$

Thus, we have shown that  $V\left(\sum_{i=1}^n X_i\right) = \sum_{i=1}^n V(X_i)$  for pairwise independent random variables  $X_1, \dots, X_n$ .



⑧ Sample space is arbitrary with an arbitrary probability distribution  
We know from Chebyshev's Inequality that

$$\text{Prob}(|X - E[X]| \geq r) \leq \frac{V(X)}{r^2} \text{ for an arbitrary sample space}$$

Claim:  $\text{Prob}(|X - E[X]| < r) \geq 75\%$

Assume that  $X$  is a random variable that may assume integer values, with some mean  $E[X]$ .

Let us consider  $E$ , the event that  $|X - E[X]| \geq r$ , where  $r = 2\sqrt{V(X)}$

$$P(E) \leq \frac{V(X)}{4V(X)} \text{ by Chebyshev's Inequality}$$

$$P(E) \leq \frac{1}{4}, \text{ by algebra}$$

$$1 - P(E) \geq 1 - \frac{1}{4}, \text{ since } P(E) \text{ is at most } \frac{1}{4}$$

$$P(\bar{E}) \geq \frac{3}{4}, \text{ since } 1 - P(E) = P(\bar{E})$$

$$\text{As such, } \text{Prob}(|X - E[X]| < r) \geq \frac{3}{4} = 75\%$$

The claim thus follows from the direct proof principle.

(9) Let  $V = [1 \dots n]$  and  $E^* = \{\{u, v\} \mid u, v \in V \text{ and } u \neq v\}$ .

$E$  then is a subset of  $E^*$ , where  $\{u, v\}$  denotes that there exists an edge between vertices  $u$  and  $v$ . As such, the sample space is  $E \subseteq E^*$ .

The probability distribution for the sample space is  $p^{|E|} \cdot (1-p)^{n-|E|}$ .

Let us define a random variable

$$X_{\{u, v\}} = \begin{cases} 1, & \{u, v\} \in E \\ 0, & \{u, v\} \notin E \end{cases}$$

As such  $P(X_{\{u, v\}} = 1) = p$  and  $P(X_{\{u, v\}} = 0) = 1-p$ .

$$E(X_{\{u, v\}}) = (p)(1) + (0)(1-p) = p.$$

Let us now consider  $X = \sum_{\substack{\{u, v\} \in E^* \\ u \neq v}} X_{\{u, v\}}$ , which is the total number of

edges in the graph. As such, since  $X_{\{u, v\}}$  is constant at  $p$  for all

$u \neq v$ , and there are  $\binom{n}{2}$  total possibilities for  $\{u, v\}$ ,

$$\begin{aligned} E\left(\sum_{\substack{\{u, v\} \in E^* \\ u \neq v}} X_{\{u, v\}}\right) &= \sum_{\substack{\{u, v\} \in E^* \\ u \neq v}} E(X_{\{u, v\}}) \text{ by linearity of expectation.} \\ &= \binom{n}{2} (p) \end{aligned}$$



b) Let  $X_{\{a,b,c\}} = 1$ , if  $\{a,b\}, \{b,c\}, \{a,c\} \in E$  and  $a \neq b \neq c$   
 0, if not otherwise

$$P(X_{\{a,b,c\}} = 1) = \text{probability that all 3 edges exist} = p^3$$

$$P(X_{\{a,b,c\}} = 0) = 1 - p^3$$

$$E(X_{\{a,b,c\}}) = p^3 + 0(1-p^3) = p^3$$

Let us now think of the sum of all such graphs isomorphic to  $K_3$ .

$$X = \sum_{\substack{\{a,b,c\} \\ a \neq b \neq c}} X_{\{a,b,c\}}$$

$$E(X) = \sum_{\substack{\{a,b,c\} \\ a \neq b \neq c}} E(X_{\{a,b,c\}}) \text{ by linearity of expectation}$$

Since there are  $\binom{n}{3}$  ways to pick 3 distinct vertices from  $V$ ,  
 by the logic used in part a,

$$E(X) = \binom{n}{3} p^3$$

$$c) = \binom{n}{3} \left(\frac{c}{n}\right)^3$$

$$= \frac{n!}{(n-3)!3!} \cdot \frac{c^3}{n^3} = \frac{(n)(n-1)(n-2)c^3}{3!n^3}$$

$$\lim_{n \rightarrow \infty} \frac{n(n-1)(n-2)c^3}{3!n^3}$$

Since there will be a  $cn^3$  term in the numerator  
 and  $3!n^3$  term in the denominator, and  
 they will be the highest degree terms in the

numerator and denominator respectively, the limit is  $\frac{c}{3!} = \boxed{\frac{c}{6}}$