Assignment 1 – SystemC intro and L1 cache

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Task:

In this task you must build a single 32kB 8-way set-associative L1 D-cache with a 32-Byte line size. The simulator must be able to be driven with the single processor trace files. Assume a Memory latency of 100 cycles, and single cycle cache latency. Use a least-recently-used, write-back replacement strategy.

The address from the processor is of 32 bits.

Since the line size is 32-Byte, offset = 2^5 = 32 Bytes => 5 bits for offset (bit 0 to bit 4)

Cache block = 32KB/32Bytes = 1KB Cache blocks

Number of sets = 1024/8 = 128 = 2^7 => 7 bits for index (bit 5 to bit 11)

Tag field is the remaining 20 bits.

Pseudo LRU replacement bits = 7(N way -1)

#define TAG_CHECK 0xFFFFF000 #define SET_CHECK 0x00000FE0 #define OFF CHECK 0x0000001F

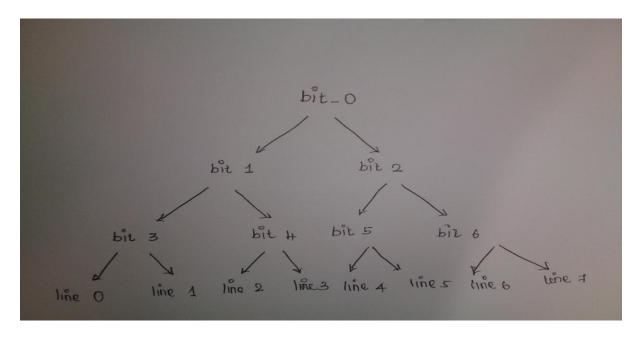
unsigned int tag = (addr & TAG_CHECK) >> 12; //addr taken from the trace file unsigned int set = (addr & SET_CHECK) >> 5; unsigned int word = (addr & OFF_CHECK) >> 2;

For the replacement strategy, pseudo LRU algorithm has been implemented. The binary tree structure and the state table for the algorithm are given below. Bit mask is used to get the next state from the present state in a particular set.

State table:

| Way | Present state | Bit mask | Next State | |
|-------|---------------|----------|------------|--|
| Way_0 | xxx0x00 | 0x0B | 1-11 | |
| Way_1 | xxx1x00 | 0x0B | 0-11 | |
| Way_2 | xx0xx10 | 0x13 | 101 | |
| Way_3 | xx1xx10 | 0x13 | 001 | |
| Way_4 | x0xx0x1 | 0x25 | -11-0 | |
| Way_5 | x1xx0x1 | 0x25 | -01-0 | |
| Way_6 | 0xxx1x1 | 0x45 | 10-0 | |
| Way_7 | 1xxx1x1 | 0x45 | 00-0 | |

State diagram:



Results:

| Trace | CPU | Reads | RHit | RMiss | Writes | Whit | Wmiss | HitRate |
|-------|-----|-------|-------|-------|--------|-------|-------|---------|
| DBG | 0 | 40 | 19 | 21 | 60 | 26 | 34 | 45 |
| RND | 0 | 33031 | 18659 | 14372 | 32505 | 18306 | 14199 | 56.4 |
| FFT | 0 | 86298 | 7652 | 78646 | 43195 | 10882 | 32313 | 14.31 |

Observations:

- Calculating the tag field and set field correctly from the address is the main task to get the appropriate hit rate for the architecture
- The hit rate gets messed up unless the LRU implementation is correct
- For pseudo LRU implementation, next state has to be updated every time irrespective of read/write hit/miss.