```
boolean primarySeekForDelete(deleteSeekRecord, deleteKey)
 2
 3
       while(true)
 4
       {
 5
         loop until a null node is reached OR promote flag is set on the node OR deleteKey is
         if keyfound
 6
 7
 8
           if the key is marked
 9
             restart primary seek
10
           else
11
             update deleteSeekRecord with pnode, node, lastUnmarkedEdge
12
             return true
13
         }
14
         else
15
         {
16
           if lastRightNode's key has not changed
17
             return false
18
19
             restart primary seek
20
         }
21
       }
22
     }
23
     //if special case secondary seek returns true
24
     boolean secondarySeekForDelete(node,nRChild,secDeleteSeekRecord)
25
26
       if special case
27
         return true
28
       else
29
         keep going left until the secondary node is located
         once the secondary node is located, populate the secDeleteSeekRecord
30
31
         with secondary node, its left & right child, its parent and the
         lastSecondaryUnmarkedEdge
32
         return false
33
     }
34
35
    delete pseudocode
36
     set CLEANUP and isSimpleDelete flags to false
37
    while(true)
38
39
       if (primarySeekForDelete (deleteSeekRecord, deleteKey) returned false)
         return false
40
41
       else
         use the values of pnode, node and lastUnmarkedEdge from the deleteSeekRecord
42
43
       if (!CLEANUP)
44
45
         try CAS (node->1Child, <nlChildAddr, 0, 0>, <nlChildAddr, 1, 0>)
46
47
         if CAS FAILED
48
           help
49
           continue from top of primary seek's while loop
50
         if CAS SUCCEEDED
51
           set CLEANUP to true and set storedNode = node
52
53
       if(storedNode != node) //Someone removed the node for me. So DONE
```

```
54
        set "deleteFlag" on node's rChild using BTS
 55
        if complex delete
 56
 57
          while(true) //secondary seek
 58
 59
            nRChild = node->rChild
 60
            if(nRChild != NULL)
              isSplCase = secondarySeekForDelete (node, nRChild, secDeleteSeekRecord)
              use the values of secondary node, its lchild, its rchild, its parent,
 62
              SeclastUnmarkedEdge from secDeleteSeekRecord
 63
            else
 64
              set isSimpleDelete flag to true
 65
              break from while loop
            if node key is unmarked
 66
 67
            {
 68
              try CAS(rnode->1Child, <NULL, 0, 0>, <nodeAddr, 0, 1>)
 69
              if CAS failed
 70
 71
                if promoteFlag is set
                  if address does not match with node's address
 72
 73
                    restart primary seek. assert (node->secFlag == DONE)
 74
                    break from while loop
 75
                else
 76
                   if address != NULL //restart secondary seek
 77
                     continue from top of secondary seek's while loop
 78
                  else
 79
                     assert(rnode->1Child's deleteFlag is set)
 80
                    help operation at secondaryLastUnmarkedEdge
                    if secondaryLastUnmarkedEdge does not exist, then help node->rChild
 81
                     //simplehelp(node,nrChild)
                     continue from top of secondary seek's while loop
 82
 83
              }
 84
              set promote flag on rnode->rChild using BTS
 85
              promote key using a simple write. Node's key changes from <0,kN> to <1,kRN>
 86
 87
            if(!isSplCase)
 88
 89
              try CAS(rpnodeLChild,<rnode,0,0>,<rnodeRChild,0,0>) //remove secondary node
              if CAS FAILED, help operation at secondaryLastUnmarkedEdge
 90
                if secondaryLastUnmarkedEdge doesn't exist, override CASinvariant and help
 91
                node->rChild //simplehelp(node,nrChild)
 92
                continue from top of secondary seek's while loop
 93
              if CAS SUCCEEDED, set node->secDoneFlag to true
 94
            }
 95
            else
 96
            {
                try CAS (nodeRChild, <rnode, 1, 0>, <rnodeRChild, 1, 0> //no problem if CAS fails
 97
 98
                set node->secDoneFlag to true
 99
            }
            oldNodeAddr = address of node
100
101
            while (true)
102
            {
              create a fresh copy of node
103
104
              newNodeKey as <0, kRN>
              newNodeLChild as <node's lChildAddr,0,0>
105
```

```
106
              newNodeRChild = <node's rChildAddr,0,0>
107
              try CAS (pnode->lChild, <node, 0, 0>, <newNode, 0, 0>)
              if CAS SUCCEEDED then DONE
108
109
              if CAS FAILED
110
                if address has changed
111
                  then someone helped me install a fresh copy.so DONE
112
                else
113
                  CAS has failed coz the edge is marked.
114
                  if lastUnmarkedEdge is not (pnode, node) then help
115
                  do primarySeekForDelete(node->key) //restart primary seek with new key
116
                  if the new key is not found then someone has installed a fresh copy. So done
117
                  if key is found and newNodeAddr != oldNodeAddr then someone has installed a
                  fresh copy. So done
118
            }
119
          }
120
        }
121
        else //simple delete
          set isSimpleDelete to true
122
123
        if(isSimpleDelete)
124
125
          try CAS(pnode->1Child, <node, 0, 0>, <node's 1/r child, 0, 0>)
          if CAS SUCCEEDED, then DONE
126
127
          if CAS FAILED
128
            if lastUnmarkedEdge is NOT (pnode, node) help
129
       }
130
      }
131
      //simplehelp(node, node's rChild)
132
      simplehelp(pnode, node)
      assert(pnode's secDoneFlag not set)
133
134
      set delete flag on node->rChild using BTS
135
      if complex delete
136
       if node->secDoneFlag is set
          create a fresh copy of node
137
138
          try CAS(pnode->rChild, <node, 1, 0>, <newNode, 1, 0>)
139
      else //simple delete
140
        try CAS(pnode->rchild, <node, 1, 0>, <node's lchild, 1, 0>)
```