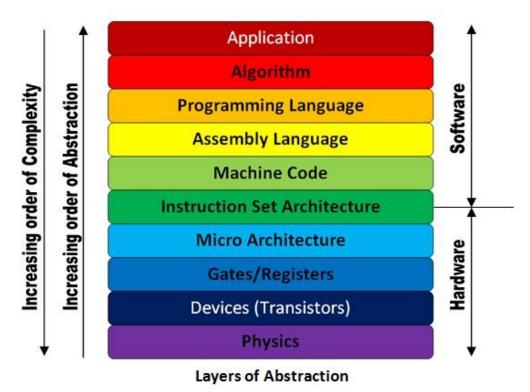
Low Level Security

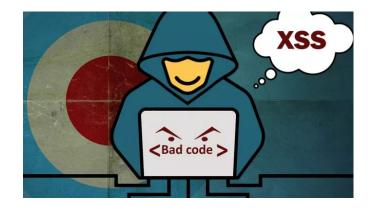
Thushara Desapriya

Why Low Level



Web Application Issues







Insecure Deserialization

Treat it Right

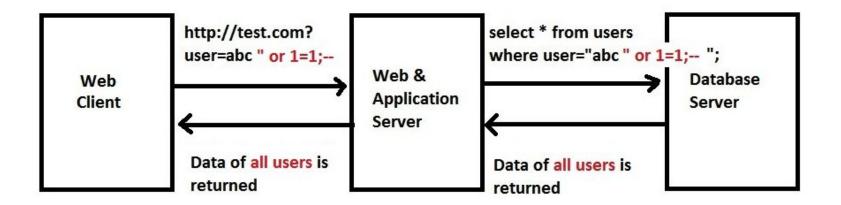
Data should not become part of the Code

Injections

Cross Site Scripting



SQL Injection



SQL Injection Example

SQL Injection

```
String SQLQuery = "SELECT Username, Password
 FROM users WHERE Username='" + Username +
    "' AND Password='" + Password +"'";
Statement stmt = connection.oreateStaement();
ResultSet rs = stmt.executeQuery(SQLQuery);
while (rs.next()) { ... }
```

SQL Injection

UserName = abc

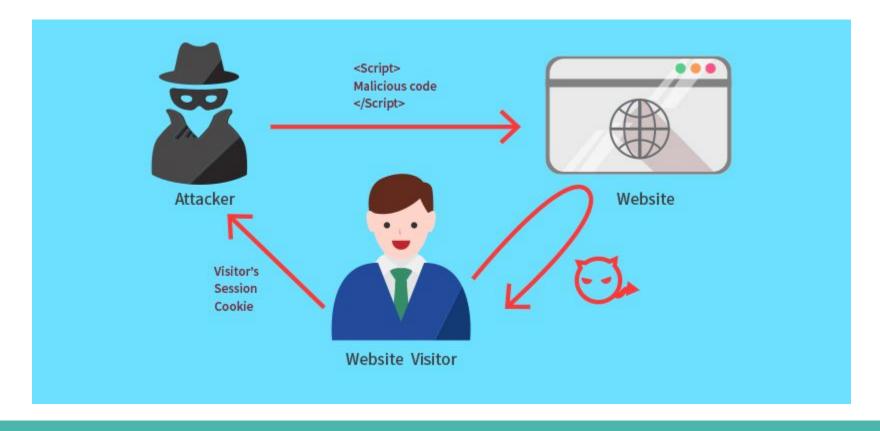
Password= xxx' or '1'='1

SELECT Username, Password FROM users WHERE

Username='abc' AND

Password='**xxx' or '1'='**1';

Cross Site Scripting - XSS

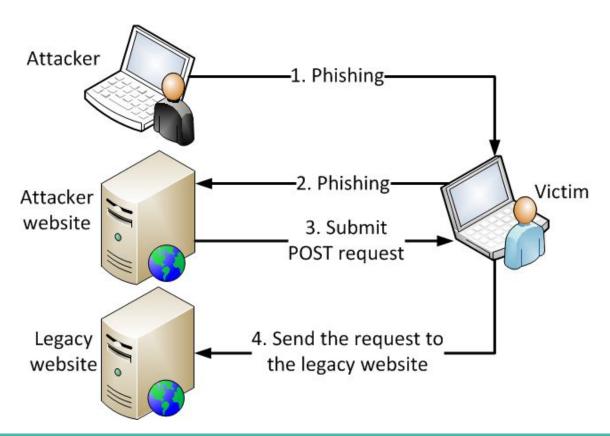


Trick the User

No Patch for Human Stupidity



CSRF - Cross-Site Request Forgery



Clickjacking



User provides username and password.
All these clicks are hijacked by the invisible frame.

Unvalidated Redirects



Buffer Overflow

Buffer overflows were understood and partially publicly documented as early as 1972

1988 Morris Worm

Shutdown over 6000 systems in just a few short hours Exploited fingered through exploiting a unchecked buffer initialized by gets()

Among SANS top 25 most dangerous software errors (Ref : http://www.sans.org/top25-software-errors/)

Buffer Overflow

2001 the Code Red worm exploited a buffer overflow in Microsoft's Internet Information Services (IIS) 5.0

2003 the SQL Slammer worm compromised machines running Microsoft SQL Server 2000.

2003, buffer overflows present in licensed Xbox games have been exploited to allow unlicensed software, including homebrew games, to run on the console without the need for hardware modifications, known as modchips.

The PS2 Independence Exploit also used a buffer overflow to achieve the same for the PlayStation 2.

The Twilight hack accomplished the same with the Wii, using a buffer overflow in The Legend of Zelda: Twilight Princess

buffer overflow, or buffer overrun, is an anomaly where a program, while writing data to a buffer, overruns the buffer's boundary and overwrites adjacent memory locations



A buffer overflow occurs when data written to a buffer also corrupts data values in memory addresses adjacent to the destination buffer due to insufficient bounds checking.

This can occur when copying data from one buffer to another without first checking that the data fits within the destination buffer.



```
char A[8] = "";
unsigned short B = 1979;
```

Initially, A contains nothing but zero bytes, and B contains the number 1979.

variable name						В				
value						19	979			
hex value	00	00	00	00	00	00	00	00	07	ВВ

```
strcpy(A, "excessive");
```

"excessive" is 9 characters long and encodes to 10 bytes including the null terminator, but A can take only 8 bytes. By failing to check the length of the string, it also overwrites the value of B:

variable name	Α					В				
value	'e'	'X'	'c'	'e'	's'	's'	'i'	۱۷'	258	356
hex	65	78	63	65	73	73	69	76	65	00

```
strlcpy(A, "excessive", sizeof(A));
```

Buffer Overflow - Solution

```
strlcpy(A, "excessive", sizeof(A));
```

Can be prevented easily through using bounds checking

```
#include <stdio.h>
int main(int argc, char **argv)
{
     char buf[8];
     printf("Type your message : ");
     fgets(buf,8,stdin);
     printf("My Message : %s \n", buf);
     return 0;
}
```

Writing data past the end of allocated memory can sometimes be detected by the operating system to generate a segmentation fault error that terminates the process.

Segmentation Fault: Fault raised by hardware with memory protection notifying an Operating System about a memory access violation

Buffer Overflow - When

Data typed into text fields of a GUI

Data sent via a program

Data provided in a file

Data provided via command line

Data provided in Environment Variables

Buffer Overflow - Damage

So Buffer overflow can raise a memory access violation error and halt a program execution which may cause a denial of service.

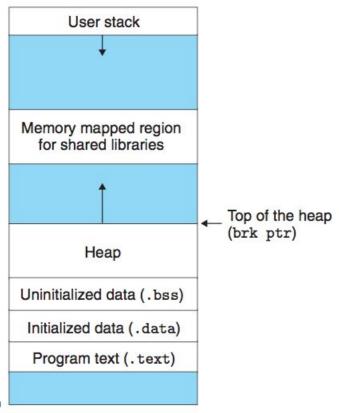
Corruption of other program data

Shutdown of the computer running the program

Corruption of Control flow structure of the program

Be able to execute code on the computer running the program

Memory Structure



Heap

 Memory Allocated to variables during runtime using malloc()

Stack

 Used to store Local variables which are only used in one function. These variables are stored in a element called as frame

Data

 Readable and writable segments containing the static, global, initialized data segments and variables

BSS

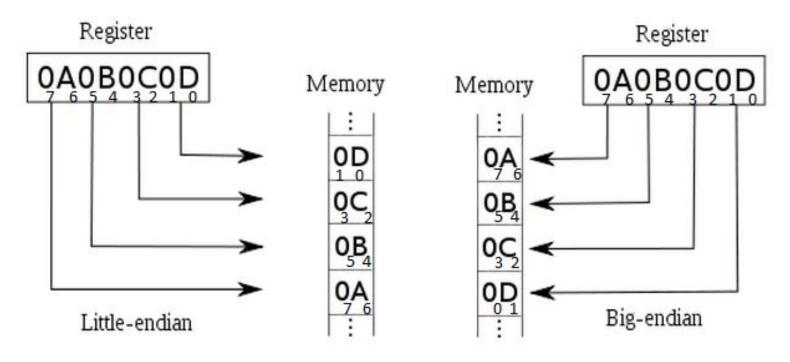
 Readable and writable segments containing the static, global, un-initialized data segments and variables

Text

 Read only segment that contains the compiled executable of the program

Big-endian and Little-endian

Big-endian and little-endian are terms that describe the order in which a sequence of bytes are stored in computer memory



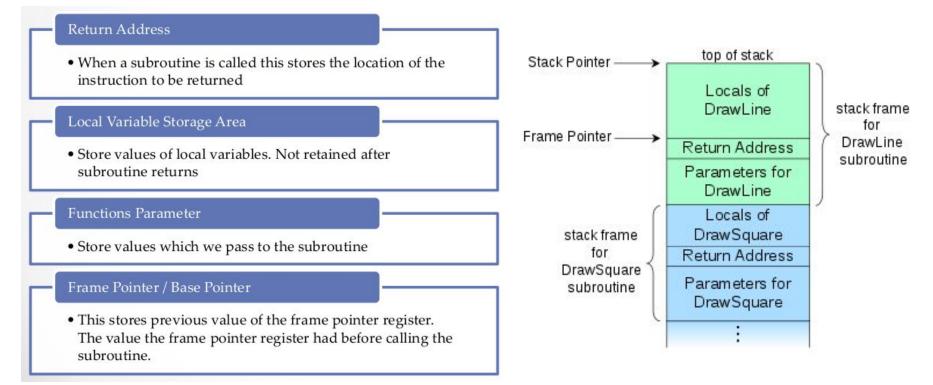
Registers

Registers are either 32 or 64 bit high-speed storage locations directly inside the CPU, designed for high-speed access. For the purposes of discussion, registers can be grouped into the four categories of Data, Segment, Index, and Control

REGISTER SET INTEL 80386 ARCHITECTURE								
Category	Register	Function						
Data	EAX (accumulator) EBX (base) ECX (counter) EDX (data)	Used for arithmetic and data movement. Each register can be addressed as either a 16 or 32 bit value. EBX can hold the address of a procedure or variable.						
Segment	CS (code segment) DS (data segment) SS (stack segment) ES (extra segment) FS & GS	Used as base locations for program instructions, data, and the stack. All references to memory involve a segment register used as a base location.						
Index	EBP (base pointer) ESP (stack pointer) ESI (source index) EDI (destination index)	Contain the offsets of data and instructions. The term offset refers to the distance of a variable or instruction from its base segment. The stack pointer contains the offset of the top of the stack						
Control	EIP (instruction pointer) EFLAGS	The instruction pointer always contains the offset of the next instruction to be executed within the current code segment.						

Stack Frame

Stack has frame which stores routines data



Stack Buffer Overflow (Stack Smashing)

By overwriting a local variable that is located near the vulnerable buffer on the stack, in order to change the behavior of the program

By overwriting the return address in a stack frame. Once the function returns, execution will resume at the return address as specified by the attacker - usually a user-input filled buffer

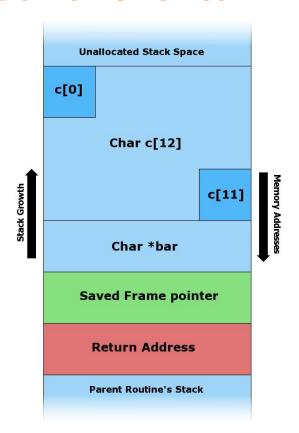
By overwriting a function pointer or exception handler, which is subsequently executed

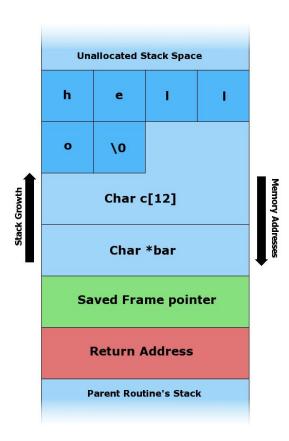
By overwriting a local variable (or pointer) of a different stack frame, which will be used by the function which owns that frame later

Stack Buffer Overflow

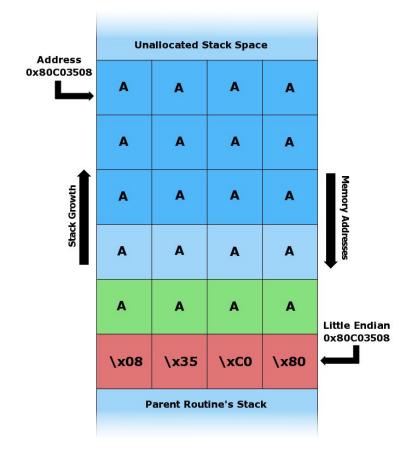
```
#include <string.h>
void foo (char *bar)
   char c[12];
   strcpy(c, bar); // no bounds checking
int main (int argc, char **argv)
   foo(argv[1]);
   return 0;
```

Stack Buffer Overflow





Stack Buffer Overflow



Heap Overflow

Different from stack-based overflows.

Memory on the heap is dynamically allocated by the application at run-time and typically contains program data.

Exploitation is performed by corrupting this data in specific ways to cause the application to overwrite internal structures such as linked list pointers.

The canonical heap overflow technique overwrites dynamic memory allocation linkage (such as malloc metadata) and uses the resulting pointer exchange to overwrite a program function pointer

Read Overflow

Similar to writing pass the end of a buffer, read overflow could permit reading past the end of buffer leaking secret information

The heartbleed bug

- SSL Server should accept heartbeat message and echoes it back
- Heartbeat message specifies the length of its echo-back portion, but the buggy SSL software did not check whether the length was accurate
- Therefore an attacker was able to request longer length and read pass the contents of the buffer leaking passwords, crypto keys etc

Counter Measures

Libsafe project

Re-implemented functions like strcpy(), strcat() and gets()

Stack Guard

Achieved by a compiler extension that adds so called canary values before the EIP saved at the function

Function Arguments

Return Address

Canary Value

Frame Pointer

Local Variables

Remote Code Execution (RCE)

```
<?php
 if (isset($_GET['domain'])) {
   echo '';
   $domain = _GET['domain'];
   $lookup = system("nslookup {$domain}");
   echo($lookup);
   echo '';
?>
```

Notice how the 'domain' parameter is taken in from the GET request, and immediately interpolated into a command string.

Questions

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