

Introduction to syntax analysis

Study Guide

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2.3 Parse Tree and Ambiguity

A **parse tree** (also called a **syntax tree** or **derivation tree**) is a **tree representation of the syntactic structure** of a string according to a grammar.

It shows **how a start symbol of a grammar derives (produces)** a given input string.

Structure of a Parse Tree

- **Root node** → represents the **start symbol** of the grammar.
- **Internal nodes** → represent **non-terminals**.
- **Leaf nodes** → represent **terminals (tokens)** of the input string.
- The tree visually shows the **derivation steps** according to production rules.

Example

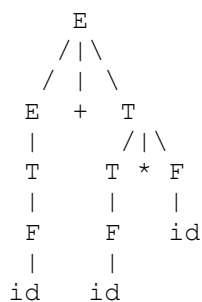
Consider the grammar:

$$\begin{aligned} E &\rightarrow E + T \mid T \\ T &\rightarrow T * F \mid F \\ F &\rightarrow (E) \mid id \end{aligned}$$

and the input string:

`id + id * id`

One possible parse tree is:



This shows that the expression is parsed as:

$$E \rightarrow E + T \rightarrow T + T \rightarrow F + T \rightarrow id + T \rightarrow id + T * F \rightarrow id + F * F \rightarrow id + id * id$$

Ambiguity in Grammars

A grammar is **ambiguous** if there exists at least one string that can have **more than one distinct parse tree (or derivation)**.

Why It Happens

Ambiguity arises when the grammar does not clearly define **operator precedence** or **associativity** rules.

Example of Ambiguous Grammar

$E \rightarrow E + E \mid E * E \mid (E) \mid id$

For the string:

`id + id * id`

There are **two possible parse trees**:

1. **(+ first)** $\rightarrow (id + id) * id$
2. **(first)*** $\rightarrow id + (id * id)$

This means the grammar is **ambiguous** because the same input has **two different structures** (and hence two possible meanings).

How to Remove Ambiguity

To eliminate ambiguity, we rewrite the grammar to enforce **operator precedence and associativity**.

Example (Disambiguated Grammar)

$E \rightarrow E + T \mid T$ // '+' has lowest precedence
 $T \rightarrow T * F \mid F$ // '*' has higher precedence
 $F \rightarrow (E) \mid id$ // parentheses and operands

Now:

- Multiplication ***** binds tighter than **+**
- Both operators are **left-associative**

This grammar is **unambiguous** for arithmetic expressions.

3. Summary of Key Concepts

- **Parse Tree:**

A hierarchical tree showing how a string is derived from a grammar's start symbol using production rules.

- Root → start symbol
- Internal nodes → non-terminals
- Leaves → terminals

Example: For $id + id * id$, the parse tree shows the order of operations and structure of the expression.

- **Ambiguity:**

A grammar is **ambiguous** if a single string can have **more than one valid parse tree** (e.g., $id + id * id$ can be parsed as either $(id + id) * id$ or $id + (id * id)$).

- **Removing Ambiguity:**

Redefine grammar rules to enforce **operator precedence** and **associativity** (e.g., separate non-terminals for $+$, $*$, and operands).

Next Steps

- Explore other basic functions of Removal left factoring and left recursion .

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