

## LAB 7 - Q2

### **Module 1:** A simple counter

#### **Task 1:**

SPEC1 AG! (p.n=4) : From now on p.n state will never be 4

SPEC2 AG (p.n=1  $\rightarrow$  AF p.n=4) : Whenever we will get state 1, in future we will definitely get state 4

#### **Task 2:**

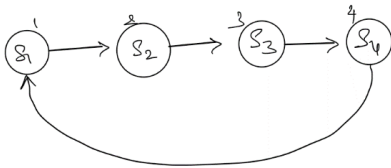
SPEC1 AG! (p.n=4): is false, as we will eventually reach 4. Counterexample is shown in the below screenshot

SPEC2 AG (p.n=1  $\rightarrow$  AF p.n=4): is true as it is a counter whenever we have 1 we will eventually reach 4

#### **Task 3:**

- Number of state variables = 1 i.e n  
p.n : {1,2,3,4}
- Reachable states:  
reachable states: 4 ( $2^2$ ) out of 4 ( $2^2$ )  
p.n = 4  
p.n = 2  
p.n = 3  
p.n = 1
- There is only 1 state available to pick i.e. "p.n = 1"
- Simulation is shown in the below screenshots
- Traces have been saved

#### **Transition Diagram:**



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```

NuSMV > show_vars
Number of Input Variables: 0
Number of State Variables: 1
1: {1, 2, 3, 4}
Number of Frozen Variables: 0
p.n : {1, 2, 3, 4}

Number of bits: 2 (0 frozen, 0 input, 2 state)
NuSMV > print_reachable_states -v
#####
system diameter: 4
reachable states: 4 (2^2) out of 4 (2^2)
----- State 1 -----
p.n = 4
----- State 2 -----
p.n = 2
----- State 3 -----
p.n = 3
----- State 4 -----
p.n = 1
-----
#####
NuSMV > pick_state -1
***** AVAILABLE STATES *****
===== State =====
0) -----
p.n = 1

There's only one available state. Press Return to Proceed.

Chosen state is: 0
NuSMV > simulate -i -k 3
***** Simulation Starting From State 1.1 *****
***** AVAILABLE STATES *****
===== State =====
0) -----
p.n = 2

0) -----
p.n = 3

There's only one available state. Press Return to Proceed.

Chosen state is: 0
NuSMV > show_traces -v -o trace_module3.txt
<!-- ##### Trace number: 1 ##### -->
NuSMV > check_ctlspec -p "AG !(p.n=4)"
unknown command 'check'
NuSMV > check_ctlspec -p "AG !(p.n=4)"
-- specification AG !(p.n = 4) is false
-- as demonstrated by the following execution sequence
Trace Description: CTL Counterexample
Trace Type: Counterexample
-> State: 2.1 <-
p.n = 1
-> State: 2.2 <-
p.n = 2
-> State: 2.3 <-
p.n = 3
-> State: 2.4 <-
p.n = 4
NuSMV > check_ctlspec -p "AG !(p.n=1 -> AF p.n=4)"
-- specification AG !(p.n = 1 -> AF p.n = 4) is false
-- as demonstrated by the following execution sequence
Trace Description: CTL Counterexample
Trace Type: Counterexample
-> State: 3.1 <-
p.n = 1
NuSMV > check_ctlspec -p "AG (p.n=1 -> AF p.n=4)"
-- specification AG (p.n = 1 -> AF p.n = 4) is true

```

### Module 2: Semaphores

#### Task 1:

SPEC1 - "AG !(proc1.state = critical & proc2.state =critical)" : It will never happen that both the process are in critical state

SPEC2 - "AG(proc1.state = entering -> AF (proc1.state = critical))": It is possible that process 1 is entering state and may never reach critical

#### Task 2:

SPEC1 - "AG !(proc1.state = critical & proc2.state =critical)" : This specification is true as both the processes cannot be in critical state simultaneously

SPEC2 - "AG(proc1.state = entering -> AF (proc1.state = critical))" is false as it is not always true that if the process is in entering state it will go to critical stage . The counterexample is shown below

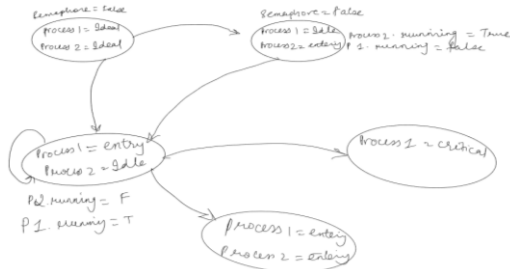
#### Task 3:

- Number of Input Variables: 1  
1: {main, proc2, proc1}
- Number of State Variables: 3  
2: {idle, entering, critical, exiting}  
1: boolean
- Number of Frozen Variables: 0  
semaphore : boolean
- proc1.state : {idle, entering, critical, exiting}
- proc2.state : {idle, entering, critical, exiting}
- \_process\_selector\_ : {main, proc2, proc1}
- Number of bits: 7 (0 frozen, 2 input, 5 state)

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- Reachable States:  
system diameter: 5  
reachable states: 12 ( $2^{3.58496}$ ) out of 32 ( $2^5$ )
- Pick state – Only 1 available state which is 0
- Simulation is shown in below screenshots
- Traces have been saved

### Transition Diagram



```
NuSMV > show_vars
Number of Input Variables: 1
1: {main, proc2, proc1}
Number of State Variables: 3
1: {idle, entering, critical, exiting}
1: boolean
Number of Frozen Variables: 0
semaphore : boolean

proc1.state : {idle, entering, critical, exiting}
proc2.state : {idle, entering, critical, exiting}
_process_selector_ : {main, proc2, proc1}

Number of bits: 7 (0 frozen, 2 input, 5 state)
NuSMV > print_reachable_states -v
#####
system diameter: 5
reachable states: 12 ( $2^{3.58496}$ ) out of 32 ( $2^5$ )
----- State 1 -----
semaphore = FALSE
proc1.state = entering
proc2.state = entering
----- State 2 -----
semaphore = FALSE
proc1.state = idle
proc2.state = entering
----- State 3 -----
semaphore = FALSE
proc1.state = entering
proc2.state = idle
----- State 4 -----
semaphore = FALSE
proc1.state = idle
proc2.state = idle
----- State 5 -----
semaphore = TRUE
proc1.state = entering
proc2.state = exiting
----- State 6 -----
semaphore = TRUE
proc1.state = idle
proc2.state = exiting
----- State 7 -----
semaphore = TRUE
proc1.state = entering
proc2.state = critical
----- State 8 -----
semaphore = TRUE
proc1.state = idle
proc2.state = critical
----- State 9 -----
semaphore = TRUE
proc1.state = exiting
proc2.state = entering
----- State 10 -----
semaphore = TRUE
proc1.state = critical
proc2.state = entering
----- State 11 -----
semaphore = TRUE
proc1.state = exiting
proc2.state = idle
----- State 12 -----
semaphore = TRUE
proc1.state = critical
proc2.state = idle
#####
NuSMV > pick_state -i
***** AVAILABLE STATES *****
===== State =====
0) =====
semaphore = FALSE
proc1.state = idle
proc2.state = idle

There's only one available state.
```

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```
NuSMV > simulate -i -k 4
***** Simulation Starting From State 1.11 *****

***** AVAILABLE STATES *****

===== State =====
semaphore = FALSE
proc1.state = entering
proc2.state = entering

This state is reachable through:
0) -----
   _process_selector_ = proc2
   running = FALSE
   proc2.running = TRUE
   proc1.running = FALSE

===== State =====
proc2.state = idle

This state is reachable through:
1) -----
   _process_selector_ = proc2
   running = FALSE
   proc2.running = TRUE
   proc1.running = FALSE

2) -----
   _process_selector_ = main
   running = TRUE
   proc2.running = FALSE

===== State =====
semaphore = TRUE
proc1.state = critical

This state is reachable through:
3) -----
   _process_selector_ = proc1
   running = FALSE
   proc2.running = FALSE
   proc1.running = TRUE
```

Choose a state from the above (0-3): 2

Chosen state is: 2

```
***** AVAILABLE STATES *****

===== State =====
semaphore = FALSE
proc1.state = entering
proc2.state = entering

This state is reachable through:
0) -----
   _process_selector_ = proc2
   running = FALSE
   proc2.running = TRUE
   proc1.running = FALSE

===== State =====
proc2.state = idle

This state is reachable through:
1) -----
   _process_selector_ = proc2
   running = FALSE
   proc2.running = TRUE
   proc1.running = FALSE

2) -----
   _process_selector_ = main
   running = TRUE
   proc2.running = FALSE

===== State =====
semaphore = TRUE
proc1.state = critical

This state is reachable through:
3) -----
   _process_selector_ = proc1
   running = FALSE
   proc2.running = FALSE
```

proc1.running = TRUE

Choose a state from the above (0-3): 3

Chosen state is: 3

```
***** AVAILABLE STATES *****

===== State =====
semaphore = TRUE
proc1.state = critical
proc2.state = entering

This state is reachable through:
0) -----
   _process_selector_ = proc2
   running = FALSE
   proc2.running = TRUE
   proc1.running = FALSE

===== State =====
proc2.state = idle

This state is reachable through:
1) -----
   _process_selector_ = proc1
   running = FALSE
   proc2.running = FALSE
   proc1.running = TRUE

2) -----
   _process_selector_ = main
   running = TRUE
   proc1.running = FALSE

3) -----
   _process_selector_ = proc2
   running = FALSE
   proc2.running = TRUE

===== State =====
```

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```

===== State =====
proc1.state = exiting

This state is reachable through:
4) -----
_process_selector_ = proc1
running = FALSE
proc2.running = FALSE
proc1.running = TRUE

Choose a state from the above (0-4): 4
Chosen state is: 4

***** AVAILABLE STATES *****

===== State =====
semaphore = FALSE
proc1.state = idle
proc2.state = idle

This state is reachable through:
0) -----
_process_selector_ = proc1
running = FALSE
proc2.running = FALSE
proc1.running = TRUE

===== State =====
semaphore = TRUE
proc1.state = exiting
proc2.state = entering

This state is reachable through:
1) -----
_process_selector_ = proc2
running = FALSE
proc2.running = TRUE
proc1.running = FALSE

===== State =====
proc2.state = idle

This state is reachable through:
2) -----
_process_selector_ = proc2
running = FALSE
proc2.running = TRUE
proc1.running = FALSE

3) -----
_process_selector_ = main
running = TRUE
proc2.running = FALSE

Choose a state from the above (0-3): 0
Chosen state is: 0
NuSMV > show_traces -v -o trace_module4.txt
<!-- ##### Trace number: 1 ##### -->
NuSMV > check_ctlspec -p "AG ! (proc1.state = critical & proc2.state = critical)"
-- specification AG ! (proc1.state = critical & proc2.state = critical) is true
NuSMV > check_ctlspec -p "AG (proc1.state = entering -> AF (proc1.state = critical))"
-- specification AG (proc1.state = entering -> AF (proc1.state = critical)) is false
-- as demonstrated by the following execution sequence
Trace Description: CTL Counterexample
Trace Type: Counterexample
-> State: 2.1 <-
semaphore = FALSE
proc1.state = idle
proc2.state = idle
-> Input: 2.2 <-
_process_selector_ = proc1
running = FALSE
proc2.running = FALSE
proc1.running = TRUE
-- Loop starts here
-> State: 2.2 <-
proc1.state = entering
-> Input: 2.3 <-
_process_selector_ = proc2
proc2.running = TRUE
proc1.running = FALSE
-> State: 2.3 <-

```

### Module 3: Spin

#### Task 1:

check\_ctlspec -p "AG (request -> AF state = busy)": Whenever we have a request we will at some point of time get busy

check\_ltlspec -p "G (request -> AF state = busy)": This is not a valid LTL formula as AF is a CTL property

check\_ctlspec -p "G (request -> F state = busy)": This is not a valid CTL property. However, it is a valid LTL property which states that whenever we have a request we will get busy

check\_ctlspec -p "EG (request -> AG state = busy)": states that sometime when we have request we will always be busy

#### Task 2:

"AG (request -> AF state = busy)": is true

"G (request -> AF state = busy)": Not a valid formula

"G (request -> F state = busy)": is true

"EG (request -> AG state = busy)": Is not valid as we will eventually come out of busy state. Counterexample is shown below:

#### Task 3:

- Number of Input Variables: 0  
Number of State Variables: 2  
1: boolean  
1: {ready, busy}  
Number of Frozen Variables: 0  
request : boolean  
  
state : {ready, busy}
- Reachable states:  
reachable states: 4 ( $2^2$ ) out of 4 ( $2^2$ )  
----- State 1 -----  
request = TRUE

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```

state = busy
----- State 2 -----
request = TRUE
state = ready
----- State 3 -----
request = FALSE
state = busy
----- State 4 -----
request = FALSE
state = ready

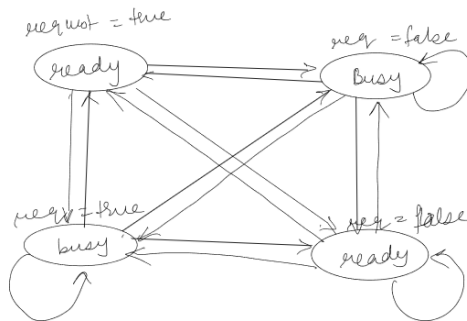
```

- There are 2 available states:
  - 0) -----
 

```
request = TRUE
state = ready
```
  - 1) -----
 

```
request = FALSE
state = busy
```
- Simulation is shown in screenshot
- Traces have been saved

### Transition Diagram



```

NuSMV > read_model -i module5.smv
NuSMV > go
NuSMV > show_vars
Number of Input Variables: 0
Number of State Variables: 2
  1: {ready, busy}
  1: boolean
Number of Frozen Variables: 0
  request : boolean

state : {ready, busy}

Number of bits: 2 (0 frozen, 0 input, 2 state)
NuSMV > print_reachable_states -v
#####
system diameter: 2
reachable states: 4 (2^2) out of 4 (2^2)
----- State 1 -----
request = TRUE
state = busy
----- State 2 -----
request = TRUE
state = ready
----- State 3 -----
request = FALSE
state = busy
----- State 4 -----
request = FALSE
state = ready
#####
NuSMV > pick_states -i
unknown command 'pick'
NuSMV > pick_state -i

***** AVAILABLE STATES *****

===== State =====
0) -----
request = TRUE
state = busy

===== State =====
1) -----
request = FALSE
state = ready

===== State =====
2) -----
request = FALSE
state = busy

===== State =====
3) -----
state = ready

NuSMV > simulate -i -k 4
***** Simulation Starting From State 1.1 *****

***** AVAILABLE STATES *****

===== State =====
0) -----
request = TRUE
state = busy

===== State =====
1) -----
request = FALSE
state = ready

===== State =====
2) -----
request = FALSE
state = busy

===== State =====
3) -----
state = ready

Choose a state from the above (0-1): 0
Chosen state is: 0

```

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```
Choose a state from the above (0-3): 3
Chosen state is: 3
***** AVAILABLE STATES *****

===== State =====
0) -----
request = TRUE
state = busy

===== State =====
1) -----
state = ready

===== State =====
2) -----
request = FALSE
state = busy

===== State =====
3) -----
state = ready

Choose a state from the above (0-3): 2
Chosen state is: 2
***** AVAILABLE STATES *****

===== State =====
0) -----
request = TRUE
state = busy

===== State =====
1) -----
state = ready

===== State =====
2) -----
request = FALSE
state = busy

===== State =====
3) -----
state = ready

NuSMV > show_traces -v -o trace_module4.txt
<!-- ##### Trace number: 1 ##### -->
NuSMV > check_ctlspec -p "AG (request -> AF state = busy)"
-- specification AG (request -> AF state = busy) is true
NuSMV > check_ltlspec -p "G (request -> AF state = busy)"

file <command-line>: line 10: : unexpected LTL operator
Parsing error: expected an "CTL" expression.
NuSMV > check_ltlspec -p "G (request -> AF state = busy)"

file <command-line>: line 10: : unexpected CTL operator
Parsing error: expected an "LTL" expression.
NuSMV > check_ltlspec -p "G (request -> AF state = busy)"

file <command-line>: line 10: : unexpected CTL operator
Parsing error: expected an "LTL" expression.
NuSMV > check_ltlspec -p "G (request -> F state = busy)"
-- specification G (request -> F state = busy) is true
NuSMV > check_ctlspec -p "EG (request -> AG state = busy)"
-- specification EG (request -> AG state = busy) is false
-- as demonstrated by the following execution sequence
Trace Description: CTL Counterexample
Trace Type: Counterexample
-> State: 2.1 <-
    request = TRUE
    state = ready
NuSMV >
```

### Module 4: Traffic lights

**Task 1:** ! F (tl1.state = g & tl2.state = g) : Both the greens lights of 2 states cannot be green at same time

**Task 2:** ! F (tl1.state = g & tl2.state = g) the property holds true

#### **Task 3:**

- Number of variables  
Number of Input Variables: 1  
1: {main, tl2, tl1}  
Number of State Variables: 2  
2: {r, y, g}  
Number of Frozen Variables: 0  
tl1.state : {r, y, g}  
tl2.state : {r, y, g}  
\_process\_selector\_ : {main, tl2, tl1}  
Number of bits: 6 (0 frozen, 2 input, 4 state)
- 5 out of 9 are reachable states
- Only 1 state available to pick
- Simulation is shown below
- Traces were saved

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```

NuSMV > real_model -i module6.smv
unknown command 'real_model'
NuSMV > read_model -i module6.smv

file module6.smv: line 1: cannot open input file module6.smv
NuSMV > read_model -i module6.smv
NuSMV > go
WARNING *** Processes are still supported, but deprecated. ***
WARNING *** In the future processes may be no longer supported. ***

WARNING *** The model contains PROCESSES or ISAs. ***
WARNING *** The HRC hierarchy will not be usable. ***
NuSMV > show_vars
Number of Input Variables: 1
  1: {main, tl2, tl1}
Number of State Variables: 2
  2: {r, y, g}
Number of Frozen Variables: 0
  tl1.state : {r, y, g}

  tl2.state : {r, y, g}

  _process_selector_ : {main, tl2, tl1}

Number of bits: 6 (0 frozen, 2 input, 4 state)
NuSMV > print_reachable_states -v
#####
system diameter: 3
reachable states: 5 (2^2.32193) out of 9 (2^3.16993)
----- State 1 -----
  tl1.state = r
  tl2.state = g
----- State 2 -----
  tl1.state = r
  tl2.state = r
----- State 3 -----
  tl1.state = r
  tl2.state = y
----- State 4 -----
  tl1.state = g
  tl2.state = r
----- State 5 -----
  tl1.state = y
  tl2.state = r

```

```

NuSMV > pick_state -i
***** AVAILABLE STATES *****

===== State =====
0) -----
  tl1.state = r
  tl2.state = r

There's only one available state. Press Return to Proceed.

Chosen state is: 0
NuSMV > simulate -i -k 1
***** Simulation Starting From State 2.1 *****

***** AVAILABLE STATES *****

===== State =====
tl1.state = r
tl2.state = y

This state is reachable through:
0) -----
  _process_selector_ = tl2
  running = FALSE
  tl2.running = TRUE
  tl1.running = FALSE

===== State =====
tl2.state = r

This state is reachable through:
1) -----
  _process_selector_ = tl1
  running = FALSE
  tl2.running = FALSE
  tl1.running = TRUE

2) -----
  _process_selector_ = main
  running = TRUE
  tl1.running = FALSE

```

```

  tl1.running = FALSE

3) -----
  _process_selector_ = tl2
  running = FALSE
  tl2.running = TRUE

===== State =====
tl1.state = y

This state is reachable through:
4) -----
  _process_selector_ = tl1
  running = FALSE
  tl2.running = FALSE
  tl1.running = TRUE

Choose a state from the above (0-4): 4
Chosen state is: 4
NuSMV > show_traces -v -o trace_module6.txt
<!-- ##### Trace number: 2 ##### -->
NuSMV > check_ltlspec -p "!(tl1.state = g & tl2.state = g)"
-- specification !(tl1.state = g & tl2.state = g) is true
NuSMV > _

```