

# Formal Languages and Abstract Machines

## Take Home Exam 2

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### 1 Context-Free Grammars (10 pts)

a) Give the rules of the Context-Free Grammars to recognize strings in the given languages where  $\Sigma = \{a, b\}$  and  $S$  is the start symbol.

$L(G) = \{w \mid w \in \Sigma^*; |w| \geq 3;$  (2/10 pts)  
the first and the second from the last symbols of  $w$  are the same}

$$\begin{aligned} S &\rightarrow A \mid Abb \\ A &\rightarrow aA \mid bA \mid a \mid b \end{aligned}$$

$L(G) = \{w \mid w \in \Sigma^*; \text{ the length of } w \text{ is odd}\}$  (2/10 pts)

$$\begin{aligned} S &\rightarrow Aa \mid Ab \\ A &\rightarrow Aaa \mid Aab \mid Aba \mid Abb \mid e \end{aligned}$$

$L(G) = \{w \mid w \in \Sigma^*; n(w, a) = 2 \cdot n(w, b)\}$  where  $n(w, x)$  is the number of  $x$  symbols in  $w$  (3/10 pts)

$$\begin{aligned} A &\rightarrow BBa \mid BaB \mid aBB \mid e \\ B &\rightarrow bS \mid Sb \end{aligned}$$

b) Find the set of strings recognized by the CFG rules given below: (3/10 pts)

$S \rightarrow X \mid Y$   
 $X \rightarrow aXb \mid A \mid B$   
 $A \rightarrow aA \mid a$   
 $B \rightarrow Bb \mid b$   
 $Y \rightarrow CbaC$   
 $C \rightarrow CC \mid a \mid b \mid \varepsilon$

$a^*b^* \mid (a|b)^* ba (a|b)^*$

## 2 Parse Trees and Derivations

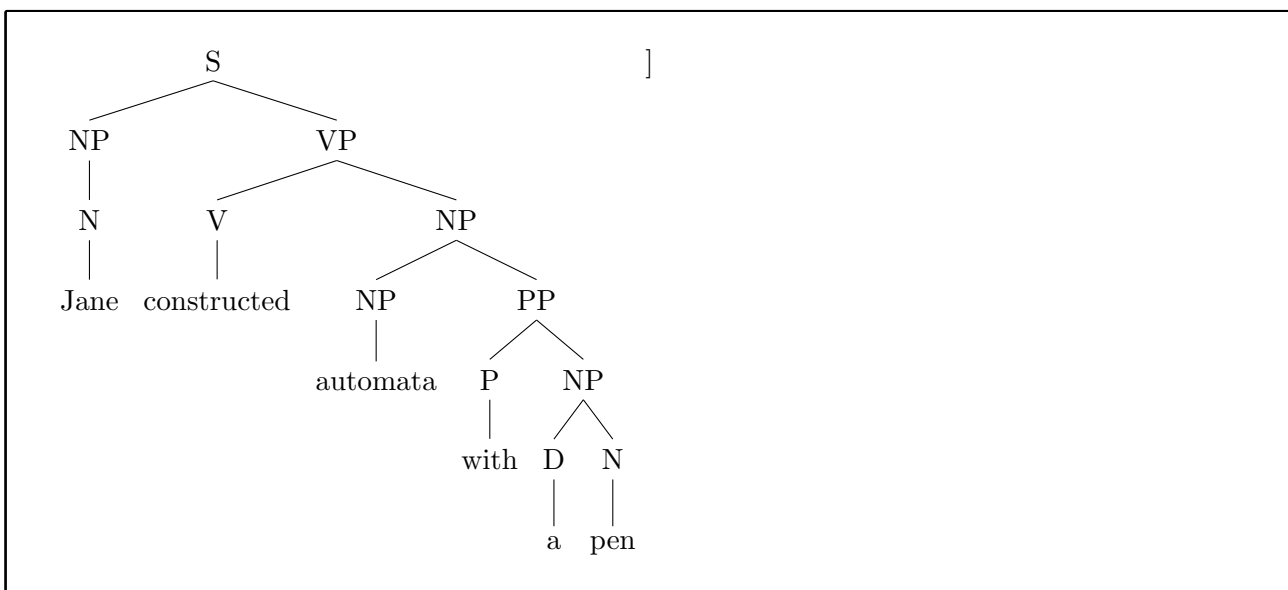
(20 pts)

Given the CFG below, provide parse trees for given sentences in **a** and **b**.

S  $\rightarrow$  NP VP  
 VP  $\rightarrow$  V NP | V NP PP  
 PP  $\rightarrow$  P NP  
 NP  $\rightarrow$  N | D N | NP PP  
 V  $\rightarrow$  wrote | built | constructed  
 D  $\rightarrow$  a | an | the | my  
 N  $\rightarrow$  John | Mary | Jane | man | book | automata | pen | class  
 P  $\rightarrow$  in | on | by | with

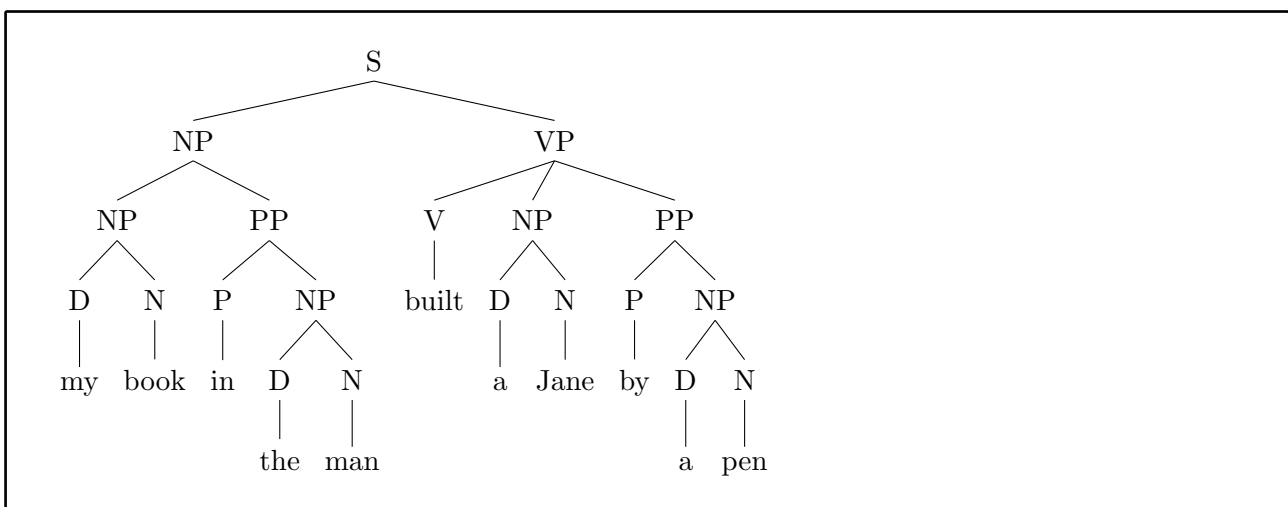
a) Jane constructed automata with a pen

(4/20 pts)



b) my book in the man built a Jane by a pen

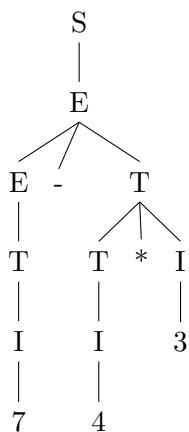
(4/20 pts)



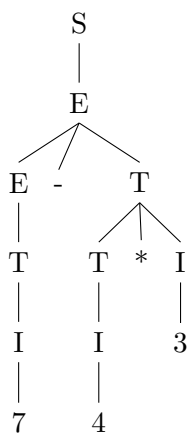
Given the CFG below, answer **c**, **d** and **e**

$$\begin{aligned}
 S &\rightarrow E \\
 E &\rightarrow E + T \mid E - T \mid T \\
 T &\rightarrow T * I \mid T / I \mid I \\
 I &\rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 6 \mid 7 \mid 8 \mid 9
 \end{aligned}$$

**c)** Provide the left-most derivation of  $7 - 4 * 3$  step-by-step and plot the final parse tree matching that derivation (4/20 pts)

$$S \rightarrow E - T \rightarrow T - T \rightarrow I - T \rightarrow 7 - T \rightarrow 7 - T * I \rightarrow 7 - 4 * I \rightarrow 7 - 4 * 3$$


**d)** Provide the right-most derivation of  $7 - 4 * 3$  step-by-step and plot the final parse tree matching that derivation (4/20 pts)

$$S \rightarrow E \rightarrow E - T \rightarrow E - T * I \rightarrow E - T * 3 \rightarrow E - T * 3 \rightarrow E - I * 3 \rightarrow E - 4 * 3 \rightarrow T - 4 * 3 \rightarrow I - 4 * 3 \rightarrow 7 - 4 * 3$$


e) Are the derivations in **c** and **d** in the same similarity class?

(4/20 pts)

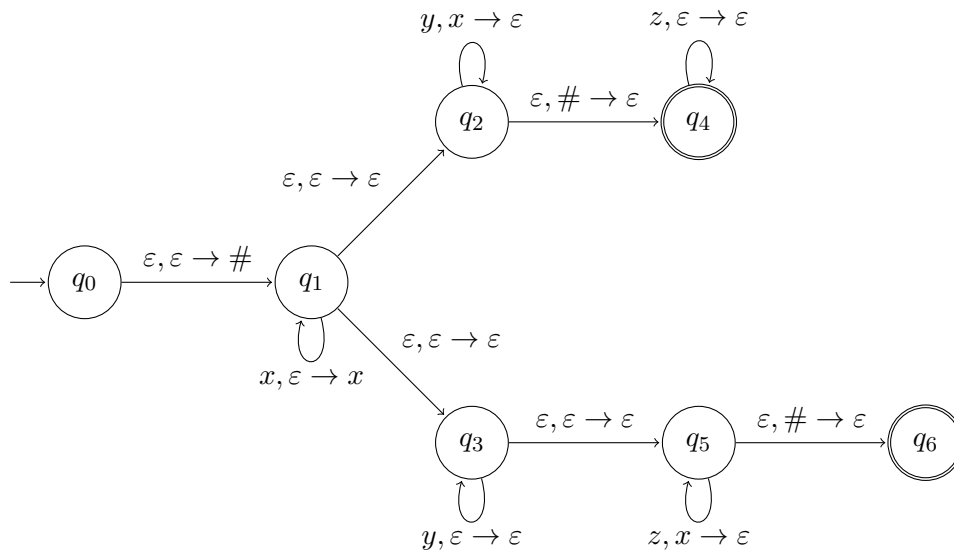
Yes, because one of them precedes the other.

### 3 Pushdown Automata

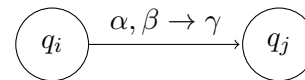
(30 pts)

a) Find the language recognized by the PDA given below

(5/30 pts)



where the transition  $((q_i, \alpha, \beta), (q_j, \gamma))$  is represented as:



answer here ...

b) Design a PDA to recognize language  $L = \{x^n y^{m+n} x^m \mid n, m \geq 0; n, m \in \mathbb{N}\}$

(5/30 pts)

answer here ...

c) Design a PDA to recognize language  $L = \{x^n y^m \mid n < m \leq 2n; n, m \in \mathbb{N}^+\}$  (10/30 pts)

Do not use multi-symbol push/pop operations in your transitions.

Simulate the PDA on strings  $xy$  (with only one rejecting derivation) and  $xyyyy$  (accepting derivation) with transition tables.

answer here ...

**d)** Given two languages  $L'$  and  $L$  as  $L' = \{w \mid w \in L; |w| = 4n + 2 \text{ for } n \in \mathbb{N}\}$  (10/30 pts)  
If  $L$  is a CFL, show that  $L'$  is also a CFL by constructing an automaton for  $L'$  in terms of another automaton that recognizes  $L$ .

answer here ...



## 4 Closure Properties

(20 pts)

Let  $L_1$  and  $L_2$  be context-free languages which are not regular, and let  $L_3$  be a regular language. Determine whether the following languages are necessarily CFLs or not. If they need to be context-free, explain your reasoning. If not, give one example where the language is a CFL and a counter example where the language is not a CFL.

a)  $L_4 = L_1 \cap (L_2 \setminus L_3)$

(10/20 pts)

No.

$L_2 - L_3 = L_2 \cap L_3'$  since  $L_3'$  is regular because regular languages are closed under complement and intersection of them is context free.

$L_1 \cap L_3$  not certainly context free.

$L_1 = q^k z^k x^l$  and  $L_2 - L_3 = z^k x^k$

b)  $L_5 = (L_1 \cap L_3)^*$

(10/20 pts)

Yes.

All regular languages are subset of context-free so that  $L_1 \cap L_3$  is also context-free.

Context-free languages are closed under Kleene Star so that  $L_5$  is CFL.

## 5 Pumping Theorem

(20 pts)

a) Show that  $L = \{a^n m^n t^i \mid n \leq i \leq 2n\}$  is not a Context Free Language using Pumping Theorem for CFLs.

(10/20 pts)

$uvxyz = ammtt$  and  $vxy = am \rightarrow$   
 $uv^2xy^2z = aaammmtt \rightarrow$   
 $n > i$   
so that this language is not Context-Free.

b) Show that  $L = \{a^n b^{2n} a^n \mid n \in \mathbb{N}^+\}$  is not a Context Free Language using Pumping Theorem for CFLs.

(10/20 pts)

$uvxyz = abba$  and  $vxy = abb \rightarrow$   
 $uv^2xy^2z = aabbba \rightarrow$   
 $n = 2, 2n = 3, n = 1$   
so that this language is not Context-Free.

## 6 CNF and CYK

(not graded)

a) Convert the given context-free grammar to Chomsky Normal Form.

$$S \rightarrow XSX \mid xY$$

$$X \rightarrow Y \mid S$$

$$Y \rightarrow z \mid \varepsilon$$

answer here ...

b) Use the grammar below to parse the given sentence using Cocke–Younger–Kasami algorithm. Plot the parse trees.

S → NP VP	VP → book   include   prefer
S → X1 VP	VP → Verb NP
X1 → Aux NP	VP → X2 PP
S → book   include   prefer	X2 → Verb NP
S → Verb NP	VP → Verb PP
S → X2 PP	VP → VP PP
S → Verb PP	PP → Prep NP
S → VP PP	Det → that   this   the   a
NP → I   she   me   Houston	Noun → book   flight   meal   money
NP → Det Nom	Verb → book   include   prefer
Nom → book   flight   meal   money	Aux → does
Nom → Nom Noun	Prep → from   to   on   near   through
Nom → Nom PP	

book the flight through Houston

Empty parse table:

<div> <div>1:5 → 1:1 2:5 1:5 → 1:2 3:5 1:5 → 1:3 4:5 1:5 → 1:4 5:5</div> </div>				
<div> <div>1:4 → 1:1 2:4 1:4 → 1:2 3:4 1:4 → 1:3 4:4</div> <div>2:5 → 2:2 3:5 2:5 → 2:3 4:5 2:5 → 2:4 5:5</div> </div>				
<div> <div>1:3 → 1:1 2:3 1:3 → 1:2 3:3</div> <div>2:4 → 2:2 3:4 2:4 → 2:3 4:4</div> <div>3:5 → 3:3 4:5 3:5 → 3:4 5:5</div> </div>				
<div>1:2 → 1:1 2:2</div>		<div>2:3 → 2:2 3:3</div>	<div>3:4 → 3:3 4:4</div>	<div>4:5 → 4:4 5:5</div>
1:1	2:2	3:3	4:4	5:5
book	the	flight	through	Houston

rest of the answer here ...

## 7 Deterministic Pushdown Automata

(not graded)

Provide a DPDA to recognize the given languages, the DPDA must read its entire input and finish with an empty stack.

a)  $a^*bc \cup a^n b^n c$

answer here ...

**b)**  $(aa)^*c \cup a^nb^nc$

answer here ...