2. (i) Let 
$$S_n := \frac{n}{n^2 + 1} + \frac{n}{n^2 + 2} + \dots + \frac{n}{n^2 + n}$$
.

Define  $T_n := \frac{n}{n^2 + 1} + \frac{n}{n^2 + 1} + \dots + \frac{n}{n^2 + 1}$ 

and  $R_n := \frac{n}{n^2 + n} + \frac{n}{n^2 + n} + \dots + \frac{n}{n^2 + n}$ .

Note that  $R_n \le S_n \le T_n \quad \forall n \in \mathbb{N}$ . (Why?)

Also,  $T_n = \frac{n^2}{n^2 + 1}$  and  $R_n = \frac{n^2}{n^2 + n}$ .

Observe that  $\lim_{n\to\infty} T_n = \lim_{n\to\infty} R_n = 1$ . (Why?)

Thus, by Sandwich Theorem,  $\lim_{n\to\infty} S_n$  exists and is equal to 1.

2. (ii) To find:  $\lim_{n\to\infty} \frac{n!}{n^n}$ .

Observe the following for n > 2:

$$a_n = \frac{n!}{n^n} = \frac{1}{n} \cdot \frac{2}{n} \cdot \dots \cdot \frac{n-1}{n} < \frac{1}{n} \cdot 1 \cdot \dots \cdot 1 = \frac{1}{n}$$

Thus,  $a_n < \frac{1}{n}$  for n > 2. Moreover,  $a_n > 0$  for all  $n \in \mathbb{N}$ .

$$\therefore 0 < a_n < \frac{1}{n} \quad \forall n > 2.$$

As  $\lim_{n\to\infty}\frac{1}{n}=0$ , we have it that  $\lim_{n\to\infty}a_n=0$ , by Sandwich Theorem.

2. (iii) 
$$\lim_{n\to\infty} \left(\frac{n^3+3n^2+1}{n^4+8n^2+2}\right)$$
  
Argue from  $a_n=\left(\frac{n^3+3n^2+1}{n^4+8n^2+2}\right)<\left(\frac{n^3+3n^2+1}{n^4}\right)=\frac{1}{n}+\frac{3}{n^2}+\frac{1}{n^4}$  and  $a_n>0$ .

2. (iv) 
$$\lim_{n \to \infty} (n)^{1/n}$$
.

Define  $h_n := n^{1/n} - 1$ .

Then,  $h_n > 0 \quad \forall n \in \mathbb{N}$ . (Why?)

Observe the following for n > 2:

$$n = (1 + h_n)^n > 1 + nh_n + \binom{n}{2}h_n^2 > \binom{n}{2}h_n^2 = \frac{n(n-1)}{2}h_n^2.$$
 Thus,  $h_n < \sqrt{\frac{2}{n-1}} \quad \forall n > 2.$ 

Using Sandwich Theorem, we get that  $\lim_{n\to\infty}h_n=0$  which gives us that  $\lim_{n\to\infty}n^{1/n}=1$ .

Where did we use that  $h_n > 0$ ?

$$2. (v) \lim_{n \to \infty} \left( \frac{\cos \pi \sqrt{n}}{n^2} \right)$$

For all  $n \in \mathbb{N}$ , we have that  $-1 \le \cos(\pi \sqrt{n}) \le 1$ .

Thus, 
$$\frac{-1}{n^2} \le \frac{\cos \pi \sqrt{n}}{n^2} \le \frac{1}{n^2} \quad \forall n \in \mathbb{N}.$$

Use Sandwich Theorem to argue that  $\lim_{n\to\infty}\frac{\cos\pi\sqrt{n}}{n^2}=0.$ 

2. (vi) 
$$\lim_{n\to\infty} (\sqrt{n}(\sqrt{n+1}-\sqrt{n}))$$

Observe that

$$a_n = \sqrt{n}(\sqrt{n+1} - \sqrt{n}) = \sqrt{n}(\sqrt{n+1} - \sqrt{n}) \cdot \frac{(\sqrt{n+1} + \sqrt{n})}{(\sqrt{n+1} + \sqrt{n})} = \frac{\sqrt{n}}{\sqrt{n+1} + \sqrt{n}}.$$

Thus, 
$$a_n < \frac{\sqrt{n}}{\sqrt{n} + \sqrt{n}} = \frac{1}{2}$$
.

Also, 
$$a_n > \frac{\sqrt{n}}{\sqrt{n+1} + \sqrt{n+1}} = \frac{1}{2}\sqrt{\frac{n}{n+1}} = \frac{1}{2}\sqrt{1 - \frac{1}{n+1}} \ge \frac{1}{2}\left(1 - \frac{1}{\sqrt{n+1}}\right)$$
.

Therefore, we have shown that 
$$\frac{1}{2}\left(1-\frac{1}{\sqrt{n+1}}\right) < a_n < \frac{1}{2}$$
.

Use Sandwich Theorem to argue that  $\lim_{n\to\infty} a_n = \frac{1}{2}$ .

4. (i) Determine whether  $\left\{\frac{n}{n^2+1}\right\}_{n\geq 1}$  is increasing or decreasing.

Let  $a_n$  denote the sequence.

$$a_{n+1} - a_n = \frac{(n+1)}{(n+1)^2 + 1} - \frac{n}{n^2 + 1} = \frac{(n+1)(n^2 + 1) - n((n+1)^2 - 1)}{((n+1)^2 + 1)(n^2 + 1)}$$
 $a_{n+1} - a_n = \frac{-n^2 - n + 1}{((n+1)^2 + 1)(n^2 + 1)} < 0.$ 

 $\therefore a_{n+1} < a_n$ , that is,  $a_n$  is a decreasing sequence.

4. (ii) 
$$a_n = \frac{2^n 3^n}{5^{n+1}}$$
.  
 $a_{n+1} = \frac{2^{n+1} 3^{n+1}}{5^{n+2}} = \frac{6}{5} a_n$ .  
 $\implies a_{n+1} - a_n = \frac{1}{5} a_n > 0$ .

Thus,  $a_n$  is an increasing sequence.

4. (iii) 
$$a_n = \frac{1-n}{n^2}$$
 for  $n \ge 2$ .  
 $a_{n+1} - a_n = \frac{1-(n+1)}{(n+1)^2} - \frac{1-n}{n^2} = \frac{n-1}{n^2} - \frac{n}{(n+1)^2}$ 

$$a_{n+1} - a_n = \frac{(n^2 - n - 1)}{n^2(n+1)^2}$$

The numerator factors as  $(n-\phi)(n+1/\phi)$  where  $1 < \phi < 2$ . Thus, for  $n \ge 2$ , the numerator is positive. Thus, the given sequence is increasing.

5. (i) 
$$a_1 = 1$$
,  $a_{n+1} = \frac{1}{2} \left( a_n + \frac{2}{a_n} \right) \quad \forall n \ge 1$ .

Claim 1.  $a_n > 0 \quad \forall n \in \mathbb{N}$ .

Proof. This can be easily seen via induction. Details are left to the reader.

Claim 2.  $a_n^2 > 2 \quad \forall n \ge 2$ .

*Proof.* We shall prove this via induction. The base case n=2 is immediate as  $a_n=3/2$ .

Assume that it holds for n = k.

$$a_{k+1}^2 - 2 = \frac{1}{4} \left( a_k + \frac{2}{a_k} \right)^2 - 2 = \frac{(a_k^2 - 2)^2}{4a_k^2}$$

 $(a_k^2-2) \neq 0$  by induction hypothesis and thus,  $a_{k+1}^2-2>0$ . Therefore, by principle of mathematical induction, we have proven our claim.

Claim 3.  $a_{n+1} < a_n \quad \forall n \geq 2$ .

Proof. Observe that  $a_{n+1} - a_n = \frac{2 - a_n^2}{2a_n}$ .

The quantity on the right is negative, using Claim 1 and Claim 2.

Thus,  $a_{n+1} < a_n$ .

We have shown that the sequence is *eventually* monotonically decreasing. Also, it is bounded below, by Claim 1. Thus, the limit  $\lim_{n\to\infty} a_n$  exists. Let  $L(\in \mathbb{R})$  denote this limit.

Note: We are assuming that an *eventually* monotonic bounded sequence is also convergent.

Thus,  $\lim_{n\to\infty} a_{n+1} = \lim_{n\to\infty} \frac{1}{2} \left( a_n + \frac{2}{a_n} \right)$ .

We had shown that  $\lim_{n\to\infty} a_{n+1} = \lim_{n\to\infty} a_n$ . Using that and other limit properties, we get

that  $L^2=2$ . Thus, L must be  $\sqrt{2}>0$ . (Why not  $-\sqrt{2}$ ?)

5. (ii)  $a_1 = \sqrt{2}$ ,  $a_{n+1} = \sqrt{2 + a_n} \quad \forall n \ge 1$ .

Claim 1.  $a_n > 0 \quad \forall n \in \mathbb{N}$ .

*Proof.* This can be easily seen via induction. Details are left to the reader.

Claim 2.  $a_n < 2 \quad \forall n \in \mathbb{N}$ .

*Proof.* We shall prove this via induction. The base case n=1 is immediate as  $\sqrt{2} < 2$ . Assume that it holds for n=k.

$$a_{k+1}^2 - 4 = (\sqrt{a_k + 2})^2 - 4 = a_k + 2 - 4 = a_k - 2.$$

But  $a_k - 2 < 0$  by induction hypothesis. Thus,  $a_{k+1}^2 < 4$  or  $a_{k+1} < 2$ . By principle of mathematical induction, we have proven the claim.

Claim 3.  $a_n < a_{n+1} \quad \forall n \in \mathbb{N}$ .

Proof. 
$$a_{n+1}^2 - a_n^2 = 2 + a_n - a_n^2 = (2 - a_n)(1 + a_n) > 0$$
.

The last inequality is by the help of Claims 1 and 2.

Thus, we have  $a_{n+1}^2 > a_n^2$ . Using Claim 1, we can conclude that  $a_{n+1} > a_n$ . By Claims 1 and 2, we have it that the sequence is bounded. By Claim 3, we have it that the sequence is monotone. Therefore, the sequence must converge. Let the limit be  $L(\in \mathbb{R})$ .

Taking limit on both sides of the recursive definition gives us  $L = \sqrt{2+L}$ . Thus,  $L^2 = 2 + L$  or (L-2)(L+1) = 0.

Note that L cannot be -1. (Why?)

$$\therefore L=2.$$

5. (iii) 
$$a_1 = \sqrt{2}$$
,  $a_{n+1} = 3 + \frac{a_n}{2} \quad \forall n \ge 1$ .

Claim 1.  $a_n < 6$   $n \in \mathbb{N}$ .

*Proof.* We shall prove this via induction. The base case n=1 is immediate as 2 < 6.

Assume that it holds for 
$$n = k$$
.  
 $a_{k+1} = 3 + \frac{a_n}{2} < 3 + \frac{6}{2} = 6$ .

By principle of mathematical induction, we have proven the claim.

Claim 2. 
$$a_n < a_{n+1} \quad \forall n \in \mathbb{N}$$
.  
Proof.  $a_{n+1} - a_n = 3 - \frac{a_n}{2} = \frac{6 - a_n}{2} > 0 \implies a_{n+1} > a_n$ .

Thus,  $(a_n)$  is a monotonically increasing sequence that is bounded above. Therefore, it must converge. Using the same method as earlier gives this limit to be 6.

#### 10. To show:

 $\{a_n\}_{n\geq 1}$  is convergent  $\iff \{a_{2n}\}_{n\geq 1}$  and  $\{a_{2n+1}\}_{n\geq 1}$  converge to the same limit.

*Proof.* ( $\Longrightarrow$ ) Let  $b_n := a_{2n}$  and  $c_n := a_{2n+1}$ . We are given that  $\lim_{n \to \infty} a_n = L$ . We must show that  $\lim_{n \to \infty} b_n = \lim_{n \to \infty} c_n$ .

Let  $\epsilon > 0$  be given. By hypothesis, there exists  $n_0 \in \mathbb{N}$  such that  $|a_n - L| < \epsilon$  for  $n \geq n_0$ .

Note that 2n > n and 2n + 1 > n for all  $n \in \mathbb{N}$ . Thus, we have that  $|b_n - L| < \epsilon$  and  $|c_n - L| < \epsilon$  for all  $n > n_0$ .

Thus, 
$$\lim_{n\to\infty} b_n = \lim_{n\to\infty} c_n = L$$
.

 $(\Leftarrow)$  Let  $(b_n)$  and  $(c_n)$  be as defined before. We are given that

 $\lim_{n\to\infty} b_n = \lim_{n\to\infty} c_n = L$ . We must show that  $(a_n)$  converges.

Let  $\epsilon > 0$  be given. By hypothesis, there exists  $n_1, n_2 \in \mathbb{N}$  such that

$$|b_n - L| < \epsilon \text{ for all } n \ge n_1 \tag{1}$$

and 
$$|c_n - L| < \epsilon$$
 for all  $n \ge n_2$ . (2)

Choose  $n_0 = \max\{2n_1, 2n_2 + 1\}.$ 

Let  $n \ge n_0$  be even. Then,  $n \ge 2n_1$  or  $n/2 \ge n_1$  and  $a_n = b_{n/2}$ . By (1), we have it that  $|a_n - L| < \epsilon$ .

Similarly, let  $n \ge n_0$  be odd. Then,  $n \ge 2n_2 + 1$  or  $(n-1)/2 \ge n_2$  and  $a_n = c_{(n-1)/2}$ . By (2), we have it that  $|a_n - L| < \epsilon$ .

Thus, we have shown that  $|a_n - L| < \epsilon$  whenever  $n \ge n_0$ . This is precisely what it means for  $(a_n)$  to converge to L.

- 1. (i) We shall show that the statement is false with the help of a counterexample.
- Let a = -1, b = 1, c = 0. Define f and g as follows:

$$f(x) = x$$
 and  $g(x) =$ 
$$\begin{cases} 1/x & ; x \neq 0 \\ 0 & ; x = 0 \end{cases}.$$

It can be seen that  $\lim_{x\to 0} f(x) = 0$  but  $\lim_{x\to c} [f(x)g(x)] = \lim_{x\to 0} 1 = 1$ .

(ii) We shall prove that the given statement is true.

We are given that g is bounded. Thus,  $\exists M \in \mathbb{R}^+$  such that  $|g(x)| \leq M \quad \forall x \in (a, b)$ .

Let  $\epsilon > 0$  be given. We want to show that there exists  $\delta > 0$  such that

$$|f(x)g(x) - 0| < \epsilon$$
 whenever  $0 < |x - c| < \delta$ .

Let  $\epsilon_1 = \epsilon/M$ . As  $\lim_{x \to \infty} f(x) = 0$ , there exists  $\delta > 0$  such that

$$0<|x-c|<\delta \implies |f(x)|<\epsilon_1.$$

Thus, whenever  $0 < |x - c| < \delta$ , we have it that

$$|f(x)g(x)-0|=|f(x)||g(x)|\leq |f(x)|\cdot M<\epsilon_1\cdot M=\epsilon.$$



(iii) We shall prove that the given statement is true.

Let  $\epsilon > 0$  be given.

Let 
$$l := \lim_{x \to c} g(x)$$
.  
Let  $\epsilon_1 = \epsilon/(|l| + \epsilon)$ .

By hypothesis, there exists  $\delta_1 > 0$  such that  $0 < |x - c| < \delta_1 \implies |g(x) - I| < \epsilon$ . Also, there exists  $\delta_2 > 0$  such that  $0 < |x - c| < \delta_2 \implies |f(x)| < \epsilon_1$ .

Let 
$$\delta = \min\{\delta_1, \ \delta_2\}$$
. Then, whenever  $0 < |x - c| < \delta$ , we have that:  $|f(x)g(x)| = |f(x)g(x) - lf(x) + lf(x)| \le |f(x)||(g(x) - l)| + |l||f(x)| < |f(x)|\epsilon + |l||f(x)| = |f(x)|(\epsilon + |l|) < \epsilon_1(\epsilon + |l|) = \epsilon$ . Thus, we have it that  $0 < |x - c| < \delta \implies |f(x)g(x) - 0| < \epsilon$ .

2. We are given that  $\lim_{x\to\alpha}f(x)$  exists. Let it be  $c(\in\mathbb{R})$ . Note that it's **not** necessary that  $c = f(\alpha)$ .

Let us evaluate  $\lim_{h\to 0} f(\alpha+h)$ . Let  $(h_n)$  be an arbitrary sequence of real numbers such that  $h_n \neq 0$  and  $h_n \to 0$ . We need to find  $\lim_{n \to \infty} f(\alpha + h_n)$ .

Consider the sequence  $(x_n)$  of real numbers defined as  $x_n := \alpha + h_n$ . Thus,  $x_n \neq \alpha$  and  $x_n \to \alpha$ . By hypothesis, we must have that  $\lim_{n \to \infty} f(x_n) = c$ .

Thus, by definition of  $x_n$ , we must have that  $\lim_{n\to\infty} f(\alpha+h_n)=c$ . This gives us that  $\lim_{h\to 0} f(\alpha+h_n)=c.$ 

Similar consideration will give  $\lim_{n\to\infty} f(\alpha - h_n) = c$  as well.

Using the limit theorems for functions, we have that:

$$\lim_{h\to 0} [f(\alpha+h)-f(\alpha-h)] = \lim_{h\to 0} f(\alpha+h) - \lim_{h\to 0} f(\alpha-h) = c-c = 0.$$

Converse of 2.

The converse of 2 does **not** hold. That is, given  $f: \mathbb{R} \to \mathbb{R}$  and  $\alpha \in \mathbb{R}$  such that  $\lim_{h \to 0} [f(a+h) - f(a-h)] = 0$ , it is not necessary that  $\lim_{x \to \alpha} f(x)$  exists. We shall demonstrate this will the help of a counterexample.

Let  $f: \mathbb{R} \to \mathbb{R}$  be defined as follows:

$$f(x) = \begin{cases} 0 & \text{if } x \in \mathbb{Q} \\ 1 & \text{if } x \notin \mathbb{Q} \end{cases}$$

It can be easily observed that  $f(x) = f(-x) \quad \forall x \in \mathbb{R}$ .

Let 
$$\alpha = 0$$
.

Then, 
$$\lim_{h\to 0} [f(\alpha+h) - f(\alpha-h)] = \lim_{h\to 0} [f(h) - f(-h)] = \lim_{h\to 0} [0] = 0.$$

However,  $\lim_{x\to 0} f(x)$  does **not** exist.

3. (i) The function is continuous everywhere except at x = 0. *Proof.* For  $x \neq 0$ , we have it that f is a composition of continuous functions. Thus, it is continuous.

To show that f is discontinuous at x=0: Consider the sequence  $(x_n)$  where  $x_n=\frac{2}{(4n+1)\pi}$ . Then,  $x_n\to 0$  but  $f(x_n)=1 \quad \forall n\in \mathbb{N}$  and thus,  $f(x_n)\to 1\neq f(0)$ . Thus, f is discontinuous at x=0, by definition.

3. (ii) The function is continuous everywhere.

*Proof.* For  $x \neq 0$ , it simply follows from the fact that product and composition of continuous functions is continuous.

To show continuity at x = 0:

Let  $(x_n)$  be any sequence of real numbers such that  $x_n \to 0$ . We must show that  $f(x_n) \to 0$ .

Let  $\epsilon > 0$  be given.

Observe that 
$$|f(x_n) - 0| = \left| x_n \sin \left( \frac{1}{x_n} \right) \right| \le |x_n|$$
.

Now, we shall use the fact  $x_n \to 0$ . By this hypothesis, there must exist  $n_1 \in \mathbb{N}$  such that  $|x_n| = |x_n - 0| < \epsilon \quad \forall n > n_1$ .

Choosing  $n_0 = n_1$ , we have it that  $|f(x_n) - 0| \le |x_n| < \epsilon \quad \forall n > n_0$ .



3. (iii) The function can be rewritten as: 
$$f(x) = \begin{cases} x & \text{if } 1 \le x < 2 \\ 1 & \text{if } x = 2 \\ \sqrt{6-x} & \text{if } 2 < x \le 3 \end{cases}$$

We claim that the function is continuous on  $[1,2) \cup (2,3]$  and discontinuous at 2. Given  $x \in [1,2)$  and any sequence  $(x_n)$  in the domain such that  $x_n \to x$ , there must exist  $n \in n_0$  such that  $x_n \in [1,2) \quad \forall n \geq n_0$ . Thus,  $f(x_n) = x_n \quad \forall n \geq n_0$ . It can now be easily shown that  $f(x_n) \to x = f(x)$ . (We have essentially used the continuity of the function  $x \mapsto x$ .) Thus, f is continuous on [1,2).

Similarly, we can argue that f is continuous on (2,3]. Again, this will follow from the fact that the function  $x \mapsto \sqrt{6-x}$  is continuous on its domain.

Now, we show that f is discontinuous at 2. Consider the sequence  $x_n := 2 - 1/n$ . It is clear that  $x_n \to 2$ .

Observe that  $1 \le x_n < 2$ . Thus,  $f(x_n) = 2 - 1/n$ . This gives us that  $f(x_n) \to 2 \ne f(2)$ .

4. We are given that f(x+y)=f(x)+f(y) for all  $x,y\in\mathbb{R}$ . Thus, we can let x=y=0. This gives us that:

$$f(0+0) = f(0) + f(0) \implies f(0) = 2f(0) \implies f(0) = 0.$$

As f is continuous at 0, we have it that  $\lim_{h\to 0} f(h) = f(0) = 0$ .

Now, we will show that f is continuous at every  $c \in \mathbb{R}$ .

Substituting x = c in the original equation gives us: f(c + y) = f(c) + f(y). As this is true for every  $y \in \mathbb{R}$ , we have that:  $\lim_{y \to 0} f(c + y) = \lim_{y \to 0} [f(c) + f(y)]$ .

We know that  $\lim_{y\to 0} f(c) = f(c)$  (constant sequence) and  $\lim_{y\to 0} f(y) = 0$  (shown above).

Thus, we can write:

 $\lim_{y\to 0} f(c+y) = \lim_{y\to 0} f(c) + \lim_{y\to 0} f(y) = f(c)$ . This is precisely what it means for f to be continuous at c.

- 4. **(Optional)** Here's a sketch of how one can show that f satisfies f(kx) = kf(x), for all  $k \in \mathbb{R}$ .
- Step 1. Use induction and show that  $f(nx) = nf(x) \quad \forall n \in \mathbb{N}$ .
- Step 2. Show that  $f(nx) = nf(x) \quad \forall n \in \mathbb{N}$ .
- Step 3. Show that  $f(qx) = qf(x) \quad \forall q \in \mathbb{Q}$ .
- Step 4. Use density of rationals and continuity of f to argue that

$$f(kx) = kf(x) \quad \forall k \in \mathbb{R}.$$

Note that we didn't require continuity of f in the first 3 steps.

12. Let  $c \in \mathbb{R}$ .

Recall that given any  $a, b \in \mathbb{R}$  such that a < b, we can construct a rational number r(a,b) such that a < r(a,b) < b. Similarly, we can construct  $i(a,b) \in \mathbb{R} \setminus \mathbb{Q}$  such that a < i(a,b) < b. (Note that we have explicit constructions of these.)

Define the two sequences  $(r_n)$  and  $(i_n)$  as follows:

$$r_n := r(c, c + 1/n)$$
 and  $i_n := i(c, c + 1/n)$ .

Thus, we have it that  $r_n \to c$  and  $i_n \to c$  and also that  $r_n \neq c \neq i_n$ .

However, observe that  $f(r_n) = 1 \quad \forall n \in \mathbb{N}$  and  $f(i_n) = 0 \quad \forall n \in \mathbb{N}$ . This gives us that  $f(r_n) \to 1$  and  $f(i_n) \to 0$ .

As  $\lim_{n\to\infty} f(r_n) = 1 \neq 0 = \lim_{n\to\infty} f(i_n)$ , f cannot be continuous at c.

13. Let  $c \in \mathbb{R}$  be such that f is continuous at c. Define sequences  $(r_n)$  and  $(i_n)$  as before.

Thus,  $f(r_n) = r_n$ . As  $r_n \to c$ , we have it that  $f(r_n) \to c$ .

Similarly,  $f(i_n) \rightarrow 1 - c$ .

For f to be continuous at c, we must have it that c = 1 - c = f(c). Solving this gives us that c = 1/2.

Thus, what we have shown so far is that: f continuous at  $c \implies c = 1/2$ .

However, we must now show that f actually is continuous at 1/2.

This is done as follows: Let  $(x_n)$  be any sequence of real numbers such that  $x_n \to 1/2$ .

We claim that  $f(x_n) \to 1/2$ . If we can prove this claim, then we are done as f(1/2) = 1/2.

Note that if 
$$x_n \in \mathbb{Q}$$
, then  $|f(x_n) - 1/2| = |x_n - 1/2|$  and if  $x \notin \mathbb{Q}$ , then  $|f(x_n) - 1/2| = |1/2 - x_n| = |x_n - 1/2|$ .

Thus, we have it that  $|f(x_n)-1/2|=|x_n-1/2| \quad \forall n\in\mathbb{N}$ . Let  $\epsilon>0$  be given. As  $x_n\to 1/2$ , there exists  $n_1\in\mathbb{N}$  such that  $|x_n-1/2|<\epsilon$  for all  $n\geq n_1$ . Choose  $n_0=n_1$ . Thus,  $|f(x_n)-1/2|<\epsilon$  for all  $n\geq n_0$ .

 $\therefore \lim_{n\to\infty} f(x_n) = 1/2.$ 

