

Smart Phone

Category: Sorting

Problem link: <https://www.codechef.com/ZCOPRAC/problems/ZCO14003>

Solution

Sort the array. If you do that, the number of people who will buy the app at the price equal to the i th person's budget will be $N - i$. After sorting, we can brute-force to maximize the value $\text{budget}[i] * (N - i)$ over all valid i using a for loop.

SUPW

Category: Dynamic Programming

Code: [submission.cpp](#)

Problem link: <https://www.codechef.com/ZCOPRAC/problems/ZCO14002>

Solution

For each day, we can recursively decide whether to work on that day or not. However, if we haven't worked on a given day.

Step 1: Define a subproblem

Let `solve(idx, days)` be the minimum number of total minutes you need to work from day `idx` till day `N - 1`, such that at no point a situation arises where you go three days without doing the SUPW duty, given that you haven't worked for the past `days` days.

Step 2: Base cases

Consider the simplest cases where you know the solution to the problem.

Assume the number of minutes given for day `i` is `minutes[i]`, where $0 \leq i < N$.

Consider the case of the $(N - 1)$ th day which is the last day.

Then `solve(N - 1, 0) = 0` because you already worked yesterday and working again on the last day is not necessary because it will only add to the number of minutes.

Similarly `solve(N - 1, 1) = 0`

However, `solve(N - 1, 2) = minutes[N - 1]` because you didn't work for the past two days, and thus you will have to work today.

Step 3: Recurrence relation

Dynamic programming is careful brute-force. We'll try out all combinations, but ensure that we do not compute the same result again.

If `days < 2`, we have a choice regarding whether to work on the given day `idx` or not. So we compute the answer for both the cases and return their minimum

```
if days < 2:
    return min( solve(idx + 1, days + 1), solve(idx + 1, 0) + minutes[i] )
```

Otherwise (when `days` is 2), we don't have a choice, so we'll have to work on day `idx`

```
else:  
    return solve(idx + 1, 0) + minutes[i]
```

Memoize

Store the return value of the `solve()` function in a memo table (perhaps an array or a vector in this case) and use it directly if already computed.

Endnote

Exercise: Argue why the base cases described above can be generalized to the following base case without affecting the recurrence:

$$\text{solve}(N, 0) = \text{solve}(N, 1) = \text{solve}(N, 2) = 0$$

Video Game

Category: Implementation Problem link: <https://www.codechef.com/ZCOPRAC/problems/ZC014001>

Solution

Store the number of boxes in each position in an array/vector. Then just follow the instructions to update the number of boxes in each position. And then just output the elements of the array/vector.