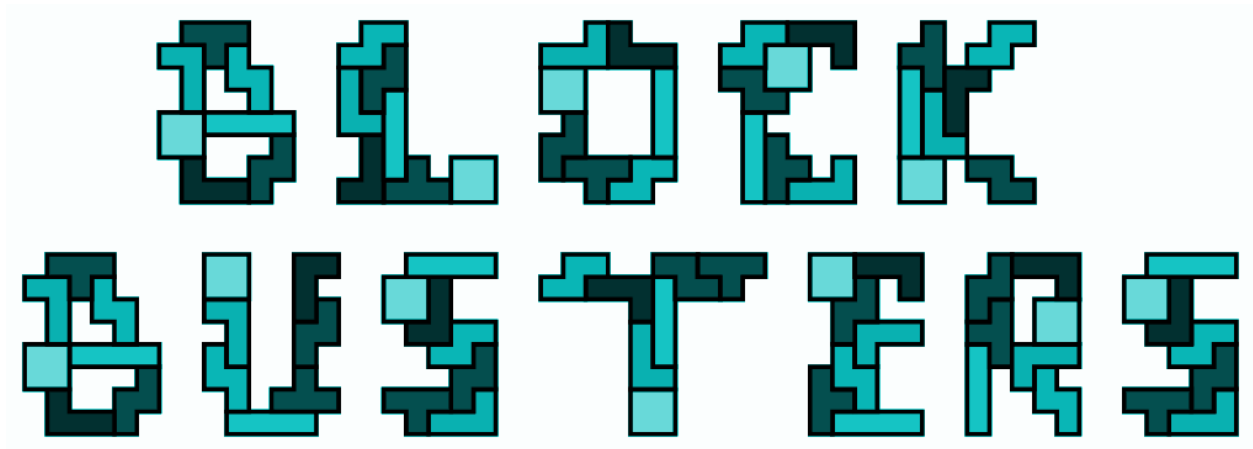


CS F363 Compiler Construction

Group-17



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1. Top - Level Design Schema:

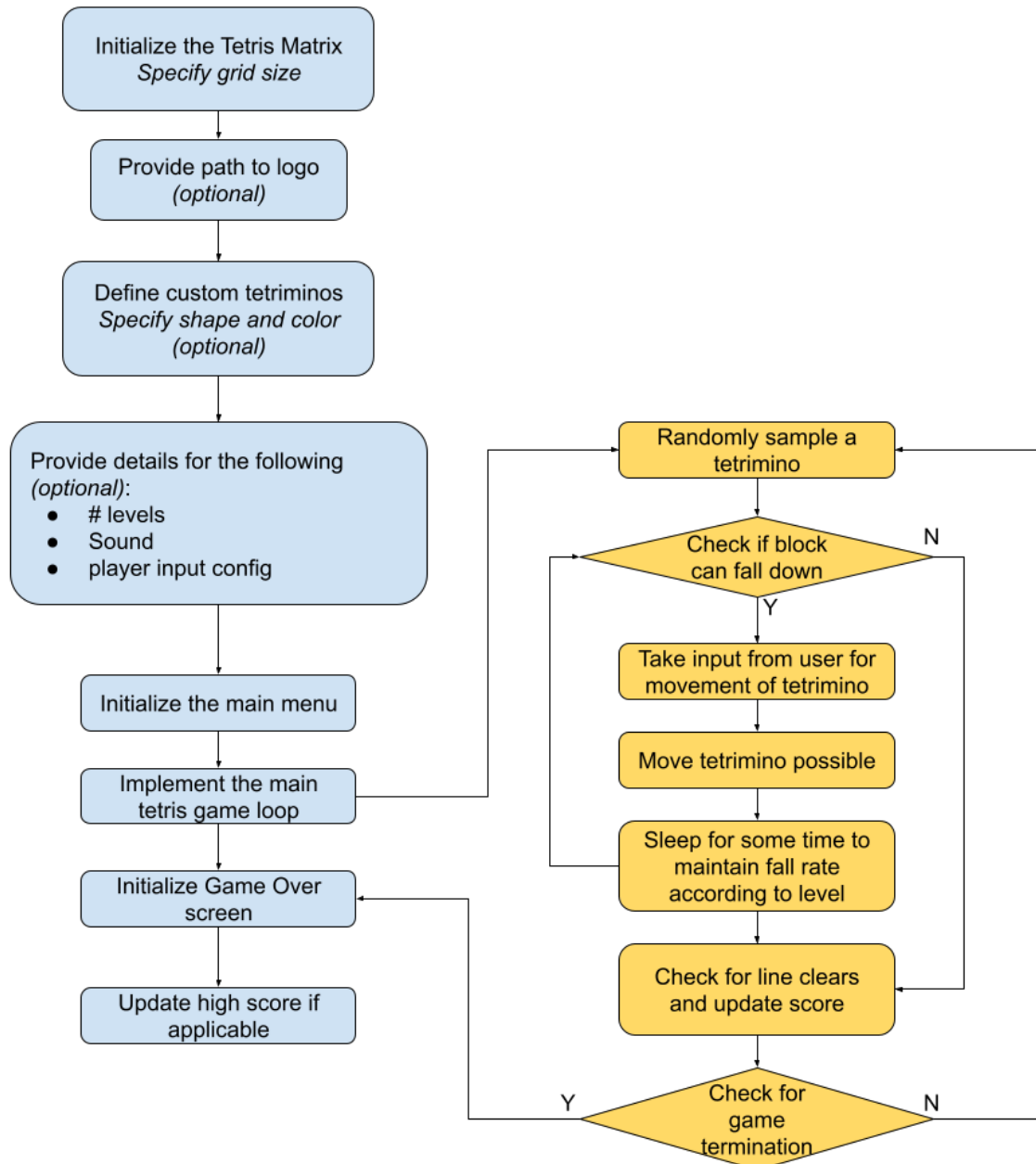
The overall program structure of the Tetris game language will be based on python. The designed language will be called Tetris Language (Tlang) and a file written by a tetris game programmer in Tlang will be saved with a `.tl` extension.

The `.tl` file which is the source program will be converted into a high level language code i.e. python, the target program, by a compiler. Finally, the generated target code must be executed by the programmer to run the tetris game. The target code will make use of an existing engine code which has certain parts of the tetris code implemented in python.

1.1 User primitives:

- Game Logo:
 - The programmer has an option to modify the logo of the game he wishes to build. The programmer must give the path of the logo(as `.png/.jpeg` file)
- Grid Size of Tetris Matrix:
 - The programmer can control the height and width of the tetris game window. A check is done during interpretation so that the height and width entered by the user satisfies the requirement $8 < \text{gridX} < \text{gridY}$.
- Buffer Size of Tetris Matrix:
 - The programmer can set the size of the buffer zone.
- Custom Tetriminos:
 - The programmer can introduce a custom tetrimino in addition to the already available default 7 tetriminos by the Tlang. The custom tetrimino must be limited to a size of 4x4 matrix cells. The following must be mentioned by the programmer when introducing a new tetrimino:
 - Shape: a 4x4 matrix with 1s denoting occupied and 0s denoting unoccupied blocks by the tetrimino
 - Color: A list of three values in RGB
 - Rotation: The possible configurations of the tetrimino when rotated
- Number of Levels:
 - The number of levels that the game will offer to the player can be defined by the programmer, subject to the constraint $0 < \# \text{ levels} < 5$. A greater level indicates greater falling speed of the tetrimino.
- Sound Effects:
 - The programmer can control whether the game will have the background music on or not during the game. The music itself can be chosen from a list of options given to the programmer.
- Player Input:
 - The programmer can configure which keyboard key, from a fixed set of keys, corresponds to the rotate and lateral movement of the tetriminos. The user can also attach additional functionality such as updating scores based on keypresses.
- Print Functionality:
 - The programmer can use a print statement for debugging.

Apart from the above mentioned features which can be used, a programmer attempting to create a Tetris game using TLang is expected to implement the main game loop using the functionalities provided by the language design. Implementations of some requisites of the Tetris Game which cannot be modified, like the game termination condition, are given to the programmer predefined in the engine file written in python. Below is a flowchart which depicts the ideal flow of code for a programmer using TLang:



The following are some of the functions defined in the language and implemented in the engine file are readily available to the programmer:

- `initialize_window(height, width, buffer_height):`
 - The tetris window will be initialized corresponding to the arguments in the function call.
 - Default (if the function is not called by the programmer): default size of tetris matrix
- `set_logo(logo_path):`
 - The logo will be set to the image whose path is passed as a string.
 - Default: original tetris logo.
- `add_block(block):`
 - The custom tetrimino defined by the user will be added to the set of tetriminos that will be considered while sampling a tetrimino. The argument should be of the *block* datatype which is given in the language definition.
 - Default: O, I, T, L, J, S and Z tetriminos
- `set_levels(levels):`
 - Set the no. of levels
 - Default: 1 level (fall speed = 1.0 sec/line)
- `set_sound(sound_path):`
 - The background will be set to a sound file whose path is passed as a string.
 - Default: no background music.
- Input keyboard config:
 - Built-in functions like `on_left_button()`, `on_W_button()` allow the programmer to add custom functionality when some input is received.
- `main_menu():`
 - The main menu window of the tetris game will be displayed. The user will select the level here and the number of levels offered will be displayed according to the number of levels set by the programmer using `set_levels`.
- `sample_tetrimino():`
 - A tetrimino from the set(default + programmer defined) will be sampled randomly.
- `block_fall():`
 - Returns true if the current tetrimino can fall further down, false if it has dropped
- `user_input():`
 - Checks if user has given any input on the keyboard

- `move_tetrimino()`:
 - Moves the tetrimino according to the user input(if given) only if possible. If the movement desired by the player is not possible, it is just moved down.
- `sleep()`:
 - A delay is introduced in order to maintain the fall rate according to the level chosen by the player.
- `line_clear()`:
 - Clear all rows that are full and update the score.
- `is_gameover()`:
 - Returns true if the game termination condition is satisfied. If not, returns false
- `gameover_window()`:
 - The game over screen is displayed which displays the score and updates the high score if applicable.

1.2 Programmable features:

- Variable declaration and assignment:

The programmer can declare different types of variables using the “**var**” keyword without having to explicitly specify the type. Type of the variable is handled by the python backend. Examples:

Declaration

```
var name1 := True;
var name2 := 42;
var name3 := "string";
var name4 := [ ];
```

Assignment

```
name := value
```

- Operations:

The programmer can perform various arithmetic, boolean, relational and bitwise operations as mentioned in the table below.

Type of Operation	Operators Provided
Arithmetic	<code>*</code> , <code>/</code> , <code>%</code> , <code>+</code> , <code>-</code>
Boolean	<code>not</code> , <code>and</code> , <code>or</code>
Relational	<code><</code> , <code><=</code> , <code>></code> , <code>>=</code> , <code>==</code> , <code>!=</code>
Bitwise	<code>!</code> , <code>&</code> , <code><<</code> , <code>>></code> , <code>^</code> , <code> </code>

- Function declaration and call:

The programmer can declare functions using the “**function**” keyword and specify the arguments enclosed by parentheses and separated by comma. Compound statements are enclosed within the “**begin**” and “**end**” keywords. Sample function declaration and call code is shown below.

Declaration

```
function fname(arg1, arg2) begin
    <stmt 1>
    <stmt 2>
    .
    .
    .
end
```

Call

```
fname(val1, val2);
```

- Conditional Statements:

Our language provides the flexibility of using conditional statements to the programmers in the following manner using the “**if**”, “**then**”, “**else**” keywords.

Usage

```
if <expr> then <stmt>;
if <expr> then <stmt> else <stmt>;
```

- Loop Statements:

The programmer can make use of loops for controlling the flow of the code. Our language provides “**while**” and “**for**” loops that can be used in the following manner.

Usage

```
while <expr> do <stmt>;
do <stmt> while <expr>;
for i := 0 to 10 do <stmt>;
```

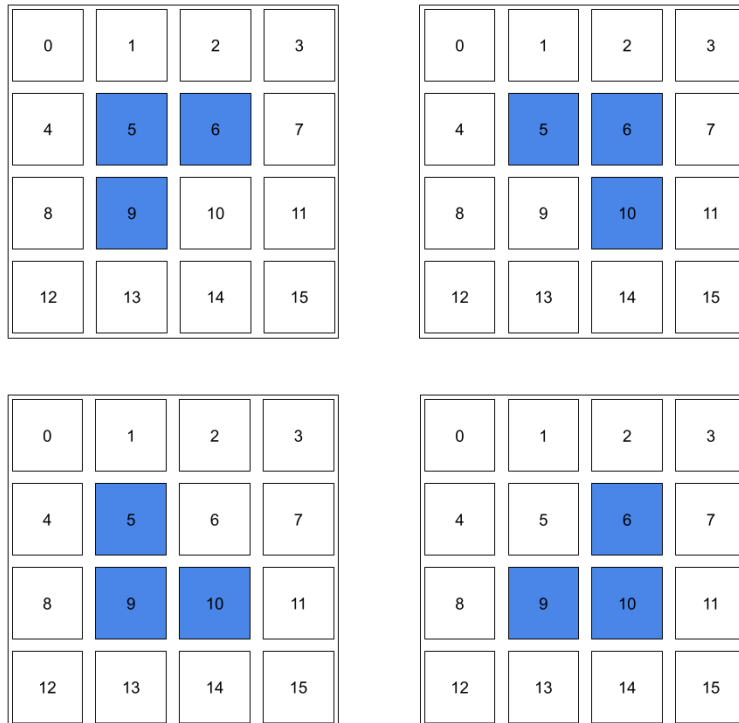
- Custom data types:

The language provides functions for creating common game objects such as new blocks and sound. Here's a code snippet demonstrating the creation of a new red colored block.

Usage

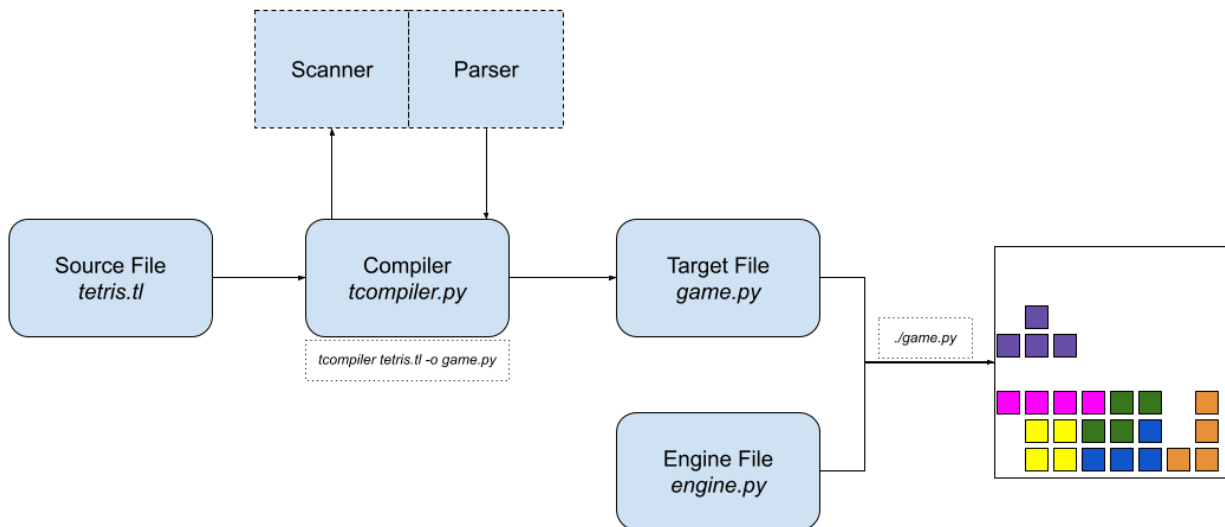
```
var newBlock := createBlock([[5,6,9],
                           [5,6,10],
                           [6,9,10],
                           [5,9,10]],
                           [255,0,0]);
```

`createBlock` is an inbuilt function in our language. It takes the configuration and color as inputs and returns a new block. The configuration parameter is a list of different orientations the block can be in. The orientation is defined as a list of indices of a 4x4 matrix (customizable!) which the block occupies.



1.3 Pipeline Schema:

The code will be written by the programmer in the TLang(.tl) and will be converted to the target code in python via a compiler(*tcompiler.py*) written in python. The compiler will use the SLY python library for scanning and parsing the source file and will eventually output an executable python file *game.py*. Finally, the game can be executed by running the *game.py* which will use the *engine.py* file.



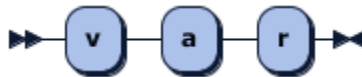
The engine file consists of an implementation of the tetris game in parts which the programmer can use. The programmer is not expected to program the basic logic of the tetris game that cannot be modified for e.g. the game termination condition.

In our project, a valid python code is generated in the parse step itself i.e. *game.py*. The action for a recognised BNF expression by the parser in the source code is to first convert the expression into valid python syntax and next to write this syntax into the executable *game.py*. Once parsing is complete, the implementations in *engine.py* are imported into the *game.py* file which can now be directly executed by a player to run the tetris game.

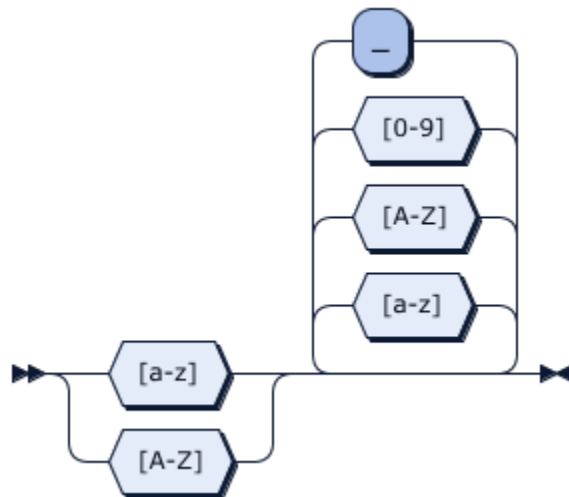
2. Scanner Design:

The overall structure of the Scanner will use the Lexer implementation from the [SLY Python library](#). We have defined the various possible tokens and their pattern, which is a description of the form that the lexemes of a token may take. The following is a subset of the pattern-action pairs.

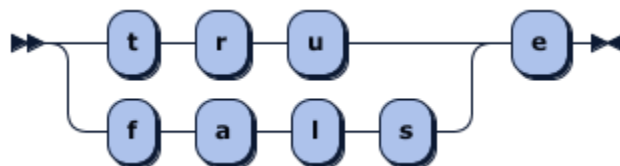
```
"var" { return VAR; }
```



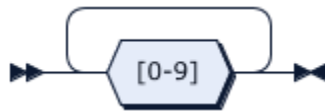
```
[a-zA-Z][a-zA-Z0-9_]* { return IDENT; }
```



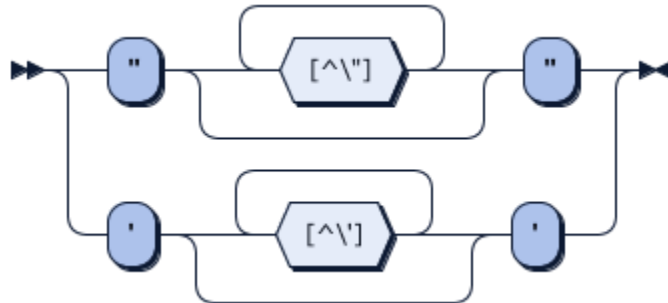
```
"true"|"false" { return BOOL; }
```



```
[0-9]+ { return NUM; }
```



```
\"[^"]*" | \'[^\']*\'          { return STRING; }
```



*more diagrams for the pattern-action pairs can be found in the [lex-diagram](#) folder.

2.1 Lexer implementation using Python library SLY (Sly Lex Yacc) :

Our scanner utilizes the predefined lexer class of the `SLY` library. Tlang scanner inherits from the `SLY` scanner. We store the names of all the tokens in the set `tokens`. Then we define the regular expression, to map suitable lexemes for each of these tokens. Tokens are specified by writing a regular expression rule compatible with the `re` module. In the next step, we invoke the `tokenize()` method which is a generator function that produces a stream of `Token` instances. `Token` contains type name and value as attributes. A special token `ignore` is used to skip over whitespaces and tabs. A literal character is a single character that is returned “as is” when encountered by the lexer. We are not using any literals in our lexer. We are also storing the line number by using the `ignore_newline()` method and the column number by using the `index` attribute for a token for the purpose of debugging. Due to the large amount of tokens and lexemes that they map on to, there exists some overlapping between the lexemes and the tokens, which can be solved using the special cases method provided by the `SLY` library.

A good way to visualize these discrepancies between the patterns of lexemes, is to look over some examples:

Consider a tokenizer, that maps the lexemes following the regex : `[a-zA-Z][a-zA-Z0-9_]*` to the token `IDENT` and the literal `true` to token `BOOL`.

<pre>var true123 := true;</pre>	<pre>type='VAR', value='var' type='IDENT', value='true123' type='ASSIGN', value=':=' type='IDENT', value='true' type='EOL', value=';'</pre>
---------------------------------	---

The above example illustrates how an incorrectly designed lexer may scan the lexemes and generate the tokens. An overlap between `IDENT` and `BOOL` tokens occurs here, where the literal `true` in the input expression does not map to `BOOL` but rather to `IDENT`.

<pre>var true123 := true;</pre>	<pre>type='VAR', value='var' type='BOOL', value='true' type='NUM', value='123' type='ASSIGN', value=':=' type='IDENT', value='true' type='EOL', value=';'</pre>
---------------------------------	---

One might think the logical solution to the above error is to rearrange the order of declaration of tokens, and define the `BOOL` token before the `IDENT` token. However, after implementing the solution, we are faced again with a rather intriguing problem. Since the literal `true` is mapped

to `BOOL` before `IDENT`, while defining the variable name as `true123` the scanner considers the `true` portion of the variable name as the `BOOL` token and the rest as `NUM`, when it should actually be the case where the entire variable name (`true123` in this case) should map to the token `IDENT`. This is where the versatile error handling capabilities of SLY are best utilized.

SLY allows us to create special cases, or rather remapping of tokens, where we can map the literals like `true` and `false` to token `BOOL`, such that while parsing, the lexemes get matched to the new token type, and not to the ones defined earlier on. An example of how one can code this, using the library in python, is shown below :

```
IDENT = r'[a-zA-Z][a-zA-Z0-9_]*'
...
IDENT['true'] = BOOL
IDENT['false'] = BOOL
```

Introducing changes to codes like this, finally gives us the required output :

<pre>var true123 := true;</pre>	<pre>type='VAR', value='var' type='IDENT', value='true123' type='ASSIGN', value=':=' type='BOOL', value='true' type='EOL', value=';'</pre>
---------------------------------	--

A demonstration of running the scanner on the terminal along with the output :

*check [scanner.py](#) for more details

The screenshot shows a code editor with a dark theme. The top part displays the scanner.py file with lines 45 to 51, where identifiers 'var', 'true', 'false', 'if', 'not', 'and', and 'or' are mapped to tokens VAR, BOOL, IF, NOT, AND, and OR respectively. Below the code editor, the 'TERMINAL' tab is active, showing the command `python3.8 scanner.py` and its output, which matches the example shown in the previous block.

```
45 IDENT['var'] = VAR
46 IDENT['true'] = BOOL
47 IDENT['false'] = BOOL
48 IDENT['if'] = IF
49 IDENT['not'] = NOT
50 IDENT['and'] = AND
51 IDENT['or'] = OR

PROBLEMS 59 OUTPUT TERMINAL GITLENS COMMENTS DEBUG CONSOLE

~/Des/S/CC/CS_F363-Compiler-Construction/A/scanner > python3.8 scanner.py
var true123 := true;
type='VAR', value='var'
type='IDENT', value='true123'
type='ASSIGN', value=':='
type='BOOL', value='true'
type='EOL', value=';'
```

For error handling, SLY ensures that if a bad character is encountered while lexing, tokenizing will stop. However, we can add an `error()` method to handle lexing errors that occur when illegal characters are detected. The error method receives a `Token` where the `value` attribute contains all remaining untokenized text.

2.2 Scanner test plan and test cases:

Since our scanner runs through the command line, we can test it using a simple bash script. We can write various instructions in the Tetris language organized into files in the input folder. We then provide the expected output in a different folder following the same naming convention. We then loop over all input files and store the compiled output in the output folder. Then we can use the diff command to check whether the generated code matches the expected output.

The command to compile all input files:

```
cd input
find . -type f | xargs -I {} -n 1 bash -c '../tcompiler {} > ../output/{}'
```

The command to check the resultant output with expected output:

```
cd ../output
find . -type f | xargs -I {} -n 1 bash -c '
  if diff {} ../expected/{} > /dev/null; then
    printf "TEST PASSED: {}\\n"
  else
    printf "TEST FAILED: {}\\n"
  fi
'
```

*see [test folder](#) for details on gitlab

Sample test cases:

Input	Expected Output
score := score + 1;	score=score+1
var name := "Test";#	name="Test"

Combining the two commands into a single bash script, gives the following result on the given test cases:

```
./runner.sh
TEST PASSED: variable_assignment
TEST PASSED: variable_declaration
```


3. Division of Labour between the Scanner and the Parser:

Suppose we are parsing the following string `x := 42`. The first step of parsing is to break the text into tokens where each token has a type and value. The Scanner is used to do this. It generates a stream of tokens containing the type and value of the token. The output of running the Scanner on the given string might be `[('IDENT', 'x'), ('ASSIGN', ':='), ('NUM', '42')]`. The Scanner can also use RegEx pattern matching to assign a type to a group of characters. This is useful when we want to parse some text whose exact value we might not know. The identifier `x` and the number `42` were matched using the patterns `[a-zA-Z][a-zA-Z0-9_]*` and `[0-9]+` respectively. The scanner is also responsible for ignoring whitespaces and comments.

Once we have a stream of tokens ready, the Parser can be used to assign structure to tokens. For example, if the Parser sees the tokens `IDENT`, `ASSIGN` and the next token is of type `NUM`, it can deduce the input to be a declaration statement. The way a parser achieves this is using a stack. It stores all the tokens it has seen so far in a stack. The token on the top determines what happens next. If the next token is a part of a rule in the grammar using the other tokens on the stack, the stack is reduced using the provided action, else the token is shifted onto the stack. If no further match can be found the parser throws an error. The stack must be empty when the parser has finished parsing all the tokens. The grammar of the language is usually defined in BNF notation. A successful parse usually results in an Abstract Syntax Tree. This AST is a machine readable form of the input tokens. What kind of AST is generated is completely controlled by the action described in the pattern-action pairs.

Appropriate tree traversal algorithms can be used to traverse the tree. Linters are an example of a tool which finds common bugs like using a variable before it is defined. Interpreters read the AST, and execute instructions based on the current node of the tree. ASTs can also be compiled to other languages. In our project, we use appropriate actions to generate a valid python code in the parse step itself.

To summarize:

- The scanner handles low-level character-by-character analysis.
- The parser embeds a scanner and handles higher level language analysis.

The scanner extracts the tokens but knows nothing about the grammar. The parser handles the grammar, but knows nothing about the original string. This keeps the code modular and simple.

4. Division and Distribution of Roles and Responsibilities Among the Team:

- Abhineet Karn , Abhinav Srivastava :
 - Working with the SLY tool to connect the engine and the TLang code written by the programmer
- Aryan Tyagi , Hrishikesh Govindrao Kusneniwar :
 - Working on the scanner and parser, and the *tcompiler.py* file and testing the scanner
- Hardik Nilesch Shah , Hitarth Daxeshbhai Kothari :
 - Working on the engine file i.e. implementations of parts of the tetris game in python