Gaurav Sengar
Section: D
Roll No: 03

Tutorial-6

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Qu-1. What do you mean by Mininum spanning tree? What is the application of MST.

Soln: A minimum spanning tree or minimum weight spanning tree is a subset of the edges of a connected, edge-weighted undirected graph that connects all the vertices together, without any cycle and with the minimum possible total edge weight.

Application:

- Designing Local Area Network
- Laying pipelines connecting offshore drilling sites, refineries and consumer markets
- Suppose you want to construct highways or railroads spanning several cities then we use the concept of MST.
- To reduce cost, you use the concept of MST to connect the houses.

Qu-2 Please analyse the time and space complexity of Prim's, Kruskal, Dijkstra and Bellman ford algorithm.

Soln:

Algorithm	Time complexity	Space complexity
Prim's	$O(V^2)$	$O(V+E)$
kruskal	$O(E \log V)$	$O(\log(E))$
Dijkstra's	$O(V+E)$	$O(V+E)$
Bellman ford	$O(VE)$	$O(V)$

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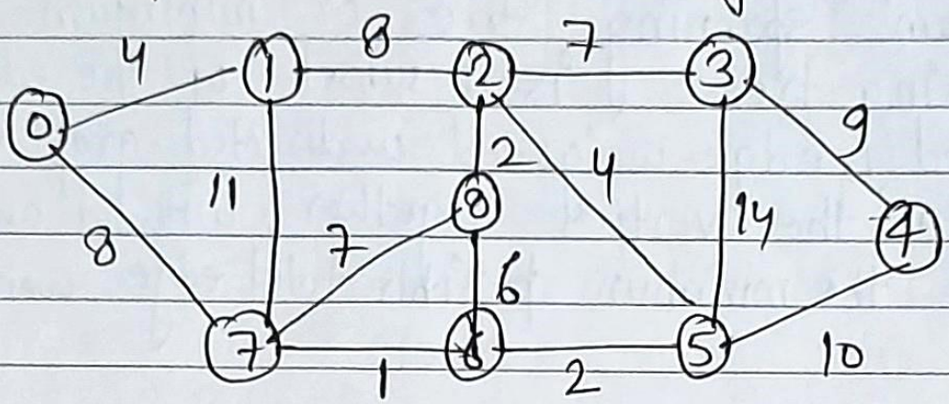
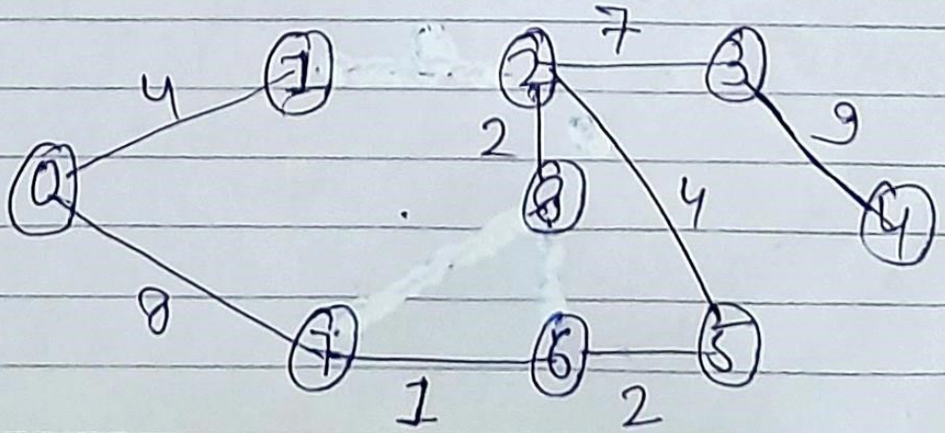
Que.3. Apply Prims and kruskal algorithm on the graph, to compute MST and its weight.



Sol^n: Kruskal.

Path.	weight
$7 \to 6$	1
$6 \to 5$	2
$2 \to 8$	2
$0 \to 1$	4
$2 \to 5$	4
$8 \to 6$	6
$2 \to 3$	7
$7 \to 8$	7
$1 \to 2$	8
$3 \to 4$	9
$5 \to 4$	10
$1 \to 7$	11
$3 \to 5$	14

$\to 0 \to 7$ 8



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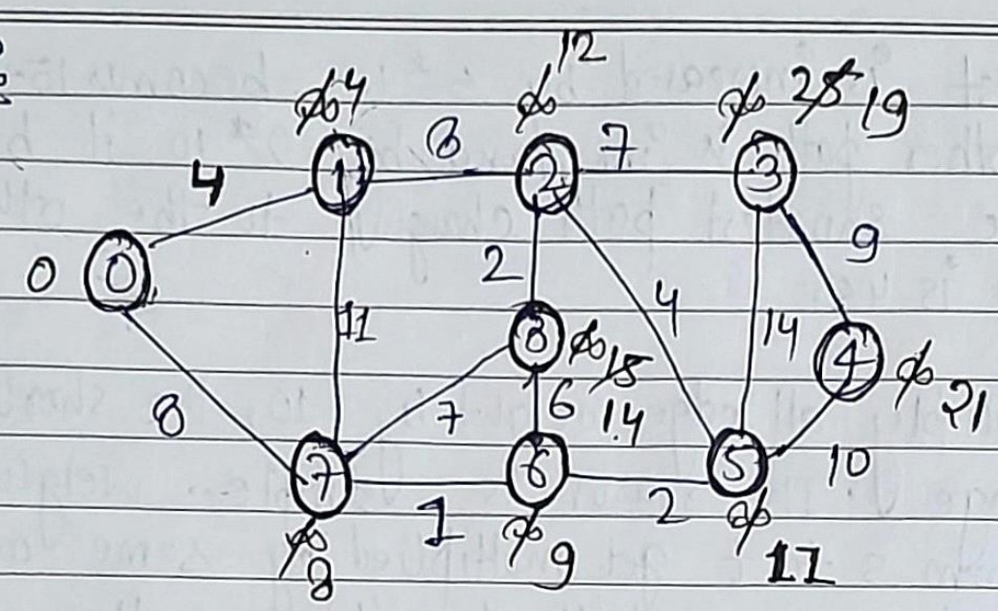
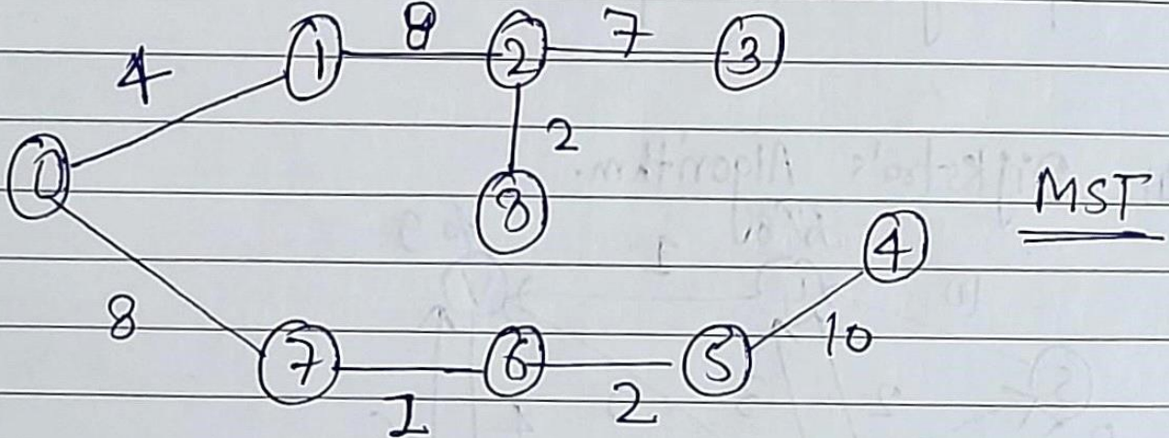
Prims:





MST

Qu-4 Given a weighted graph. You are also given the shortest path from a source vertex 's' to a given destination vertex 'f'. Does the shortest path remains same in the modified graph in the following case.

- If weight of every edge is incorpased by 10 units
- If weight of every edge is multiplied by 10 units

Sol^{n}: (i) The shortest path may change. The reason is that there may me different no. of edges in different paths from 's' to 't'. for example let shortest path be of weight 15 and has 5 edges. let there be another path with 2 edges and total weight is 25. The weight

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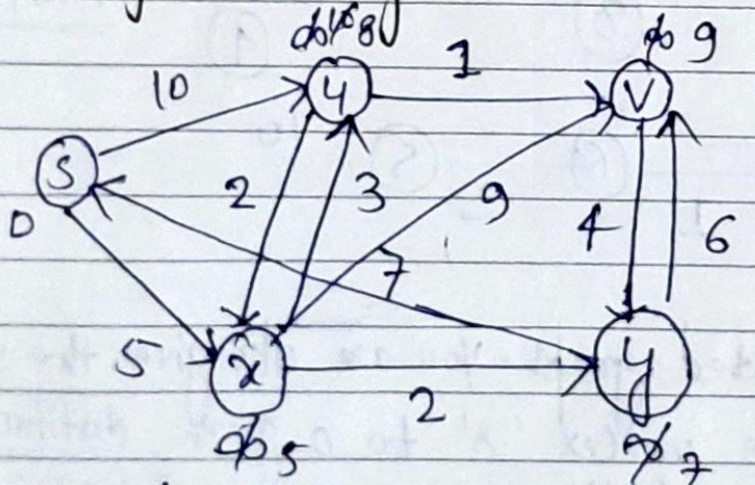
of the shortest is increased by 5*10 becomes 15+50. Weight of other path x increased by 2*10 it becomes 25+20, so the shortest path changes to the other path whose weight is 45.

(ii) If we multiply all edge weight by 10, the shortest path doesn't change. The reason is simple. weight of all paths from s to t get multiplied by some amount. The no. of edges on a path doesn't matter.

Qu.5

Sol^n: (a) Dijkstra's Algorithm.



node	shortest distance from source node
u	8
x	5
v	7
y	9

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Bellman ford Algorithm

1st → (s)0 (t)10 (x)5 (y)∞ (z)∞

2nd → (s)0 (t)10 (x)5 (y)∞ (z)11

3rd → (s)0 (t)8 (x)5 (y)7 (z)9

4th → (s)0 (t)8 (x)5 (y)7 (z)9

final Graph