

Definition of subproblem used:

$DP[i][j]$ as the max success probability if reaction i has j units of catalyst.

$comp[i][j]$ stores the units of catalyst assigned to the i th reaction such that we can get the max success probability.

base cases:

$DP[k][0]=0$ for $1 \leq k \leq N$ because its given $e[k][0]=0$

$DP[1][j]=e[1][j]$ for $1 \leq j \leq C$ as if there is only one reaction, all the j units of catalyst must be assigned to it.

$comp[1][j]=j$ for $1 \leq j \leq C$ as for only 1 reaction all the units of catalyst be set to that reaction.

Recursive formulation:

for $2 \leq x \leq N$, $x \leq y \leq C$:

if ($x==y$) : $DP[x][y]=D[x-1][y-1]*e[x][1]$ (as all essential reactions must have one catalyst each)

$comp[x][y]=1$ as x th reaction will be assigned 1 catalyst.

else : $DP[x][z]=\max(DP[x-1][y-z]*e[x][z])$ where $1 \leq z \leq (y-1)$

$comp[x][y]=tempz$ where $tempz$ =value of z for which maximum is found above
(as this means x th reaction must have $tempz$ catalyst)