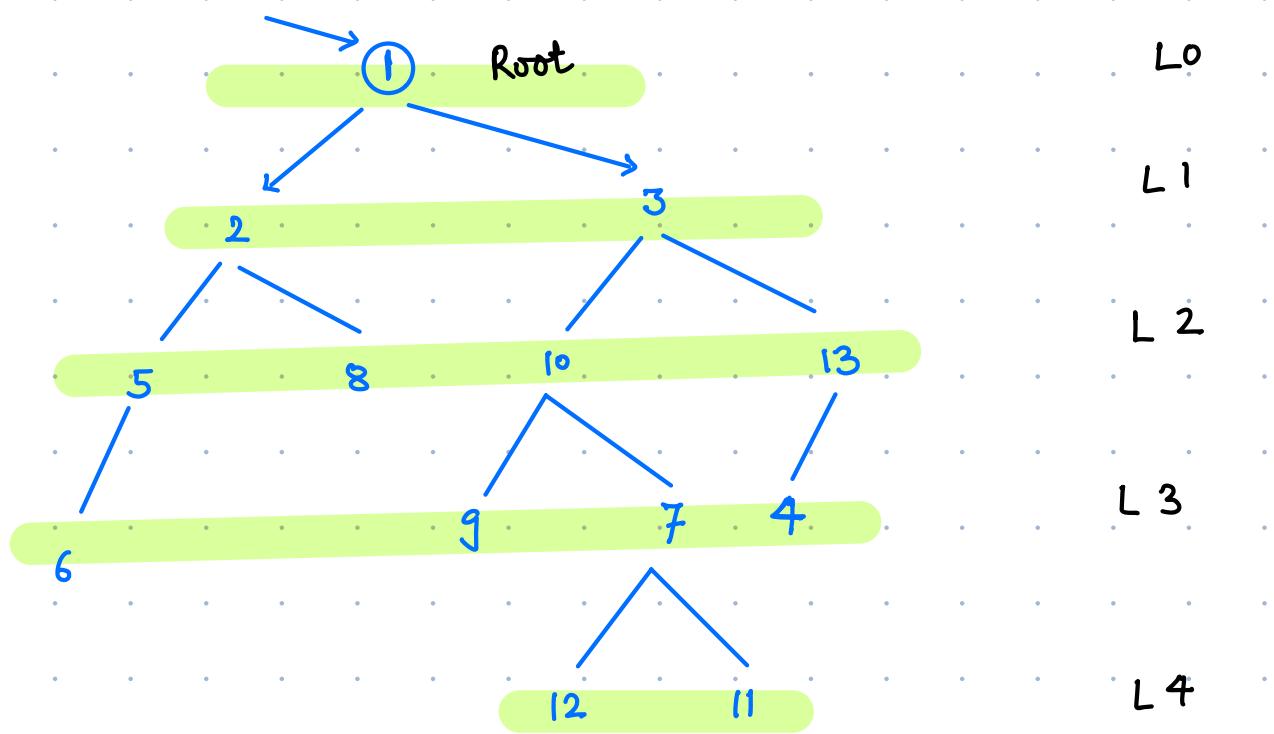
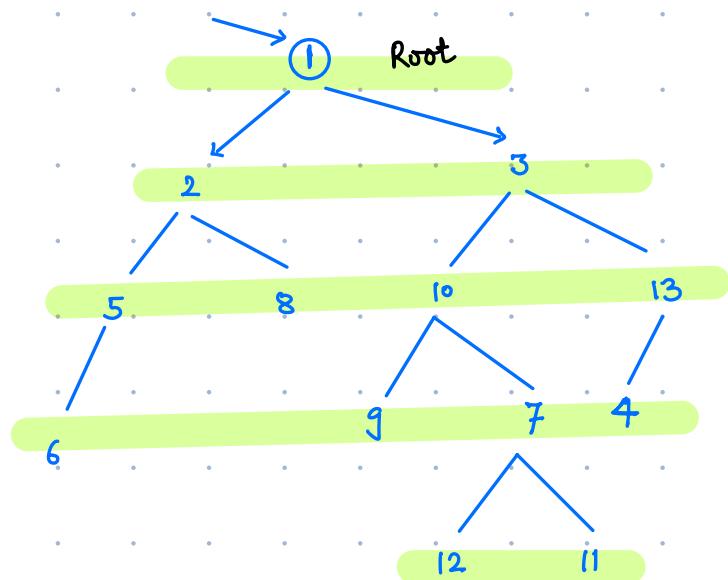


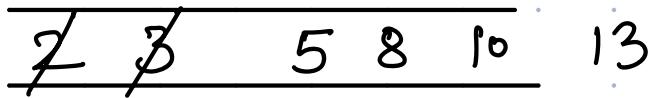
Q.



1
2 3
5 8 10 13
6 9 7 4
12 11

1. Process all nodes of current level before moving to next level
2. DS Used = Queue





Pseudocode :

$q.$ enqueue (root)

while ($q.$ isEmpty() == false) {

$x = q.$ dequeue();

print ($x.$ data);

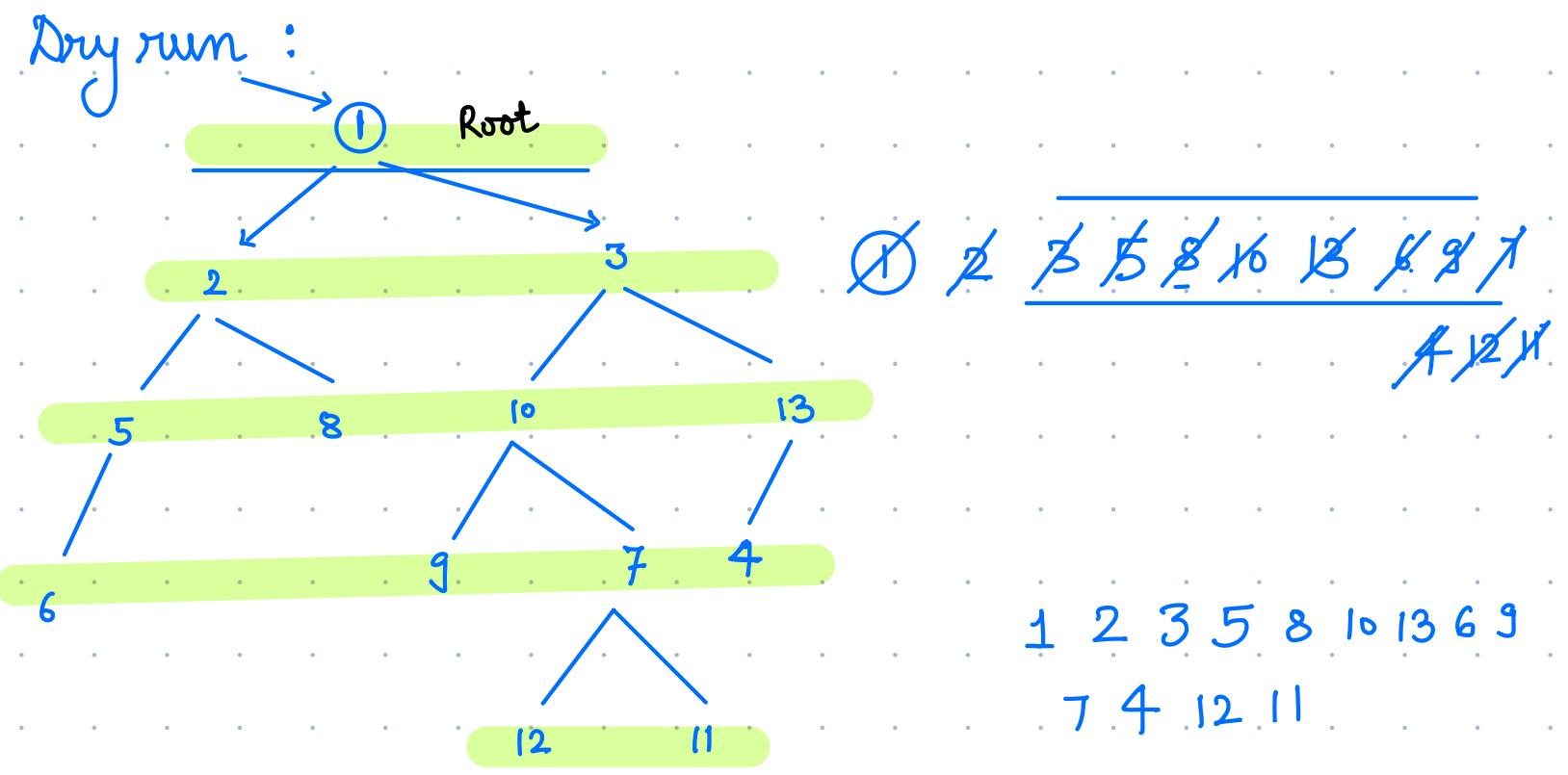
if ($x.$ left != NULL) { $q.$ enqueue ($x.$ left) }

if ($x.$ right != NULL) { $q.$ enqueue ($x.$ right) }

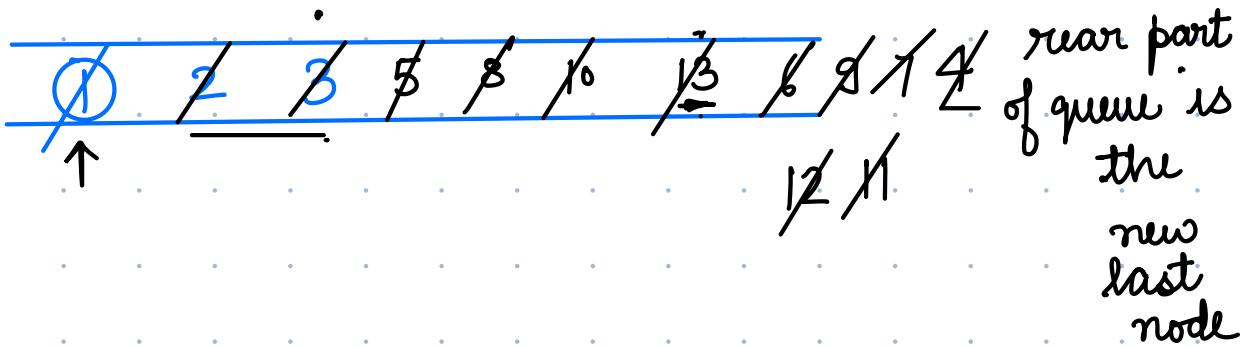
}

TC: $2N = O(N)$

SC: $O(N)$



11 ~~4~~ ~~13~~ ~~3~~ ~~1~~ = last



- ✓ When you remove last node, next level is completely ready
- ✓ Whenever you put/update last node, add a new line.

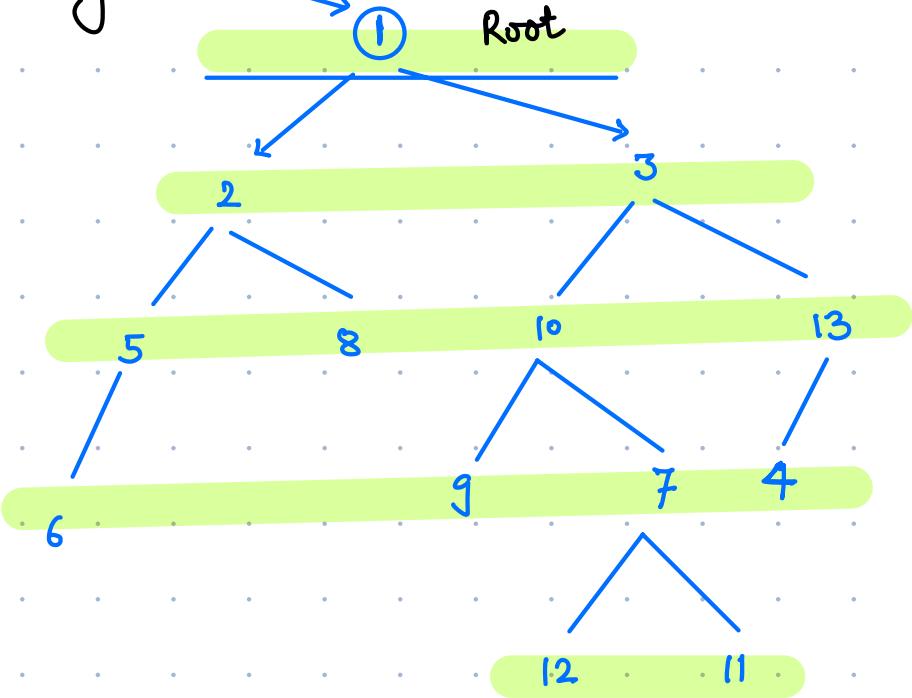
Pseudocode :

```
last = root
q. enqueue (root)
while ( q.isEmpty() == false) {
    x = q.dequeue();
    print (x.data);
    if (x.left != NULL) { q.enqueue(x.left) }
    if (x.right != NULL) { q.enqueue(x.right) }
    if ( x == last && q.isEmpty() == false) {
        print ("\n")
        last = q.rear();
    }
}
```

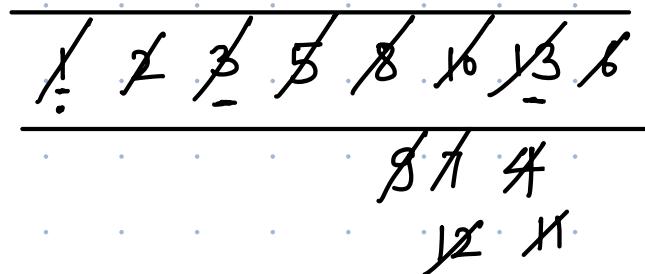
TC: $2N = O(N)$

SC: $O(N)$

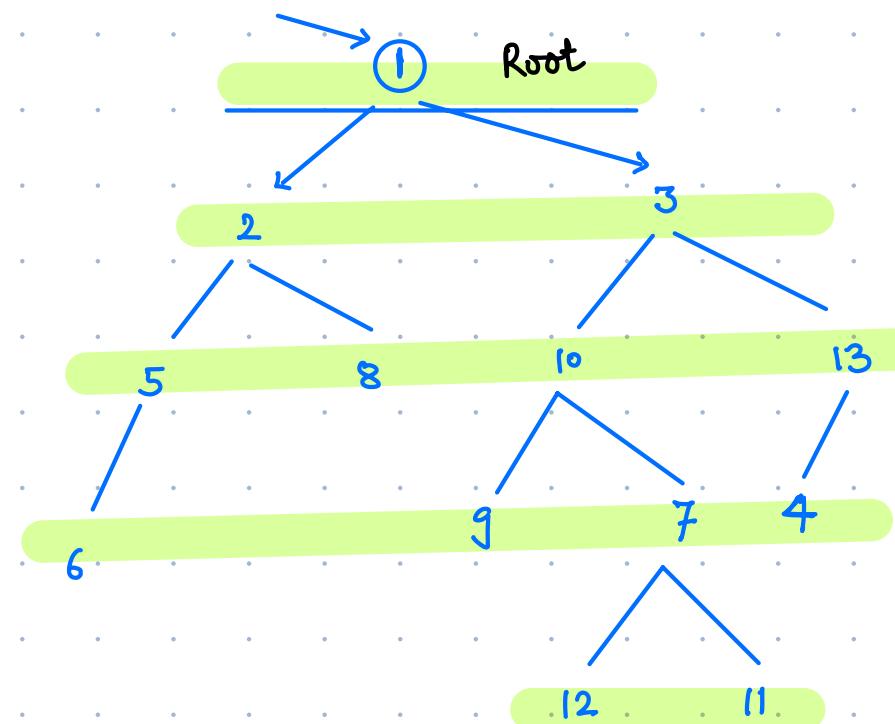
Dry Run



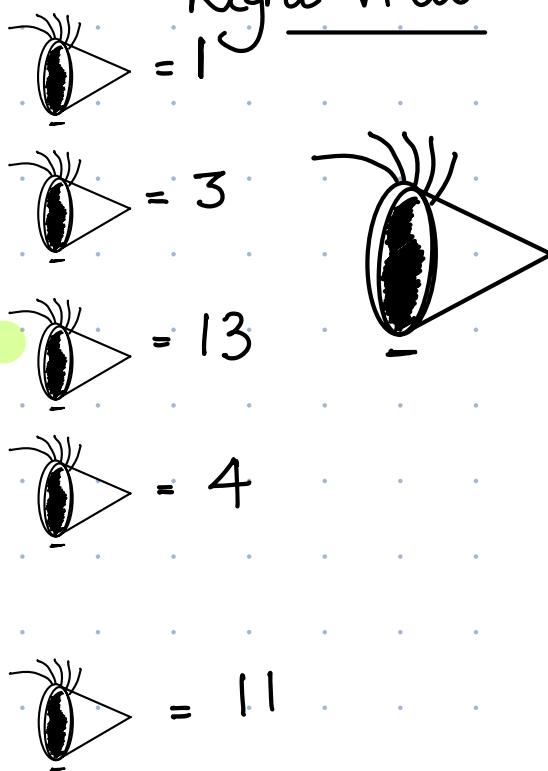
last = 1 2 3 13 4 !!



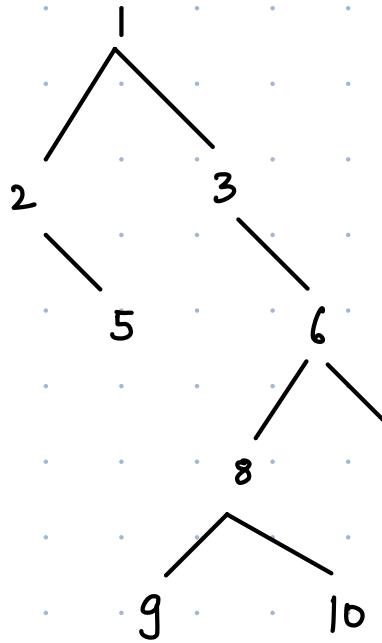
1 ←
2 3 ←
5 8 10 13 ←
6 9 7 4 ←
12 11



Right view

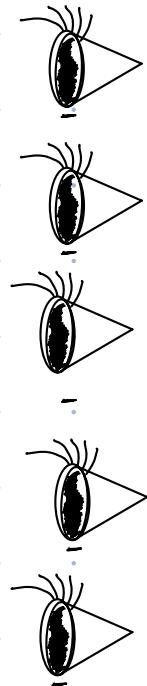


Last Node



1 3 6 7 10

1
3
6
7
10



Pseudocode :

```

last = root
q. enqueue (root)
while (q. isEmpty() == false) {
    x = q. dequeue();
  
```

```

        if (x.left != NULL) { q. enqueue(x.left) }
        if (x.right != NULL) { q. enqueue(x.right) }
    
```

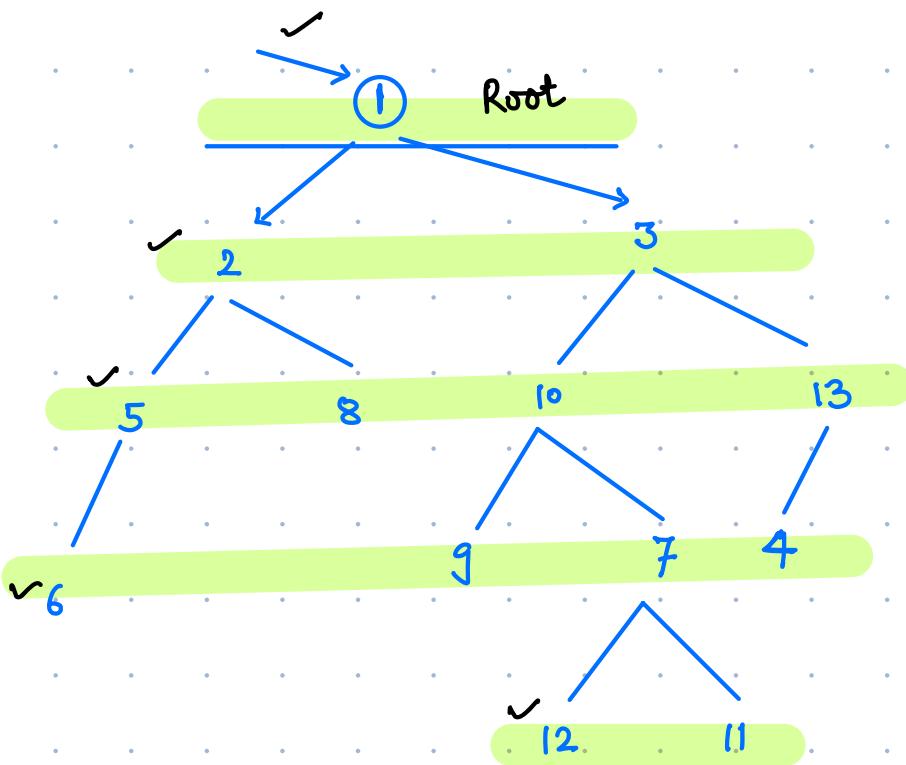
```

if (x == last && q.isEmpty() == false) {
    print(x.data);
    print("\n");
    last = q.rear();
}

```

TC: $2N = O(N)$
SC: $O(N)$

Left view



last \rightarrow

next node

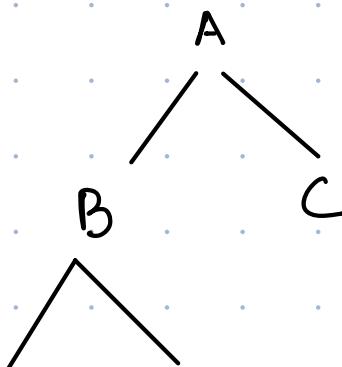
\downarrow
first
node

1 2 5 6 12

Types of binary tree —

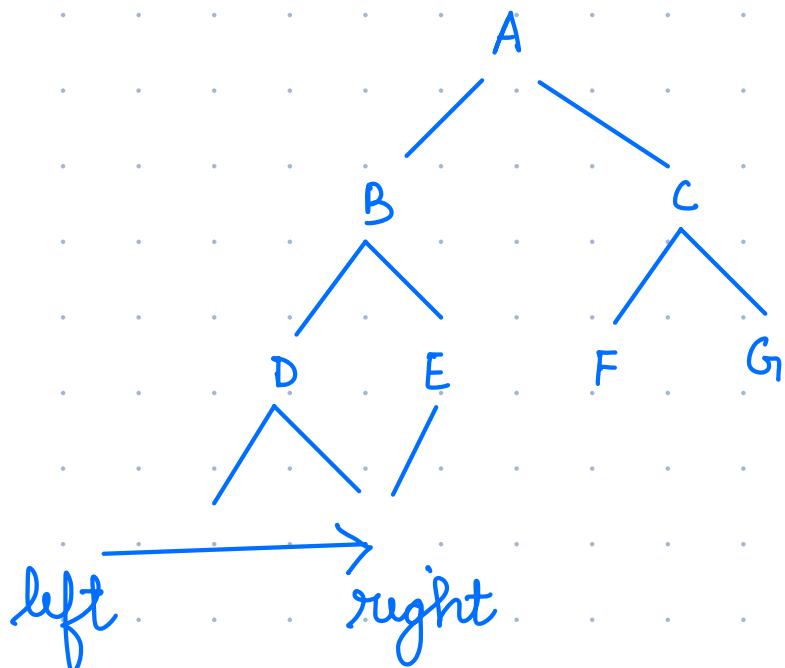
Proper binary tree (Strict BT)

0 or 2 children (never has 1 child)



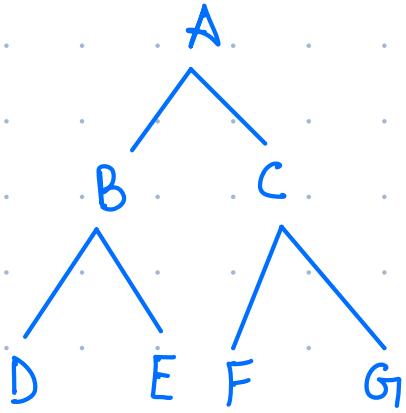
Complete Binary tree

All level of tree completely filled, except last level.



Perfect binary tree

- All internal nodes have exactly 2 children
- all the leaf nodes are at same level

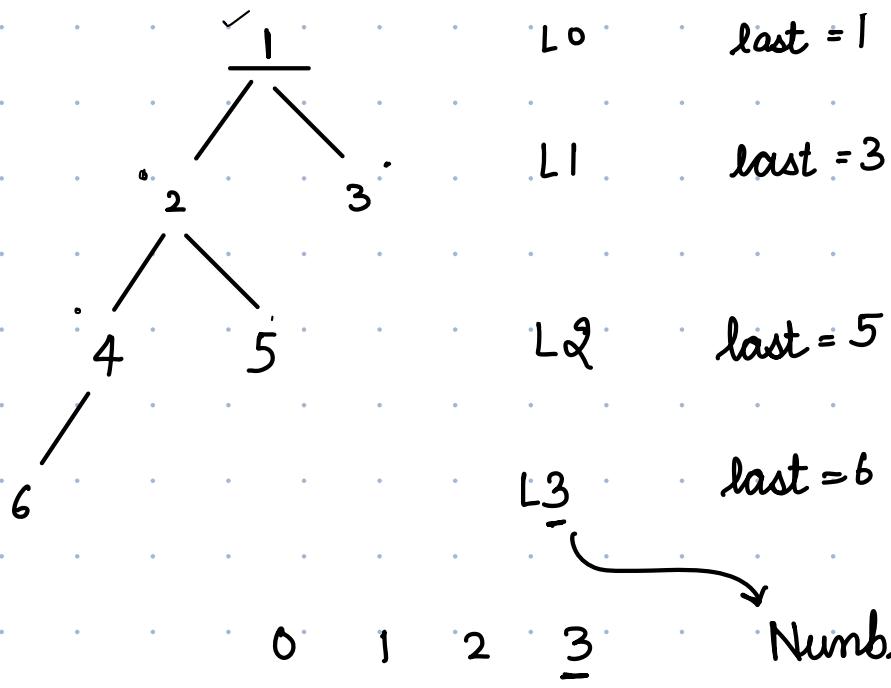
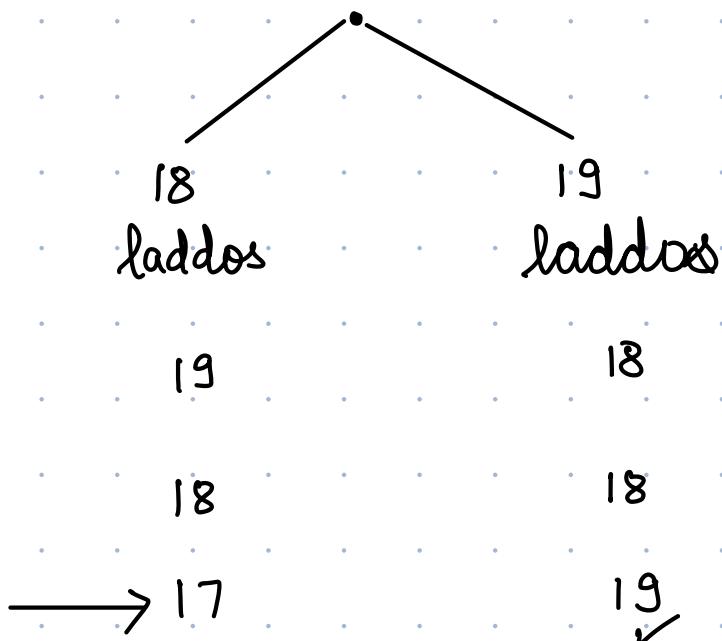


Q.

Check height balanced tree

For all nodes if

$$|h_l - h_r| = 0 \text{ or } 1$$



Height of tree = Height of root node

function calcHeight (root) {
 if (root is leaf) { return 0; }
 ~ hl = calcHeight (root.left);
 ~ hr = calcHeight (root.right);
 return Max(hl, hr) + 1 ← Post Order
}

function heightbalanced (root) {
 if (root == null) {
 return true;
 }
 hl = calcHeight (root.left);
 hr = calcHeight (root.right);
 if (abs(hl - hr) > 1) {
 return false;
 }
 return heightbalanced (root.left);
 && heightbalanced (root.right);
}

At any node —

$$|h_l - h_r| > 1$$

global variable → false

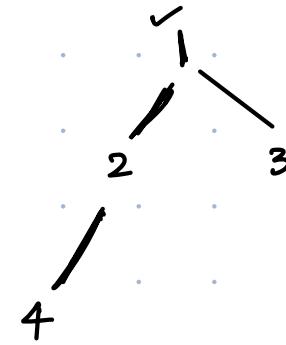
boolean ishb = true

```
function height (root, —) {  
    if (root is leaf) { return 0; }  
    if (root == null) { return -1 }  
    l = height (root.left)  
    r = height (root.right)  
    if (abs (l - r) > 1) ishb=false  
    return max (l, r) + 1  
}
```

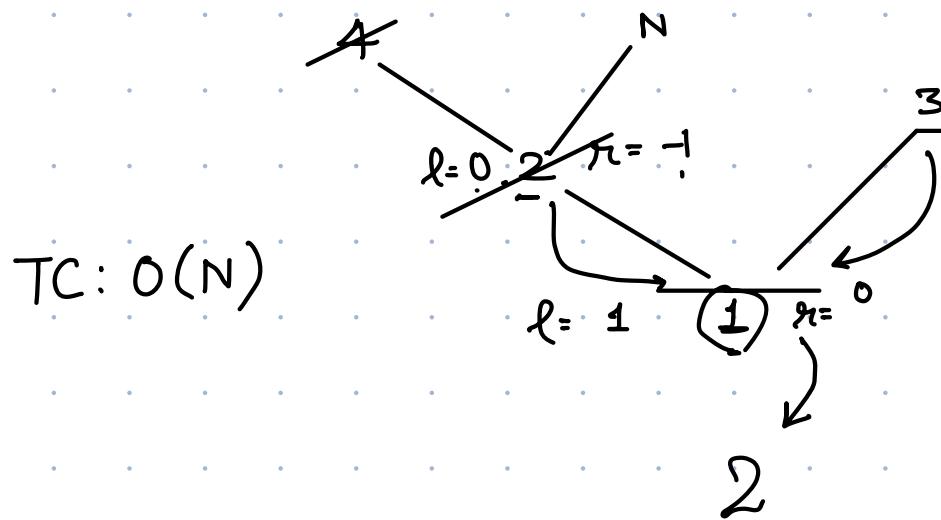
Dry Run →

boolean ishb = true

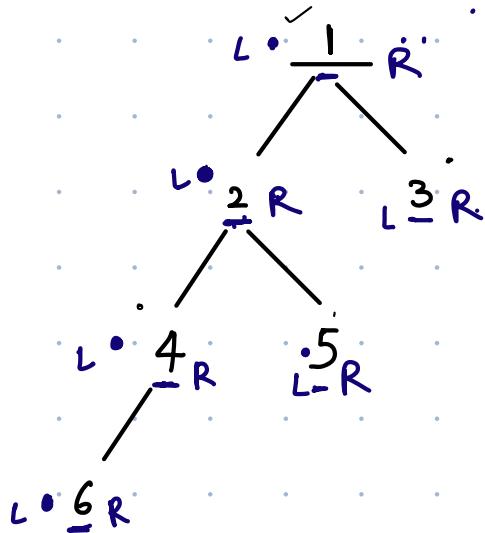
function height (root, →) {
 if (root is leaf) { return 0; }
 if (root == null) { return -1 }
 l = height (root.left)
 r = height (root.right)
 • if (abs (l - r) > 1) ishb=false
 return max (l, r) + 1
}



ishb = true



Q Consider inorder & post order — you have to construct binary tree



6 4 2 5 1 3

Inorder

Go to left

Print

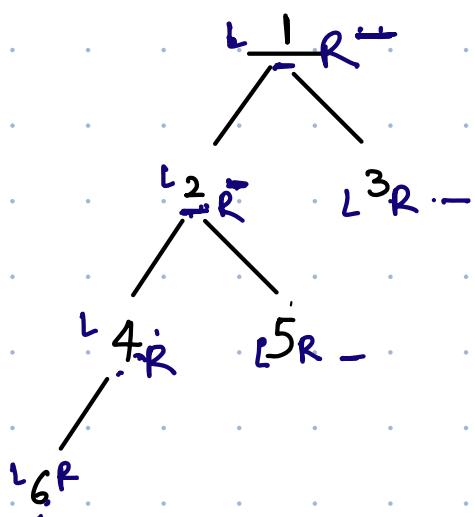
Go to right

PostOrder

Go to left

Go to right

Print

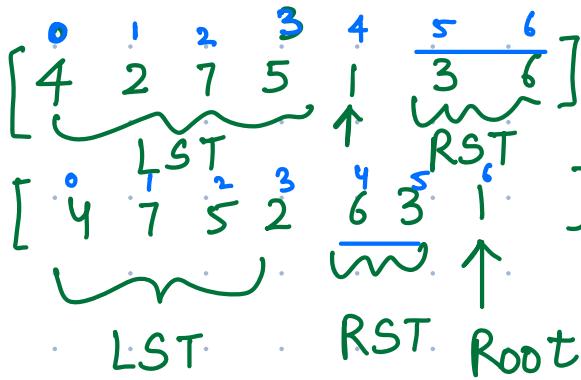


6 4 5 2 3 1

PostOrder

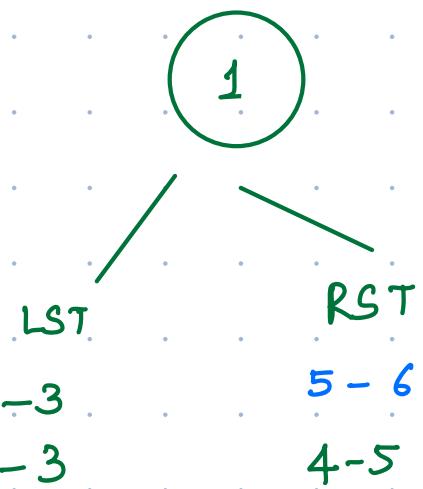
Idea:

In:



Post:

RST
Root



Inorder

LST RST

4 2 7 5 1 3 6

PostOrder

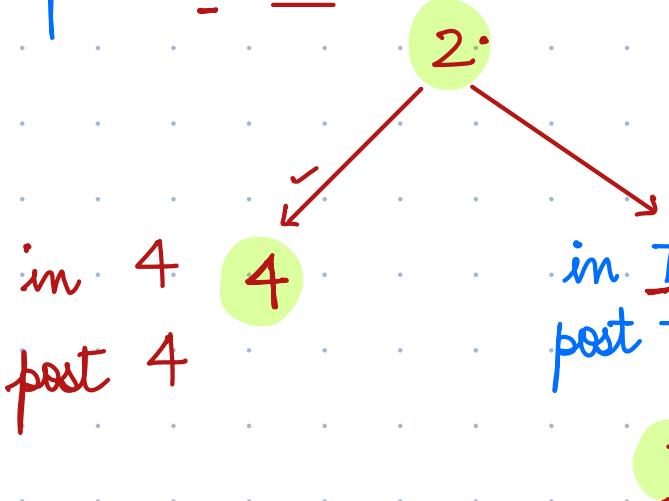
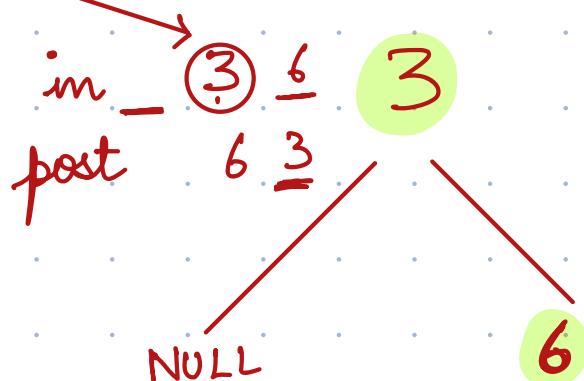
4 7 5 2 6 3 1

LST

RST

↑ Root

in 4 2 7 5
post 4 7 5 2




```
root.right = build tree (subarray (inorder,  
                                 rootIndex-1, end  
                               ),  
                         subarray (postOrder,  
                                 rootIndex  
                               ,  
                               end - 1  
                             ))
```

return root

}