

2 D ARRAY

Act
as if what
you do makes
a difference.
It does.

—WILLIAM JAMES



Good
Morning ☺

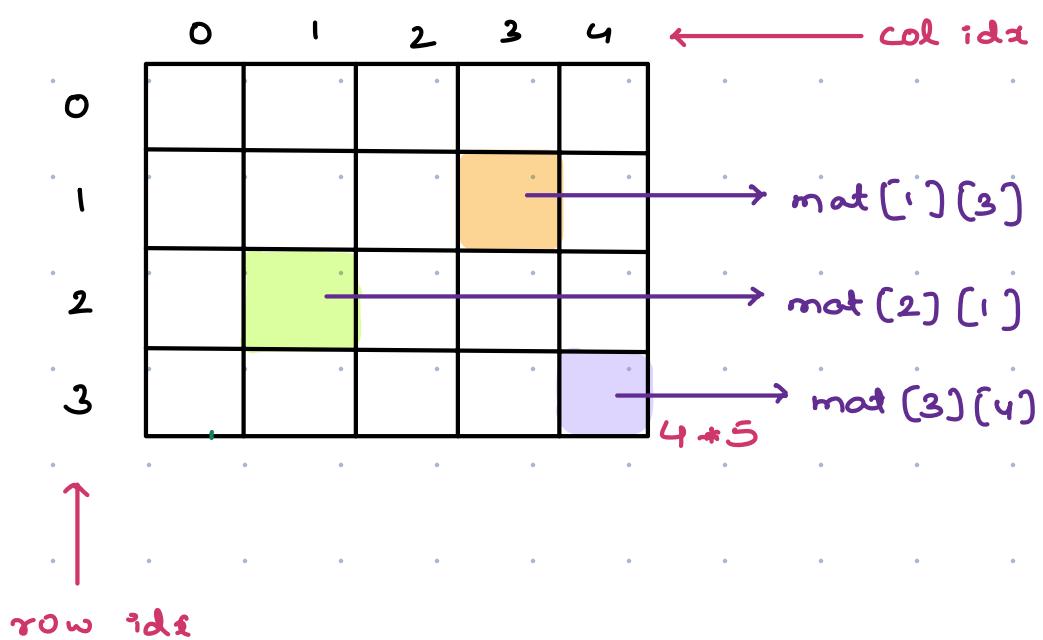
Content

Starting at 7:05

01. Matrix basics
02. Print row wise sum
03. Print col wise sum
04. Print diagonal of sq. matrix
05. Print all diagonal of rectangular matrix
06. Transpose of matrix
07. Rotate matrix by 90°

Matrix / 2D Array → Data in the grid format
 → Data represented in form of
 tables
 → Rows & Columns

Declare → int mat [n] [m]
 Data type Name of matrix rows columns



0	1	2	...	j	.	.	.	M-1
0				(0,j)				
1				(1,j)				
2				(2,j)				
i	(i,0)	(i,1)	(i,2)	(i,j)		(i,m-1)
:				⋮				
N-1				(n-1,j)				

$\rightarrow A[0][M-1]$

$\rightarrow A[N-1][M-1]$

$n \times m$

If we want to move on i' row

↳ col will change from 0 to m-1

If we want to move on j' col

↳ row will change from 0 to n-1

Q Given $\text{mat}[N][M]$. print row wise sum.

	0	1	2	3	<u>sum</u>
0	4	3	1	7	15
1	6	2	3	4	15
2	5	3	2	7	17

Idea → Iterate on a particular row & calculate the sum.
Row = 0 to $(n-1)$

```
void printrowsum( int [ ] [ ] A )
```

```
int N = A.length; // rows
```

```
int M = A[0].length; // cols
```

Tc: $O(N \times M)$

```
for ( r=0 ; r < N ; r++ ) {
```

Sc: $O(1)$

```
    int sum=0
```

```
    for ( c=0; c < M ; c++ ) {
```

```
        sum = sum + A[r][c];
```

```
    }
```

```
    System.out.println(sum);
```

* Not optimal but another way

```
void printrowsum( int [ ] [ ] A)
```

```
    int N = A.length; // rows
```

```
    int M = A[0].length; // cols
```

```
    int [ ] ans = new int [n]
```

```
    for ( r=0 ; r < N ; r++ ) {
```

```
        int sum=0
```

```
        for ( c=0 ; c < M ; c++ ) {
```

```
            sum = sum + A[r][c];
```

```
        ans[r] = sum;
```

TC : O(N*M)

SC : O(N)

```
    for ( i=0 ; i < ans.length ; i++ ) {
```

```
        println( Ans[i] );
```

Q Given mat[n][m] , print col wise sum

	0	1	2	3
0	4	3	1	7
1	6	2	3	4
2	5	3	2	7

sum 15 8 6 18

```
void print_col_sum( int [][] A )
{
    int N = A.length;      // rows
    int M = A[0].length;   // cols
    Tc: O(N*M)

    for (c=0; c < M; c++) {
        int sum=0
        Sc: O(1)

        for ( r=0; r < N; r++) {
            sum = sum + A[r][c];
        }
    }
    Println(sum);
}
```

Q3 Given a square matrix, print its diagonal.

→ Principal Diagonal

→ Top left to bottom right

→ Anti Diagonal

→ Top right to Bottom left corner

	0	1	2	3
0				
1				
2				
3				
0	1	2	3	4
1	5	6	7	8
2	9	10	11	12
3	13	14	15	16

Principal diagonal

0	0
1	1
2	2
3	3

4

4 → stop

```

for (i=0 → n)
    for (j=0 → n) {
        if (i==j) print A[i][j] :
    }

```

$Tc: O(N \times M)$

3
3

Principal diagonal

```
i=0, j=0
```

```

while (i < N && j < N)
    println (A[i][j])
    i++;
    j++;

```

3

$Tc: O(n)$
 $Sc: O(1)$

Anti-diagonal

```
i=0, j=n-1
```

```

while (i < N && j ≥ 0)
    println (A[i][j])
    i++;
    j--;

```

3

Anti diagonal

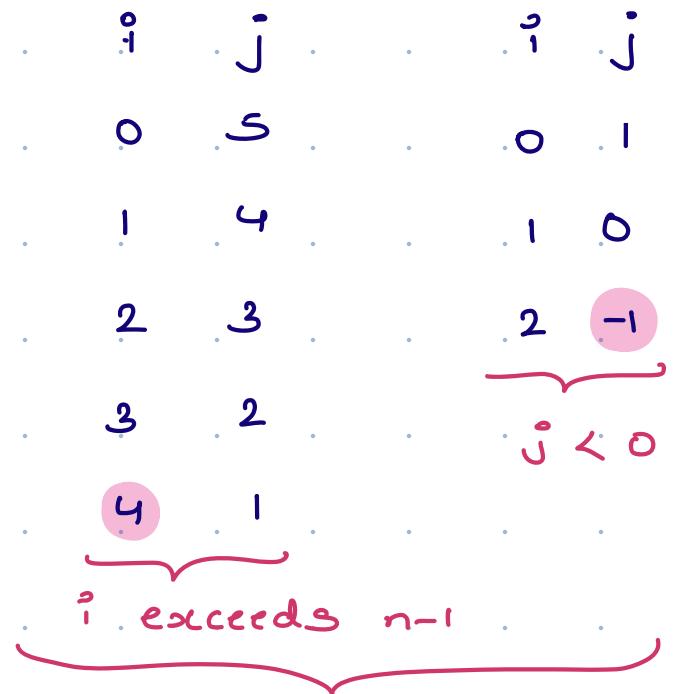
row	col
0	3
1	2
2	1
3	0

4 -1 stop

03. Given a matrix $mat[N][M]$, print all the diagonals from right to left. (Anti-diagonals)

Note:- Diagonals should start from 0' row & can also start from last col.

0	1	2	3	4	5
0	(0,0)				(0,5)
1	(1,0)			(1,4)	(1,5)
2			(2,3)		(2,5)
3		(3,2)			(3,5)



Obs :- $i \leq n$ & $j \geq 0$

$\Delta[3][5]$

0	1	2	3	4	5
0	1	2	3	4	5
1	6	7	8	9	10
2	11	12	13	14	15

Output

1
2 6
3 7 11
4 8 12
5 9 13
10 14
15

Idea → Print all diagonals start from 0th row

print all diagonals start from last col

void print anti diagonals (int [][] A) N*m

// print all diagonals from 0th row

for (c=0 ; c<m ; c++) {

i=0, j=c

// start point of a diagonal

while (i < n && j ≥ 0)

print (A[i][j]);

i++;

j--;

TC: O(n*m)

SC: O(1)

println();

// print diagonals from last col

for (r=1 ; r < n ; r++)

i=r, j=m-1

while (i < n && j ≥ 0)

print (A[i][j]);

i++;

j--;

println();

}

Q4. Given a mat $[N][N]$, calculate transpose of matrix

Expected sc: $O(1)$

interchanging rows & columns

0	1	2	3	4
1	2	3	4	5
6	7	8	9	10
11	12	13	14	15
16	17	18	19	20
21	22	23	24	25

interchange
rows & cols

0	1	2	3	4
1	6	11	16	21
2	7	12	17	22
3	8	13	18	23
4	9	14	19	24
5	10	15	20	25

Output

$\text{mat}[0][1] \longleftrightarrow \text{mat}[1][0]$

$\text{mat}[0][2] \longleftrightarrow \text{mat}[2][0]$

$\text{mat}[0][3] \longleftrightarrow \text{mat}[3][0]$

$\text{mat}[i][j] \xleftrightarrow{\text{swap}} \text{mat}[j][i]$

$i = j \rightarrow \text{no need to swap}$

```
void transpose ( int ( ) ( ) A)
```

```
for ( i=0; i<n; i++) {
```

```
    for ( j=i+1 ; j<n; j++)
```

```
        temp = A[i][j]
```

```
        A[i][j] = A[j][i]
```

```
        A[j][i] = temp;
```

TC: $O(n^2)$

SC: $O(1)$

3

3

3

	0	1	2
0	1	2	3
1	4	5	6
2	7	8	9

i	j
0	0
0	1
0	2
0	3
1	0

$$A[0][\cdot] \leftrightarrow A[\cdot][0]$$

$$A[0][2] \leftrightarrow A[2][0]$$

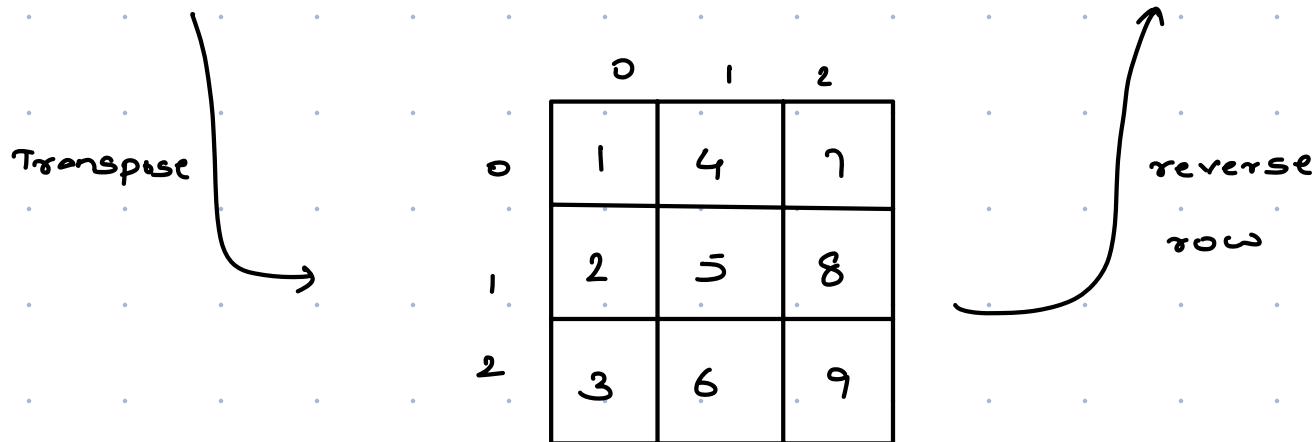
$$A[\cdot][0] \leftrightarrow A[0][\cdot]$$

Q Given a $\text{mat}[n][n]$, rotate the matrix by 90° in clockwise direction.

	0	1	2
0	1	2	3
1	4	5	6
2	7	8	9

Rotation by
 90°

	0	1	2
0	7	4	1
1	8	5	2
2	9	6	3



```
void rotateMatrix( int [][] A)
```

```
    transpose(A);
```

```
    for (r=0; r<n; r++) {
```

```
        i=0, j=n-1
```

```
        while (i < j) {
```

```
            temp = A[r][i]
```

```
A[r][i] = A[r][j] .
```

```
A[r][j] = temp ;
```

```
            i++, j--;
```

TC: $O(n^2)$

SC: $O(1)$

Doubts

+ Rectangular matrix

0	1	2	3	4
0	2	3	4	5
1	6	7	8	9
2	11	12	13	14
3	15			15

$n \times m$

1	6	11
2	7	12
3	8	13
4	9	14
5	10	15

Ans

$m \times n$

$$A[0][1] \rightarrow A[1][0]$$

for ($i=0; i < n; i++$)

 for ($j=0; j < m; j++$) {

 Ans[j][i] = $A[i][j]$

 3

 3

}

Transpose

i	-	j	
0	11	6	1
1	12	7	2
2	13	8	3
3	14	9	4
4	15	10	5

reverse

Transpose

	i	j
0	0	1
1	1	2
2	2	0

→ 0 → 1 → 2

→ 0 → 1 → 2

→ 0 → 1 → 2

for ($r=0; r < n; r++$) {

$i=0, j=n-1$

while ($i < j$) {

int temp = $A[r][i]$

$A[r][i] = A[r][j]$

$A[r][j] = temp$

$i++, j--$

y

$0 < 3$

$$i = \emptyset \quad < \quad j = 2$$

1 j = 1

$$\text{temp} = A[0][0] = 1$$

$$A[0][0] = A[0][2]$$

$$A[0][2] = 1$$

y

$1 < 3$

$$i \quad j$$

0 2

$$\text{temp} = A[1][0] = 2$$

Q : elements $\leq B$ together

$$\boxed{B=5}$$

4 6 5 8 3 2 7 7

count no. of elements that
needs to be group together = 4

minimum
swap

4 6 5 8 3 2 7 7

int bad elements \rightarrow look at no. of elements I need
to swap

bad ele count = 2

ans = 00 vs 2 = 2

= 2

ans = 2

bad = 2 - 1 = 1

ans = 1

bad = 1 + 1 = 2

ans = 1

$$\text{bad} = 2 - 1 + 1 = 2$$

$$\text{ans} = 1$$

```
int minSwaps(int arr[], int B)
```

```
    count = 0
```

```
    for (int i=0; i<n; i++) {
```

```
        if (arr[i] >= B) count++;
```

```
}
```

// Count → No. of ele we have to group together.

```
    if (count == 0 || count == 1 || count == n) return 0;
```

// Calculate ans for first window

```
    bad = 0, ans = ∞
```

```
    for (i = 0; i < count; i++) {
```

```
        if (arr[i] >= B) bad++;
```

```
}
```

TC: O(n)

SC: O(1)

```
    if (bad < ans) ans = bad;
```

// Apply sliding window technique

```
s = 1, e = count
```

```
while (e < n)
```

```
    if (arr[s-1] >= B) bad--; // dropping bad
```

```
    if (arr[e] >= B) bad++; // adding bad
```

```
    if (bad < ans) ans = bad;
```

```
s++, e++;
```

```
return ans;
```

