Two variables / pointers

"The more you sweat in peace, the less you bleed on war"

Q. given an integer sorted array A , integer K.

Find any pair (i, j) ⟶ A[i] + A[j] = K

i != j

[-5 -2 1 8 10 12 15] , K = 11

0 1 2 3 4 5 6

i j

2, 4

{-3 0 1 3 6 8 11 14 18 25}

0 1 2 3 4 5 6 7 8 9

K = 12

i = 2, j = 6

A[2] + A[6] = 12

App 1:

Run a loop and for every pair of indices (i, j)

check if sum = K

TC : O(N^2) , SC : O(1)

App2 :

A[i] + A[j] = K , A[j] = K - A[i]

for every i,

N — binary search and find K - A[i]

log N

TC : N log N

SC : O(1)

App3 : ✓ Hashing i

A[]

K - A[i]

↑

hashset

TC : O(N)

SC : O(N)

App 4:

2 variables/pointers

0	1	2	3	4	5	6
-5	-2	1	8	10	12	15
i					j	

k = 11

$-5 + (-2) = -7 < 11$

increase

$-5 + 1 = -4 < 11$

$-5 + 8 = 3 < 11$

$-5 + 10 = 5 < 11$

Not correct

---

2 variables/pointers

return true

0	1	2	3	4	5	6
-5	-2	1	8	10	12	15
		i		j		

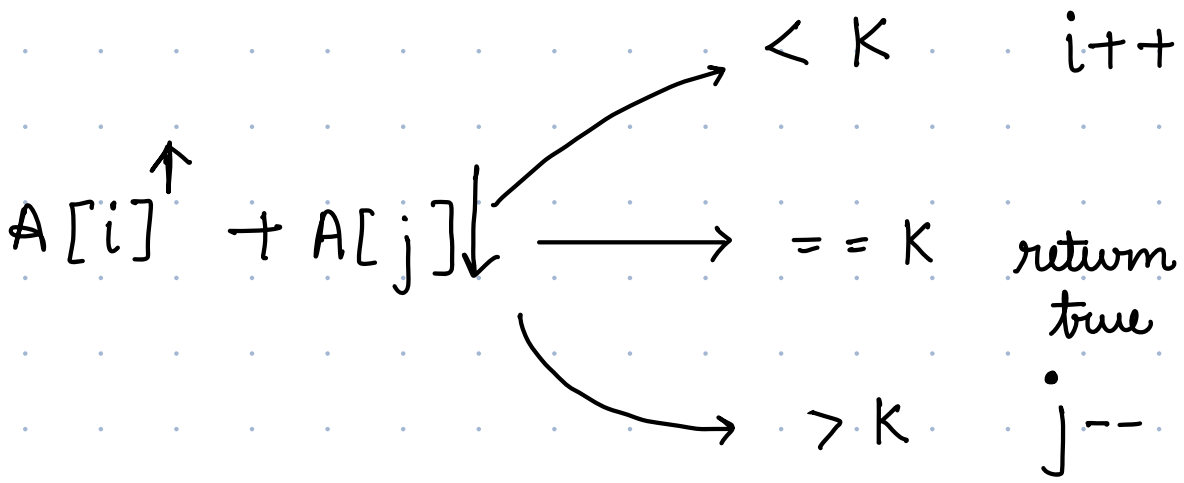
k = 11

$1 + 10 = 10 < 11$ (inc)

$11 == 11$

$13 > 11$

Core idea:



Code:

```
i = 0      j = N-1
while ( i < j ) {
    if ( A[i] + A[j] == K) { return true; }
    else if( A[i] + A[j] < K) {
                i++
    } else {
                j--
    }
}
return false
```

TC: O(N)

SC: O(1)

- Meet in the middle approach

- **Count pairs:**

A = { 1 2 3 4 5 6 8 } , K = 10

Case 1: When elements are different

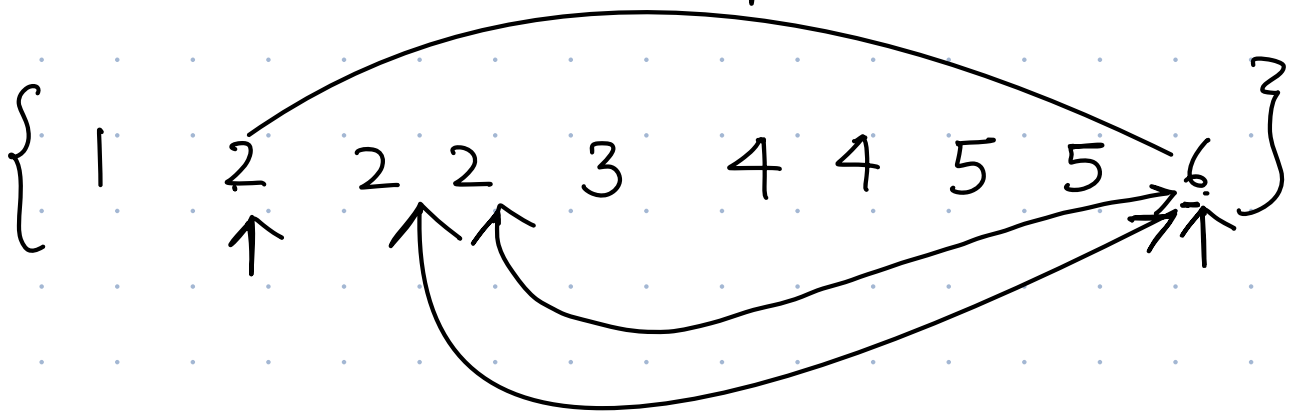
```
Code: count = 0
      i = 0    j = N-1
      while ( i < j ) {
          if ( A[i] + A[j] == K) { count++, i++, j-- }
          elseif ( A[i] + A[j] < K) {
                  i++
          } else {
                  j--
          }
      }
      return count
```

TC: O(N)
SC: O(1)

{ 1 2 3 4 5 6 7 8 }

K = 10	count
(2, 8)	1
(7, 3)	2
(4, 6)	3

Case 2: When elements are duplicate



K = 8

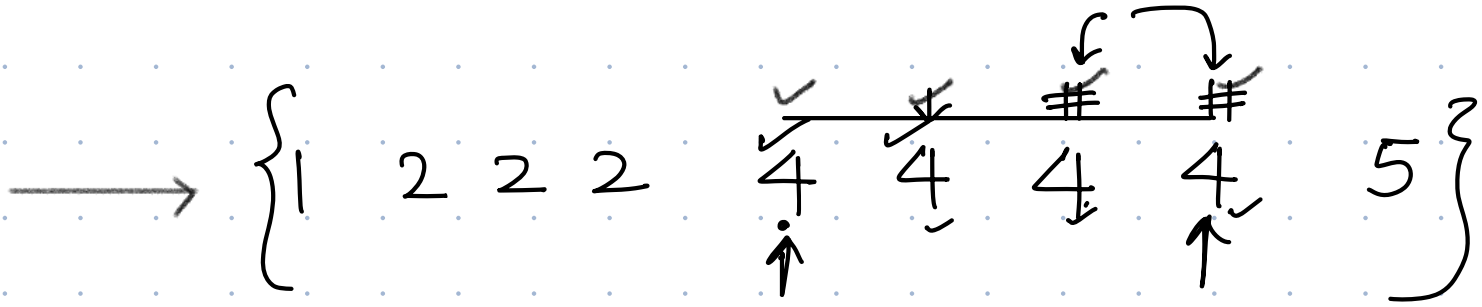
1	1
2	3
3	1
4	2
5	2
6	1

$3 \times 1 = 3$

$freq(A[i]) = m$

$freq(A[j]) = n$

$A[i] + A[j] == K \rightarrow m \times n$



K = 8

$^4C_2 \quad \frac{4 \times 3}{2} \quad \begin{matrix} 3 \\ 2 \\ 1 \end{matrix} = 6$

$A[i] == A[j]$

$\frac{4 \times 3}{2}$

out of 4, choose any 2

$^{n}C_{r} = {}^{4}C_{2}$

4 4

4 4

$A[i] + A[j]$ } K

$A[j] + A[i]$ }

Q.

Pair difference

Sorted Array , K

Pair (i, j) $\longrightarrow$ $A[j] - A[i] = K$

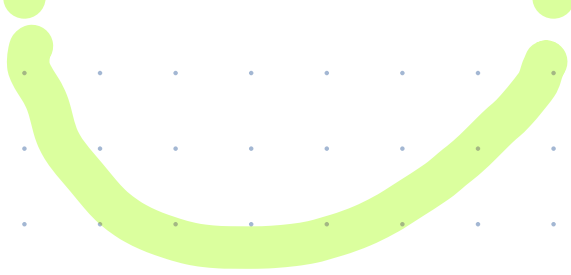
$K > 0$

- Run 2 loops
- Binary search

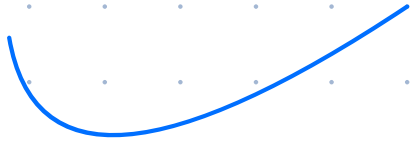
$\longrightarrow$

0	1	2	3	4	5
1	2	4	5	6	12

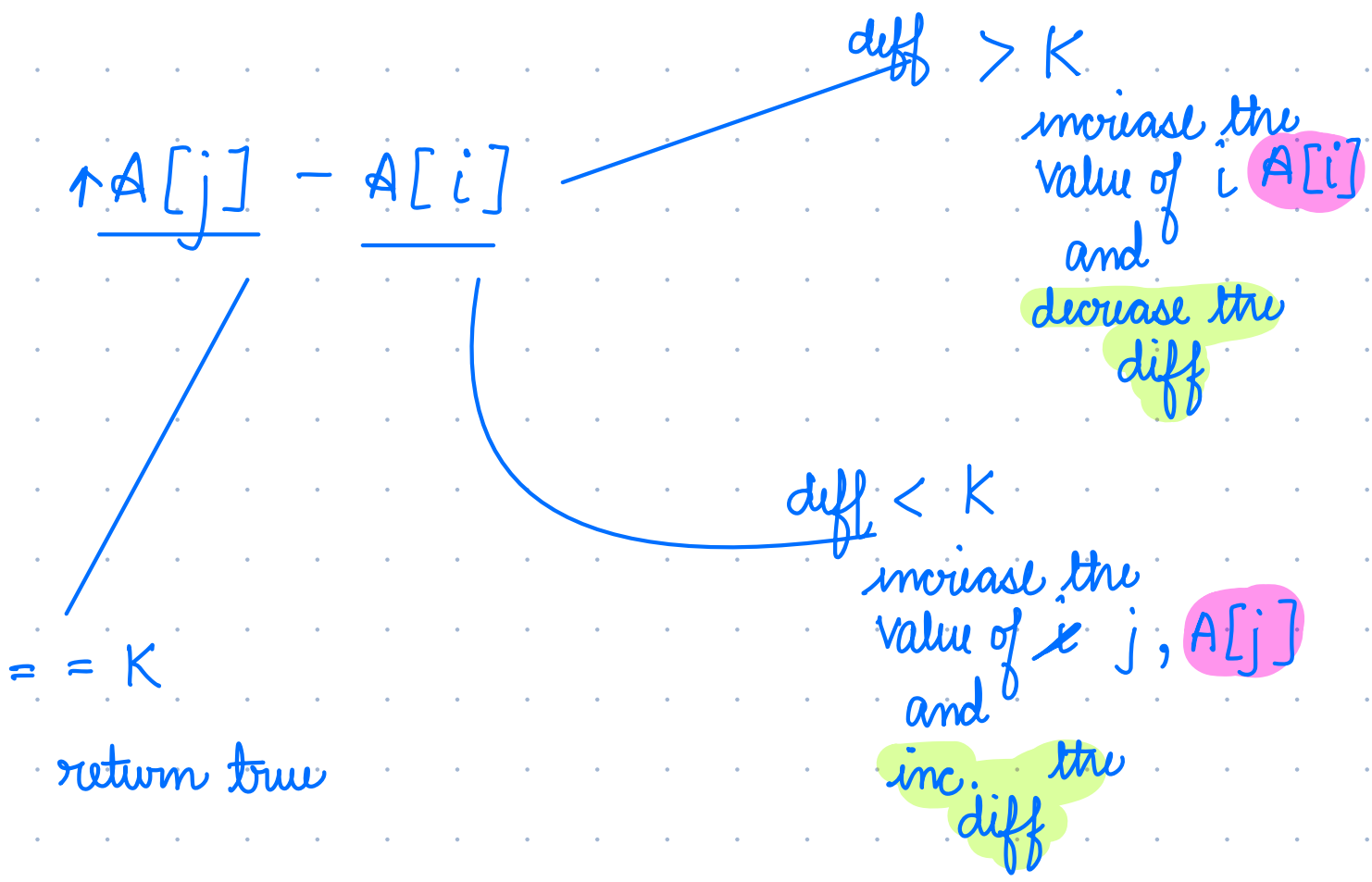
, $K = 10$

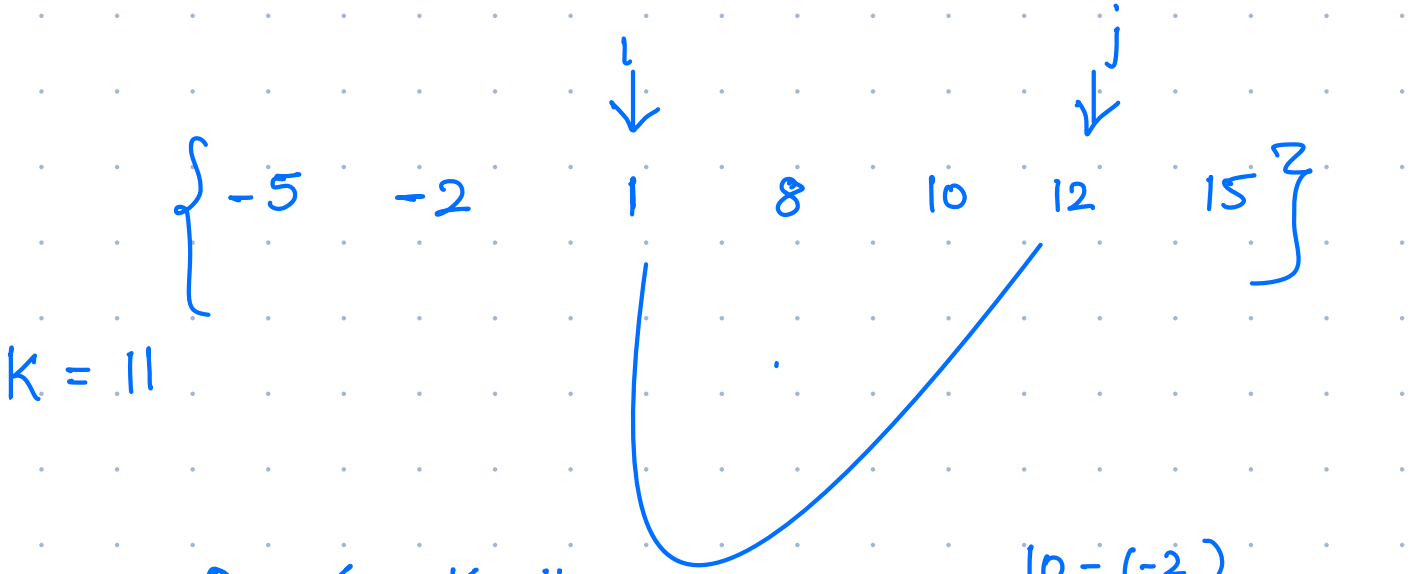


-5	-2	1	8	10	12	15



$K = 11$





$0 < K = 11$

$3 < 11$

$6 < 11$

$13 > 11$

$8 - (-2)$

$10 < 11$

$10 - (-2)$

$12 > 11$

$10 - 1$

$9 < 11$

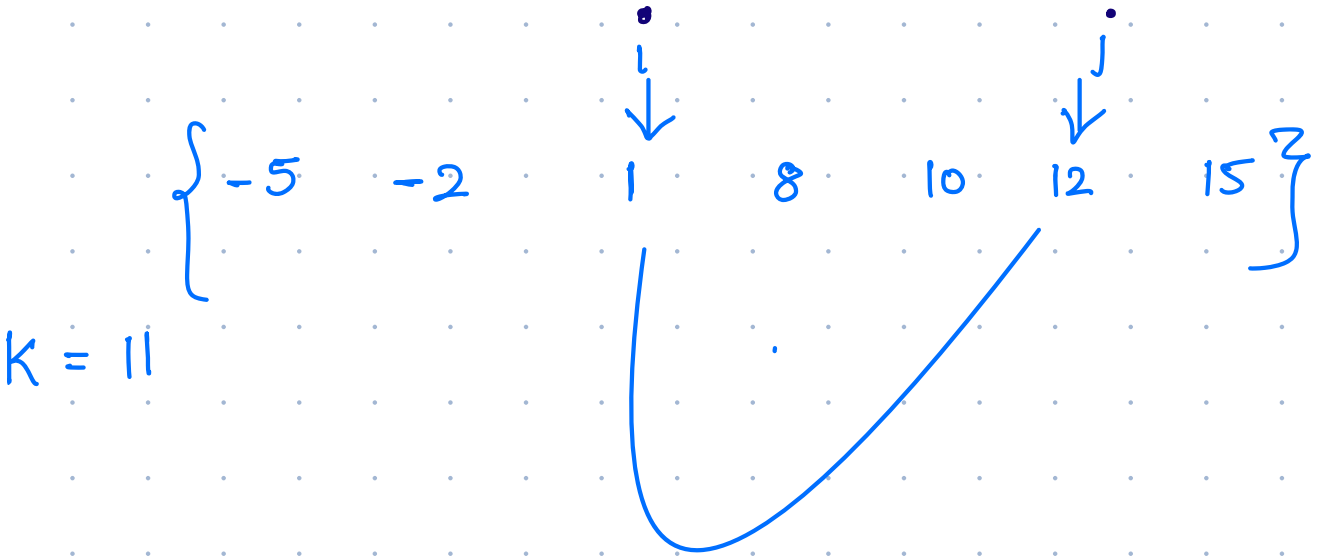
$12 - 1 = 11 == 11$

Code: i = 0, j = 0

j++

j = 1

```
while ( j < N) {
    diff = A[j] - A[i]
    if (diff == K) {
        return true
    } else if (diff < K) {
        j++
    } else {
        i++
    }
}
return false
```



Q. Check if subarray with sum = K exists.

K = 9

Al: { 1 2 5 4 3 }

3 K

12

Prefix Sum approach: 1 3 8 12 15

if diff of prefix sum = K, then exists

A2:

K = 9

{ 1 2 5 4 3 }

i ↑ 5

j ↑ 4

sum = 1 < K

3 < 9

8 < 9

12 > ✗ 9

11 > 9

9 == 9

{ 1 3 15 10 20 3 23 }, K = 33

i ↑ 10

j ↑ 3

Sum = 1 < 33

+3

4 < 33

+15

19 < 33

10

29 < 33

20

49 > 33

− 1

48 > 33

− 3

45 > 33

− 15

30 < 33

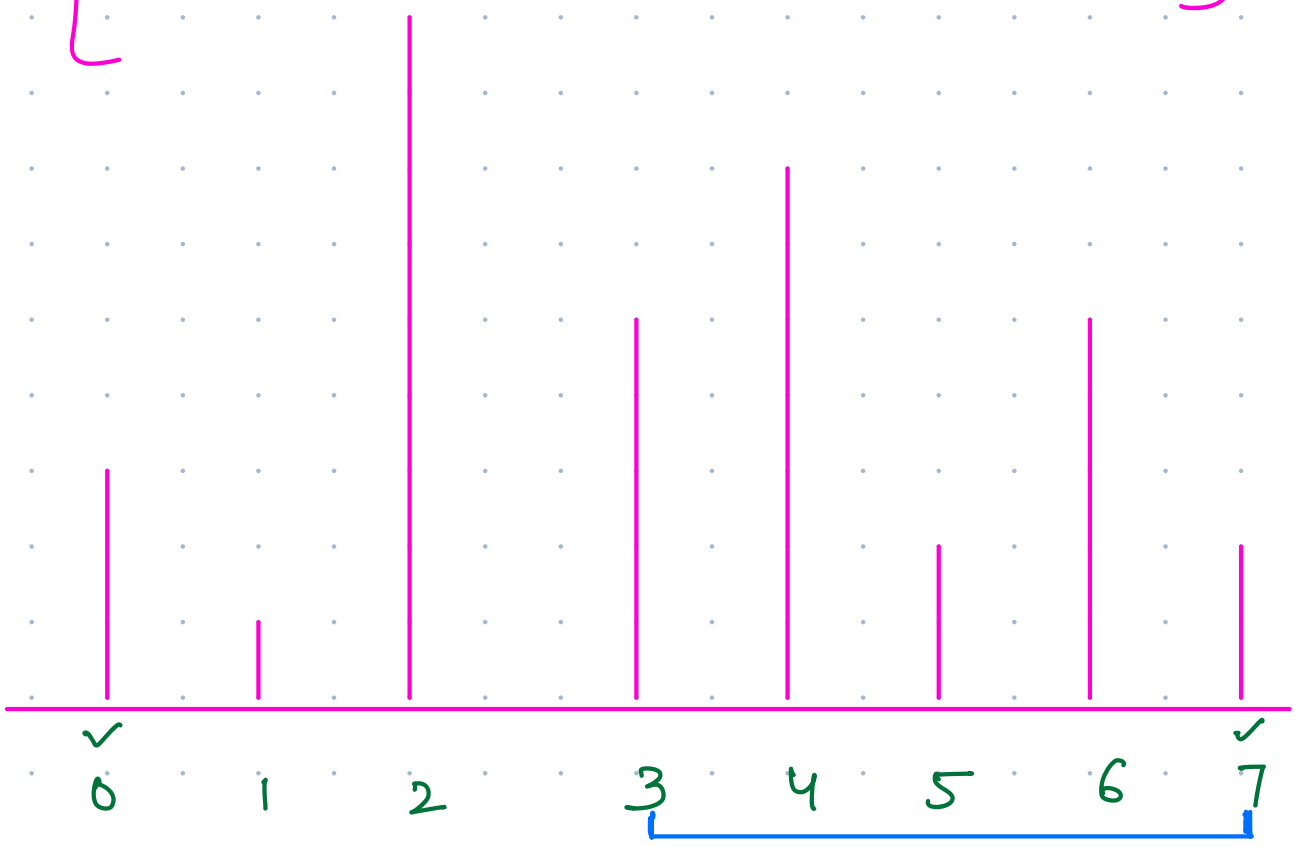
33 = = 33

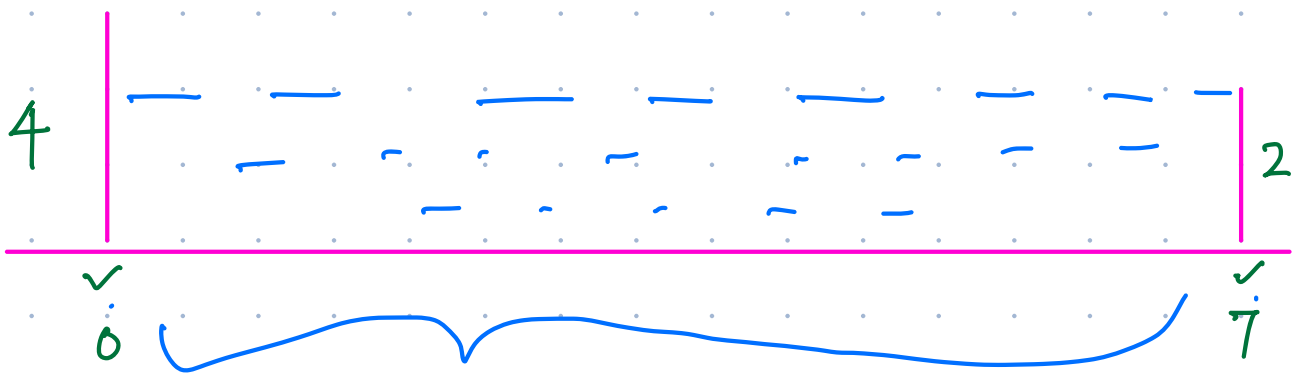
Pseudocode :

normal array

```
i=0, j=0, sum = A[0]
while ( j < N && i<=j ) {
    if (sum == K) {
        return true
    }
    else if (sum > K) {
        sum = sum - A[i]
        i++,
    } else {
        j++
        if (j==N) {break;}
        sum = sum + A[j]
    }
}
return false
```

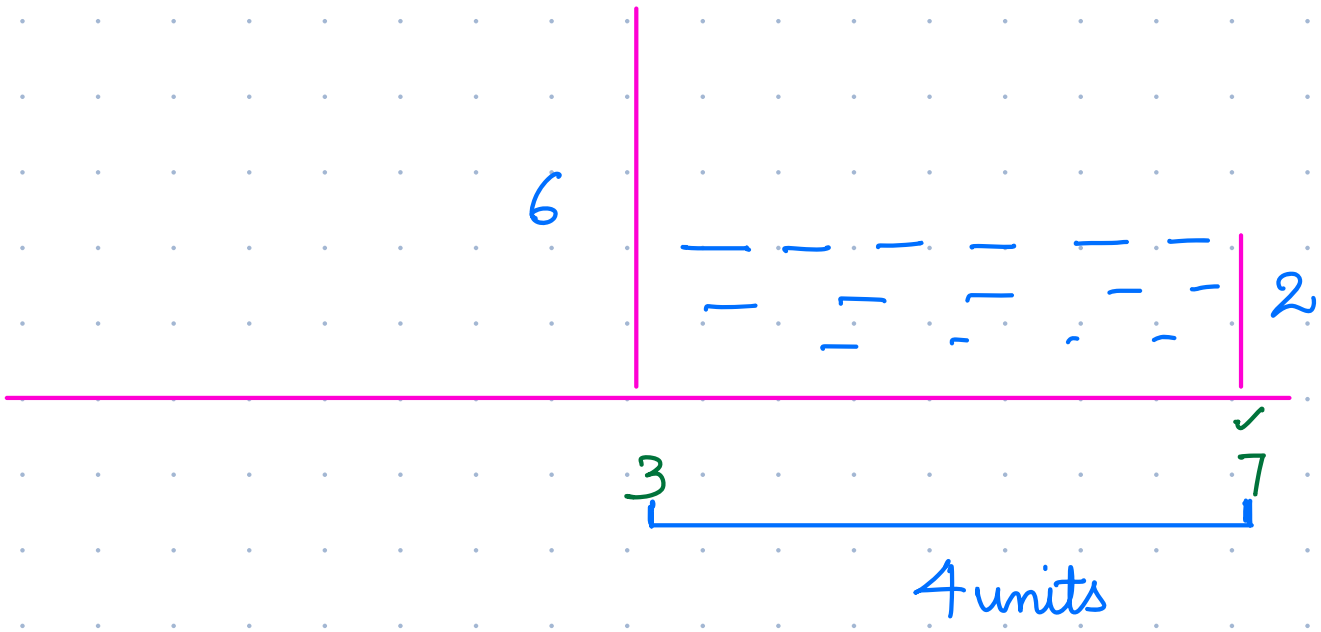
Q. {4 2 10 6 8 2 6 2}



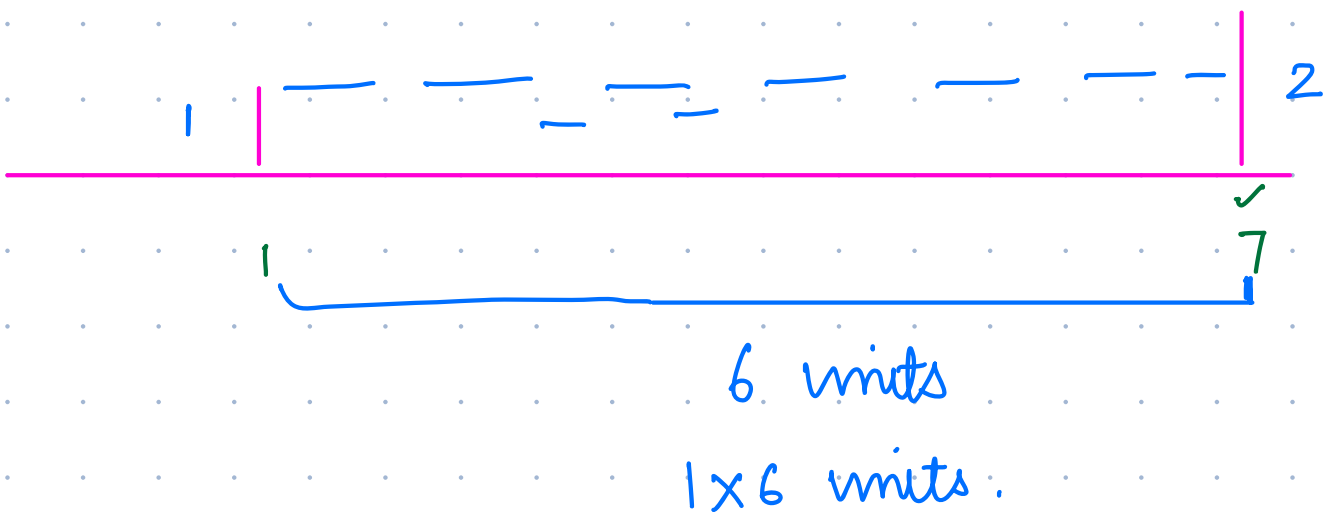


height = 2
width = 7

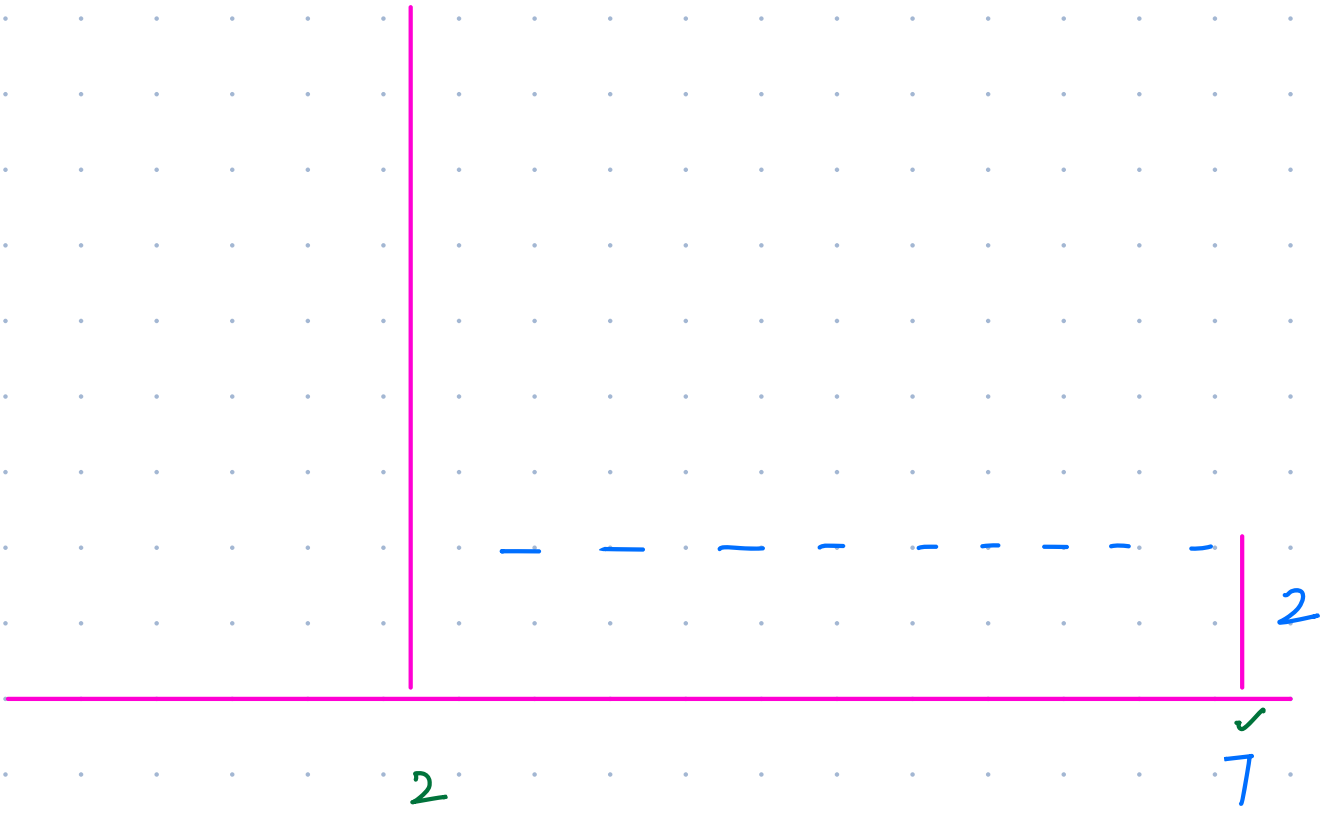
= 14



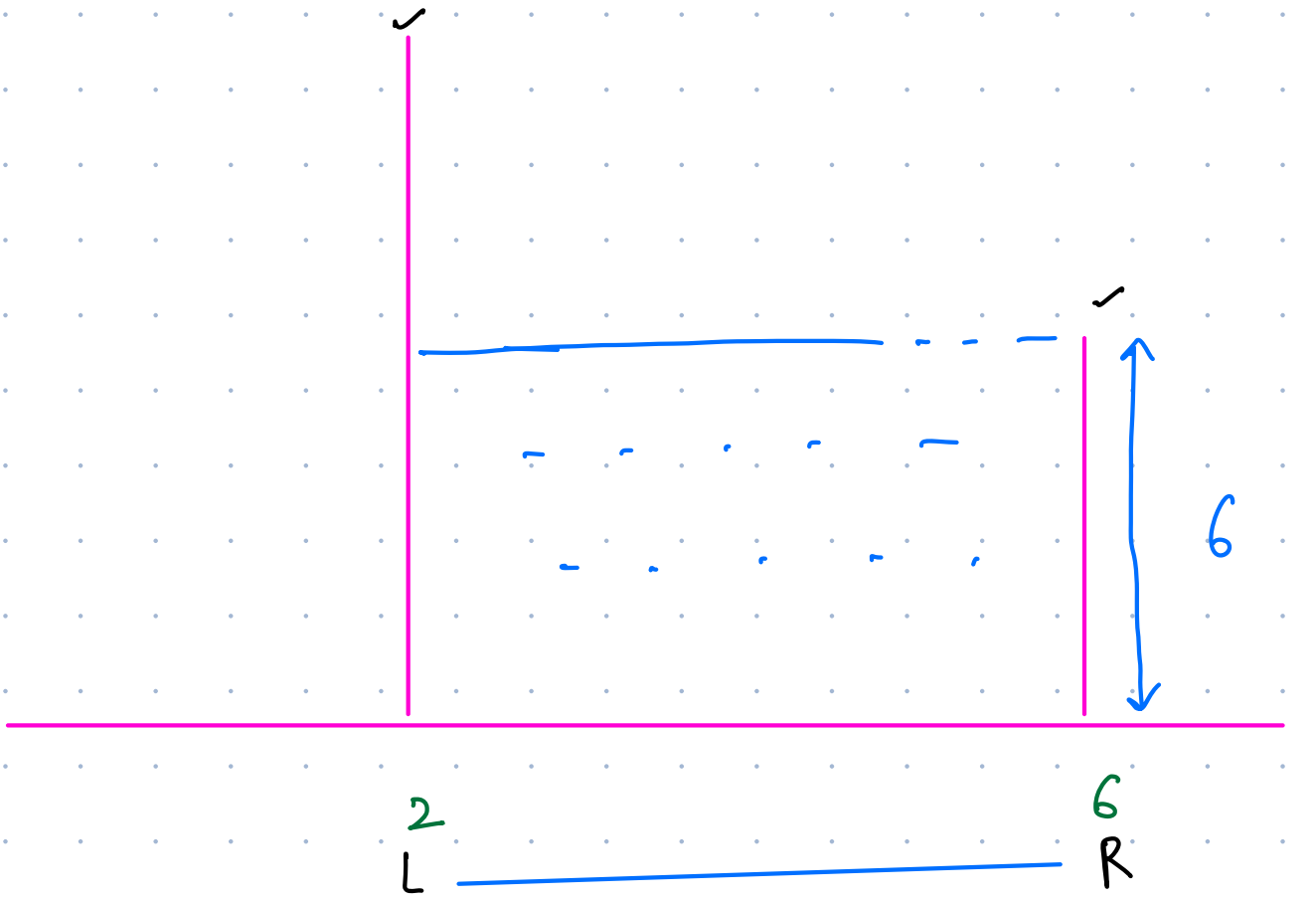
$2 \times 4 = 8$ units



1×6 units.



$5 \times 2 = 10$ units



width = R − L 4 units

24 units.

Brute force :

Run 2 loops

TC: $O(N^2)$

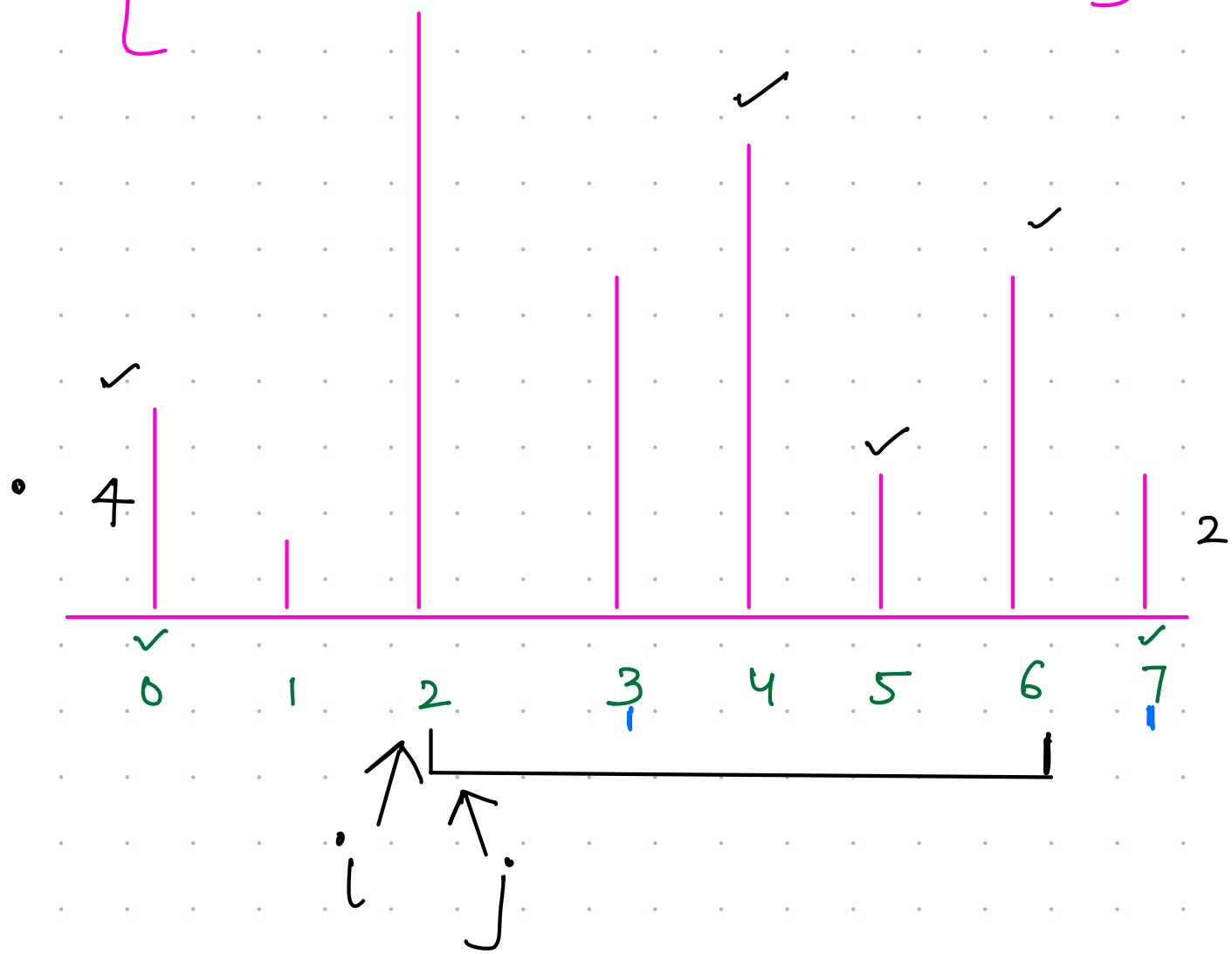
SC: $O(1)$

- Water trapped between 2 walls $= (R - L) \times \min(A[R], A[L])$

  $(R - L)$: width

  $\min(A[R], A[L])$: height

{4 2 10 6 8 2 6 2}



WC = 2 x 7 = 14

= 4 x 6 = 24

= 1 x 5 = 5

= 6 x 4 = 24

= 2 x 3 = 6

$= 8 \times 2 = 16$

$= 6 \times 1 = 6$

Pseudocode:

```
i = 0        j = N-1        maxarea = 0
while ( i < j ) {
    area = (j - i) x min(A[i], A[j])
    maxarea = max(area, maxarea);
    if( A[i] < A[j])
        i++
    else if( A[i] > A[j]){
        j--
    }else {
        i++, j--          // i
                          // j
                          // i, j
    }
}
return maxarea ;
```

TC: O(N)

SC: O(1)

Pair diff:

{ 1 3 10 13 15 20 23 }, K=33

23 - 1 = 22 ↑ < 33

20 - 1 } j--

23 - 3 } i++