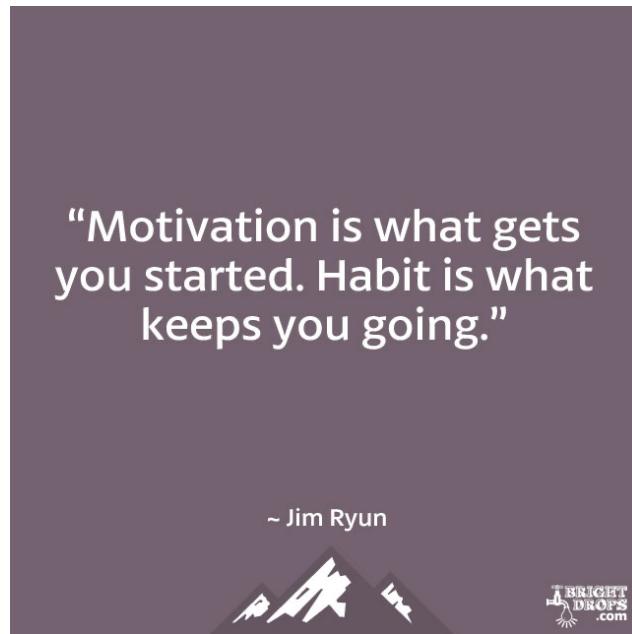


PREFIX SUM



Good
Morning

Today's content

01. Range sum query
02. Prefix sum construction
03. Sum of even indexed elements
04. Special index

01. Given arr [N] elements & Q queries

For each query : Given L & R calculate & print sum of all

elements in range [L R]. L & R both has to be calculate

Note :- L & R are array indices such that $0 \leq L \leq R \leq N$

Constraints

$$1 \leq N, Q \leq 10^5$$

$$-10^9 \leq A[i] \leq 10^9$$

$$0 \leq L \leq R \leq N$$

arr [10] =

-3	6	2	4	5	2	8	-9	3	1
0	1	2	3	4	5	6	7	8	9

$$Q = 5$$

	L	R
i = 0	4	8
1	3	7
2	1	3
3	0	4
4	7	7

$$L = Q[i][0]$$

$$R = Q[i][1]$$

Idea \rightarrow For each query, we have to iterate from L to R and add all elements \downarrow

print the sum

```
void rangesum ( int []A, int [][] Q)
```

```
for ( i=0 ; i< Q.length ; i++) {
```

```
    int L = Q[i][0]
```

```
    int R = Q[i][1]
```

```
    sum = 0;
```

```
    for ( j=L ; j<=R ; j++) {
```

```
        sum = sum + A[j];
```

```
    }
```

```
    print (sum);
```

```
}
```

```
3
```

TC : $O(Q * N)$

SC : $O(1)$

According to the constraint

Constraints

$$1 \leq N, Q \leq 10^5$$

$$-11 \leq A[i] \leq 10^9$$

$$\underline{0 \leq L \leq R < N}$$

BF Approach = $Q * N$

$$= 10^5 * 10^5 = \frac{10^{10}}{X}$$

TLE

Q : Given Indian Cricket Team score , for first 10 overs of Batting .

After every over , total score is given as :

overs →	1	2	3	4	5	6	7	8	9	10
cumulative Scores →	2	8	14	29	31	49	65	79	88	97

$$* \text{ runs scored in } 7^{\text{th}} \text{ over} = \frac{\text{Total runs after } 7^{\text{th}} \text{ over}}{\text{Total runs after } 6^{\text{th}} \text{ over}}$$

$$= 65 - 49 \Rightarrow \underline{\underline{16}}$$

$$* \text{ Runs scored from } 6^{\text{th}} \text{ to } 10^{\text{th}} \text{ over} =$$

$$\underbrace{\text{sc}[1-10]}_{97} = \text{sc}[1-5] + \text{sc}[6-10]$$

$$97 = 31 + x$$

$$x = 97 - 31 = \underline{\underline{66}}$$

* Runs scored in 10th over =

$$sc[1-10] = sc[1-9] + \underbrace{sc[10-10]}$$

$$97 = 88 + x$$

$$x = 97 - 88 = 9$$

* Runs scored from 3rd to 6th over = $sc[6] - sc[3-1]$

$$\Rightarrow 49 - 8 = 41$$

* Runs scored from 4th to 9th over $\Rightarrow sc[9] - sc[4-1]$

$$\Rightarrow 88 - 14$$

$$= 74$$

$or[10] =$	<table border="1"> <tr> <td>-3</td> <td>6</td> <td>2</td> <td>4</td> <td>5</td> <td>2</td> <td>8</td> <td>-9</td> <td>3</td> <td>1</td> </tr> <tr> <td>0</td> <td>1</td> <td>2</td> <td>3</td> <td>4</td> <td>5</td> <td>6</td> <td>7</td> <td>8</td> <td>9</td> </tr> </table>	-3	6	2	4	5	2	8	-9	3	1	0	1	2	3	4	5	6	7	8	9
-3	6	2	4	5	2	8	-9	3	1												
0	1	2	3	4	5	6	7	8	9												

$psum[10] =$	<table border="1"> <tr> <td>-3</td> <td>3</td> <td>5</td> <td>9</td> <td>14</td> <td>16</td> <td>24</td> <td>15</td> <td>18</td> <td>19</td> </tr> <tr> <td>0</td> <td>1</td> <td>2</td> <td>3</td> <td>4</td> <td>5</td> <td>6</td> <td>7</td> <td>8</td> <td>9</td> </tr> </table>	-3	3	5	9	14	16	24	15	18	19	0	1	2	3	4	5	6	7	8	9
-3	3	5	9	14	16	24	15	18	19												
0	1	2	3	4	5	6	7	8	9												

cumulative
sum
array

$Q = 5$

	L	R
0	4	8
1	3	1
2	1	3
3	0	4
4	7	7

$$\begin{aligned} & \rightarrow \text{psum}[8] - \text{psum}[4-1] = 18 - 9 = 9 \\ & \rightarrow \text{psum}[7] - \text{psum}[3-1] = 15 - 5 = 10 \\ & \rightarrow \text{psum}[3] - \text{psum}[1-1] = 9 - (-3) = 12 \\ & \rightarrow \text{psum}[4] = 14 \\ & \rightarrow \text{psum}[7] - \text{psum}[7-1] = 15 - 24 = -9 \end{aligned}$$

Range sum query

```

if ( $L \neq 0$ ) ans = psum[R] - psum[L-1]
else ans = psum[R]
    
```

* Construct psum array

$$A[] = \{ 9, 10, 32, 6, 12, 20, 1 \}$$

0 1 2 3 4 5

$$\text{psum}[] = \{ 9, 10, 42, 48, 60, 80, 81 \}$$

0 1 2 3 4 5

$\text{psum}[i] = \text{sum of all ele from } 0 \text{ to } i$

$$A[] = \{ 9, 10, 32, 6, 12, 20, 1 \}$$

0 1 2 3 4 5

$$psum[] = \{ 9, 10, 42, 48, 60, 80, 81 \}$$

0 1 2 3 4 5

$$pf[0] = A[0]$$

$$pf[1] = \underbrace{A[0]}_{pf[0] + A[1]} + A[1]$$

$$pf[2] = \underbrace{A[0] + A[1]}_{pf[1] + A[2]} + A[2]$$

$$pf[3] = \underbrace{A[0] + A[1] + A[2]}_{pf[2] + A[3]} + A[3]$$

if ($i == 0$) $pf[0] = A[0]$

else { $pf[i] = pf[i-1] + A[i]$ }

```
void rangesum (int []A, int [][]Q)
```

// construct psum[] array

$n = A.size()$

```
long [] psum = new long [n]
psum[0] = A[0]
for (i=1; i<n; i++) {
    psum[i] = psum[i-1] + A[i];
}
```

TC: $O(Q+n)$

SC: $O(n)$

// Answer all queries

```
for (i=0; i< Q.length; i++) {
    int l = Q[i][0]
    int R = Q[i][1]
    if (l == 0) { print(psum[R]) }
    else { print (psum[R] - psum[l-1]); }
```

3

3

03. Given arr[N] elements & Q queries.

For each query : Given L & R calculate & print sum of even indexed ele in given range [L R].

Note :- L & R are array indices such that $0 \leq L \leq R \leq N$

Constraints

$$1 \leq N, Q \leq 10^5$$

$$1 \leq A[i] \leq 10^9$$

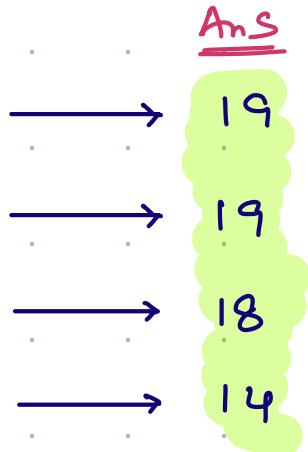
$$0 \leq L \leq R \leq N$$

Eg:- arr[10] :

2	4	3	7	9	8	6	3	4	9
0	1	2	3	4	5	6	7	8	9

Queries ↴

	0	1
0	4	8
1	3	9
2	2	7
3	0	4



Brute force $\rightarrow \{ \text{TODO} \}$

Obs \rightarrow Odd indexes are going to contribute 0 to the sum.

Eg:- arr[10] :

2	4	3	7	9	8	6	3	4	9
0	1	2	3	4	5	6	7	8	9

pre[10] =

2	2	5	5	14	14	20	20	24	24
0	1	2	3	4	5	6	7	8	9

$$pre[0] = A[0]$$

for ($i=1$; $i < n$; $i++$) {

 if ($i \% 2 == 0$) {

$$pre[i] = pre[i-1] + A[i];$$

 else {

$$pre[i] = pre[i-1] + 0;$$

// Answer all queries

for ($i=0$; $i < Q.length$; $i++$) {

$$int L = Q[i][0]$$

$$int R = Q[i][1]$$

 if ($L == 0$) { print(psum[R]) }

 else { print (psum[R] - psum[L-1]); }

* Queries on sum of all odd indexed elements



{ TODO }

HIN →

03. Given $ar[N]$ elements & Q queries.

For each query : Given L & R calculate & print **sum of even ele** in given range [L R]

Note :- L & R are array indices such that $0 \leq L \leq R \leq N$

Constraints

$$1 \leq N, Q \leq 10^5$$

$$1 \leq ar[i] \leq 10^9$$

$$0 \leq L \leq R < N$$

$ar[10]$:

2	4	3	7	9	8	6	3	4	9
0	1	2	3	4	5	6	7	8	9

$$L = 0$$

$$\text{sum} = 6$$

$$R = 3$$

$$L = 3$$

$$\text{sum} = \underline{\underline{14}}$$

$$R = 7$$

Q2. Special Index

Given an array of size N, count the no. of special indexes

Note → Special indexes are those after removing which sum of all even indexed ele = sum of all odd indexed ele.

$$A[] = \{ 4 \underset{0}{}, 3 \underset{1}{}, 2 \underset{2}{}, ? \underset{3}{}, 6 \underset{4}{}, -2 \underset{5}{} \}$$

i	A[i]	Sc	S _o	
0	{ 3 2 7 6 -2 } 0 1 2 3 4	8	8	Yes
1	{ 4 2 7 6 -2 } 0 1 2 3 4	9	8	No
2	{ 4 3 7 6 -2 } 0 1 2 3 4	9	9	Yes
3	{ 4 3 2 6 -2 } 0 1 2 3 4	4	9	No
4	{ 4 3 2 7 -2 } 0 1 2 3 4	4	10	No
5	{ 4 3 2 7 6 } 0 1 2 3 4	12	10	No
Count = 2				

Quiz 1 $A[] = \{ 4, 1, \cancel{3}, 7, 3, 10 \}$

\downarrow remove 2nd idx

$A[] = \{ 4, 1, 7, 10 \}$

$S_0 = 1 + 10 = 11$

Quiz 2 $A[] = \{ 2, 3, 1, \cancel{4}, 0, -1, 2, -2, 10, 8 \}$

\downarrow remove 3rd index

$A[] = \{ 2, 3, 1, 0, -1, 2, -2, 10, 8 \}$

$S_0 = 3 + 0 + 2 + 10 = \underline{\underline{15}}$

Quiz 3 $A[] = \{ 2, 3, 1, \cancel{4}, 0, -1, 2, -2, 10, 8 \}$

\downarrow remove 3rd index

$A[] = \{ 2, 3, 1, 0, -1, 2, -2, 10, 8 \}$

$S_0 = 2 + 1 + (-1) + (-2) + 8 \Rightarrow \underline{\underline{8}}$

Sum of even indexed ele after removal of 3rd idx =

sum of even + sum of odd idx
idx ele from . . . ele from 4 to 9
0 to 2

$$A[] = \{ 2, 3, 1, 4, 0, -1, 2, -2, 10, 8 \}$$

remove 5th index

$$A[] = \{ 2, 3, 1, 4, 0, 2, -2, 10, 8 \}$$

$S_0 =$ sum of odd idx + sum of even idx
ele from 0 to 4 ele from 6 to 9

Remove i^{th} idx

$S_0 = ?$

$S_e = ?$

$$A[] = \{ \begin{matrix} 3 & 4 \\ 0 & 1 \end{matrix} \quad \cancel{\begin{matrix} 5 \\ 2 \end{matrix}} \quad 6 & 7 & 2 \}$$

$$MA[] = \{ \begin{matrix} 3 & 4 & 6 & 7 & 2 \end{matrix} \} \quad \cancel{\begin{matrix} 5 \\ 2 \end{matrix}}$$

$$S_0 = 4 + \cancel{7} \quad \text{in } \underline{MA}$$

$$\underline{\underline{1}} \quad \underline{\underline{3}}$$

$S_0 = \text{sum of all odd index ele in } A + \text{sum of all even index ele in } A \text{ from } i+1 \text{ to } n-1$

from 0 to $i-1$

$\text{podd}[i-1]$

+

$\text{peven}[i+1 \text{ to } n-1]$

$L \quad R$

$\text{peven}[R] - \text{peven}[L-1]$

$= \text{peven}[n-1] - \text{peven}[i+1-1]$

$S_0 = \text{podd}[i-1] + \text{peven}[n-1] - \text{peven}[i]$

```
int countspecial (int A[])
```

```
int peven[n] } TODO  
int podd [n]
```

```
count = 0
```

```
for (i=0; i<n; i++) {
```

// removing i^{th} index

TC: O(n)

SC: O(n)

$$S_0 = p_{\text{even}}[n-1] - p_{\text{even}}[i]$$

```
if (i > 0) {  $S_0 = S_0 + p_{\text{odd}}[i-1]$  }
```

$$S_e = p_{\text{odd}}[n-1] - p_{\text{odd}}[i]$$

```
if (i > 0) {  $S_e = S_e + p_{\text{even}}[i-1]$  }
```

```
if ( $S_0 == S_e$ ) count++;
```

}

```
return count
```

3