

CSCS 311

# Data Communications and Networking

## Lecture 15

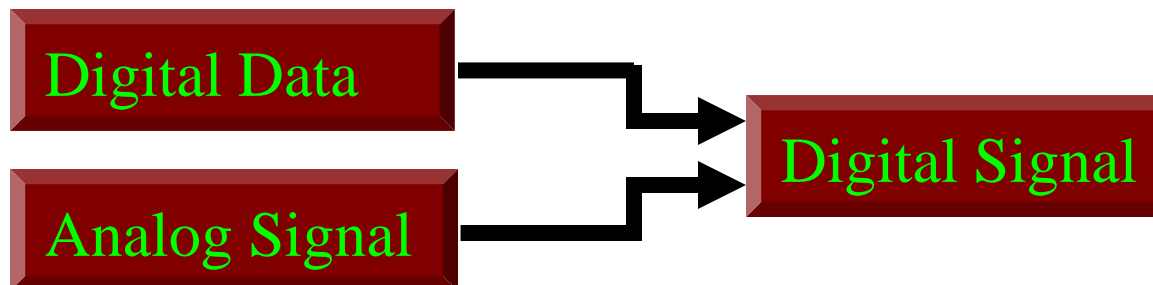
### *Lecture Focus:*

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⦿ Digital Transmission

# Digital Transmission

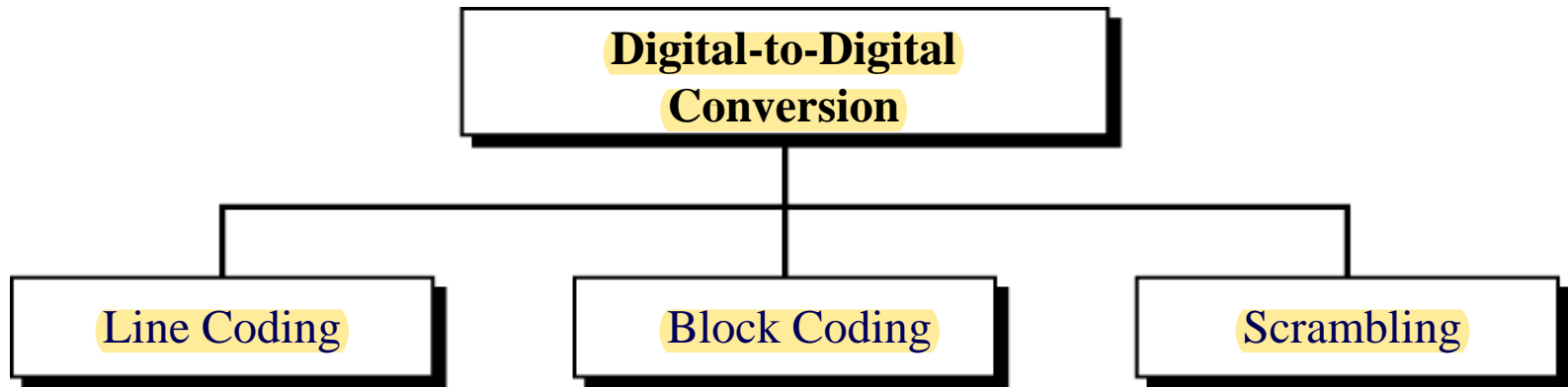
- ⦿ A computer network is designed to send information from one point to another.
- ⦿ This information needs to be converted to either a digital signal or an analog signal for transmission.
- ⦿ Here, we study the schemes and techniques used to transmit data digitally.
  - ⦿ First, we discuss **digital-to-digital** conversion techniques, methods which convert digital data to digital signals.
  - ⦿ Second, we discuss **analog-to-digital** conversion techniques, methods which change an analog signal to a digital signal.



## DIGITAL-TO-DIGITAL CONVERSION

- ⊙ Data can be either digital or analog. Signals that represent data can also be digital or analog.
- ⊙ Here, we study how we can represent digital data by using digital signals.
- ⊙ The conversion involves three techniques:
  - ⊙ Line coding
  - ⊙ Block coding
  - ⊙ Scrambling
- ⊙ Line coding is always needed.
- ⊙ Block coding and scrambling may or may not be needed.

## Digital-to-Digital Conversion



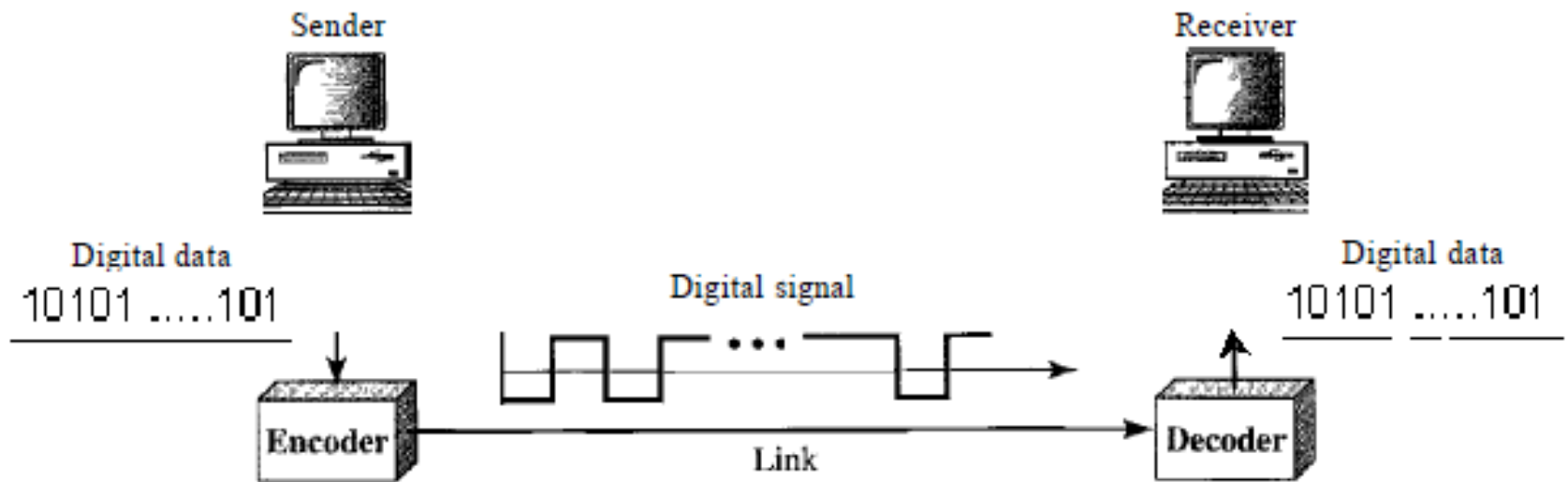
# Digital-to-Digital Conversion

## Line Coding

- ⊙ Line coding is the process of converting digital data to digital signals.
- ⊙ We assume that data, in the form of text, numbers, graphical images, audio, or video, are stored in computer memory as sequences of bits.
- ⊙ Line coding converts a sequence of bits to a digital signal.
  - ⊙ At the sender, digital data are encoded into a digital signal;
  - ⊙ At the receiver, the digital data are recreated by decoding the digital signal.

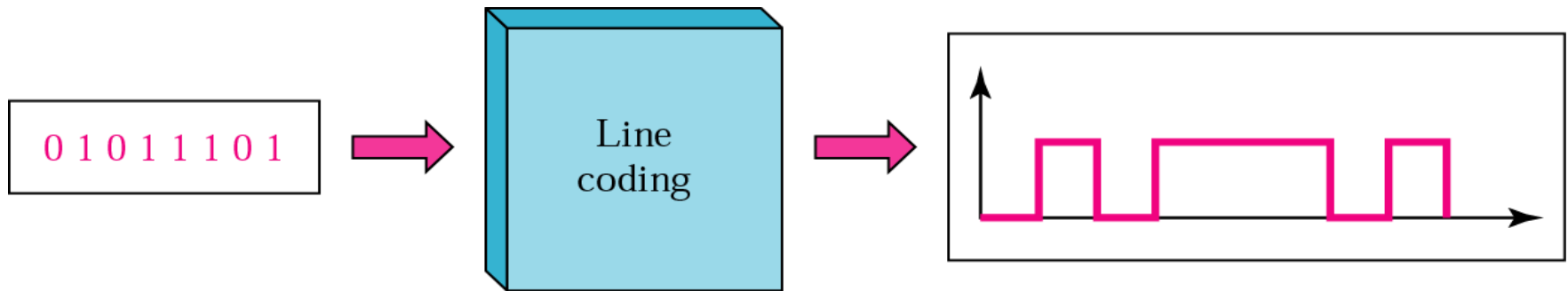
# Digital-to-Digital Conversion

## Line Coding Process



# Digital-to-Digital Conversion

## Line Coding Process



Line coding and decoding

# Characteristics of Line Coding Schemes

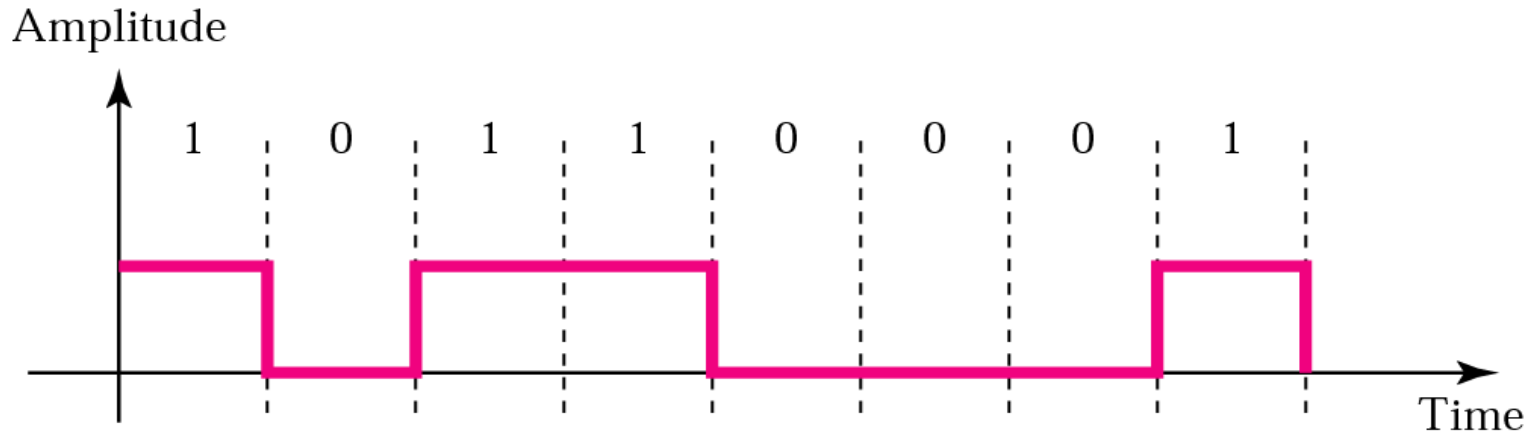
## Signal Element *Versus* Data Element

- ⊙ A data element is the smallest entity that can represent a piece of information: this is the bit.
  - ⊙ In digital data communications, a signal element carries data elements.
  - ⊙ A signal element is the shortest unit (time wise) of a digital signal.
  - ⊙ In other words, data elements are what we need to send; signal elements are what we can send.
  - ⊙ Data elements are being carried; signal elements are the carriers.
- 
- ⊙ We define a ratio  $r$  which is the number of data elements carried by each signal element.
  - ⊙ Figure shows several situations with different values of  $r$ .

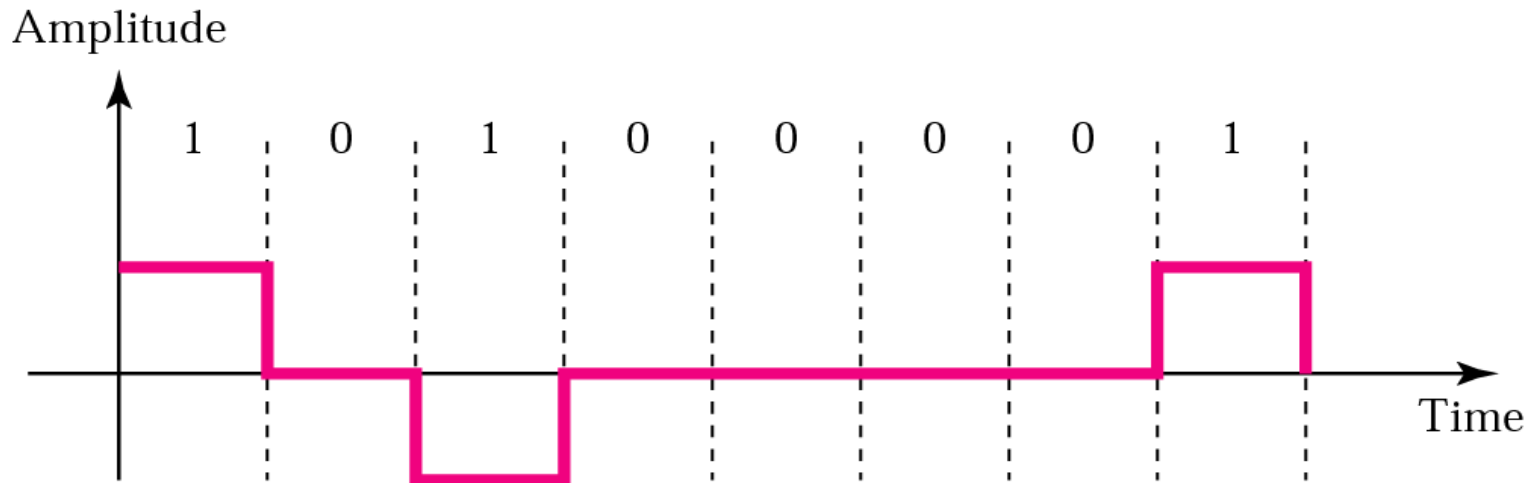


# Characteristics of Line Coding Schemes

## Signal Element Versus Data Element



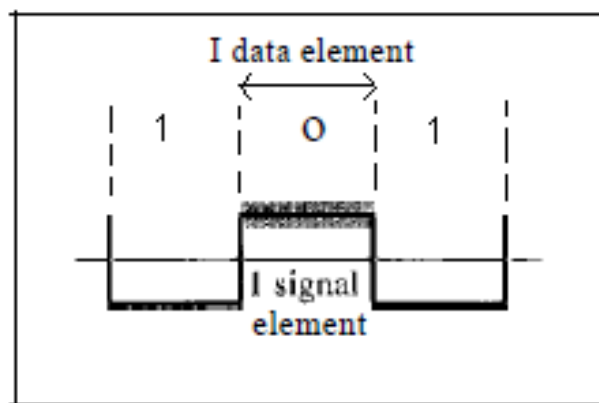
a. Two signal levels, two data levels



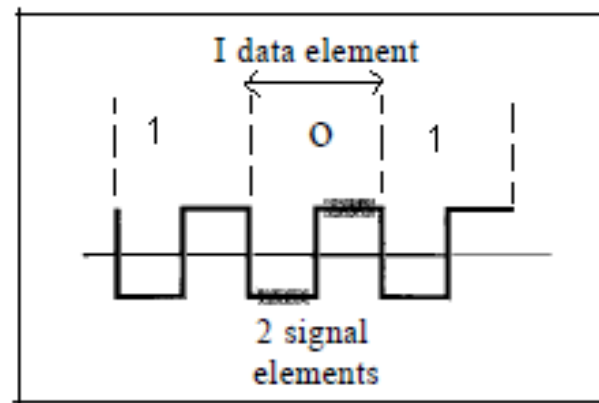
b. Three signal levels, three data levels

# Characteristics of Line Coding Schemes

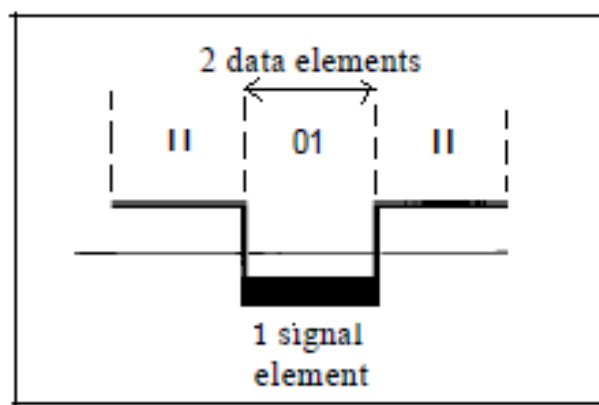
## Signal Element Versus Data Element



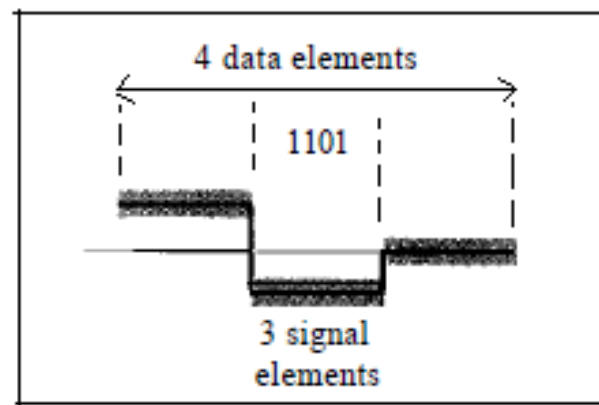
a. One data element per one signal element ( $r=1$ )



b. One data element per two signal elements ( $r = \frac{1}{2}$ )



c. Two data elements per one signal element ( $r=2$ )



d. Four data elements per three signal elements ( $r = \frac{4}{3}$ )

We need two signal elements (two transitions) to carry each data element.

# Characteristics of Line Coding Schemes

## Data Rate Versus Signal Rate

- ⊙ **Data rate:**
  - ⊙ The number of data elements (bits) sent in 1s.
  - ⊙ The unit is bits per second (bps).
- ⊙ **Signal rate:**
  - ⊙ The number of signal elements sent in 1s.
  - ⊙ The unit is the baud.
- ⊙ The data rate is sometimes called the bit rate.
- ⊙ The signal rate is sometimes called the pulse rate, the modulation rate, or the baud rate.

# Characteristics of Line Coding Schemes

## Data Rate Versus Signal Rate

- ⊙ One goal in data communications is to:
  - ⊙ Increase the data rate, and
  - ⊙ Decrease the signal rate
- ⊙ Increasing the data rate increases the speed of transmission; decreasing the signal rate decreases the bandwidth requirement.
- ⊙ We need to carry more people in fewer vehicles to prevent traffic jams. We have a limited bandwidth in our transportation system.

# Characteristics of Line Coding Schemes

## Relationship between Data Rate and Signal Rate

- ⊙ This relationship depends on the value of  $r$ .
- ⊙ It also depends on the data pattern.
  - ⊙ If a data pattern consists of all 1s or all 0s, the signal rate may be different from a data pattern of alternating 0s and 1s.
- ⊙ To derive a formula for the relationship, we need to define three cases: the **worst**, the **best**, and the **average**.
  - ⊙ The worst case is when we need the maximum signal rate.
  - ⊙ The best case is when we need the minimum signal rate.
- ⊙ We are usually interested in the average case.

# Characteristics of Line Coding Schemes

## Relationship between Data Rate and Signal Rate

- ⊙ We can formulate the relationship between data rate and signal rate as:

$$S = c \times N \times \frac{1}{r} \quad \text{baud}$$

Where:

N is the data rate (bps);

c is the case factor, which varies for each case;

S is the number of signal elements; and

r is the previously defined factor.

# Characteristics of Line Coding Schemes

## Relationship between Data Rate and Signal Rate

### Example

A signal is carrying data in which one data element is encoded as one signal element ( $r = 1$ ). If the bit rate is 100 kbps, what is the average value of the baud rate if  $c$  is between 0 and 1?

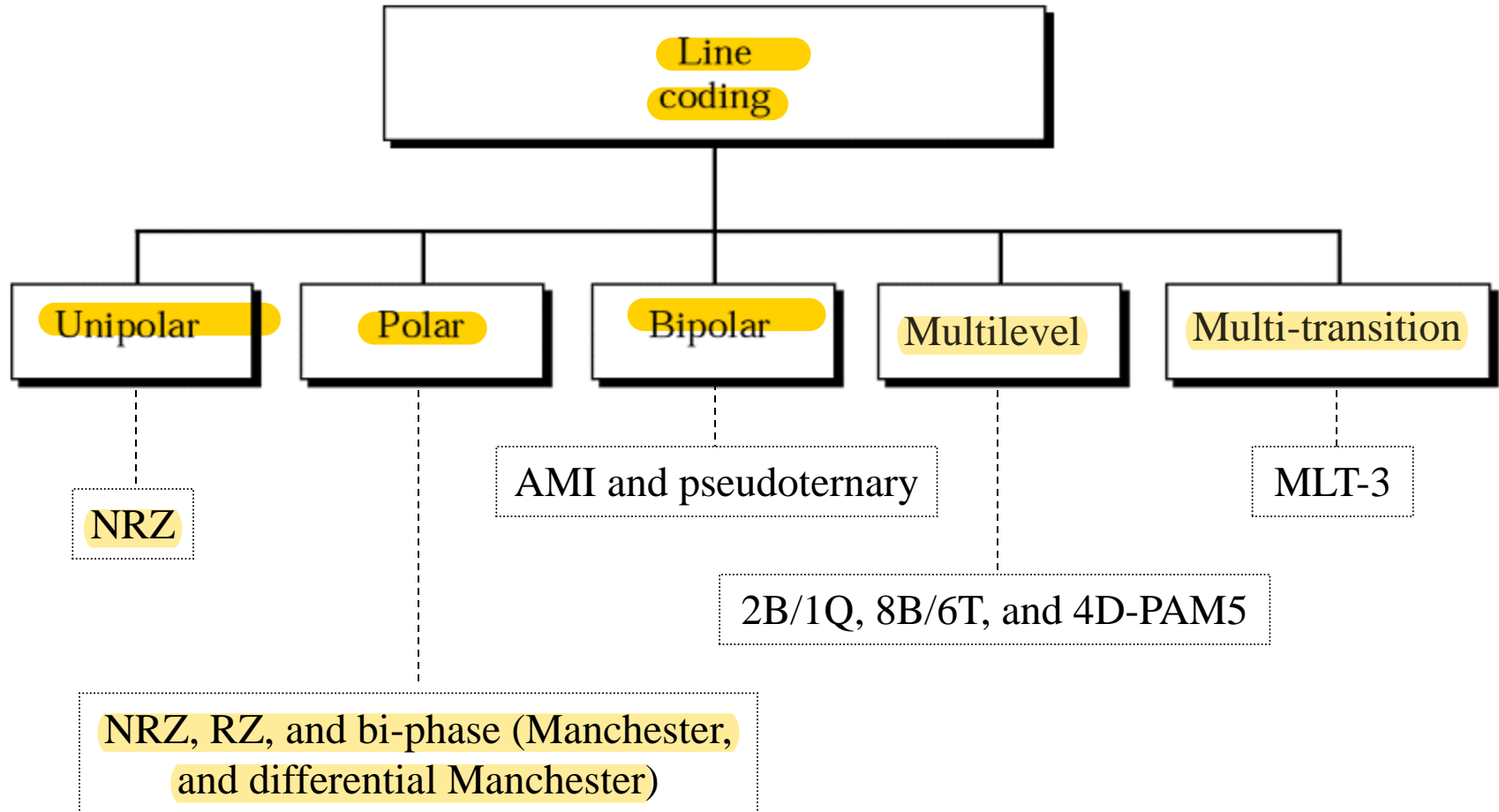
### Solution

We assume that the average value of  $c$  is  $1/2$ . The baud rate is then:

$$\begin{aligned} S &= c \times N \times 1/r \\ &= 1/2 \times 100,000 \times 1/1 \\ &= 50,000 \text{ bauds} \\ &= 50 \text{ kbaud} \end{aligned}$$

# Line Coding Schemes

## Five categories



There are several schemes in each category.



# Line Coding Schemes

## UNIPOLAR

In a unipolar scheme, all the signal levels are on one side of the time axis, either above or below.

Unipolar encoding uses only one voltage level.

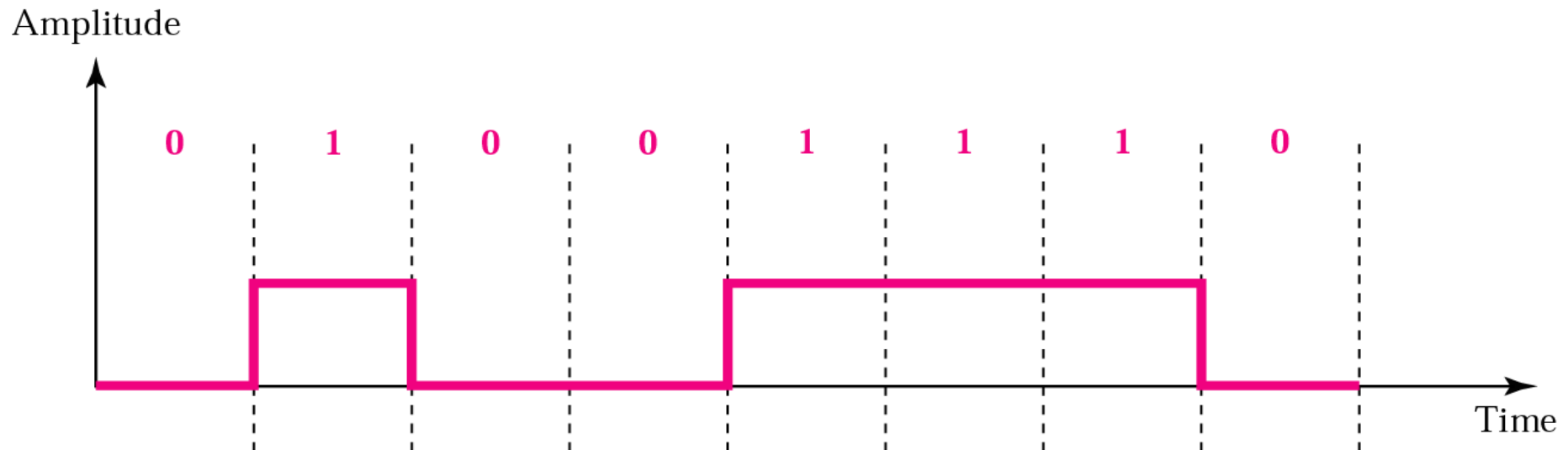
UNI = Single

# Line Coding Schemes

## UNIPOLAR

### NRZ (Non-Return-to-Zero)

- ⊙ Unipolar scheme was designed as a non-return-to-zero (NRZ) scheme in which the positive voltage defines bit 1 and the zero voltage defines bit 0.
- ⊙ It is called NRZ because the signal does not return to zero at the middle of the bit.



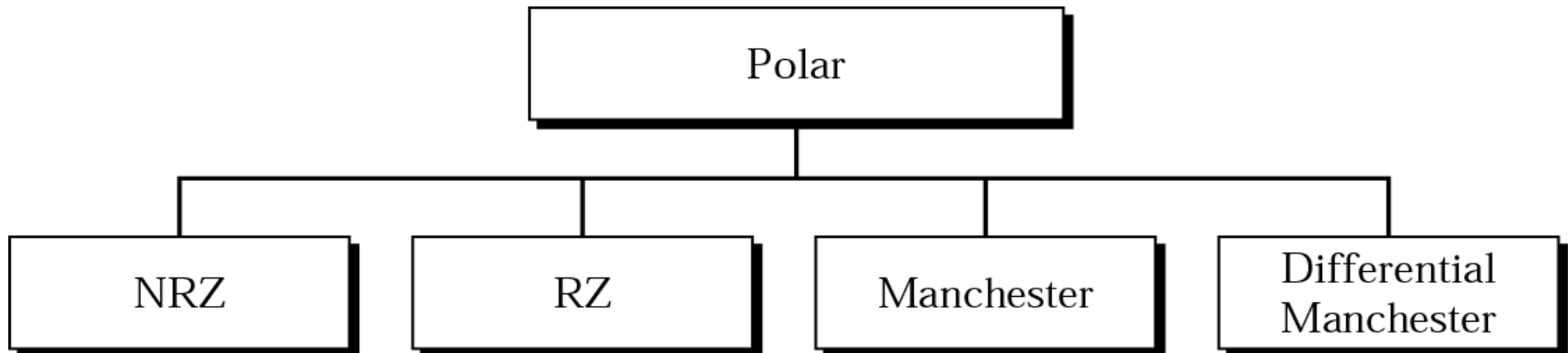
# Line Coding Schemes

## POLAR SCHEMES

- ⊙ In polar schemes, the voltages are on the both sides of the time axis.
- ⊙ For example, the voltage level for 0 can be positive and the voltage level for 1 can be negative.

Polar encoding uses two voltage levels (positive and negative).

### Types of polar encoding

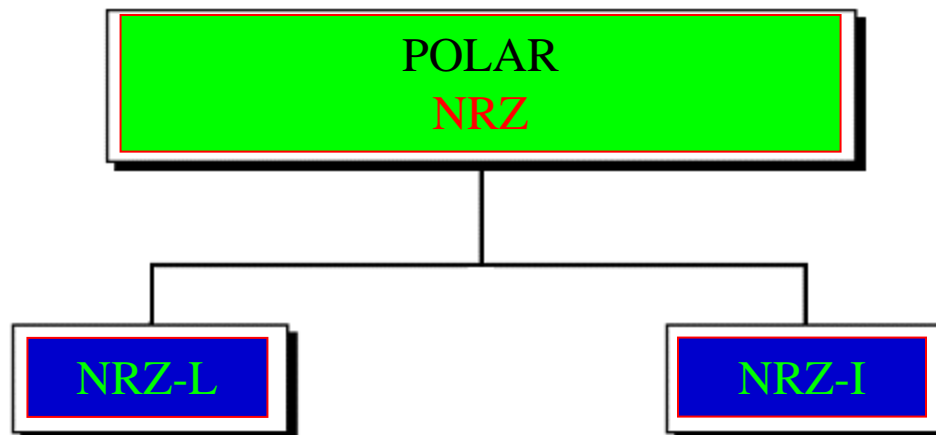


# Line Coding Schemes

## POLAR SCHEMES

### NRZ (Non-Return-to-Zero)

- ⊙ In polar NRZ encoding, we use two levels of voltage amplitude.
- ⊙ We have two versions of polar NRZ:
  - ⊙ NRZ-L
  - ⊙ NRZ-I



# Line Coding Schemes

## POLAR SCHEMES

### NRZ (Non-Return-to-Zero)

- ⊙ In NRZ-L (NRZ-Level), the level of the voltage determines the value of the bit.

In NRZ-L the level of the signal is dependent upon the state of the bit.

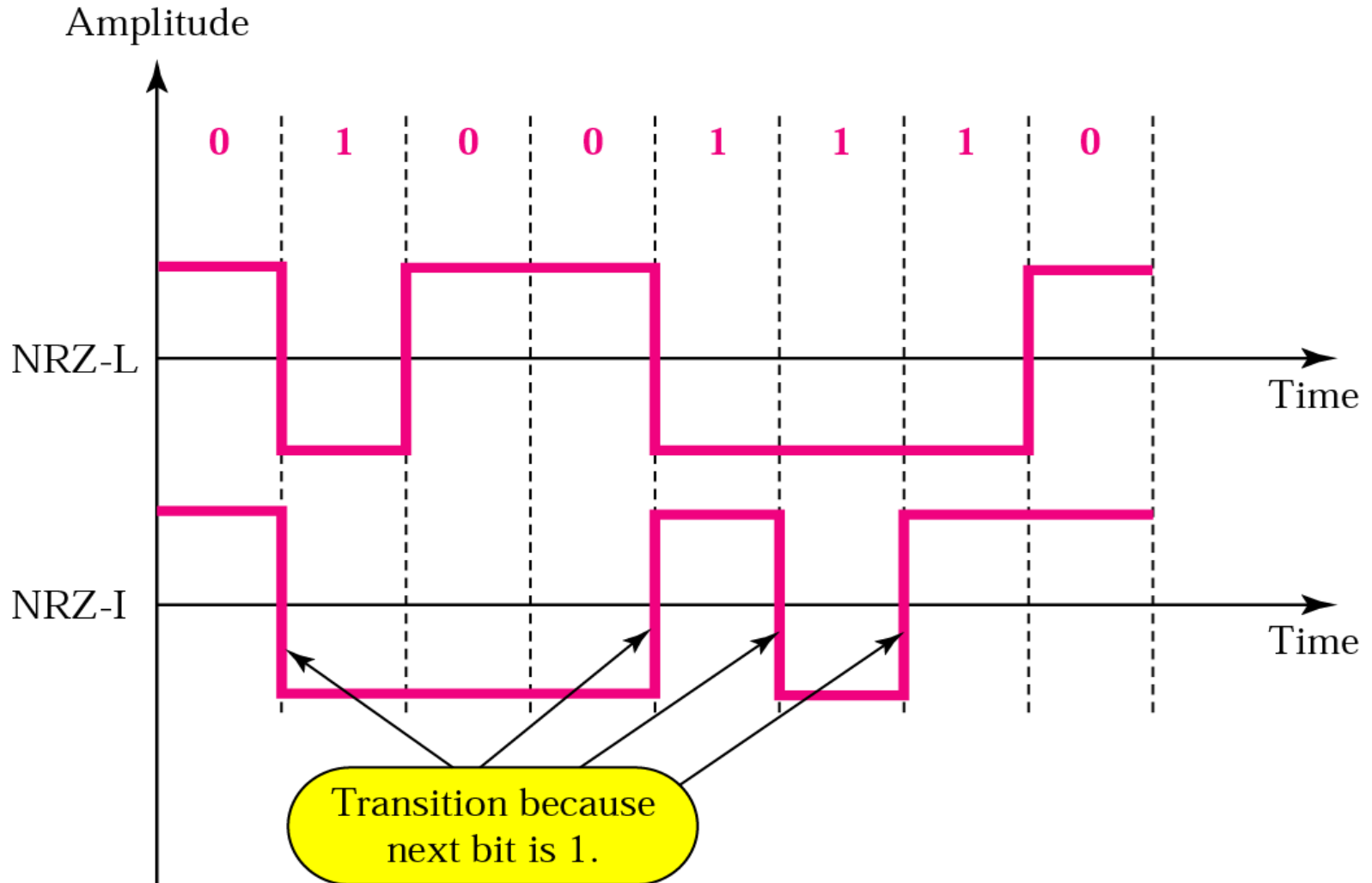
- ⊙ In NRZ-I (NRZ-Invert), the change or lack of change in the level of the voltage determines the value of the bit.
  - ⊙ If there is no change, the bit is 0.
  - ⊙ If there is a change, the bit is 1.

In NRZ-I the signal is inverted if a 1 is encountered.

# Line Coding Schemes

## POLAR SCHEMES

### NRZ (Non-Return-to-Zero)



# Line Coding Schemes

## POLAR SCHEMES

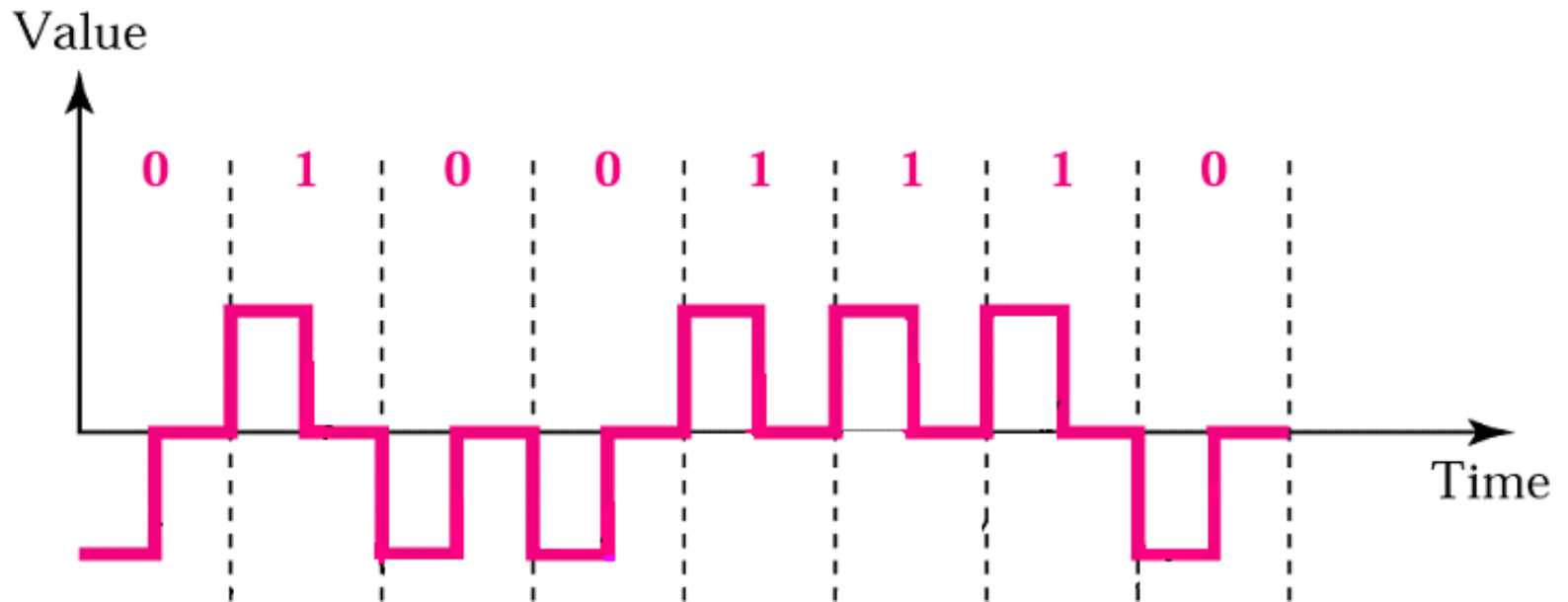
### RZ (Return-to-Zero)

- ⊙ This scheme uses three values: positive, negative, and zero.
- ⊙ In RZ, the signal changes not between bits but during the bit.
- ⊙ Signal goes to 0 in the middle of each bit. It remains there until the beginning of the next bit.
- ⊙ The main disadvantage of RZ encoding is that it requires two signal changes to encode a bit and therefore occupies greater bandwidth.
- ⊙ Another problem is the complexity: RZ uses three levels of voltage, which is more complex to create and discern.
- ⊙ As a result of all these deficiencies, the scheme is not used today.
  - ⊙ It has been replaced by the better-performing Manchester and differential Manchester schemes.

# Line Coding Schemes

## POLAR SCHEMES

### RZ (Return-to-Zero)



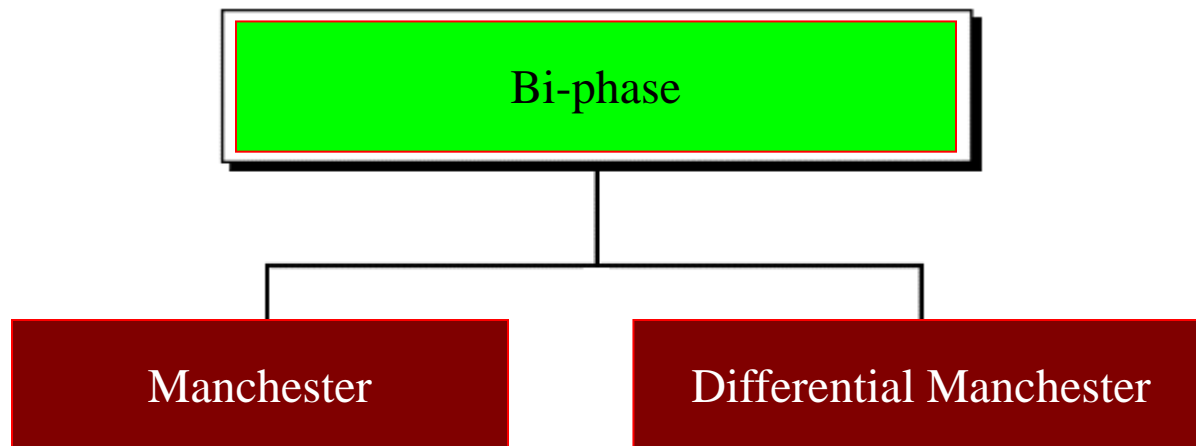


# Line Coding Schemes

## POLAR SCHEMES

### Bi-phase

- ⊙ Manchester
- ⊙ Differential Manchester



# Line Coding Schemes

## POLAR SCHEMES

### Bi-phase: Manchester

- ⊙ The idea of RZ (transition at the middle of the bit) and the idea of NRZ-L are combined into the Manchester scheme.
- ⊙ In Manchester encoding, the duration of the bit is divided into two halves.
  - ⊙ The voltage remains at one level during the first half and moves to the other level in the second half.

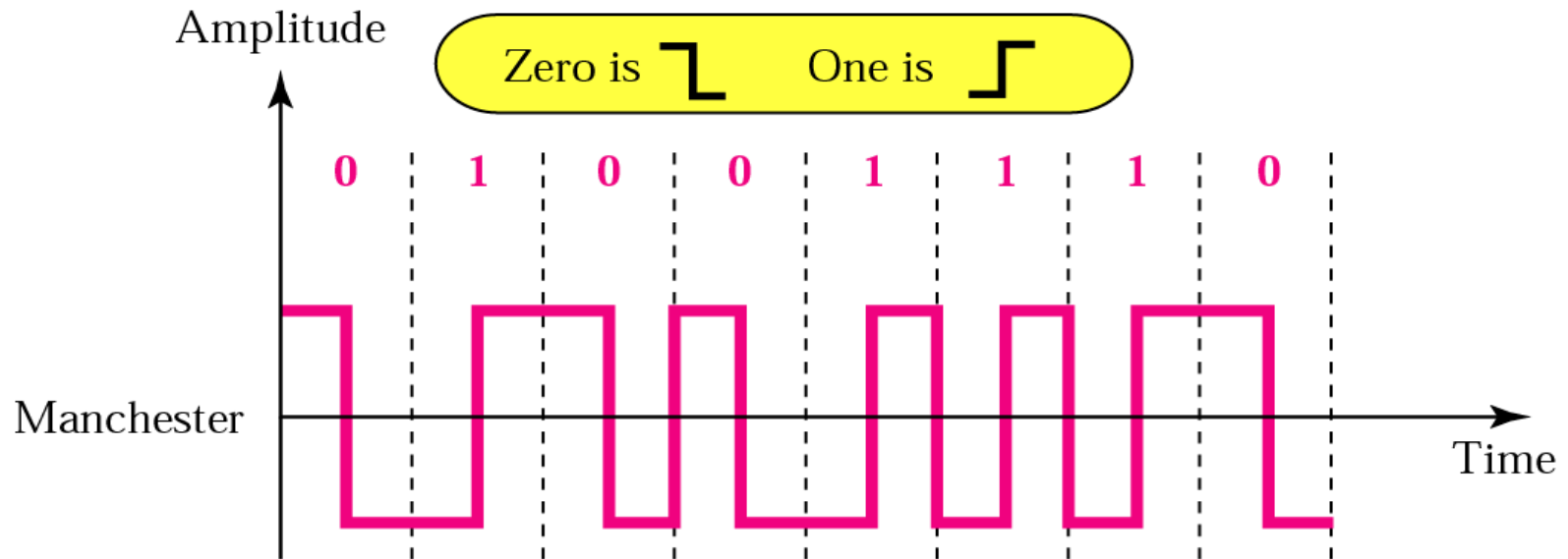
### Bi-phase: Differential Manchester

- ⊙ This scheme combines the ideas of RZ and NRZ-I.
- ⊙ There is always a transition at the middle of the bit, but the bit values are determined at the beginning of the bit.
  - ⊙ If the next bit is 0, there is a transition.
  - ⊙ If the next bit is 1, there is none.

# Line Coding Schemes

## POLAR SCHEMES

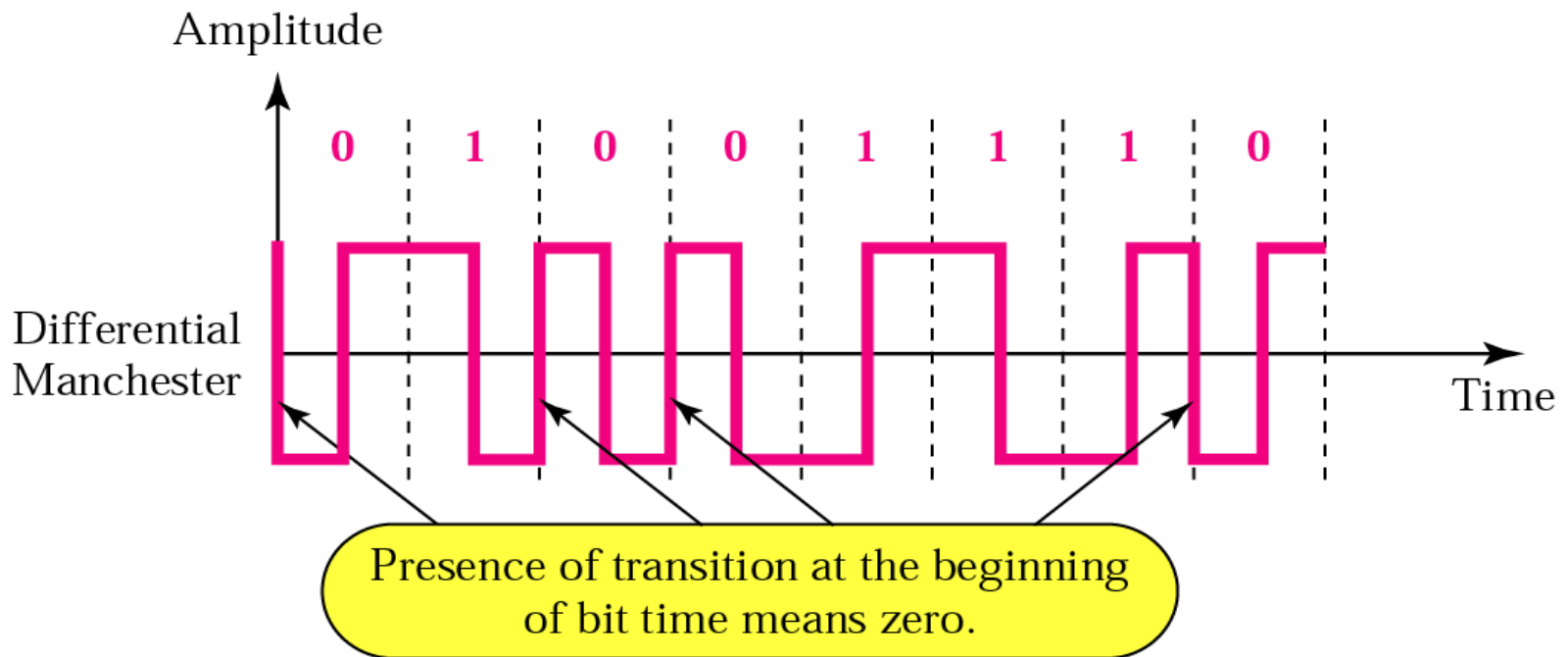
### Manchester encoding



# Line Coding Schemes

## POLAR SCHEMES

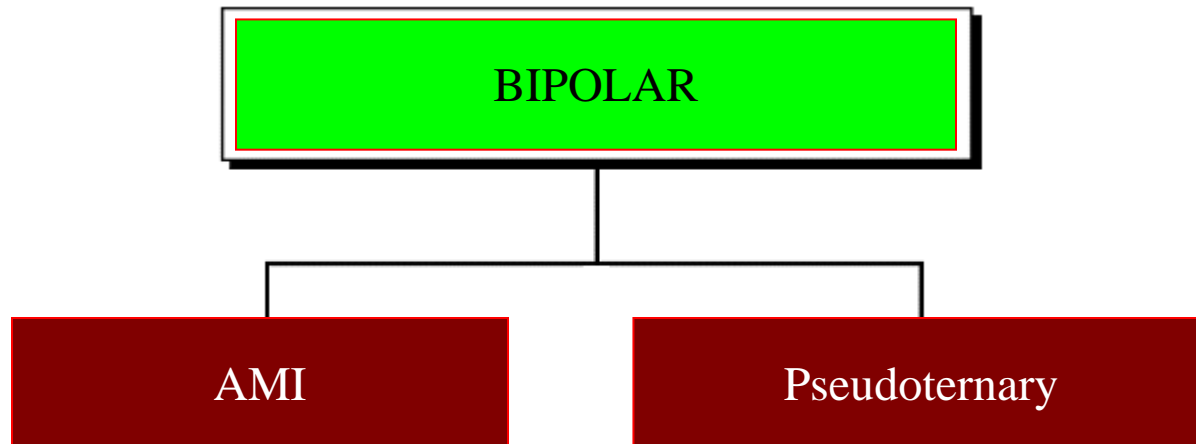
### Differential Manchester encoding



# Line Coding Schemes

## BIPOLAR SCHEMES

- ⊙ In bipolar encoding (sometimes called multilevel binary), there are three voltage levels: positive, negative, and zero.
- ⊙ The voltage level for one data element is at zero.
- ⊙ The voltage level for the other element alternates between positive and negative.



# Line Coding Schemes

## BIPOLAR SCHEMES

### AMI: Alternate Mark Inversion

- ⊙ In the term alternate mark inversion, the word mark comes from telegraphy and means 1. So AMI means alternate 1 inversion.
- ⊙ A neutral zero voltage represents binary 0.
- ⊙ Binary 1s are represented by alternating positive and negative voltages.

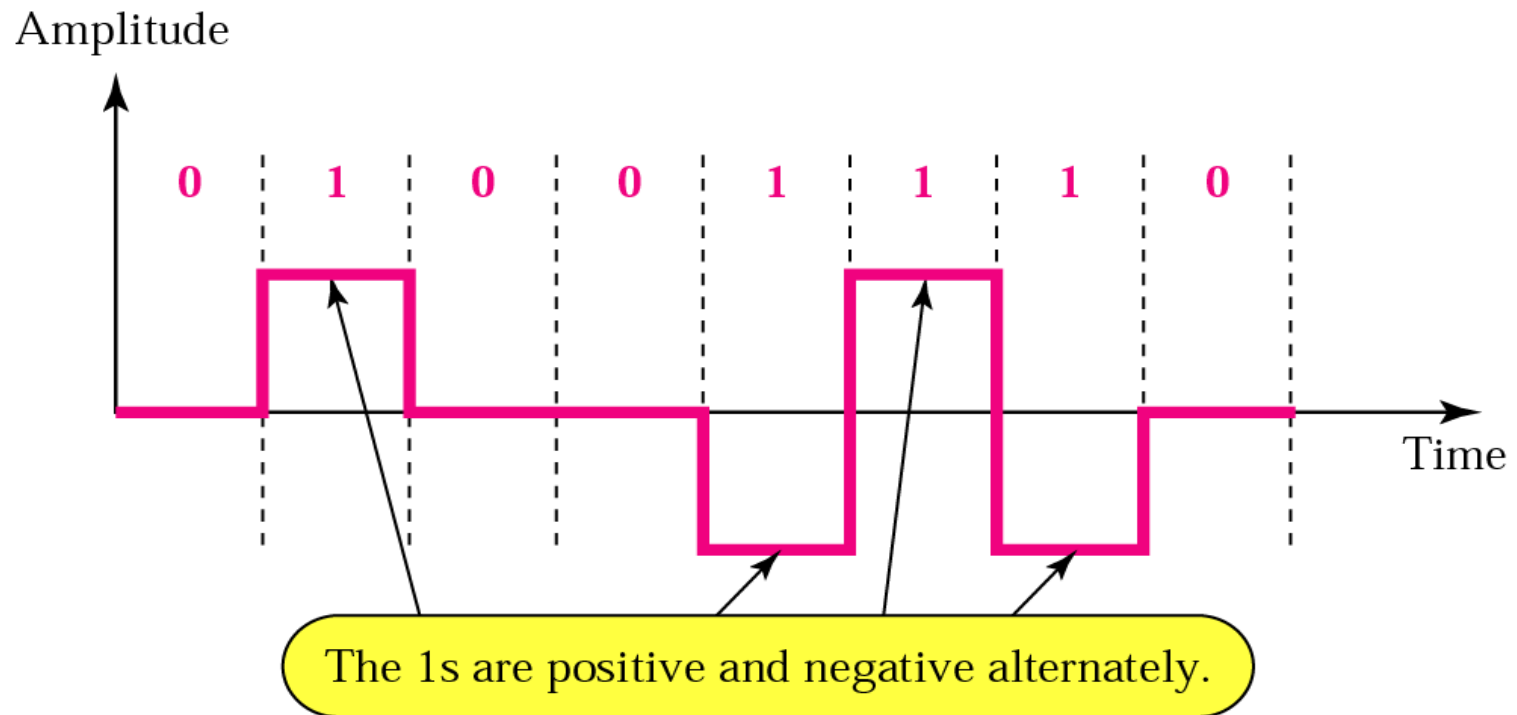
### Pseudoternary

- ⊙ A variation of AMI encoding is called pseudoternary in which the 1 bit is encoded as a zero voltage and the 0 bit is encoded as alternating positive and negative voltages.

# Line Coding Schemes

## BIPOLAR SCHEMES

**AMI:** Alternate Mark Inversion

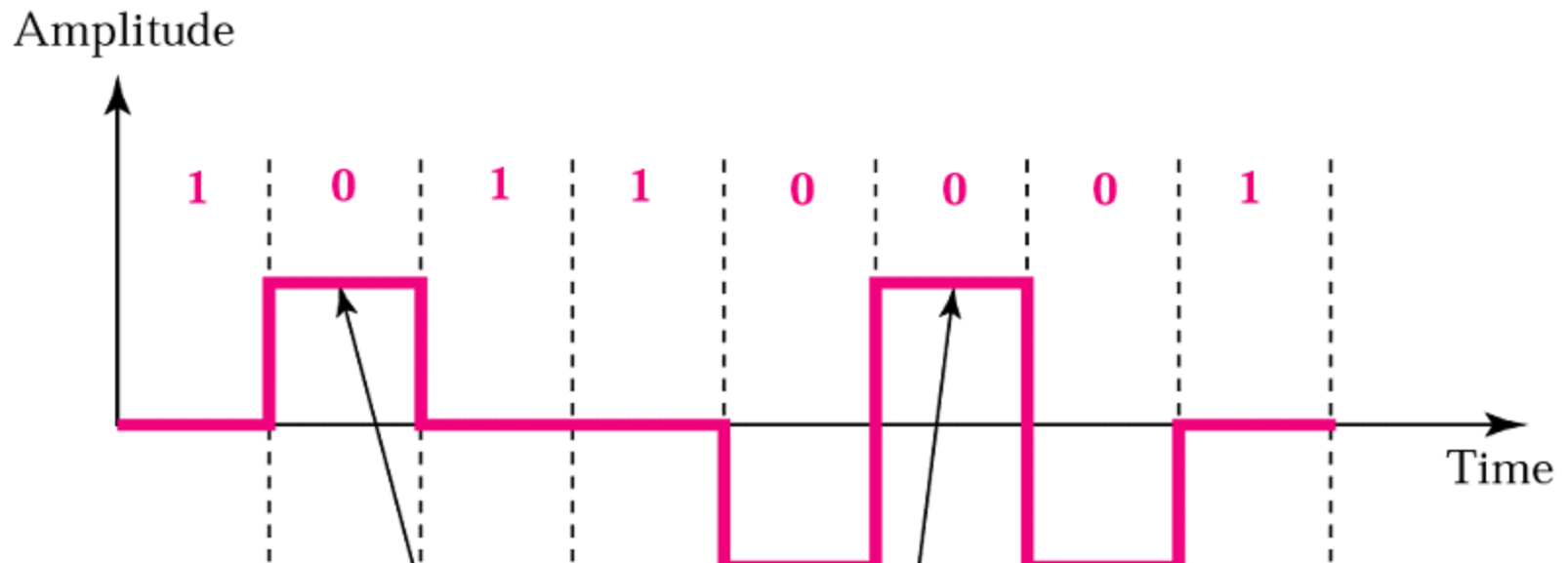


Bipolar AMI encoding

# Line Coding Schemes

## BIPOLAR SCHEMES

### Pseudoternary



The 0s are positive and negative alternately.

Bipolar Pseudoternary encoding



# Line Coding Schemes

## MULTILEVEL SCHEMES

- ⊙ The desire to increase the data speed or decrease the required bandwidth has resulted in the creation of many schemes.
    - ⊙ The goal is to increase the number of bits per baud by encoding a pattern of  $m$  data elements into a pattern of  $n$  signal elements.
  - ⊙ We only have two types of data elements (0s and 1s).
    - ⊙ It means that a group of  $m$  data elements can produce a combination of  $2^m$  data patterns.
  - ⊙ We can have different types of signal elements by allowing different signal levels.
    - ⊙ If we have  $L$  different levels, we can produce  $L^n$  combinations of signal patterns.
- 
- ⊙ If  $2^m = L^n$ , then each data pattern is encoded into one signal pattern.
  - ⊙ If  $2^m < L^n$ , data patterns occupy only a subset of signal patterns.
  - ⊙ Data encoding is not possible if  $2^m > L^n$  because some of the data patterns cannot be encoded.

# Line Coding Schemes

## MULTILEVEL SCHEMES

- ⊙ The code designers have classified these types of coding as **mBnL**, where:
  - ⊙ m is the length of the binary pattern,
  - ⊙ B means binary data,
  - ⊙ n is the length of the signal pattern, and
  - ⊙ L is the number of levels in the signaling.
- ⊙ A letter is often used in place of L:
  - ⊙ B (binary) for L = 2,
  - ⊙ T (ternary) for L = 3, and
  - ⊙ Q (quaternary) for L = 4.
- ⊙ Note that:
  - ⊙ The first two letters define the data pattern, and
  - ⊙ The second two define the signal pattern.

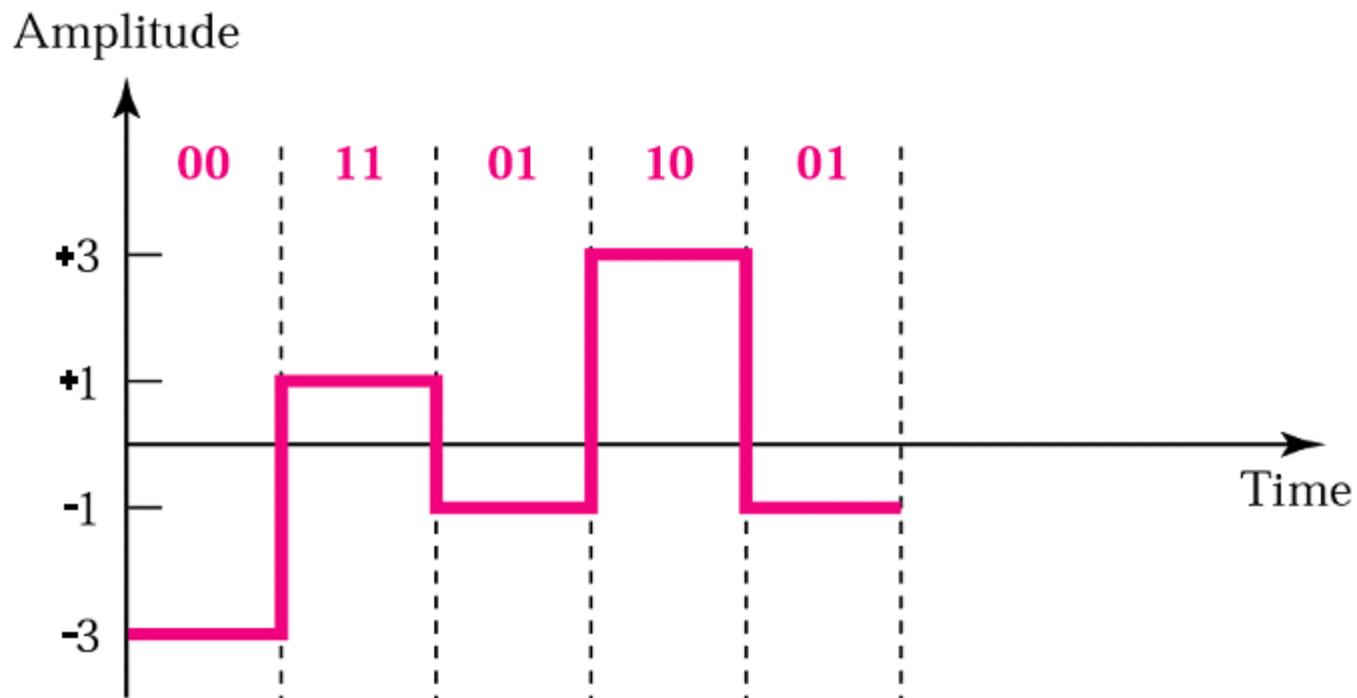
In mBnL schemes, a pattern of **m data elements** is encoded as a pattern of **n signal elements** in which  $2^m \leq L^n$ .

# Line Coding Schemes

## MULTILEVEL SCHEMES

### 2B1Q

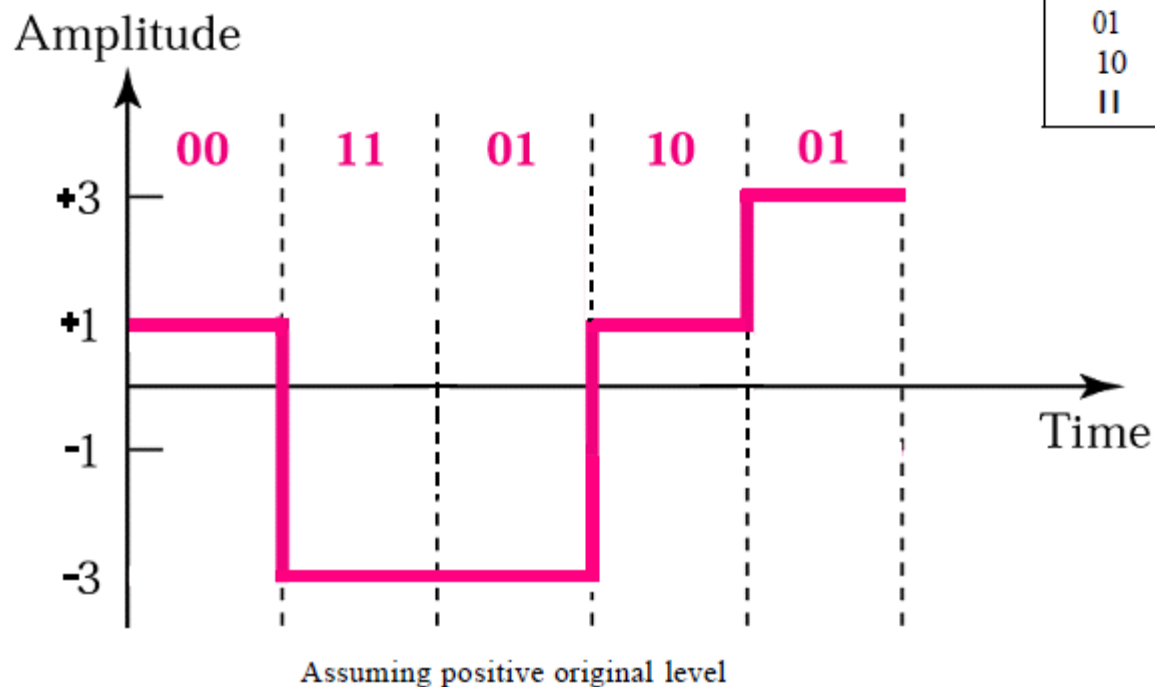
- ⊙ The first mBnL scheme, two binary, one quaternary (2B1Q), uses data patterns of size 2 and encodes the 2-bit patterns as one signal element belonging to a four-level signal.
- ⊙ In this type of encoding  $m = 2$ ,  $n = 1$ , and  $L = 4$  (quaternary).
- ⊙ Figure below shows an example of a 2B1Q signal.



# Line Coding Schemes

## MULTILEVEL SCHEMES

## 2B1Q



| Next bits | Previous level:<br>positive | Previous level:<br>negative |
|-----------|-----------------------------|-----------------------------|
|           | Next level                  | Next level                  |
| 00        | +1                          | -1                          |
| 01        | +3                          | -3                          |
| 10        | -1                          | +1                          |
| 11        | -3                          | +3                          |

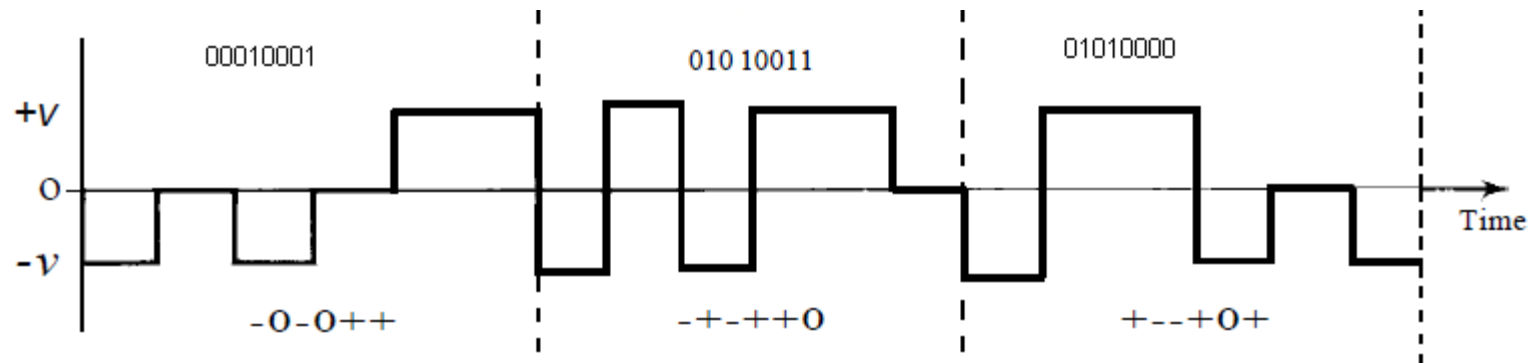
Transition table

# Line Coding Schemes

## MULTILEVEL SCHEMES

### 8B6T : Eight binary, six ternary

- ⊙ The idea is to encode a pattern of 8 bits as a pattern of 6 signal elements, where the signal has three levels (ternary).
- ⊙ In this type of scheme, we can have  $2^8 = 256$  different data patterns and  $3^6 = 478$  different signal patterns.
- ⊙ There are  $478 - 256 = 222$  redundant signal elements that provide synchronization and error detection.



# Line Coding Schemes

## Multi-Transition Schemes

### Multiline Transmission: MLT-3

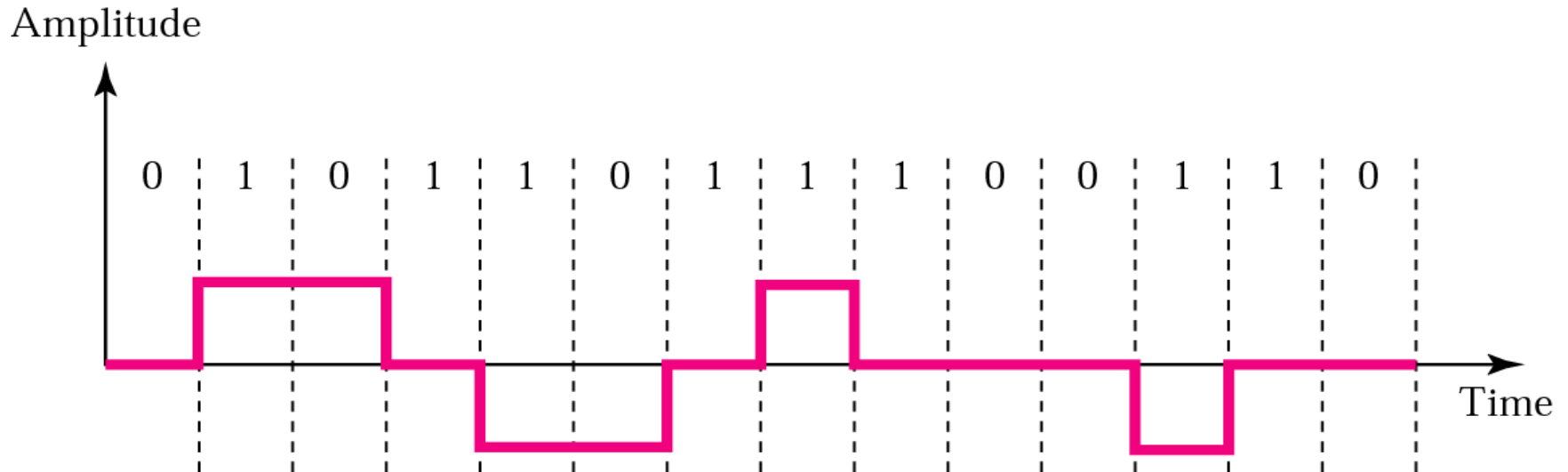
- ⊙ NRZ-I and differential Manchester are classified as differential encoding but use two transition rules to encode binary data (no inversion, inversion).
- ⊙ If we have a signal with more than two levels, we can design a differential encoding scheme with more than two transition rules.
- ⊙ Multiline transmission, three level (MLT-3) scheme uses three levels (+V, 0, and -V) and three transition rules to move between the levels.
  1. If the next bit is 0, there is no transition.
  2. If the next bit is 1 and the current level is not 0, the next level is 0.
  3. If the next bit is 1 and the current level is 0, the next level is the opposite of the last nonzero level.

# Line Coding Schemes

## Multi-Transition Schemes

### Multiline Transmission: MLT-3

1. If the next bit is 0, there is no transition.
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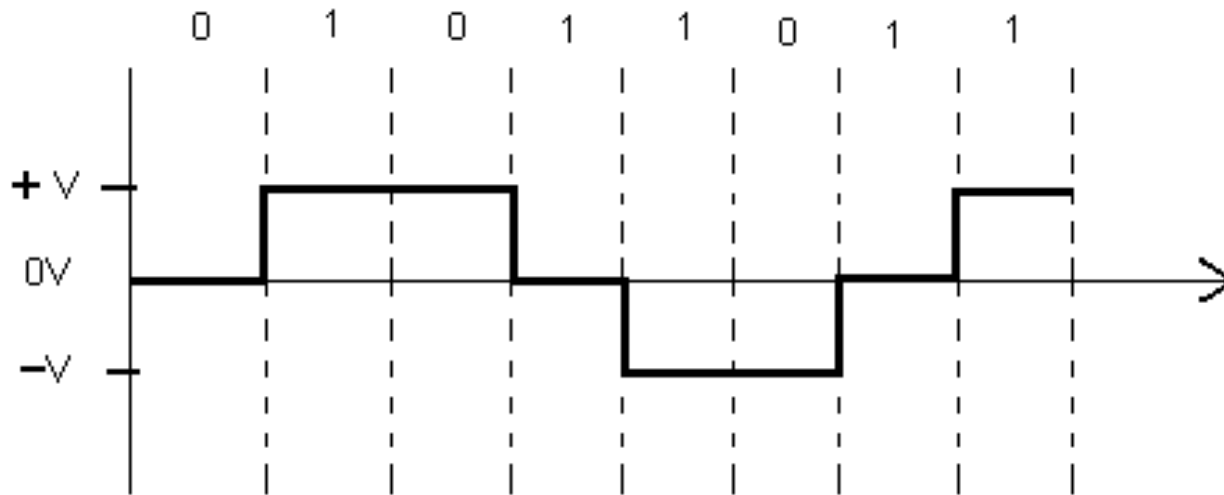


# Line Coding Schemes

## Multi-Transition Schemes

### Multiline Transmission: MLT-3

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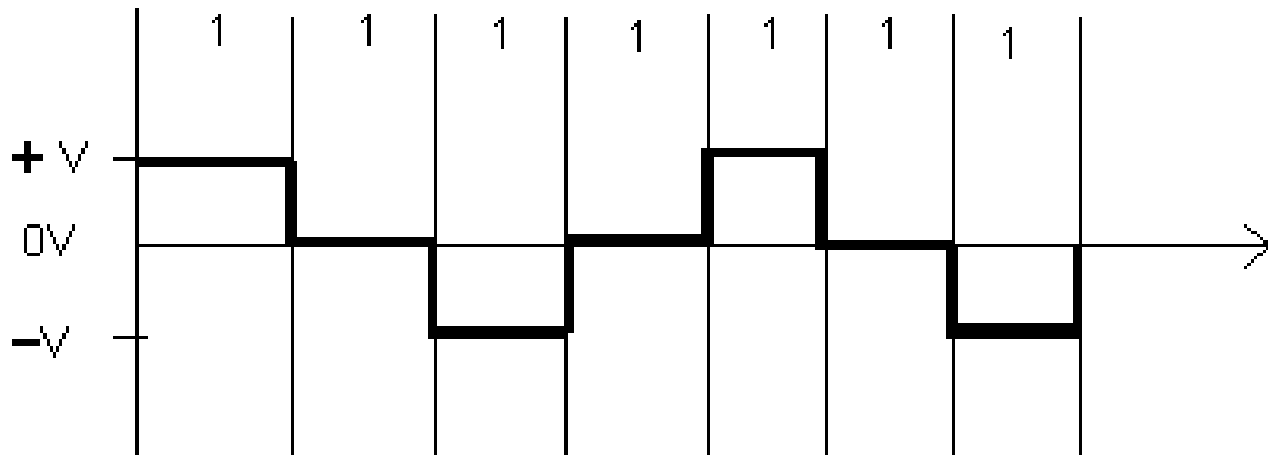


# Line Coding Schemes

## Multi-Transition Schemes

### Multiline Transmission: MLT-3

1. If the next bit is 0, there is no transition.
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3. If the next bit is 1 and the current level is 0, the next level is the opposite of the last nonzero level.



WORST CASE

# Line Coding Schemes

Draw the graph of the NRZ-L, NRZ-I, Manchester, Differential Manchester schemes using each of the following data streams, assuming that the last signal level has been positive.

- a. 00000000
- b. 11111111
- c. 01010101
- d. 00110011

2B1Q scheme:

- a. 0000000000000000
- b. 1111111111111111
- c. 0101010101010101
- d. 0011001100110011

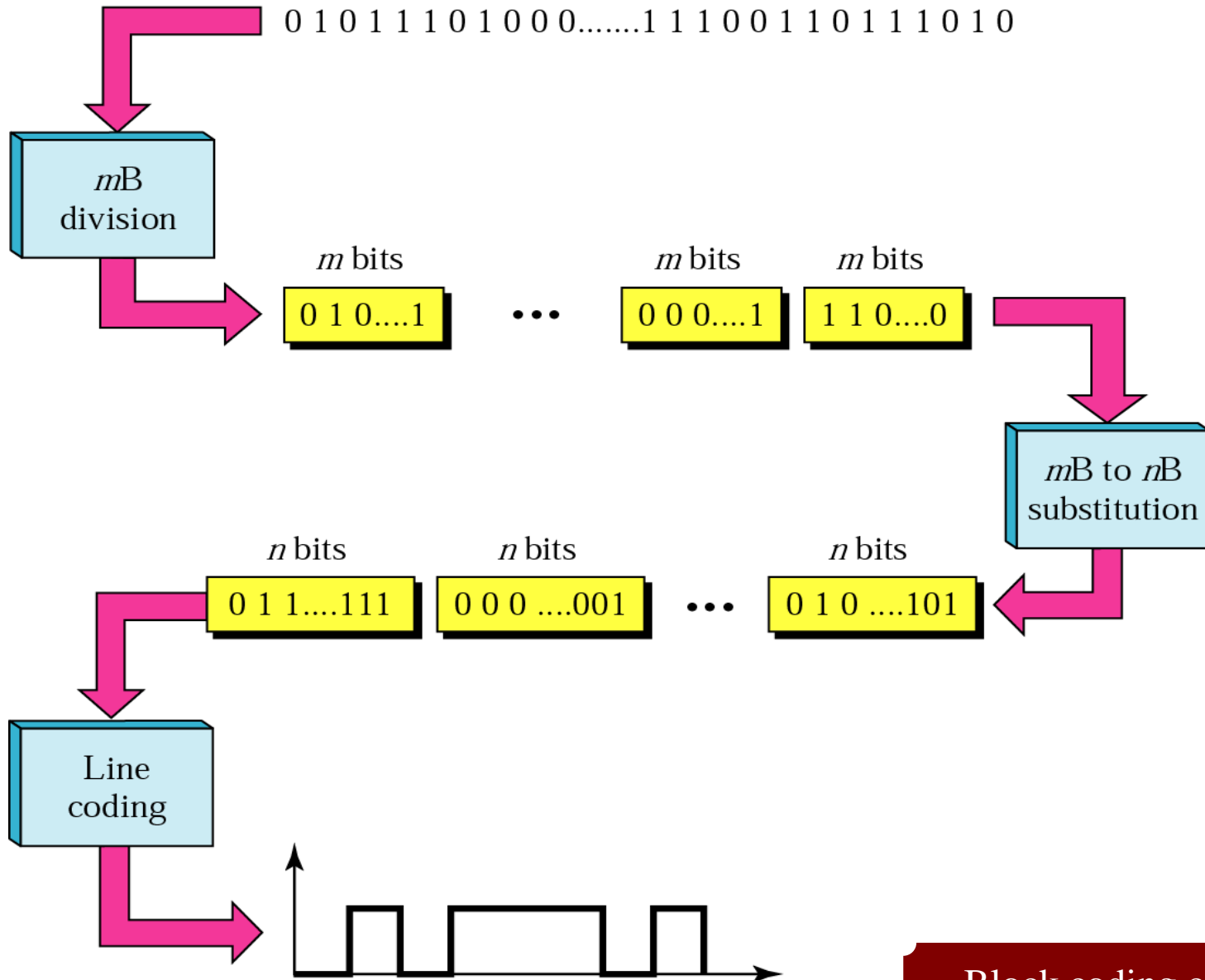
MLT-3 scheme:

- a. 00000000
- b. 11111111
- c. 01010101
- d. 00011000

## Block Coding

- ⊙ Block coding changes a block of  $m$  bits into a block of  $n$  bits, where  $n$  is larger than  $m$ .
- ⊙ Block coding is referred to as an  $mB/nB$  encoding technique.
  - ⊙ It replaces each  $m$ -bit group with an  $n$ -bit group.
- ⊙ Block coding normally involves three steps:
  - ⊙ Division, Substitution, and Combination.
- ⊙ In the division step, a sequence of bits is divided into groups of  $m$  bits.
  - ⊙ For example, in 4B/5B encoding, the original bit sequence is divided into 4-bit groups.
- ⊙ The heart of block coding is the substitution step.
  - ⊙ In this step, we substitute an  $m$ -bit group for an  $n$ -bit group.
  - ⊙ For example, in 4B/5B encoding we substitute a 4-bit code for a 5-bit group.
- ⊙ Finally, the  $n$ -bit groups are combined together to form a stream. The new stream has more bits than the original bits.

# Block Coding

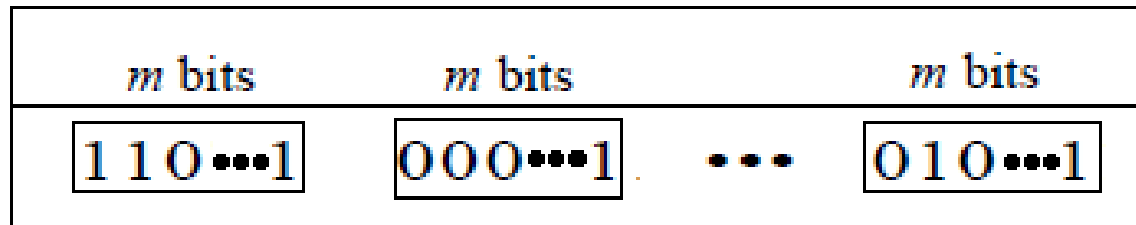


Block coding concept

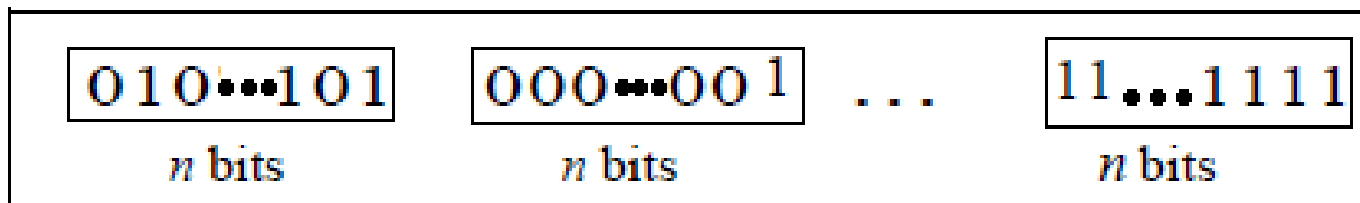
# Block Coding

## Block coding concept

Division of a stream into  $m$ -bit groups



$mB$ -to- $nB$   
substitution

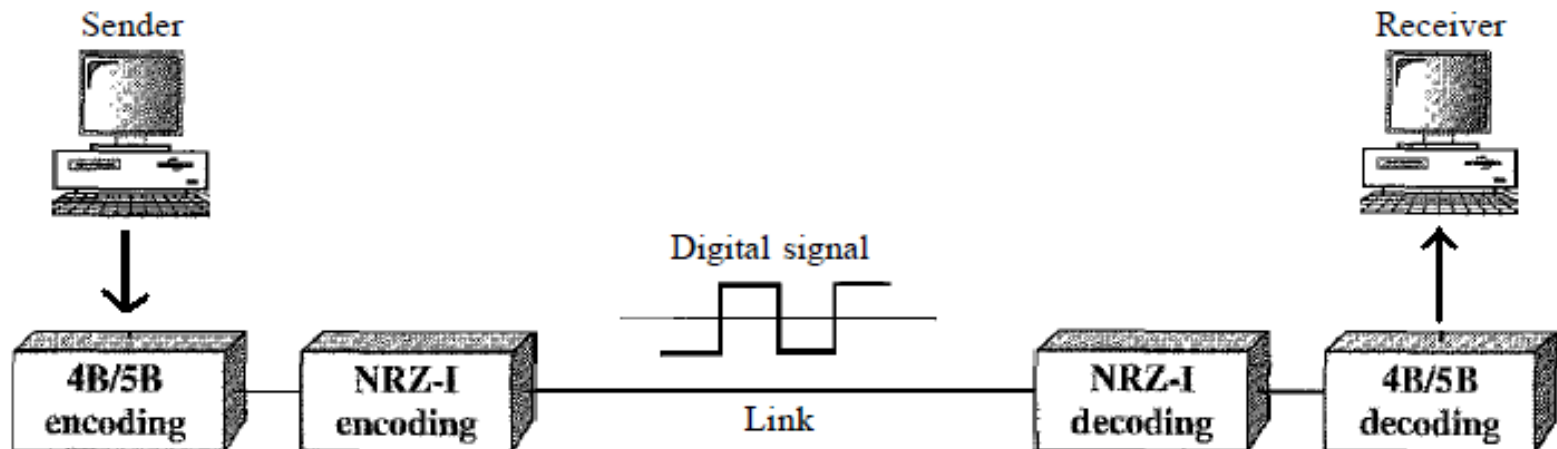


Combining  $n$ -bit groups into a stream

# Block Coding

## 4B/5B

- ⊙ The four binary/five binary (4B/5B) coding scheme was designed to be used in combination with NRZ-I.
- ⊙ Steps:
  - ⊙ At the sender side, change the bit stream, prior to encoding with NRZ-I.
  - ⊙ At the receiver, the NRZ-I encoded digital signal is first decoded into a stream of bits and then decoded.

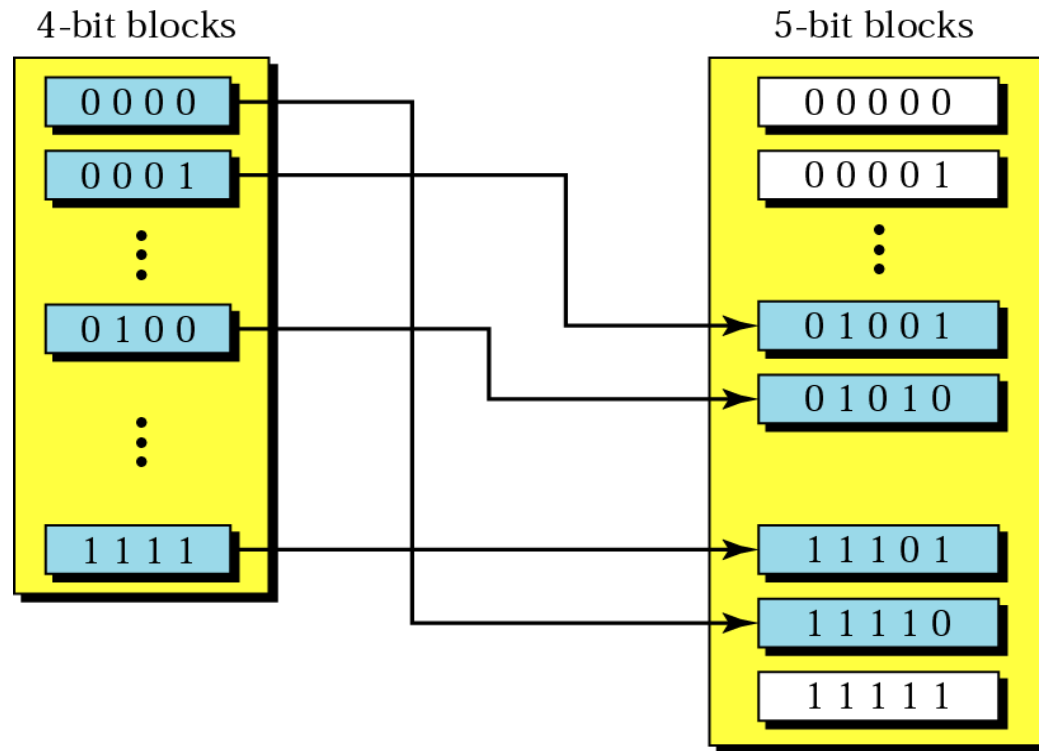


Using block coding 4B/5B with NRZ-I line coding scheme

# Block Coding

## 4B/5B

- ⊙ In 4B/5B, the 5-bit output that replaces the 4-bit input has no more than one leading zero (left bit) and no more than two trailing zeros (right bits).
- ⊙ So when different groups are combined to make a new sequence, there are never more than three consecutive 0s.



# Block Coding

## 4B/5B

- ⊙ A group of 4 bits can have only 16 different combinations while a group of 5 bits can have 32 different combinations.
- ⊙ This means that there are 16 groups that are not used for 4B/5B encoding.
- ⊙ Some of these unused groups are used for control purposes; the others are not used at all. The latter provide a kind of error detection.
- ⊙ If a 5-bit group arrives that belongs to the unused portion of the table, the receiver knows that there is an error in the transmission.



## Block Coding 4B/5B

| <i>Data Sequence</i> | <i>Encoded Sequence</i> | <i>Control Sequence</i> | <i>Encoded Sequence</i> |
|----------------------|-------------------------|-------------------------|-------------------------|
| 0000                 | 11110                   | Q (Quiet)               | 00000                   |
| 0001                 | 01001                   | I (Idle)                | 11111                   |
| 0010                 | 10100                   | H (Halt)                | 00100                   |
| 0011                 | 10101                   | J (Start delimiter)     | 11000                   |
| 0100                 | 01010                   | K (Start delimiter)     | 10001                   |
| 0101                 | 01011                   | T (End delimiter)       | 01101                   |
| 0110                 | 01110                   | S (Set)                 | 11001                   |
| 0111                 | 01111                   | R (Reset)               | 00111                   |
| 1000                 | 10010                   |                         |                         |
| 1001                 | 10011                   |                         |                         |
| 1010                 | 10110                   |                         |                         |
| 1011                 | 10111                   |                         |                         |
| 1100                 | 11 010                  |                         |                         |
| 1101                 | 11011                   |                         |                         |
| 1110                 | 11100                   |                         |                         |
| 1111                 | 11101                   |                         |                         |

## Digital-to-Digital Conversion

