

1.5 Laws of Propositional logic

$$p \rightarrow q \equiv \neg p \vee q$$

This rule can also apply to compound propositions.

$$(\neg t \wedge r) \rightarrow (\neg s \vee t) \equiv \neg(\neg t \wedge r) \vee (\neg s \vee t)$$

Laws of propositional Logic

Idempotent Laws	$p \vee p \equiv p$	$p \wedge p \equiv p$
Associative Laws	$(p \vee q) \vee r \equiv p \vee (q \vee r)$	$(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$
Commutative Laws	$p \vee q \equiv q \vee p$	$p \wedge q \equiv q \wedge p$
Distributive Laws	$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$	$p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$
Identity Laws	$p \vee F \equiv p$	$p \wedge T \equiv p$
Domination Laws	$p \wedge F \equiv F$	$p \vee T \equiv T$
Double Negation Law	$\neg\neg p \equiv p$	
Complement laws	$p \wedge \neg p \equiv F$ $\neg T \equiv F$	$p \vee \neg p \equiv T$ $\neg F \equiv T$
De Morgan's Laws	$\neg(p \vee q) \equiv \neg p \wedge \neg q$	$\neg(p \wedge q) \equiv \neg p \vee \neg q$
Absorption Laws	$p \vee (q \wedge p) \equiv p$	$p \wedge (q \vee p) \equiv p$
Conditional Identities	$p \rightarrow q \equiv \neg p \vee q$	$p \leftarrow q \equiv (p \rightarrow q) \wedge (q \rightarrow p)$

1.6 Predicates and Qualifiers

" x is an odd number" \neq proposition

- because x is not specified

Predicate: A logical statement with a truth value that is a function of 1 or more variables.

The Domain of a variable in a predicate is the set of all possible values for the variable.

"The city has a population over 1,000,000" = predicate

- because it has a variable
- the variable is "city"

• " $\forall x P(x)$ " is read "for all x , $P(x)$ "

- This statement asserts that $P(x)$ is true for every value of x in domain

• The symbol \forall is a universal quantifier.

• The statement " $\forall x P(x)$ " is a universally quantified statement.

• " $\forall x P(x)$ " = proposition

- because it is either T or F

• Arbitrary element: nothing is assumed about the element other than the fact that it is in the domain.

• Counterexample for a universally quantified statement is an element in the domain for which the predicate is false

- If the domain for variable x is empty, then the statement $\forall x P(x)$ is True because there are no elements in the domain for which $P(x)$ is false.
- " $\exists x P(x)$ " is read "there exists an x such that $P(x)$ "
 ① This statement asserts that $P(x) = T$ for at least one x in domain
- The symbol \exists is an existential quantifier.
- The statement " $\exists x P(x)$ " is called an existentially quantified statement.
 ② It is a proposition because it is either true or false
- An example for an existentially qualified statement is an element in the domain for which the predicate is true.



1.7 Quantified Statements

- The universal and existential quantifiers are generally called quantifiers.

Quantified Statement: A logical statement that includes a universal or existential quantifier.

- A variable x in the predicate $P(x)$ is called a free variable.
- The variable x in $\forall x P(x)$ is a bound variable because the variable is bound to a quantifier.

1.8 DeMorgan's Law for quantified statements

$$\neg \forall x F(x) \equiv \exists x \neg F(x)$$

$$\neg (F(a_1) \wedge F(a_2) \dots \wedge F(a_n)) \equiv \neg F(a_1) \vee \neg F(a_2) \dots \vee \neg F(a_n)$$

$$\neg \exists x P(x) \equiv \forall x \neg P(x)$$

$$\neg (P(a_1) \vee P(a_2) \vee \dots \vee P(a_n)) \equiv \neg P(a_1) \wedge \neg P(a_2) \dots \wedge \neg P(a_n)$$