

Lecture 12

Partial Orders

- A relation is a partial order if its:

1. reflexive

2. transitive

3. anti-symmetric

- If aRb for a partial order R , we say $a \leq b$.

- Partially ordered set: (A, \leq) is called a poset.

- if $x \leq y$, then x, y are comparable

- Total order: $A \leq$ is a total order if every two elements in the domain are comparable.

- x is minimal if there is no y such that $y \leq x$.

- x is maximal if there is no y such that $x \leq y$.

Strict Order

- a relation is a strict order if its:

1. transitive

2. anti-symmetric.

Algorithms

- An algorithm is a finite set of precise instructions for solving a problem.
- A problem is defined by a pair (input, desired output).

The Growth of Functions

- We use "time complexity" to evaluate algorithms.
- Big O notation:

• When f's growth rate is less than or equal to g's growth rate.

$$f(x) = \boxed{3x^2} + 2x + 4 = O(x^2)$$

analyze
dominant
part and
drop constants

$$f(x) = 10x^2 \log(x) + x + 4 = O(x^2 \log x)$$

f(x) is	<table border="0" style="width: 100%;"> <tr> <td style="padding-right: 20px;">faster than g(x)</td><td>if $\lim_{x \rightarrow \infty} \frac{f(x)}{g(x)} = +\infty$</td></tr> <tr> <td style="padding-right: 20px;">as fast as g(x)</td><td>if $\lim_{x \rightarrow \infty} \frac{f(x)}{g(x)} = \text{Some constant}$</td></tr> <tr> <td style="padding-right: 20px;">slower than g(x)</td><td>if $\lim_{x \rightarrow \infty} \frac{f(x)}{g(x)} = 0$</td></tr> </table>	faster than g(x)	if $\lim_{x \rightarrow \infty} \frac{f(x)}{g(x)} = +\infty$	as fast as g(x)	if $\lim_{x \rightarrow \infty} \frac{f(x)}{g(x)} = \text{Some constant}$	slower than g(x)	if $\lim_{x \rightarrow \infty} \frac{f(x)}{g(x)} = 0$
faster than g(x)	if $\lim_{x \rightarrow \infty} \frac{f(x)}{g(x)} = +\infty$						
as fast as g(x)	if $\lim_{x \rightarrow \infty} \frac{f(x)}{g(x)} = \text{Some constant}$						
slower than g(x)	if $\lim_{x \rightarrow \infty} \frac{f(x)}{g(x)} = 0$						

Example

$$4x^3 = O(x^3)$$

$$4x^3 = O(x^4)$$

$$4x^3 = O(x^{3.5})$$

$$4x^3 \neq O(x^{2.5})$$

- Little - Oh : "less than and not equal to"
- Big - theta : "Exactly equal to"
- Big - Omega : "Greater than or equal to"
- Little - Omega : "Greater than and not equal to"

Formal / Mathematical Approach

• f is $O(g)$ if :

$$\exists c, n_0 \quad \forall n > n_0 : f(n) \leq c \cdot g(n)$$

Example

$$3n + 7 = O(n)$$

$$f(n) \leq c \cdot g(n)$$

n	$f(n)$	$g(n)$	$\lceil \frac{f(n)}{g(n)} \rceil$	ceiling function
1	10	1	10	
10	37	10	4	
100	307	100	4	

↑
4 is potential candidate for c

$$\text{so } n > 10 : f(n) \leq 4 \cdot g(n)$$

$$\text{so } n > 10$$

$$3n + n > 3n + 10$$

$$\begin{aligned} 4n &> 3n + 10 > 3n + 7 \\ c \cdot g(n) &> f(n) \end{aligned}$$