

## 1.5 Laws of Proposition logic

$$p \rightarrow q \equiv \neg p \vee q$$

• This rule can also apply to compound propositions.

$$(\neg t \wedge r) \rightarrow (\neg s \vee t) \equiv \neg((\neg t \wedge r) \vee (\neg s \vee t))$$

## Laws of Propositional Logic

Idempotent Laws	$p \vee p \equiv p$	$p \wedge p \equiv p$
Associative Laws	$(p \vee q) \vee r \equiv p \vee (q \vee r)$	$(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$
Commutative Laws	$p \vee q \equiv q \vee p$	$p \wedge q \equiv q \wedge p$
Distributive Laws	$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$	$p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$
Identity Laws	$p \vee F \equiv p$	$p \wedge T \equiv p$
Domination Laws	$p \wedge F \equiv F$	$p \vee T \equiv T$
Double Negation Law	$\neg \neg p \equiv p$	
Complement laws	$p \wedge \neg p \equiv F$ $\neg T \equiv F$	$p \vee \neg p \equiv T$ $\neg F \equiv T$
De Morgan's Laws	$\neg (p \vee q) \equiv \neg p \wedge \neg q$	$\neg (p \wedge q) \equiv \neg p \vee \neg q$
Absorption Laws	$p \vee (q \wedge p) \equiv p$	$p \wedge (q \vee p) \equiv p$
Conditional Identities	$p \rightarrow q \equiv \neg p \vee q$	$p \leftrightarrow q \equiv (p \rightarrow q) \wedge (q \rightarrow p)$

## 1.6 Predicates and Qualifiers

" $x$  is an odd number"  $\neq$  proposition

- because  $x$  is not specified

Predicate: A logical statement with a truth value that is a function of 1 or more variables.

The Domain of a variable in a predicate is the set of all possible values for the variable.

"The city has a population over 1,000,000" = predicate

- because it has a variable
- the variable is "city"

• " $\forall x P(x)$ " is read "for all  $x$ ,  $P(x)$ "

- This statement asserts that  $P(x)$  is true for every value of  $x$  in domain

• The symbol  $\forall$  is a universal quantifier

• The statement " $\forall x P(x)$ " is a universally quantified statement.

• " $\forall x P(x)$ " = proposition

- because it is either T or F

• Arbitrary element: nothing is assumed about the element other than the fact that it is in the domain.

• Counterexample for a universally quantified statement is an element in the domain for which the predicate is false

- If the domain for variable  $x$  is empty, then the statement  $\forall x P(x)$  is True because there are no elements in the domain for which  $P(x)$  is false.
- " $\exists x P(x)$ " is read "there exists an  $x$  such that  $P(x)$ "
  - This statement asserts that  $P(x) = T$  for at least one  $x$  in domain
- The symbol  $\exists$  is an existential quantifier.
- The statement " $\exists x P(x)$ " is called an existentially quantified statement.
  - It is a proposition because it is either true or false
- An example for an existentially qualified statement is an element in the domain for which the predicate is true.

## 1.7 Quantified Statements

- The universal and existential quantifiers are generally called quantifiers.

Quantified Statement: A logical statement that includes a universal or existential quantifier.

- A variable  $x$  in the predicate  $P(x)$  is called a free variable.
- The variable  $x$  in  $\forall x P(x)$  is a bound variable because the variable is bound to a quantifier.

## 1.8 De Morgan's Law for quantified statements

$$\neg \forall x F(x) \equiv \exists x \neg F(x)$$

$$\neg (F(a_1) \wedge F(a_2) \dots \wedge F(a_n)) \equiv \neg F(a_1) \vee \neg F(a_2) \dots \vee \neg F(a_n)$$

$$\neg \exists x P(x) \equiv \forall x \neg P(x)$$

$$\neg (P(a_1) \vee P(a_2) \vee \dots \vee P(a_n)) \equiv \neg P(a_1) \wedge \neg P(a_2) \dots \wedge \neg P(a_n)$$