

8.12 Analyzing time complexity of recursive algorithms

- The time complexity of an algorithm $T(n)$ indicating the number of atomic operations performed by the algorithm on an input of size n .

Example

Factorial (n)

If ($n=0$), Return(1)

$r := \text{Factorial}(n-1)$

Return($r * n$)

if $n=0$, two atomic operations are performed
(conditional and return)

Recursive call to Factorial on input $n-1$ is $T(n-1)$.

Initial Value: $T(0) = 2$

Recurrence relation: for $n \geq 1$, $T(n) = T(n-1) + 5$

Examples of simplifying Recurrence relations

$$\bullet T(n) = 3 \cdot T\left(\frac{n}{2}\right) + 7n \rightarrow T\left(\frac{n}{2}\right) + \Theta(n)$$

↳ INCORRECT, CANNOT Remove scalar on $T\left(\frac{n}{2}\right)$ term

$$\bullet T(n) = T\left(\left\lfloor \frac{n}{2} \right\rfloor\right) + 5n^2 + 10 \rightarrow T(n) = T\left(\frac{n}{2}\right) + \Theta(n^2)$$

↳ Correct

$$\bullet T(n) = T\left(\left\lfloor \frac{n}{2} \right\rfloor\right) + T\left(\left\lfloor \frac{n}{2} \right\rfloor\right) + \Theta(n) \rightarrow T(n) = 2 \cdot T\left(\frac{n}{2}\right) + \Theta(n)$$

↳ Correct

8.13 Divide and Conquer Algorithms: Introduction and mergesort

- A divide and conquer algorithm solves a problem recursively by breaking the original input into smaller subproblems of roughly equal size. They follow 3 basic steps:

1. Break the input into smaller subproblems of the same type on smaller inputs.
2. Solve each subproblem recursively.
3. Combine the solutions of the subproblems to obtain a solution to the original problem.

Merge Sort

- Merge Sort sorts the list $(17, 3, 9, 11, 21, 4, 7, 1)$ in 4 basic steps:

1. Break list into two sub-lists: $(17, 3, 9, 11)$ and $(21, 4, 7, 1)$
2. Make a recursive call to MergeSort with input $(17, 3, 9, 11)$, which returns $(3, 9, 11, 17)$
3. Make a recursive call to MergeSort with input $(21, 4, 7, 1)$, which returns $(1, 4, 7, 21)$
4. Merge the sorted sub-lists to obtain $(1, 3, 4, 7, 9, 11, 17, 21)$.

- Merge Sort is $\Theta(n \log n)$

- Merge Sort uses queue data structure:

- New items can be added to one end called the back.
- Items can be removed from the other end of the queue called the front.

Example: Queue B: $(5, 8, 7)$

↳ Remove (B) : $(8, 7)$

↳ Add (9) : $(8, 7, 9)$

8.14 Divide and Conquer algorithms: Binary Search

• Pseudocode for recursive binary search

RecBinarySearch(low, high, A, x)

if (low = high AND $a_{low} = x$), Return (low)

if (low = high AND $a_{low} \neq x$), Return (-1)

$$mid := \left\lfloor \frac{low + high}{2} \right\rfloor$$

if ($x \leq a_{mid}$), then high := mid

if ($x > a_{mid}$), then low := mid + 1

Return [RecBinarySearch(low, high, A, x)]

- If size of list is even, then the two subranges are both $\frac{n}{2}$: $(mid - low + 1) = (high - mid) = \frac{n}{2}$
- If size of list is odd, then the size of first subrange is $(mid - low + 1) = \lceil \frac{n}{2} \rceil$, and the size of the second subrange is $(high - mid) = \lfloor \frac{n}{2} \rfloor$.