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Section: 11

# Homework 10

P1. a

PRESENT STATE	NEXT		OUTPUT
	W=0	W=1	
A	A	B	0
B	B	C	1
C	C	D	2
D	D	A	3

b.

$Y_1$	$Y_0$	$W=0$ $Y_1, Y_0$	$W=1$ $Y_1, Y_0$	$Z_1, Z_0$
0	0	00	00	00
0	1	01	01	01
1	0	10	10	10
1	1	11	11	11

c.

W	$Y_1$	$Y_0$	$Y_1$	$Y_0$	$Z_1$	$Z_0$
0	0	0	0	0	0	0
0	0	1	0	1	0	1
0	1	0	1	0	1	0
0	1	1	1	1	1	1
1	0	0	0	1	0	0
1	0	1	1	0	0	1
1	1	0	1	1	1	0
1	1	1	0	0	1	1

$Y_1$

$W$	$Y_1, Y_0$	00	01	10	11
0	0	0	0	1	1
1	0	0	1	0	1

$$Y_1 = Y_1 \bar{W} + Y_1 \bar{Y}_0 + \bar{Y}_1 Y_0 W$$

$Y_0$

$Y_1, Y_0$

$w$

0	1	1	0
1	0	0	1

$$Y_0 = Y_0 \oplus W$$

\_/\_/\_

$Z_1$   $y_1, y_0$

$w$	00	01	11	10
0	0	0	1	1
1	0	0	1	1

$Z_1$   $y_1$

$y_0$	0	1
0	0	1
1	0	1

$$Z_1 = y_1$$

$Z_i$  only depended on  $y_i$  ( $i = 0, 1$ )

$Z_0$   $y_1, y_0$

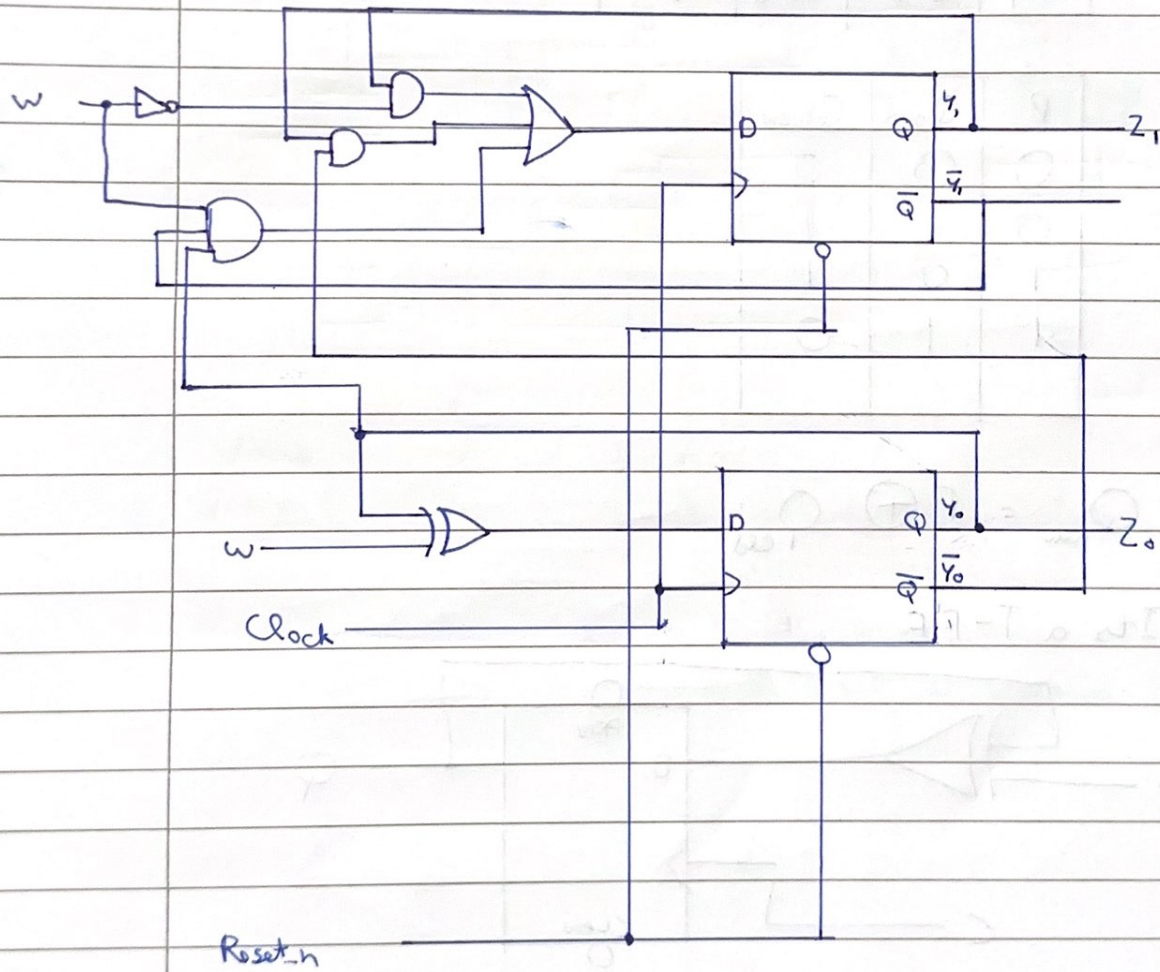
$w$	00	01	11	10
0	0	1	1	0
1	0	1	1	0

$Z_0$   $y_0$

$y_1$	0	1
0	0	0
1	1	1

$$Z_0 = y_0$$

d.



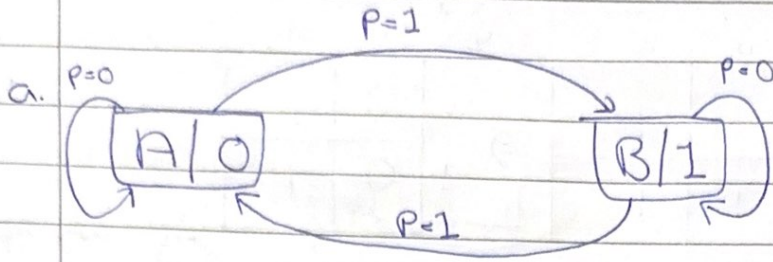


e. FSM, when  $w=0$ ; Points back to same state.

FSM when  $w=1$ ; points to the next state.

(upcount in binary). ( $A \rightarrow B$  or  $B \rightarrow C$  or  $C \rightarrow D$  or  $D \rightarrow A$ )

P2.



b.

	P=0	P=1	Q
A	A	B	0
B	B	A	1

c.

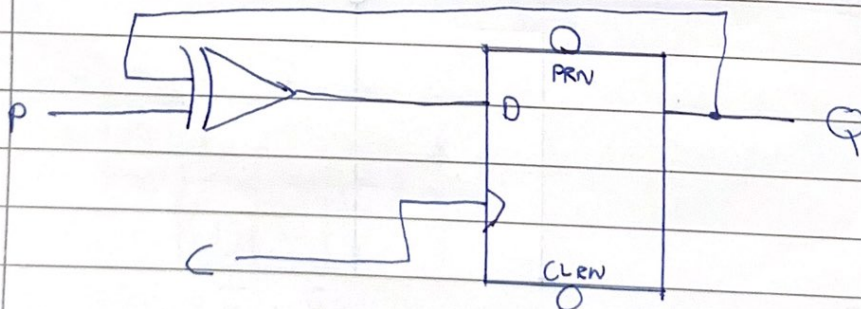
	P=0	P=1	Q
A:0	A:0	B:1	0
B:1	B:1	A:0	1

d.

P	Q <sub>OLD</sub>	Q <sub>NEW</sub>
0	0	0
0	1	1
1	0	1
1	1	0

e.  $Q_{\text{new}} = P \oplus Q_{\text{old}}$

f. Its a T-FF





P3. Counting Sequence: 000, 001, 010, 111

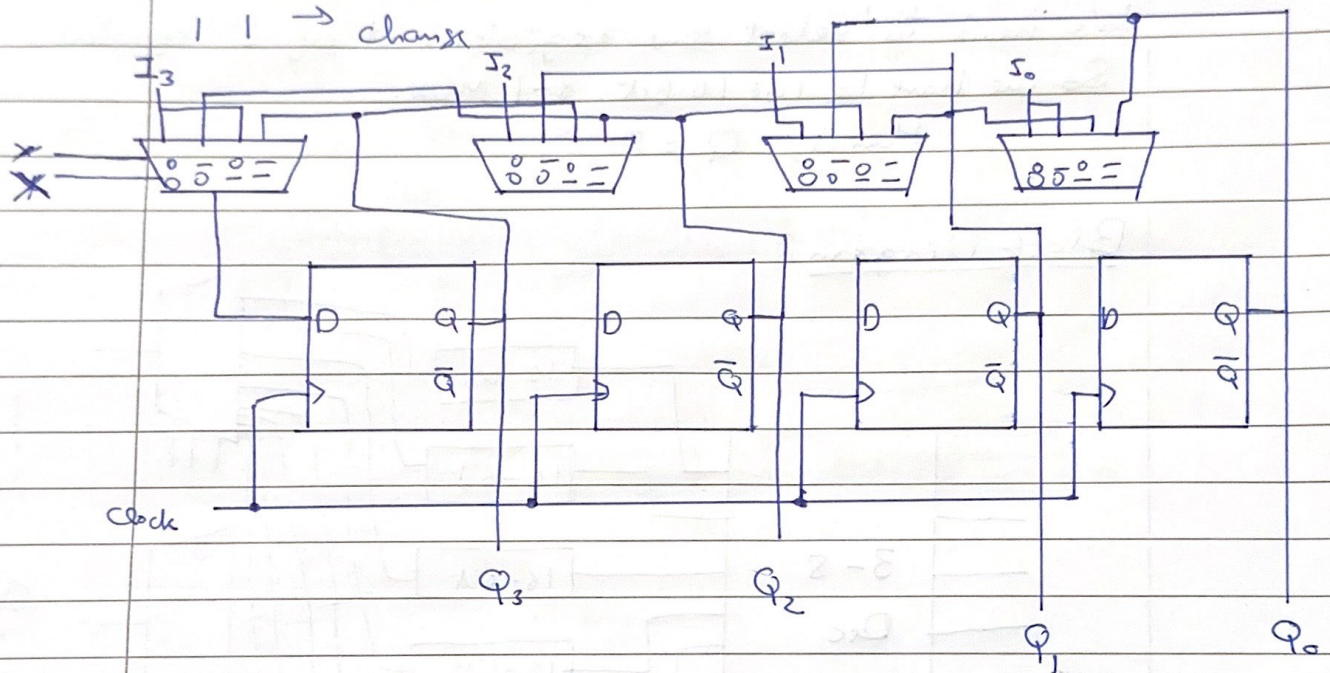
P4.	X	Y
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○ ○ → new date

○ 1 → shift left

10 → " eight

1 1  $\rightarrow$  change



P5. There are eight 16-bit registers. We can access those 8 registers using 1 3-8 decoder.

$$W = 1, X = 3, Y = 8$$

3 output ports means we need 3 MUX's.

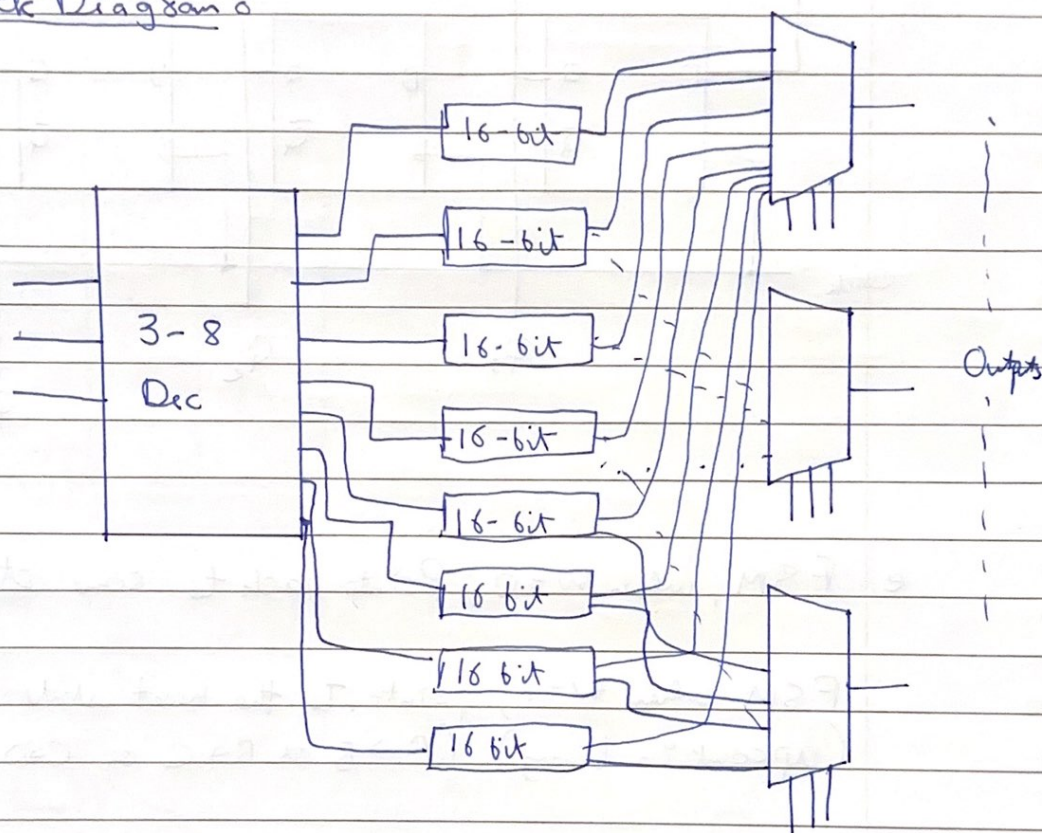
$$\text{Hence, } Z = 3.$$

We need to select one register out of 8 registers.

So we have to use 16 bit 8-1 MUX.

$$\text{Hence, } Q = 8$$

### Block Diagram :



$$W = 1, X = 3, Y = 3, Z = 3, Q = 8$$

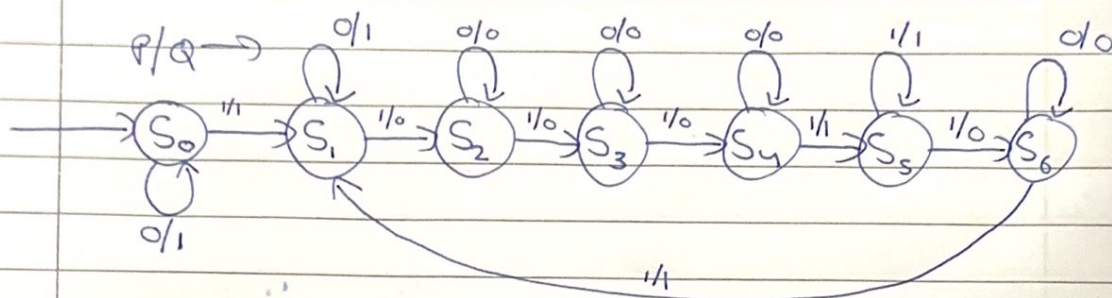


P6. Since, we are only concerned with numbers of 1's, we can ignore the 0's by changing state only when  $P=1$ .

We start in a special state  $S_0$  where  $Q=1$  because we have no 1's yet.

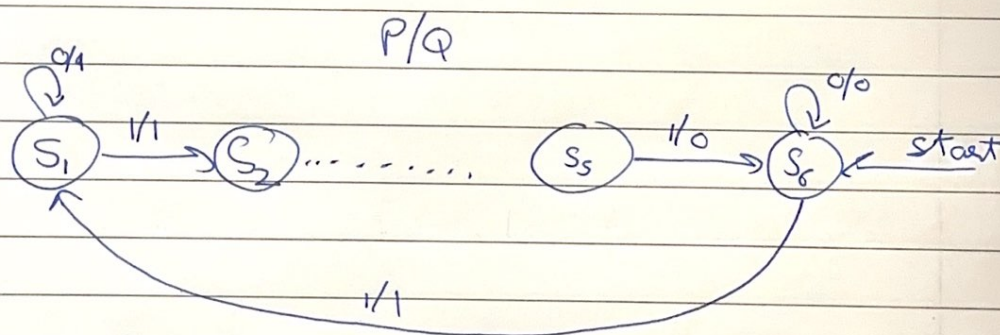
Each time we have  $P=1$ , we advance to the next state, so it's basically counting 1's.

When it gets to 6, that is div. by 2 & 3, we can restart.



Alternative with 0 defined as divisible by 2 or 3.

Six states instead of Seven.



It is a Mealy machine. This has an output value rather than a final state.

Moreover, its output values are determined both by its current state and current inputs.