

Algorithm to compute FIRST, FOLLOW and nullable

Initialize FIRST and FOLLOW to all empty sets and nullable to all false.

```
For each terminal symbol Z
        FIRST(Z)=\{Z\};
repeat
For each production X -> Y_1 Y_2 ... Y_k
        For each i from 1 to k, each j from i+1 to k
            if all the Y<sub>i</sub> are nullable
                    then nullable[X]=true
            if Y_1...Y_{i-1} are all nullable
                    then FIRST[X]=FIRST[X] U FIRST[Yi]
            if Y_{i+1}...Y_k are all nullable
                    then FOLLOW[Y<sub>i</sub>]=FOLLOW[Y<sub>i</sub>]UFOLLOW[X]
            if Y<sub>i+1</sub>..Y<sub>i-1</sub> are all nullable
            then FOLLOW[Y<sub>i</sub>]=FOLLOW[Y<sub>i</sub>] U FIRST[Y<sub>i</sub>]
until FIRST and FOLLOW nullable did not change in this iteration.
```

First - 1

example for compute nullable FIRST() and FOLLOW()

Z-> d	Y->	X->Y			
Z->XYZ	Y->c	X-> a			
Grammar					

	nullable	FIRST	FOLLOW		nullable	FIRST	FOLLOW
$\overline{\mathbf{X}}$	no			$\overline{\mathbf{X}}$	no	a	cd
\mathbf{Y}	no			\mathbf{Y}	yes	c	d
${f Z}$	no			${\bf Z}$	no	d	
		Initial					

IIIIIII

Iterative 1

	nullable	FIRST	FOLLOW
$\overline{\mathbf{X}}$	yes	ac	acd
\mathbf{Y}	yes	c	acd
Z	no	acd	

In iterative 3, because no new information is generated, the algorithm stops.



Example for generating predictive parsing table

S->E\$				
E->E+T	T->T*F	F->id		
E->E-T	T->T/F	F->num		
E->T	T->F	\mathbf{F} ->(\mathbf{E})		
Grammar				

S->E\$ E->TE'	T->FT'	F->id
E'->+TE'	T'->*FT'	
E'->-TE'	T'->/FT'	F->num
E'->	T'->	F->(E)



Compute nullable FIRST() and FOLLOW()

-	NULLABLE	FIRST	FOLLOW
S	no	(id num	
${f E}$	no	(id num)\$
E '	yes	+ =)\$
\mathbf{T}	no	(id num)+ -\$
T '	yes	* /)+-\$)+-\$)*/+-\$
${f F}$	no	(id num)*/ +-\$



•Build a predictive parsing table.

- •Enter production $X\rightarrow r$ in row X, column T of the table for each $T \in FIRST(r)$
- •If r is nullable, enter the production in row X, column T for each $T \in FOLLOW(X)$

	+, -	*,/	Id, num	()	\$
\mathbf{S}			S->E\$	S->E\$		
\mathbf{E}			E->TE'\$	E->TE'\$		
E '	E'->+TE'				E'->	E'->
\mathbf{T}			T->FT'\$	T->FT \$		
T '	T'->	T'->*FT'			T'->	T'->
${f F}$			F->id	F->(