

Semantic

<<Semantic.ppt>>

(* Attribute Grammar: synthetic/inherited attributes *)
 (* synthetic attribute *)

Grammar Rule

$number_1 \rightarrow number_2 \text{ digit}$
 $number \rightarrow digit$

$digit \rightarrow 0$

$digit \rightarrow 1$

$digit \rightarrow 2$

$digit \rightarrow 3$

$digit \rightarrow 4$

$digit \rightarrow 5$

$digit \rightarrow 6$

$digit \rightarrow 7$

$digit \rightarrow 8$

$digit \rightarrow 9$

Semantic Rules

$number_1.val = number_2.val * 10 + digit.val$
 $number.val = digit.val$

$digit.val = 0$

$digit.val = 1$

$digit.val = 2$

$digit.val = 3$

$digit.val = 4$

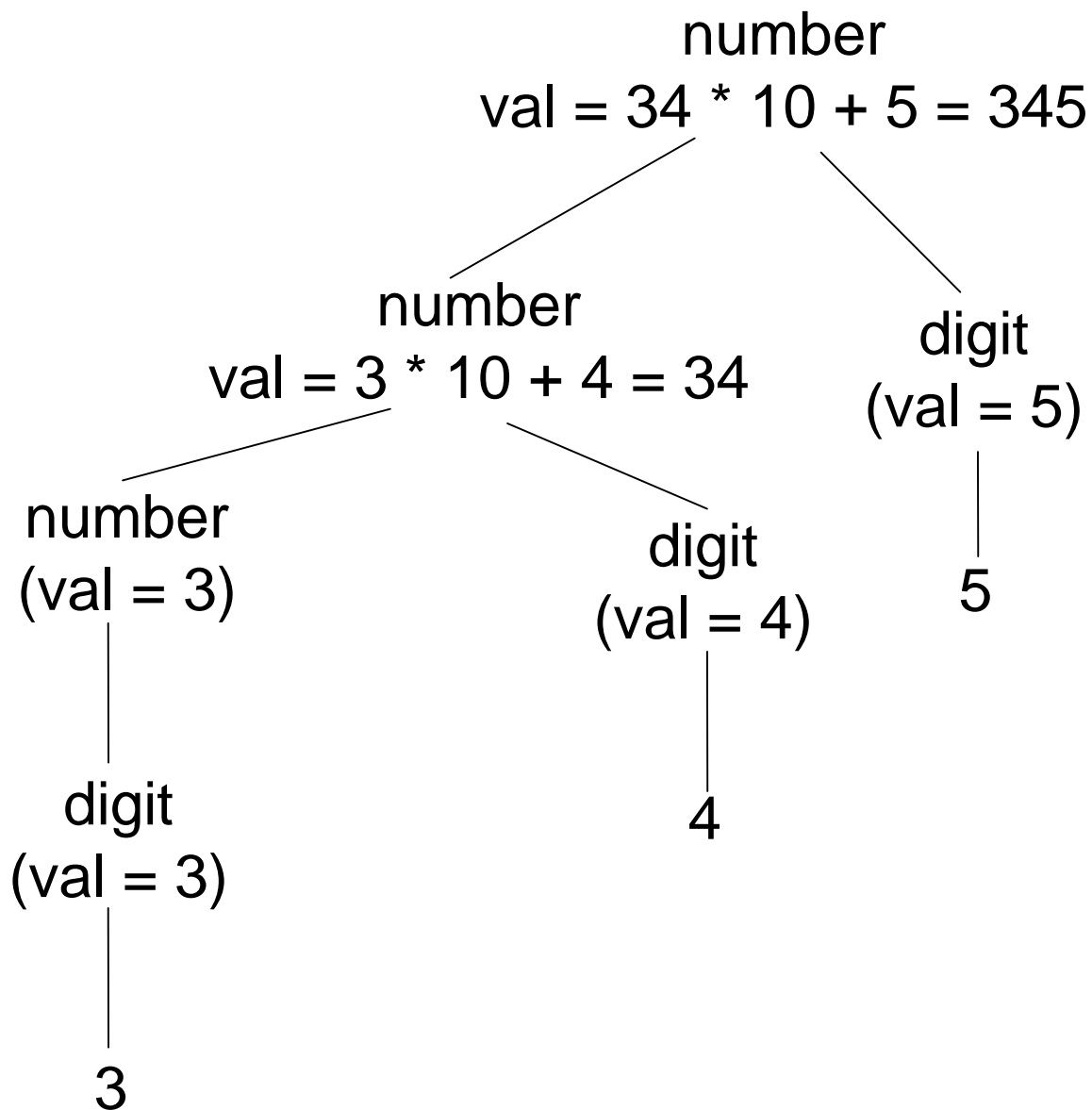
$digit.val = 5$

$digit.val = 6$

$digit.val = 7$

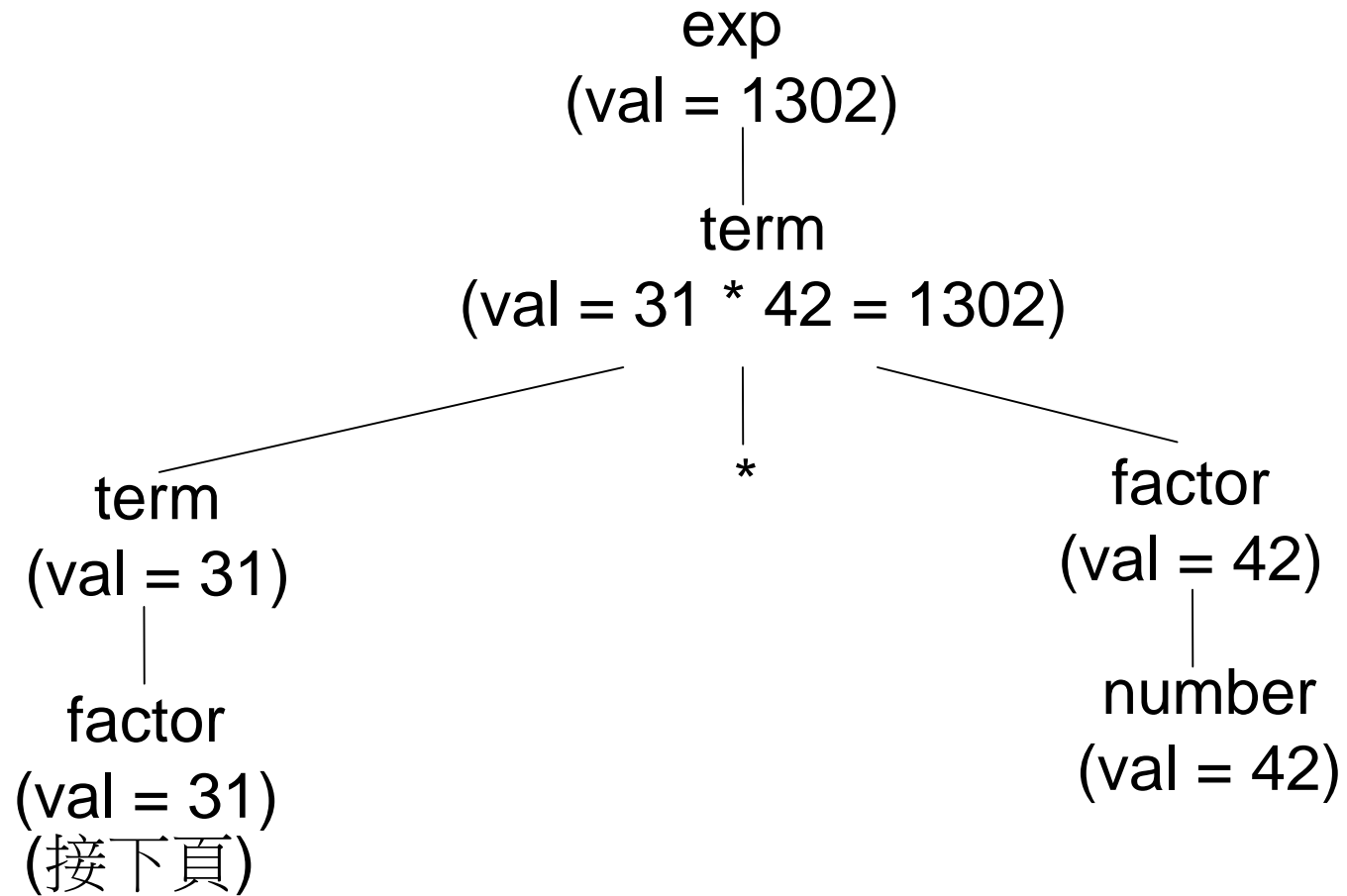
$digit.val = 8$

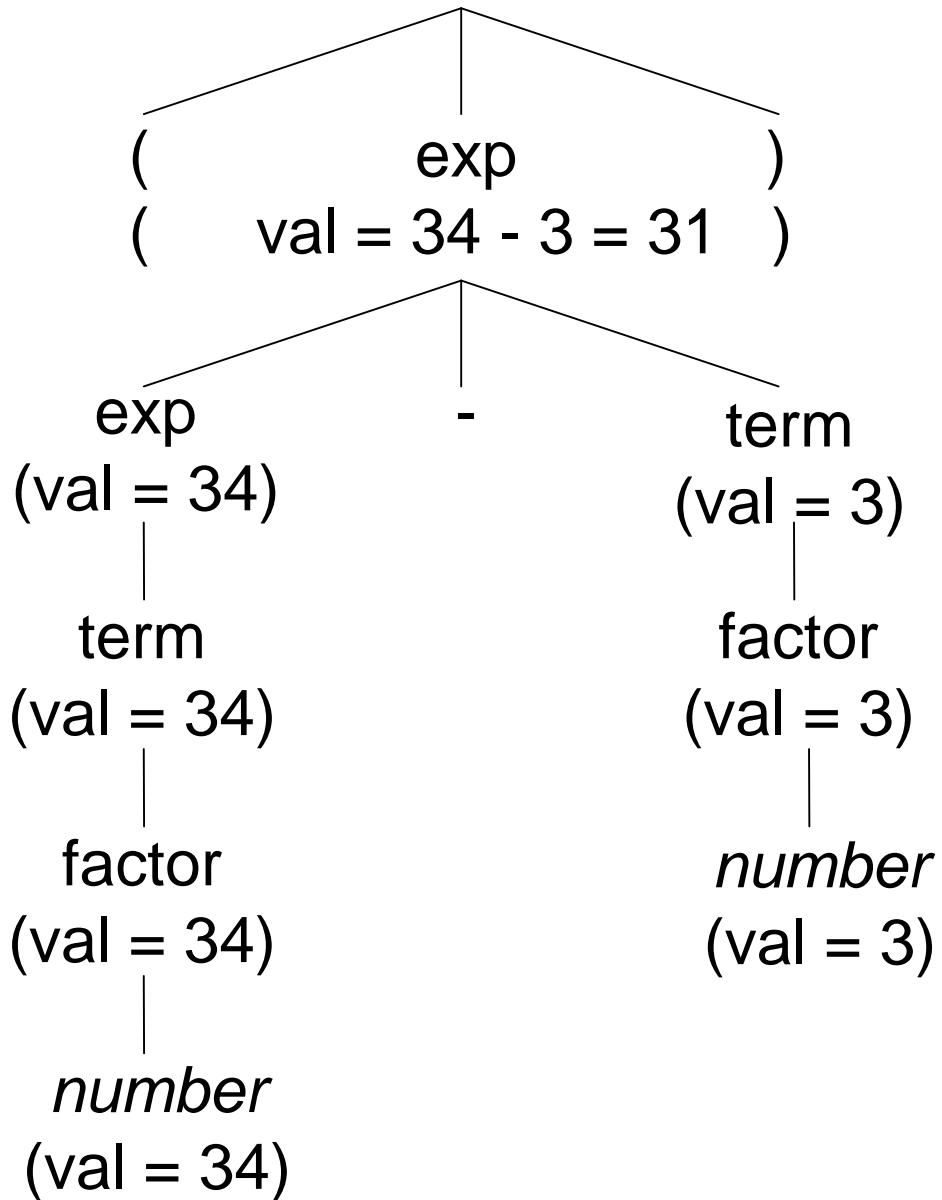
$digit.val = 9$



(* synthetic attribute *)

| Grammar Rule | Semantic Rules |
|---|---|
| $\text{exp}_1 \rightarrow \text{exp}_2 + \text{term}$ | $\text{exp}_1.\text{val} = \text{exp}_2.\text{val} + \text{term.val}$ |
| $\text{exp}_1 \rightarrow \text{exp}_2 - \text{term}$ | $\text{exp}_1.\text{val} = \text{exp}_2.\text{val} - \text{term.val}$ |
| $\text{exp} \rightarrow \text{term}$ | $\text{exp.val} = \text{term.val}$ |
| $\text{term}_1 \rightarrow \text{term}_2 * \text{factor}$ | $\text{term}_1.\text{val} = \text{term}_2.\text{val} * \text{factor.val}$ |
| $\text{term} \rightarrow \text{factor}$ | $\text{term.val} = \text{factor.val}$ |
| $\text{factor} \rightarrow (\text{exp})$ | $\text{factor.val} = \text{exp.val}$ |
| $\text{factor} \rightarrow \textbf{number}$ | $\text{factor.val} = \textbf{number.val}$ |





(* inherited attribute *)

Grammar Rule

$\text{decl} \rightarrow \text{type var_list}$
 $\text{type} \rightarrow \textit{int}$
 $\text{type} \rightarrow \textit{float}$
 $\text{var_list}_1 \rightarrow \textit{id}, \text{var_list}_2$
 $\text{var_list} \rightarrow \textit{id}$

Semantic Rules

$\text{var_list.dtype} = \text{type.dtype}$
 $\text{type.dtype} = \text{integer}$
 $\text{type.dtype} = \text{real}$
 $\textit{id.dtype} = \text{var_list}_1.dtype$
 $\text{var_list}_2.dtype = \text{var_list}_1.dtype$
 $\textit{id.dtype} = \text{var_list.dtype}$

Grammar Rule

Semantic Rules

$S \rightarrow \text{exp}$

$S.\text{etype} =$
 if $\text{exp}.\text{isFloat}$ **then** float **else** int
 $S.\text{val} = \text{exp}.\text{val}$

$\text{exp}_1 \rightarrow \text{exp}_2 / \text{exp}_3$

$\text{exp}_1.\text{isFloat} =$
 $\text{exp}_2.\text{isFloat}$ **or** $\text{exp}_3.\text{isFloat}$
 $\text{exp}_2.\text{etype} = \text{exp}_1.\text{etype}$
 $\text{exp}_3.\text{etype} = \text{exp}_1.\text{etype}$
 $\text{exp}_1.\text{val} =$
 if $\text{exp}_1.\text{etype} = \text{int}$
 then $\text{exp}_2.\text{val} \text{ div } \text{exp}_3.\text{val}$
 else $\text{exp}_2.\text{val} / \text{exp}_3.\text{val}$

(接下頁)

exp -> ***num***

exp.isFloat = ***false***

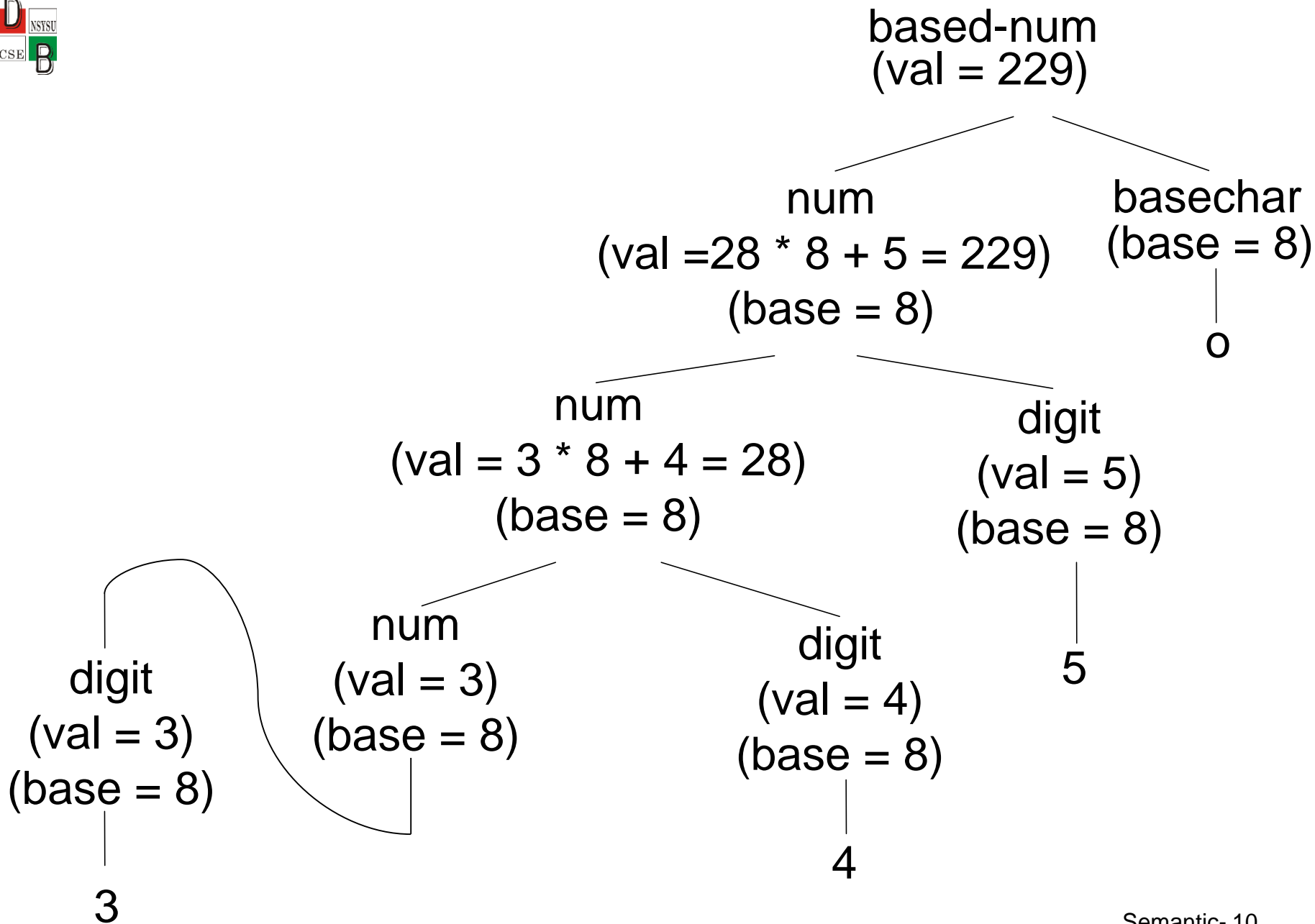
exp.val =

if exp.etype = int ***then*** ***num.val***
else Float(***num.val***)

exp -> ***num.num***

exp.isFloat = ***true***

exp.val = ***num.num.val***



| Grammar Rule | Semantic Rules |
|--------------------------------------|---|
| based-num \rightarrow num.basechar | based-num.val = num.val num.base = basechar.base |
| basechar \rightarrow o | basechar.base = 8 |
| basechar \rightarrow d | basechar.base = 10 |
| $num_1 \rightarrow num_2$ digit | $num_1.val =$ if digit.val = error or $num_2.val$ = error then error else $num_2.val * num_1.base + digit.val$ $num_2.base = num_1.base$ digit.base = $num_1.base$ |

| | |
|-------------------|---|
| num -> digit | num.val = digit.val digit.base = num.base |
| digit -> 0 | digit.val = 0 |
| digit -> 1 | digit.val = 1 |
| | |
| digit -> 7 | digit.val = 7 |
| digit -> 8 | digit.val = <i>if</i> digit.base = 8 then error else 8 |
| digit -> 9 | digit.val = <i>if</i> digit.base = 8 then error else 9 |

(* inherited attribute *)

| Grammar Rule | Semantic Rules |
|---|--|
| $\text{decl} \rightarrow \text{type var_list}$ $\text{type} \rightarrow \textit{int}$ $\text{type} \rightarrow \textit{float}$ $\text{var_list}_1 \rightarrow \text{var_list}_2, \textit{id}$ $\text{var_list} \rightarrow \textit{id}$ | $\text{var_list.dtype} = \text{type.dtype}$ $\text{type.dtype} = \text{integer}$ $\text{type.dtype} = \text{real}$ $\text{insert}(\textit{id.name}, \text{var_list}_1.dtype)$ $\text{var_list}_2.dtype = \text{var_list}_1.dtype$ $\text{insert}(\textit{id.name}, \text{var_list.dtype})$ |

(* the resulting data type; synthetic attribute *)

| 產生規則 | 語意規則 |
|---|--|
| $E \rightarrow \text{num}$ $E \rightarrow \text{num} . \text{num}$ $E \rightarrow \text{id}$ $E \rightarrow E_1 \text{ op } E_2$ | $E.\text{type} := \text{integer}$ $E.\text{type} := \text{real}$ $E.\text{type} := \text{lookup}(\text{id}.\text{entry})$ $E.\text{type} := \text{if } E_1.\text{type} = \text{integer} \text{ and } E_2.\text{type} = \text{integer} \text{ then integer}$ $\text{else if } E_1.\text{type} = \text{integer} \text{ and } E_2.\text{type} = \text{real} \text{ then real}$ $\text{else if } E_1.\text{type} = \text{real} \text{ and } E_2.\text{type} = \text{integer} \text{ then real}$ |

(接下頁)

| 產生規則 | 語意規則 |
|------|---|
| | <p><i>else if</i> $E_1.type = \text{real}$ <i>and</i> $E_2.type = \text{real}$ <i>then</i> real <i>else</i> type_error</p> |

(the resulting data type: synthetic attribute *)

| | |
|--------------------------------------|---|
| $E \rightarrow id$ | { $E.type := \text{lookup}(id.entry)$ } |
| $E \rightarrow E_1 \text{ mod } E_2$ | { $E.type :=$ if $E_1.type = \text{integer}$ and $E_2.type = \text{integer}$ then integer else type_error } |
| $E \rightarrow E_1 [E_2]$ | { $E.type :=$ if $E_2.type = \text{integer}$ and $E_1.type = \text{array}(s, t)$ then t else type_error } |
| $E \rightarrow E_1 \uparrow$ | { $E.type :=$ if $E_1.type = \text{pointer}(t)$ then t else type_error } |
| $S \rightarrow id := E$ | { $S.type :=$ if $id.type = E.type$ then void else type_error } |

(接下頁)

$S \rightarrow \text{if } E \text{ then } S_1 \quad \{ S.type := \text{if } E.type = \text{boolean} \text{ then } S_1.type \text{ Else type_error } \}$

$S \rightarrow \text{while } E \text{ do } S_1 \quad \{ S.type := \text{if } E.type = \text{boolean} \text{ then } S_1.type \text{ else type_error } \}$

$S \rightarrow S_1 ; S_2 \quad \{ S.type := \text{if } S_1.type = \text{void} \text{ and } S_2.type = \text{void} \text{ then void else type_error } \}$

| | Parsing Stack | Input | Parsing Action | Value Stack | Semantic Action |
|---|---------------------|----------|-----------------------------------|------------------------|------------------------------------|
| 1 | \$ | 3*4+5 \$ | shift | \$ | |
| 2 | \$ <i>n</i> | *4+5 \$ | reduce $E \rightarrow \mathbf{n}$ | \$ <i>n</i> | $E.val = \mathbf{n.val}$ |
| 3 | \$ E | *4+5 \$ | shift | \$ 3 | |
| 4 | \$ E^* | 4+5 \$ | shift | \$ 3 * | |
| 5 | \$ $E^* \mathbf{n}$ | +5 \$ | reduce $E \rightarrow \mathbf{n}$ | \$ 3 * <i>n</i> | $E.val = \mathbf{n.val}$ |
| 6 | \$ $E^* E$ | +5 \$ | reduce $E \rightarrow E^* E$ | \$ 3 * 4 | $E_1.val =$ $E_2.val * E_3.val$ |

(接下頁)

| | | | | | |
|----|------------|-------|---------------------------------|-------------|------------------------------------|
| 7 | \$ E | +5 \$ | shift | \$ 12 | |
| 8 | \$ $E+$ | 5 \$ | shift | \$ 12 + | |
| 9 | \$ $E + n$ | \$ | reduce $E \rightarrow n$ | \$ 12 + n | $E.val = n.val$ |
| 10 | \$ $E + E$ | \$ | reduce $E \rightarrow E + E$ | \$ 12 + 5 | $E_1.val =$ $E_2.val + E_3.val$ |
| 11 | \$ E | \$ | | \$ 17 | |

(* structure information)

| | |
|---|--|
| $P \rightarrow D$ | { offset := 0 } |
| $D \rightarrow D ; D$ | |
| $D \rightarrow \textit{id} : T$ | { enter (<i>id</i> .name, T.type, offset); offset := offset + T.width } |
| $T \rightarrow \textit{integer}$ | { T.type := integer; T.width := 4 } |
| $T \rightarrow \textit{real}$ | { T.type := <i>real</i> ; T.width := 8 } |
| $T \rightarrow \textit{array} [\textit{num}] \textit{of} T_1$ | { T.type := array(<i>num</i> .val, T_1 .type); T.width := <i>num</i> .val * T_1 .width } |
| $T \rightarrow \uparrow T_1$ | { T.type := pointer(T_1 .type); T.width := 4 } |

(* symbol table structure for implementing scoping rule.)

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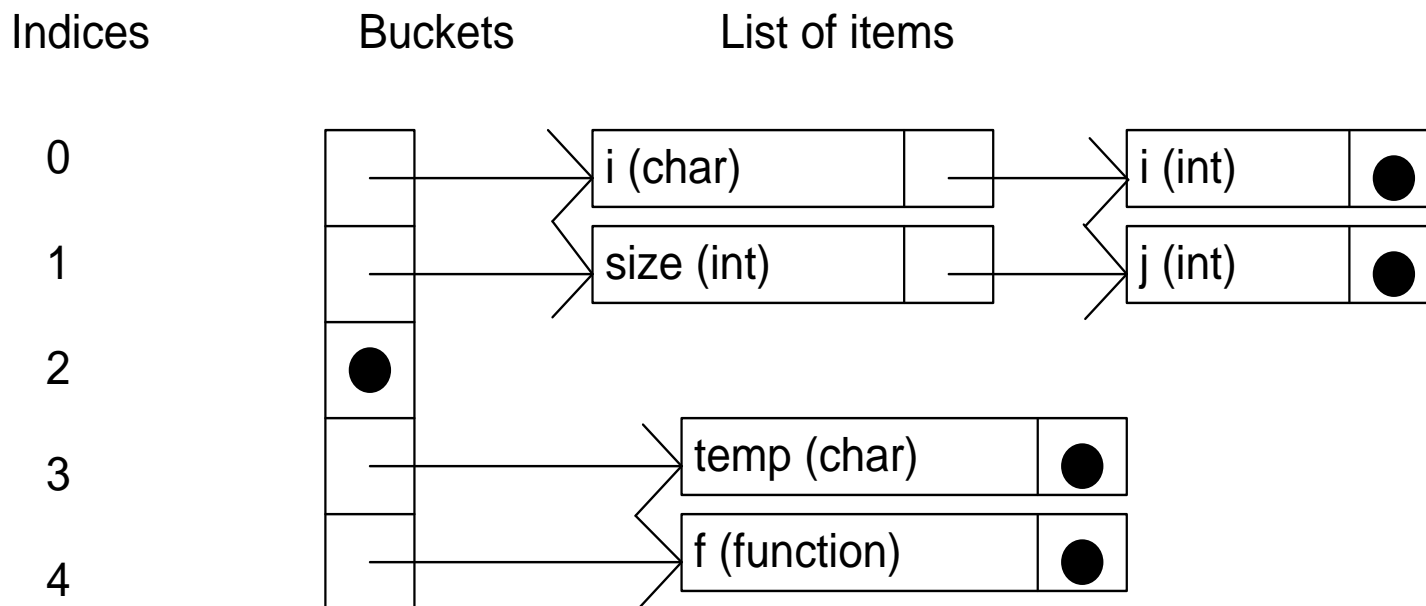
Program Ex;
var i , j : integer;
function f (size: integer): integer;
var i , temp: char;
    procedure g;
    var j: real;
    begin
        ...
    end;
    procedure h;
    var j: ^char;
    begin
        ...
    end;

```

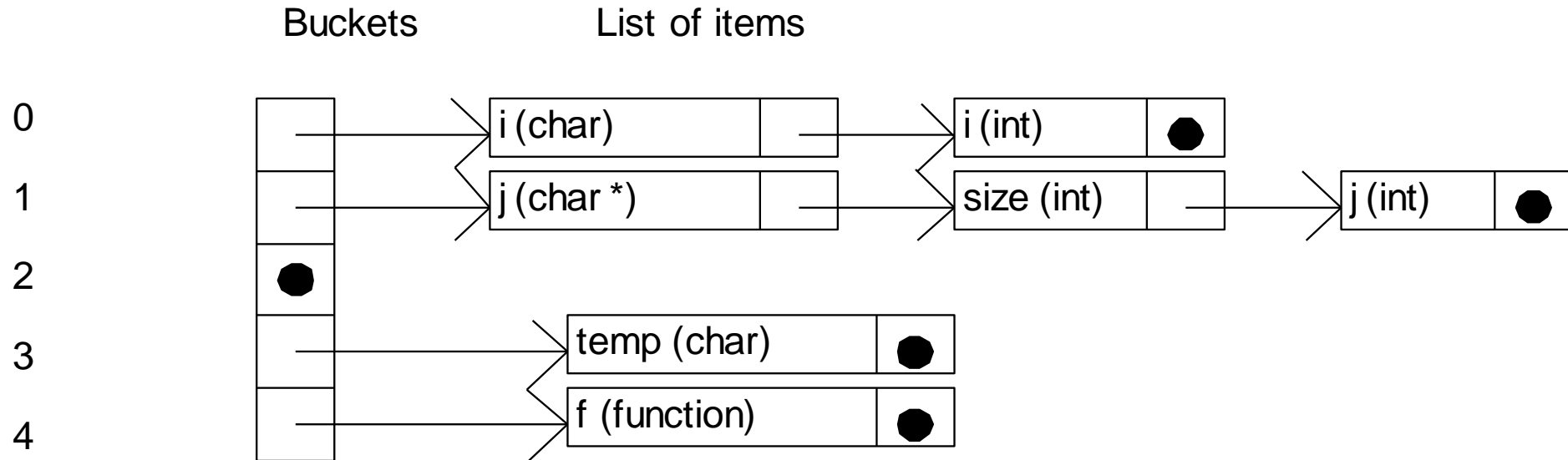
```

begin ( * f *)
    ...
end;
begin ( * main program *)
    ...
end.

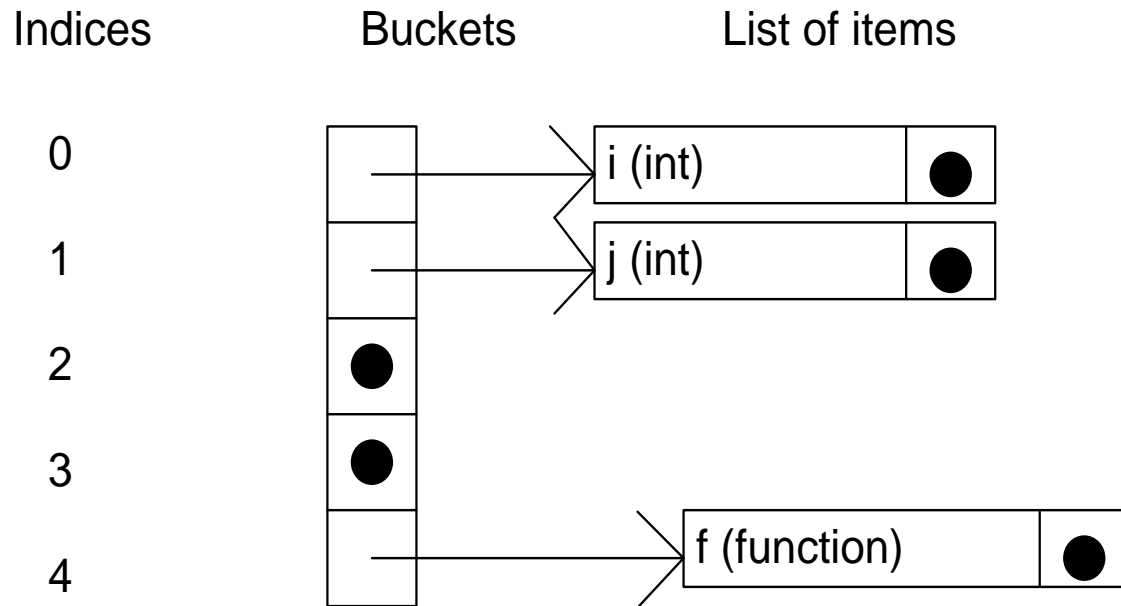
```



(a) After processing the declarations of the body of f



(b) After processing the declarations of the second nested compound statement within the body of f



(c) After exiting the body of f (and deleting its declarations)