## **Contents**

3.3 AP/Bridge

## 1 Basic

## 1.1 Default Code

```
#include <bits/stdc++.h>
#define int long long
// #pragma GCC target("popent")
// #pragma GCC optimize("03")
using namespace std;

void solve() {
}

signed main() {
  ios_base::sync_with_stdio(false);
  cin.tie(nullptr);
  int tt = 1;
  cin >> tt;
  while (t--) {
     solve();
  }
  return 0;
}
```

# 2 Math

#### 2.1 快速冪

#### | // 根據費馬小定

```
理,若 a p 互質, a^(p-2) 為 a 在 mod p 時的乘法逆元
int fast_pow(int a, int b, int mod)
{
    // a^b % mod
    int res = 1;
    while(b)
    {
        if(b & 1) res = (res * a) % mod;
        a = (a * a) % mod;
        b >>= 1;
    }
    return res;
```

#### 2.2 擴展歐幾里得

```
int gcd(int a, int b)
{
    return b == 0 ? a : gcd(b, a % b);
}
int lcm(int a, int b)
{
    return a * b / gcd(a, b);
}

pair<int, int> ext_gcd
    (int a, int b) //擴展歐幾里德 ax+by = gcd(a,b)
{
    if (b == 0)
        return {1, 0};
    if (a == 0)
        return {0, 1};
    int x, y;
    tie(x, y) = ext_gcd(b % a, a);
    return make_pair(y - b * x / a, x);
}
```

# 3 Graph

# 3.1 Tarjan SCC

```
class tarjan{
    // 1-base
    int time = 1;
    int id = 1;
    stack<int> s;
    vector<int> low;
    vector<int> dfn;
    vector<bool> in_stack;
    void dfs(int node, vector<vector<int>> &graph){
      in_stack[node] = true;
      s.push(node);
      dfn[node] = low[node] = time++;
      for(auto &j : graph[node]){
        if(dfn[j] == 0){
          dfs(j, graph);
          // 看看往下有沒有辦法回到更上面的點
          low[node] = min(low[node], low[j]);
        else if(in stack[j]){
          low[node] = min(low[node], low[j]);
      }
      vector < int > t; // 儲存這個強連通分量
      if(dfn[node] == low[node]){
        while(s.top() != node){
          t.push_back(s.top());
          in_stack[s.top()] = false;
          scc_id[s.top()] = id;
          s.pop();
        t.push_back(s.top());
        scc_id[s.top()] = id;
        in_stack[s.top()] = false;
        s.pop();
      if(!t.empty()) ans.push_back(t);
  public:
    vector<int> scc_id;
    vector<vector<int>> ans:
    // ans ans[i] 代表第 i 個強連通分量裡面包涵的點
    // scc_id[i] 代表第 i 個點屬於第幾個強連通分量
        <vector<int>> scc(vector<vector<int>> &graph){
      int num = graph.size();
      scc_id.resize(num, -1);
      dfn.resize(num, 0);
      low.resize(num, 0);
      in_stack.resize(num, false);
      for(int i = 1; i < num; i++){</pre>
        if(dfn[i] == 0) dfs(i, graph);
      return ans;
};
```

```
3.2 2 SAT
   下面的 tarjan scc 算法來解 2 sat 問題,若 事件 a 發
   生時,事件 b 必然發生,我們須在 a \rightarrow b 建立一條有向
   cses 的 Giant Pizza 來舉例子,給定 n 個人 m 個配料
   表,每個人可以提兩個要求,兩個要求至少要被滿足一個
// 3 5
// + 1 + 2
// - 1 + 3
// + 4 - 2
// 以這
   個例子來說,第一個人要求要加 配料1 或者 配料2 其中
   一項,第二個人要求不要 配料1 或者 要配料3 其中一項
// 試問能不能滿足所有人的要求,我們可以把 要加
   配料 i 當作點 i ,不加配料 i 當作點 i + m(配料數量)
// 關於第一個人的要求 我們可以看成若不加 配
   料1 則必定要 配料2 以及 若不加 配料2 則必定要 配料1
// 關於第二個人要求 可看做加了 配料
   1 就必定要加 配料3 以及 不加 配料3 就必定不加 配料1
```

```
// 以這些條件建立有像圖,並且
    找尋 scc ,若 i 以及 i + m 在同一個 scc 中代表無解
// 若要求解,則若 i 的 scc_id
     小於 i + m 的 scc_id 則 i 為 true , 反之為 false
// tarjan 的模板在上面
cin >> n >> m;
vector<vector<int>> graph(m * 2 + 1);
function < int(int) > tr = [&](int x){
  if(x > m) return x - m;
  return x + m;
for(int i = 0; i < n; i++){</pre>
 char c1, c2;
  int a, b;
  cin >> c1 >> a >> c2 >> b;
  // a 代表 a 為真, m + a 代表 a 為假
 if(c1 == '-') a += m;
if(c2 == '-') b += m;
  graph[tr(a)].push_back(b);
 graph[tr(b)].push_back(a);
tarjan t;
auto scc = t.scc(graph);
for(int i = 1; i <= m; i++){</pre>
  if(t.scc id[i] == t.scc id[tr(i)]){
    cout << "IMPOSSIBLE\n";</pre>
    return 0;
}
for(int i = 1; i <= m; i++){</pre>
  if(t.scc_id[i] < t.scc_id[tr(i)]){</pre>
    cout << '+';
  else cout << '-';</pre>
  cout << ' ';
cout << '\n';
```

# 3.3 AP/Bridge

```
// adj[u] = adjacent nodes of u
// ap = AP = articulation points
// p = parent

// disc[u] = discovery time of u

// low[u] = 'low' node of u
int dfsAP(int u, int p) {
  int children = 0;
  low[u] = disc[u] = ++Time;
  for (int& v : adj[u]) {
    if (v == p) continue; //
         we don't want to go back through the same path.
                            // if we go back is because
                                we found another wav back
    if (!disc
         [v]) { // if V has not been discovered before
      children++;
      dfsAP(v, u); // recursive DFS call
      if (disc[u] <= low[v]) // condition #1</pre>
        ap[u] = 1;
      low[u] = min(low[u],
           low[v]); // low[v] might be an ancestor of u
    } else // if v was already
          discovered means that we found an ancestor
      low[u] = min(low[u], disc[v]); // finds
            the ancestor with the least discovery time
  return children;
void AP() {
  ap = low = disc = vector<int>(adj.size());
  Time = 0:
  for (int u = 0; u < adj.size(); u++)</pre>
    if (!disc[u])
      ap[u] = dfsAP(u, u) > 1; // condition #2
// br = bridges, p = parent
```

```
vector<pair<int, int>> br;
int dfsBR(int u, int p) {
  low[u] = disc[u] = ++Time;
  for (int& v : adj[u]) {
    if (v == p) continue; //
        we don't want to go back through the same path.
                           // if we go back is because
                               we found another way back
    if (!disc
        [v]) { // if V has not been discovered before
      dfsBR(v, u); // recursive DFS call
      if (disc
          [u] < low[v]) // condition to find a bridge</pre>
        br.push_back({u, v});
      low[u] = min(low[u],
           low[v]); // low[v] might be an ancestor of u
    } else // if v was already
         discovered means that we found an ancestor
      low[u] = min(low[u], disc[v]); // finds
           the ancestor with the least discovery time
 }
}
void BR() {
  low = disc = vector<int>(adj.size());
  Time = 0;
  for (int u = 0; u < adj.size(); u++)</pre>
    if (!disc[u])
      dfsBR(u, u)
```

#### 3.4 Max flow min cut

```
#define int long long
// Edmonds-Karp Algorithm Time: O(VE^2) 實際上會快一點
// 記得在 main 裡面 resize graph
// 最小割,找
    到最少條的邊切除,使得從 src 到 end 的 maxflow 為 \theta
// 枚舉所有邊 i -> j , src 可
    以到達 i 但無法到達 j , 那這條邊為最小割裡的邊之一
class edge{
  public:
    int next;
    int capicity;
    int rev;
    bool is_rev;
    edge(int _n, int _c, int _r, int _ir) :
        next(_n), capicity(_c), rev(_r), is_rev(_ir){};
};
vector<vector<edge>> graph;
void add_edge(int a, int b, int capacity){
  graph[a].push_back
      (edge(b, capacity, graph[b].size(), false));
  graph[b].
      push_back(edge(a, 0, graph[a].size() - 1, true));
int dfs(int now, int end
   , int flow, vector<pair<int, int>> &path, int idx){
  if(now == end) return flow;
  auto &e = graph[now][path[idx + 1].second];
  if(e.capicity > 0){
    auto ret = dfs(e.next
        , end, min(flow, e.capicity), path, idx + 1);
    if(ret > 0){
      e.capicity -= ret;
      graph[e.next][e.rev].capicity += ret;
      return ret;
   }
  }
  return 0;
vector<pair<int, int>> search_path(int start, int end){
  vector<pair<int, int>> ans;
  queue < int > q;
  vector
      <pair<int, int>> parent(graph.size(), {-1, -1});
  q.push(start);
```

```
while(!q.empty()){
    int now = q.front();
    q.pop();
    for(int i = 0; i < (int)graph[now].size(); i++){</pre>
      auto &e = graph[now][i];
      if(e.
          capicity > 0 and parent[e.next].first == -1){
        parent[e.next] = {now, i};
        if(e.next == end) break;
        q.push(e.next);
      }
   }
  if(parent[end].first == -1) return ans;
  int now = end;
  while(now != start){
    auto [node, idx] = parent[now];
    ans.emplace_back(node, idx);
    now = node;
  ans.emplace_back(start, -1);
  reverse(ans.begin(), ans.end());
  return ans;
int maxflow(int start, int end, int node_num){
  int ans = 0;
  while(1){
    vector < bool > visited(node_num + 1, false);
    auto tmp = search_path(start, end);
    if(tmp.size() == 0) break;
    auto flow = dfs(start, end, 1e9, tmp, 0);
    ans += flow;
  return ans;
}
```

# 4 String

# 4.1 Hash

```
vector<int> Pow(int num){
  int p = 1e9 + 7;
  vector<int> ans = {1};
  for(int i = 0; i < num; i++)</pre>
    ans.push_back(ans.back() * b % p);
  return ans;
}
vector<int> Hash(string s){
  int p = 1e9 + 7;
  vector<int> ans = {0};
  for(char c:s){
    ans.push_back((ans.back() * b + c) \% p);
  return ans;
}
// 閉區間[l, r]
int query
    (vector<int> &vec. vector<int> &pow. int l. int r){
  int p = 1e9 + 7;
  int length = r - l + 1;
       (vec[r + 1] - vec[l] * pow[length] % p + p) % p;
}
```

#### 4.2 Zvalue

# 5 Geometry

### 5.1 Static Convex Hull

```
#define mp(a, b) make_pair(a, b)
#define pb(a) push_back(a)
#define F first
#define S second
template < typename T>
pair<T, T> operator -(pair<T, T> a, pair<T, T> b){
    return mp(a.F - b.F, a.S - b.S);
}
template < typename T>
T cross(pair<T, T> a, pair<T, T> b){
    return a.F * b.S - a.S * b.F;
}
template < typename T >
vector<pair
    <T, T>> getConvexHull(vector<pair<T, T>>& pnts){
    sort(pnts.begin
         (), pnts.end(), [](pair<T, T> a, pair<T, T> b)
    { return
         a.F < b.F || (a.F == b.F && a.S < b.S); });
    auto cmp = [&](pair<T, T> a, pair<T, T> b)
    { return a.F == b.F && a.S == b.S; };
    pnts.erase(unique
        (pnts.begin(), pnts.end(), cmp), pnts.end());
    if(pnts.size()<=1)</pre>
        return pnts;
    int n = pnts.size();
    vector<pair<T, T>> hull;
    for(int i = 0; i < 2; i++){</pre>
        int t = hull.size();
        for(pair<T, T> pnt : pnts){
             while(hull.size() - t >= 2 &&
    cross(hull.back() - hull[hull.size() -
                 2], pnt - hull[hull.size() - 2]) <= 0){
                 hull.pop_back();
             hull.pb(pnt);
        hull.pop_back();
        reverse(pnts.begin(), pnts.end());
    return hull:
}
```

## 6 Data Structure

### 6.1 Sparse Table

```
class Sparse_Table{
 // 0-base
 // 要改成找最大把min換成max就好
 private:
  public:
    int spt[500005][22][2];
   Sparse Table(vector<int> &ar){
     int n = ar.size();
     for (int i = 0; i < n; i++){</pre>
         spt[i][0][0] = ar[i];
         // spt[i][0][1] = ar[i];
     for (int j = 1; (1 << j) <= n; j++) {</pre>
       for (int i = 0; (i + (1 << j) - 1) < n; i++) {</pre>
         spt[i][j][0] = min(spt[i + (1 <<
         }
     }
    int query_min(int l, int r)
     if(l>r) return INT_MAX;
     int j = (int)__lg(r - l + 1);
     ///j = 31 - \_builtin_clz(r - l+1);
      return min
         (spt[l][j][0], spt[r - (1 << j) + 1][j][0]);
   int query_max(int l, int r)
     if(l>r) return INT_MAX;
```

### 6.2 Segement Tree

```
template < typename T>
class segment_tree
  // 1-base4
 private:
 public:
    template < typename F>
    class node{
     public:
      int lb, rb;
      F num, tag;
      node<F> *left, *right;
      node(){
        tag = 0;
        right = nullptr, left = nullptr;
      T rv(){
       return num + tag * (rb - lb + 1);
      void pull(){
        if(left) left -> tag += tag;
        if(right) right -> tag += tag;
        num = rv();
        tag = 0;
     }
    };
    node < T > * root;
    node<T> *build(vector<T> &save, int l, int r){
      node<T> *temp = new node<T>;
      temp -> lb = l;
      temp -> rb = r;
      if (l == r)
      {
        temp -> num = save[l];
        return temp;
      int mid = (l + r) / 2;
      temp -> left = build(save, l, mid);
      temp -> right = build(save, mid + 1, r);
      node<T> *left_node, *right_node;
left_node = temp -> left;
      right_node = temp -> right;
      temp ->
           num = left_node -> num + right_node -> num;
      return temp:
    T query(int l, int r, node<T> *t){
      t -> pull();
      if(l == t -> lb and r == t -> rb)
         return t -> num:
      int mid = (t -> lb + t -> rb) / 2;
      if(r <= mid) return query(l, r, t -> left);
      else if(l > mid) return query(l, r, t -> right);
      else return query(l, mid
          , t -> left) + query(mid + 1, r, t -> right);
    void modify_node(int index, T delta, node<T> *t){
      if(t \rightarrow lb == t \rightarrow rb)
        t -> num += delta;
       return;
      int mid = (t -> lb + t -> rb) / 2;
         > mid) modify_node(index, delta, t -> right);
      else modify_node(index, delta, t -> left);
      t -> num += delta;
    void modify_scope
        (int lb, int rb, int delta, node<T> *t){
      t -> tag += delta;
       return;
      int mid = (t -> lb + t -> rb) / 2;
      if(t -> left and rb <=</pre>
           mid) modify_scope(lb, rb, delta, t -> left);
```

```
4
      else if(t -> right and lb >
          mid) modify_scope(lb, rb, delta, t -> right);
      else{
        modify_scope(lb, mid, delta, t -> left);
        modify_scope(mid + 1, rb, delta, t -> right);
      if(t -> left and t -> right) t ->
          num = t -> left -> rv() + t -> right -> rv();
};
signed main()
  int n, q;
  cin >> n >> q;
  vector<int> save(n + 1, 0);
  for(int i = 1; i <= n; i++){</pre>
   cin >> save[i];
  segment_tree<int> s;
  // init [1, n]
  s.root = s.build(save, 1, n);
  // modify [a, b] add c
  s.modify_scope(a, b, c, s.root);
  // query [a, b]
  s.query(a, b, s.root)
```