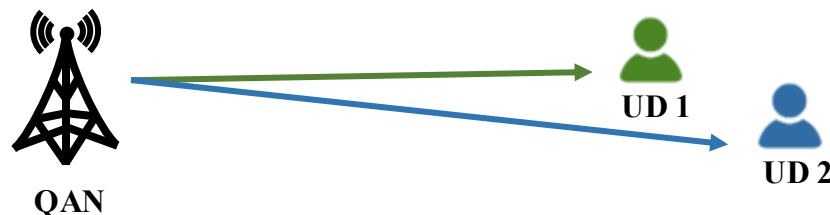


Data Structures Programming Project #1

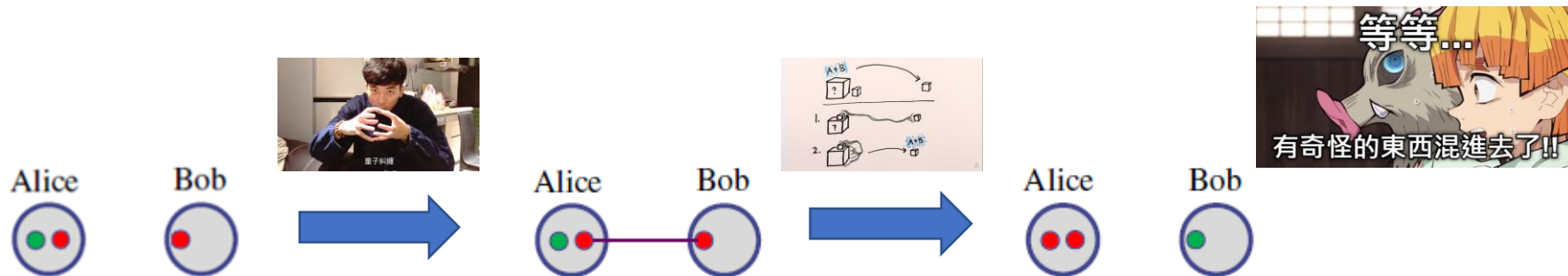
Entanglement Distribution in FSO QNs

- A quantum access node (QAN) with user devices (UDs) in free-space optical quantum networks (QNs)
- Nodes has quantum memory to store qubits
- Nodes are interconnected with light-of-sight (LoS) free-space optical (FSO) quantum channels
- Entangled pairs are distributed from QAN to UD
- Applications: teleportation, symmetric encryption



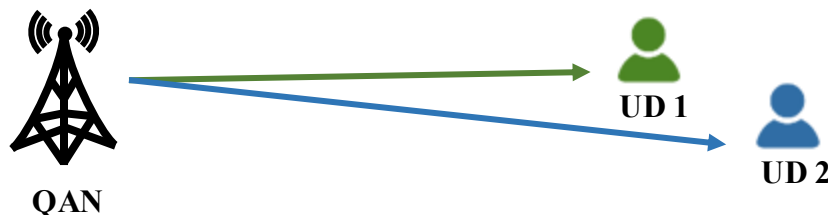
Distribution and Teleportation in FSO QNs

- **Entanglement distribution:**
Create an entangled pair between two nodes
- **Precondition:**
Two nodes each with **a unit of quantum memory** are interconnected with a **quantum channel**
- Teleportation starts as an entangled pair is ready



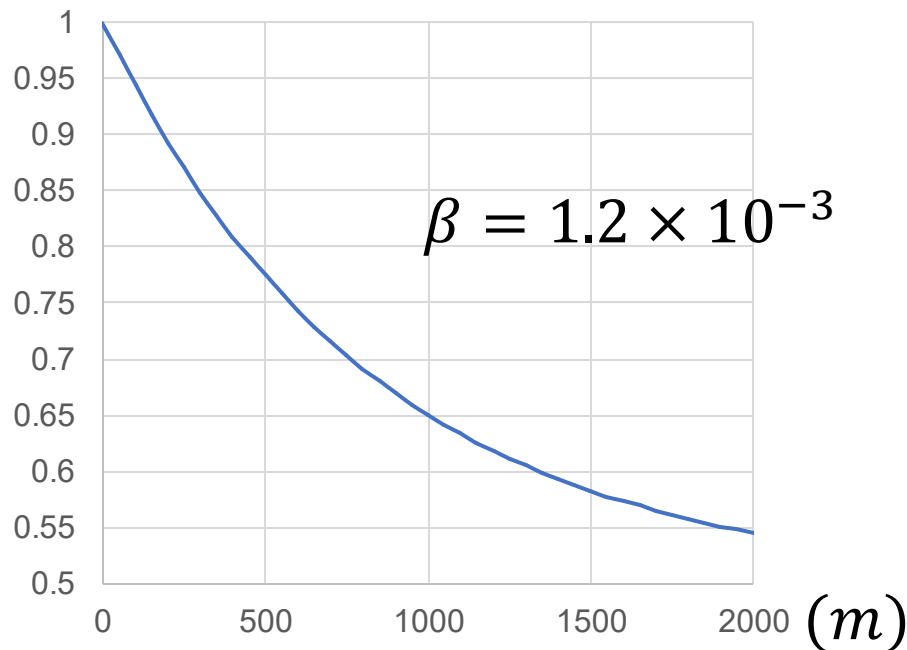
UD's Request and Requirements

- Each UD has an **all-or-nothing** request with a **profit, a fidelity threshold, and a demand expected key rate**
- **Either reject or accept it**
- Once accepted, the resulting fidelity should **meet the fidelity threshold** and the received key rate should be **no less than the demand expected key rate**

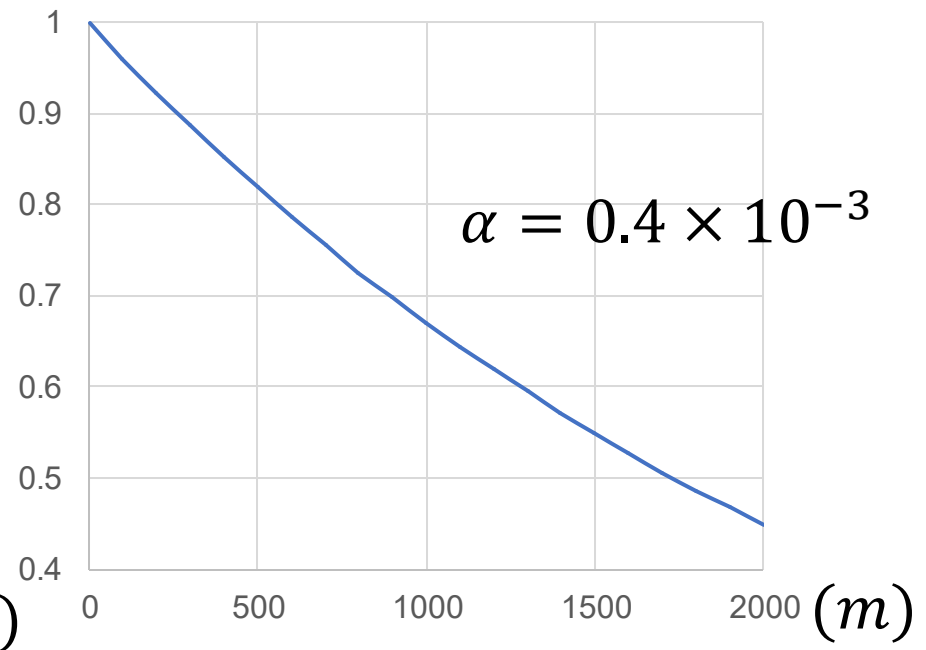


Fidelity and Probability of Entangling

- The **success probability** of establishing an entangled pair and the resulting **fidelity** are related to the **FSO distance l (m)** (here we adopt **binary state**)



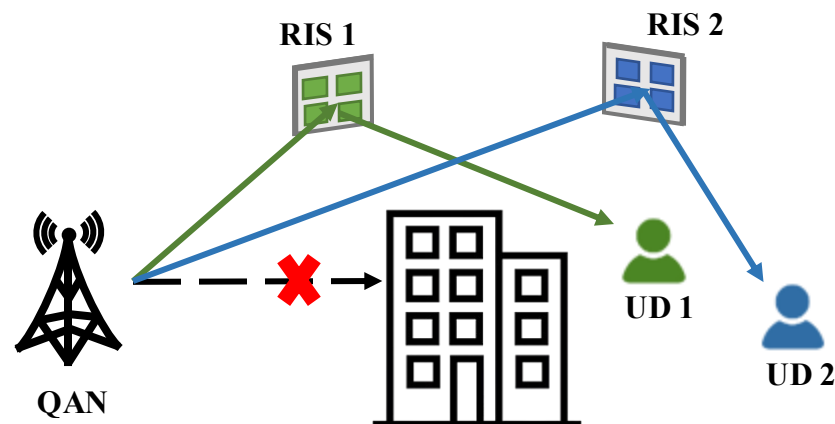
Fidelity: $\frac{1}{2} + \frac{1}{2} e^{-\beta \cdot l}$



Probability: $e^{-\alpha \cdot l}$

NLoS Challenge and RIS

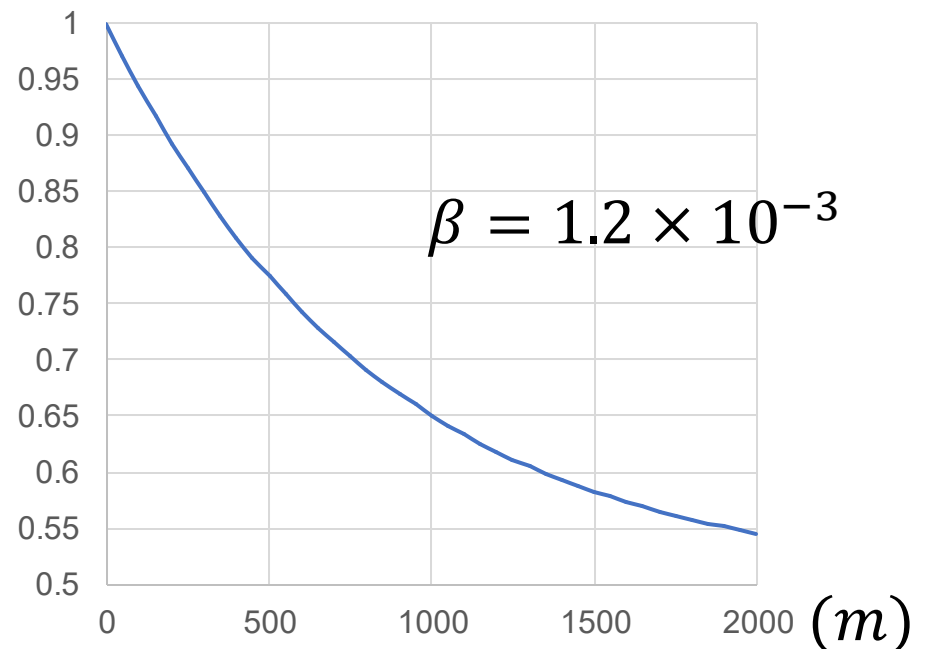
- Two nodes may be **non-line-of-sight** due to blockage by buildings or other obstacles
- A **reconfigurable intelligent surfaces (RIS)** reflects the signal via **phase shift** so that a blocked UD can still receive it (i.e., virtual LoS)
- Each RIS can be assigned to **at most one UD**



Penalty of RISs – Longer Optical Path

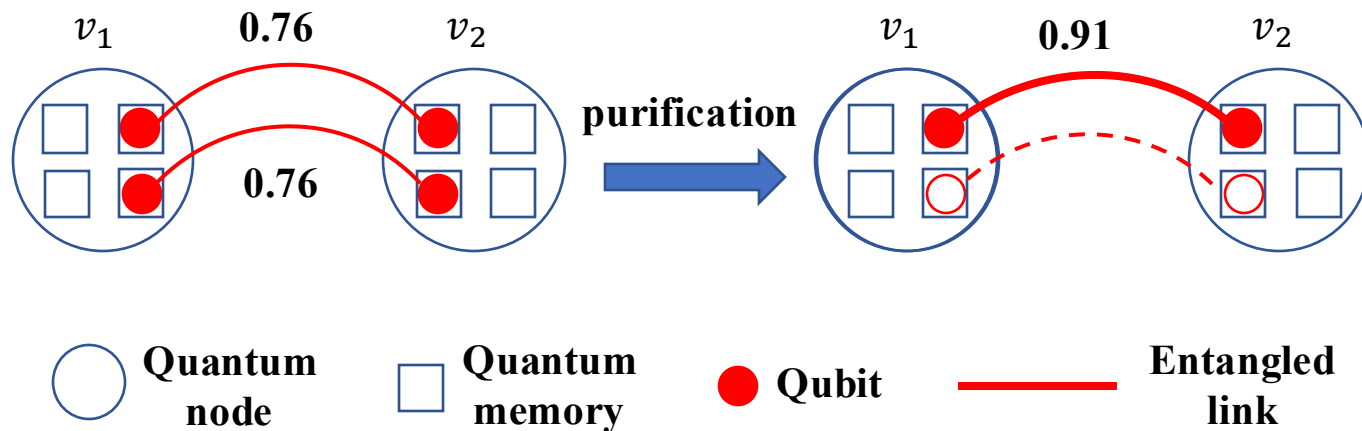
- An inherent **penalty of RIS** is the longer optical path, leading to degraded fidelity
- The fidelity could be terrible (i.e., **cannot meet the fidelity threshold**) when **the FSO distance l is overlength**

$$\text{Fidelity: } \frac{1}{2} + \frac{1}{2} e^{-\beta \cdot l}$$



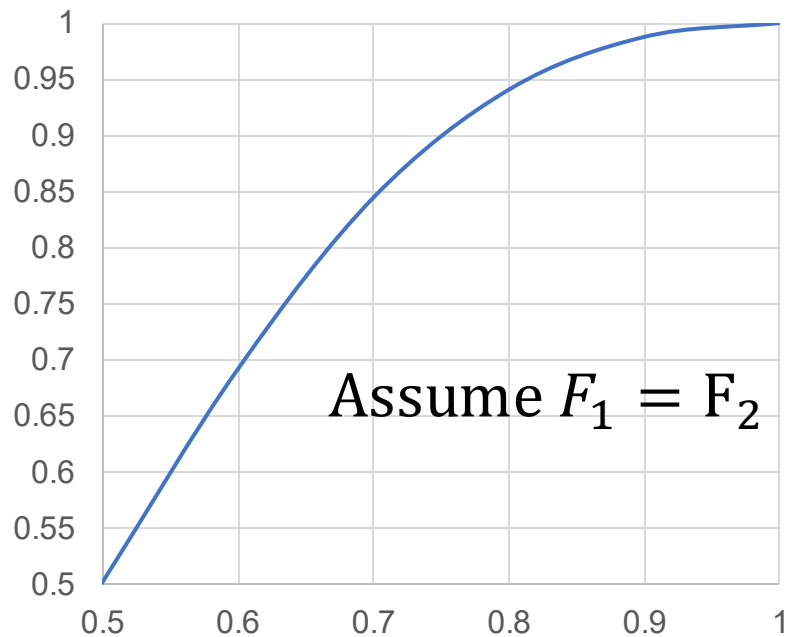
Remedy of Fidelity Loss – Purification

- An inherent **penalty of RIS** is the longer optical path, leading to degraded fidelity
- **Entanglement purification** improves fidelity at the cost of consuming an additional entangled pair

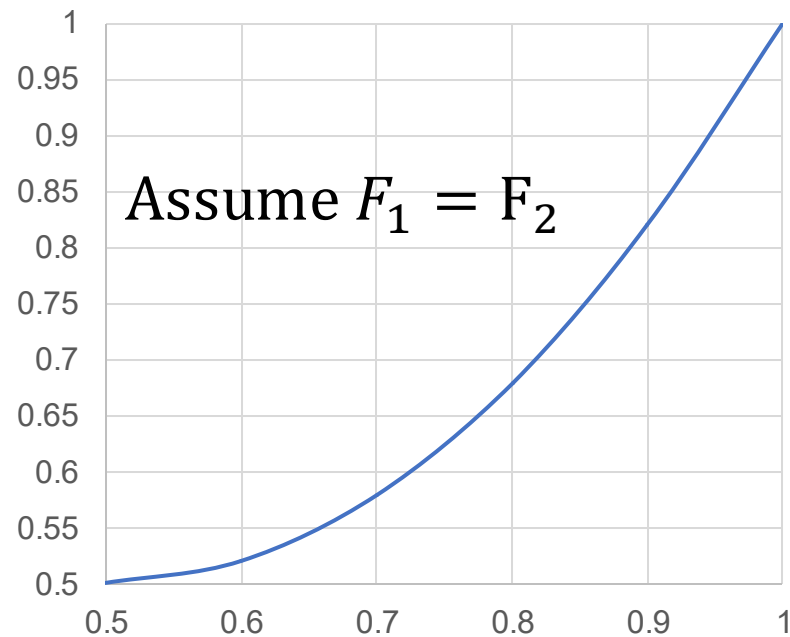


Fidelity and Probability of Purification

- Purification can **raise the fidelity** with **a probability**
- Note that we adopt **binary state**



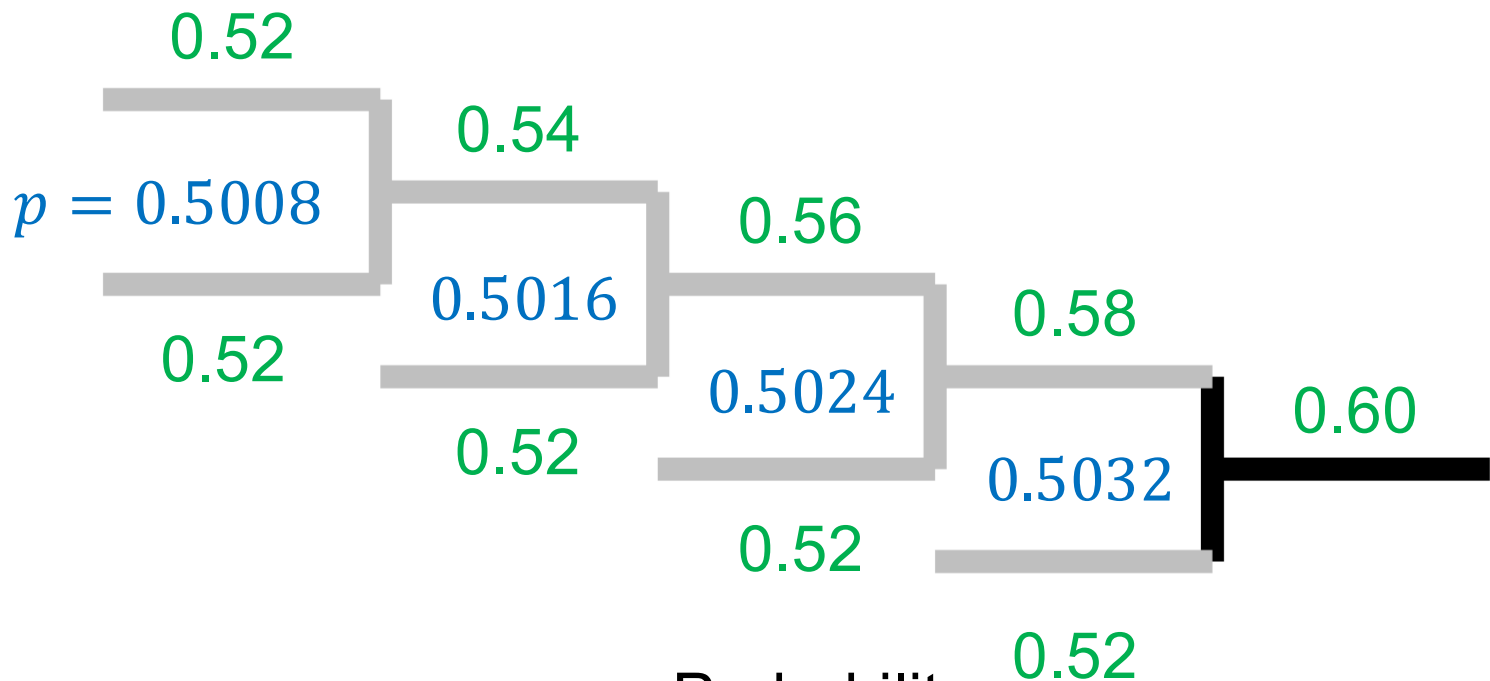
Fidelity:
$$\frac{F_1 \cdot F_2}{F_1 \cdot F_2 + (1 - F_1) \cdot (1 - F_2)}$$



Probability:
$$F_1 \cdot F_2 + (1 - F_1) \cdot (1 - F_2)$$

Fidelity and Probability of Purification

- We can conduct **multiple rounds of purification**
- For simplicity, here we adopt **pumping**



Fidelity:
$$\frac{F_1 \cdot F_2}{F_1 \cdot F_2 + (1 - F_1) \cdot (1 - F_2)}$$

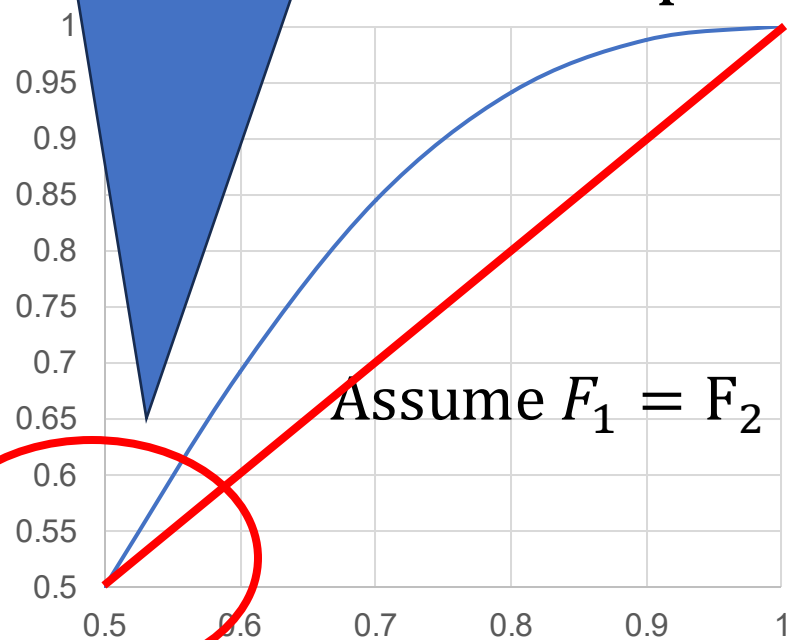
Probability:
$$F_1 \cdot F_2 + (1 - F_1) \cdot (1 - F_2)$$

Fidelity and Probability of Purification

It is difficult to fix the fidelity when the fidelity is low

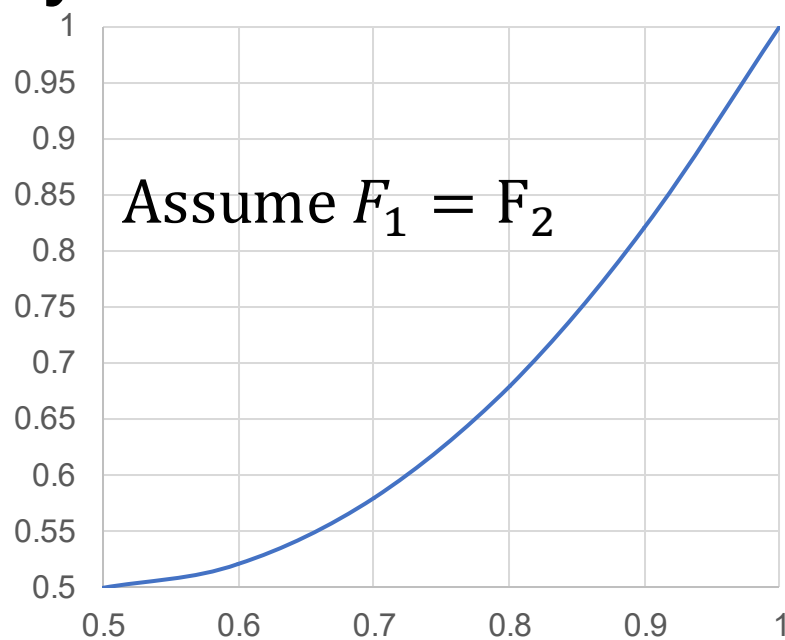
the fidelity with a probability

that we adopt binary state



Assume $F_1 = F_2$







Fidelity:
$$\frac{F_1 \cdot F_2}{F_1 \cdot F_2 + (1 - F_1) \cdot (1 - F_2)}$$



Assume $F_1 = F_2$

Probability:
$$F_1 \cdot F_2 + (1 - F_1) \cdot (1 - F_2)$$

Summary of Operations

	Fidelity	Probability
Entangling	Distance  Fidelity 	Distance  Probability 
Purification		

Summary of Formulas (1/4)

- UD: u
- Indirect path via RIS or direct path to the QAN: r
- The success probability of **e**ntangling from the QAN to UD u via path r :

$$P_{\textcolor{red}{e}}(u, r) = e^{-\alpha \cdot l(u, r)}$$

- The resulting fidelity:

$$F_{\textcolor{red}{e}}(u, r) = \frac{1}{2} + \frac{1}{2} e^{-\beta \cdot l(u, r)}$$

Summary of Formulas (2/4)

- Fidelity of two entangled pairs: F_1 and F_2
- The success probability of **p**urification on two entangled pairs with fidelity F_1 and F_2 :

$$P_{\text{p}}(F_1, F_2) = F_1 \times F_2 + (1 - F_1) \times (1 - F_2)$$

- The resulting fidelity:

$$F_{\text{p}}(F_1, F_2) = \frac{F_1 \cdot F_2}{F_1 \cdot F_2 + (1 - F_1) \cdot (1 - F_2)}$$

Summary of Formulas (3/4)

- # rounds of pumping purification: $n \in \mathbb{Z}_0^+$
- The success probability of **executing n rounds of pumping purification** between the QAN and UD u via path r :

$$P(u, r, n) = \begin{cases} P_p(F(u, r, n-1), F_e(u, r)) \cdot P(u, r, n-1), & \text{if } n \geq 1 \\ 1, & \text{if } n = 0 \end{cases}$$

- The resulting fidelity:

$$F(u, r, n) = \begin{cases} F_p(F(u, r, n-1), F_e(u, r)), & \text{if } n \geq 1 \\ F_e(u, r), & \text{if } n = 0 \end{cases}$$

Summary of Formulas (4/4)

- The fidelity threshold of UD u : $F_{th}(u) \geq 0$
- The expected number of key rate of UD u : $d(u) \geq 0$
- The minimum number of required entangled pairs to generate one purified entangled pair that meets the fidelity threshold $F_{th}(u)$:
$$n(u, r) = \min\{n \in \mathbb{Z}_0^+ \mid F(u, r, n) \geq F_{th}(u)\}$$
- Satisfying $d(u)$ requires the QAN to **roughly** provide entangled pairs to UD u at a distribution rate of
$$s(u, r) = \frac{d(u)}{P_e(u, r)} \cdot \frac{n(u, r) + 1}{P(u, r, n(u, r))}$$

Programming Project #1: Entanglement Distribution Optimization

- **Given:**
 - A QAN with a limited entanglement generation rate
 - Multiple RISs, each has its coverages (i.e., a subset of UD's)
 - Multiple UD's, each has an all-or-nothing request with a profit
 - Requirements: fidelity threshold and demand expected key rate
 - Each node has a coordinate position
 - Assume that quantum memory/channels are always sufficient
- **Goal:** maximize the total profit
- **Constraints:**
 - Each UD's request is either accepted or rejected
 - The resulting fidelity and received key rate of accepted requests must meet the requirements
- The grade is proportional to total weighted received

The Competition

- The grade is proportional to **the total profit**
- **Basic: 60 (deadline)**
 - A baseline solution (see the following pages)
 - (Meet the coding style requirements)
- **Performance ranking** (decided after the deadline)
 - [0%, 30%) (bottom): +0
 - [30%, 50%): + 5
 - [50%, 75%): + 10
 - [75%, 85%): + 15
 - [85%, 90%): + 20
 - [90%, 95%): + 25
 - [95%, 100%] (top): + 30
- **Homework assistant** (superb deadline)
 - +10

The Competition

- The grade
- **Basic: 60** (t
 - A baseline
 - (Meet the
- **Performance** (deadline)
 - [0%, 30%
 - [30%, 50%
 - [50%, 75%
 - [75%, 85%
 - [85%, 90%
 - [90%, 95%
 - [95%, 100%
- **Homework assistant** (superb deadline)
 - +10



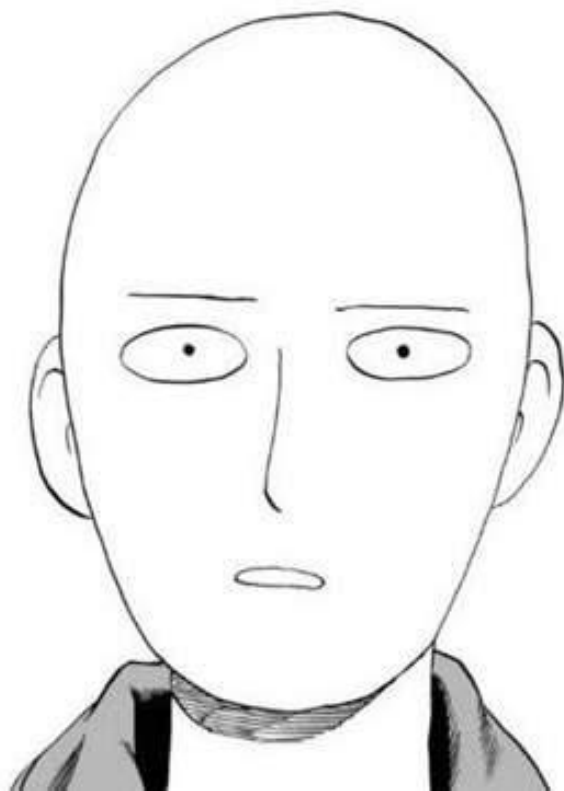
The Competition

- The grade
- **Basic: 60** (t)
 - A baseline
 - (Meet the
- **Performance**
 - [0%, 30%
 - [30%, 50%
 - [50%, 75%
 - [75%, 85%
 - [85%, 90%
 - [90%, 95%
 - [95%, 100%
- **Homework assistant** (superb deadline)
 - +10



相信你們在做完作業以後

也變強了



禿了

Baseline Algorithm

- For each UD (from ID 0 to ID $n-1$)
 - Calculate $s(u, r)$ for each available path r that can serve UD u
 - Note that $s(u, r)$ should not be greater than the residual entanglement generation rate of the QAN
 - Note that each path is either indirect via an available RIS or direct toward the QAN
 - Each available RIS is assigned to at most one UD
 - If such a path exists, accept UD u 's request and select the available path with the minimum value of $s(u, r)$ to serve UD u
 - If there is a tie, select the path via the RIS with a smaller ID
 - Once accepted, receive its profit and reduce the residual entanglement generation rate of the QAN by $s(u, r)$

Input Sample: use scanf

Format:

#UDs #RISs alpha beta Ent_Gen_Rate

QAN_X_Position QAN_Y_Position #Covered_UDs

Covered_UD1_ID Covered_UD2_ID ...

UD_ID X_Position Y_Position Profit Exp_Rate Fidelity_Th

...

RIS_ID X_Position Y_Position #Covered_UDs

Covered_UD1_ID Covered_UD2_ID ...

...

Ex:

2 1 0.4 1.2 55

500 500 0

0 1000 1000 10 100 0.7

1 0 0 8 100 0.7

0 1000 500 1

0

Output Sample: use printf

Format:

#AccUDs

Acc_UD_ID Used_RIS_ID_(if no RIS: -1)

...

Ex:

1

0 0

Note

- Superb deadline: 10/9 Thu
- Deadline: 10/16 Thu
- Pass the test of our online judge platform
- Submit your code to E-course2
- Demonstrate your code remotely or in person with TA
- C Source code (i.e., only .c)
- Show a good programming style