

Another Org-based Presentation

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Fully automated cyclic planning for
large-scale manufacturing domains.

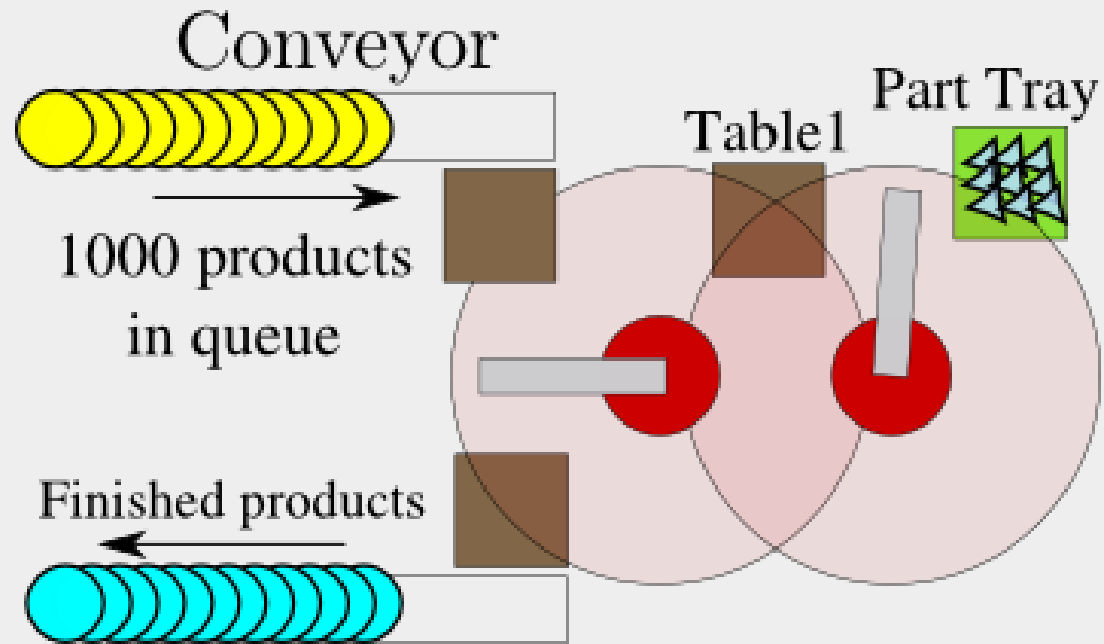
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1 Motion Planning in Cell Assembling System

- It's **NOT** about spatial search, all about **actions** !
- Require professional human resource w/o automated planner

2 CELL-ASSEMBLY



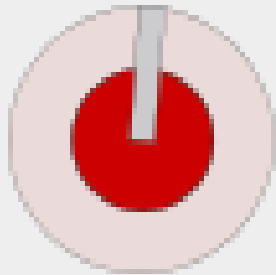
gripper + woodworking + logistics.

- : many products, with
: parts

- to complete each product, multiple operations –



Assembly

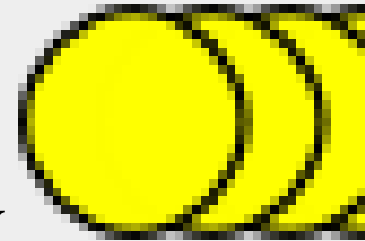


Robot arms &
range of motion

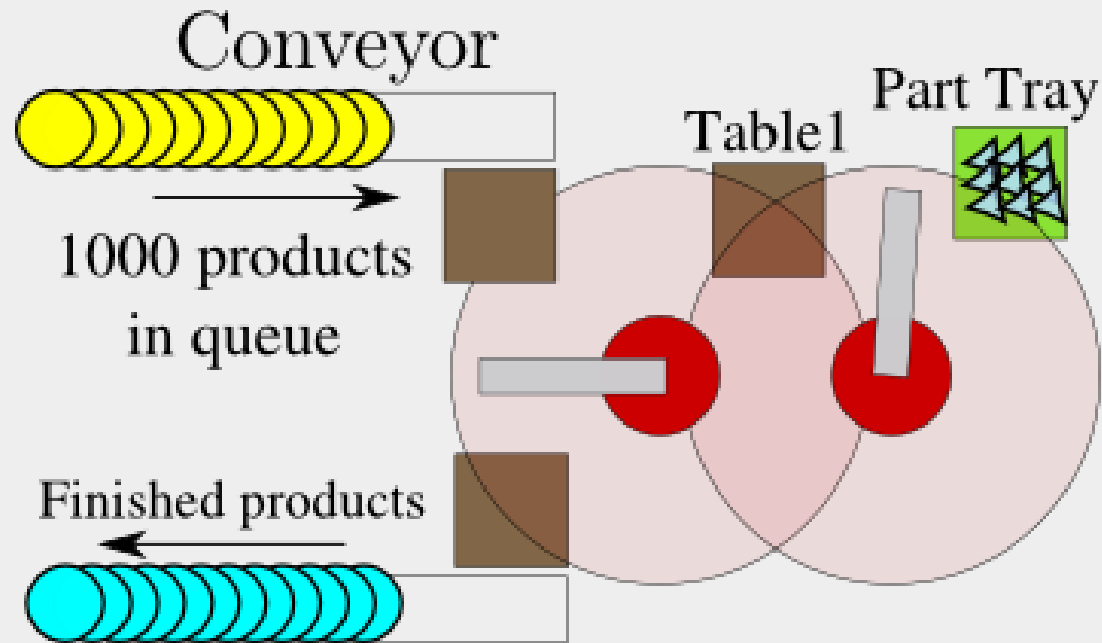
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without collision.

moves to carry



2.1 Planner's Task



Assembling recipes are provided in the problem

All products use the same recipe (Identical)

... (we tried to relax it in the KEPS paper!)

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Primary Task : **optimizing the arm motion**

2.2 High degree of symmetry

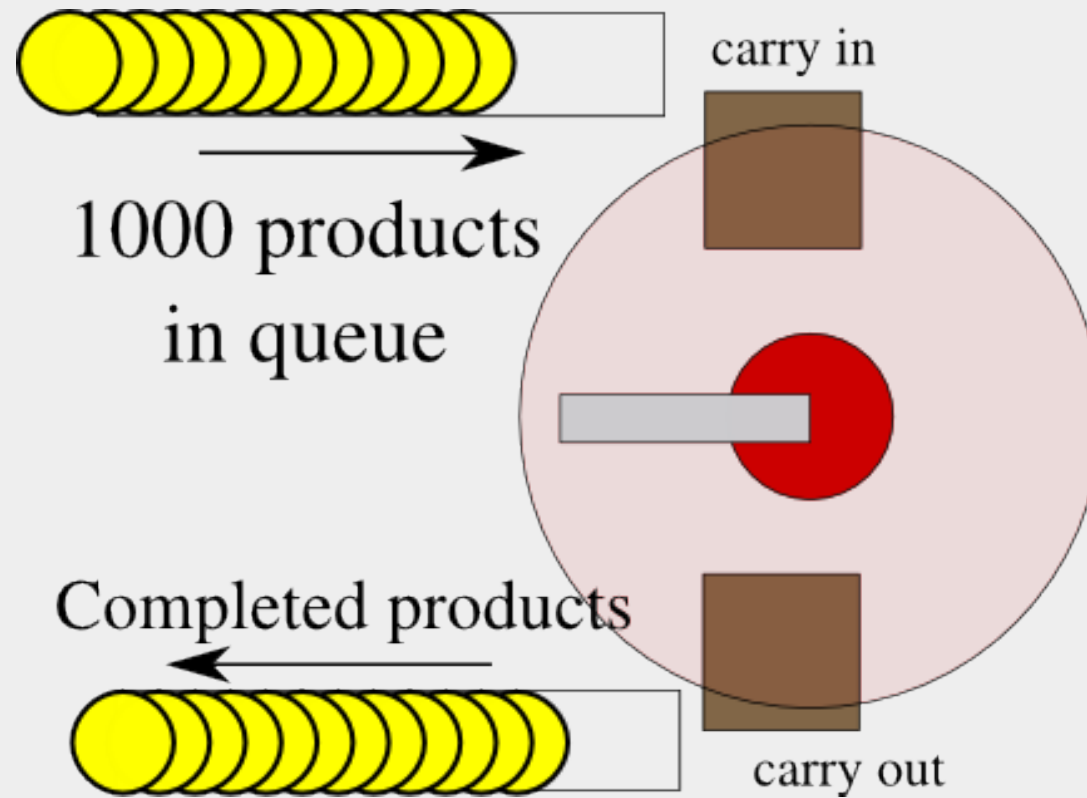


Figure 1: Simplest problem with many product

2.3 No Existing Planner Can Solve This Practical Problem!

- symmetry breaking
- $\langle - \rangle$ ANOTHER METHOD

.

Problem	# of products	Standard Planner				
		FD/LM _{cut} + scheduler	FD/LAMA + scheduler	yahsp	DAE	CPT4
	N					
CELL-ASSEMBLY	4	fail	892	807	774	fail
2a	16	fail	fail	fail	fail	fail
(2 arms, 5 jobs)	64	fail	fail	fail	fail	fail
(τ = base)	256	fail	fail	fail	fail	fail
	1024	fail	fail	fail	fail	fail
CELL-ASSEMBLY	4	249	256	607	332	fail
	16	fail	fail	fail	fail	fail
2b (1a, 5j)	64	fail	fail	fail	fail	fail
	256	fail	fail	fail	fail	fail
	1024	fail	fail	fail	fail	fail

3 Summary of Contribution

Automated Framework to Form a Loop Structure .

First attempt: form and find the *best* cyclic plan

General domain/problem analysis method (owner/lock)

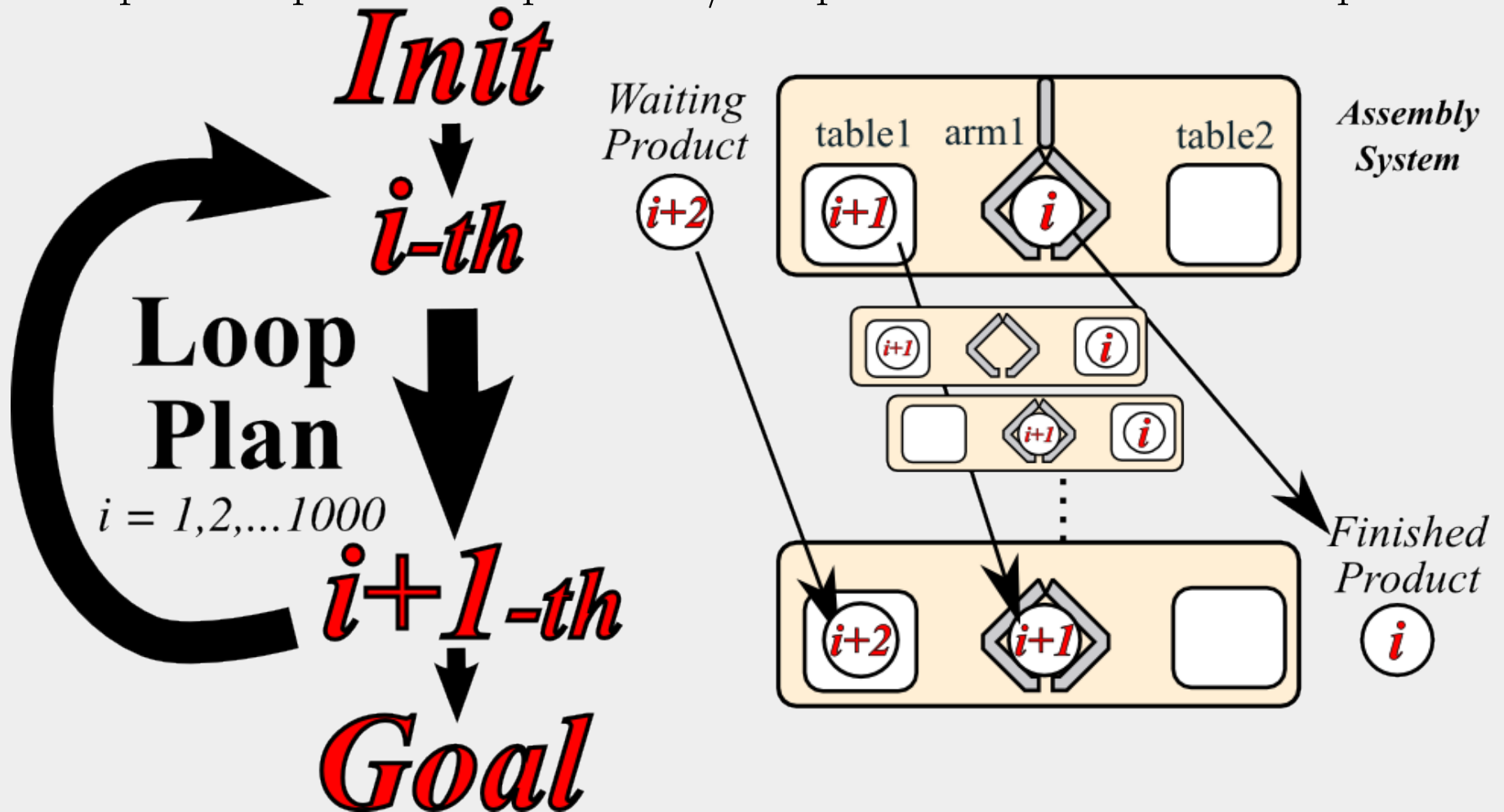
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the basic method is applicable to other domains

Solved EXTREMELY large PDDL instances (inc. IPC domain)

.
This is *clearly beyond* state-of-the-art planners

4 Strategy: Cyclic Planning

One loop – completes one product / all products to the next steps



4.1 Loop unrolling <-> Loop *rolling*

length : 1000 x 3 = **3000**

```
paint(A[0],table1)
weld(A[1],table2)
move(A[0],table1,table2)

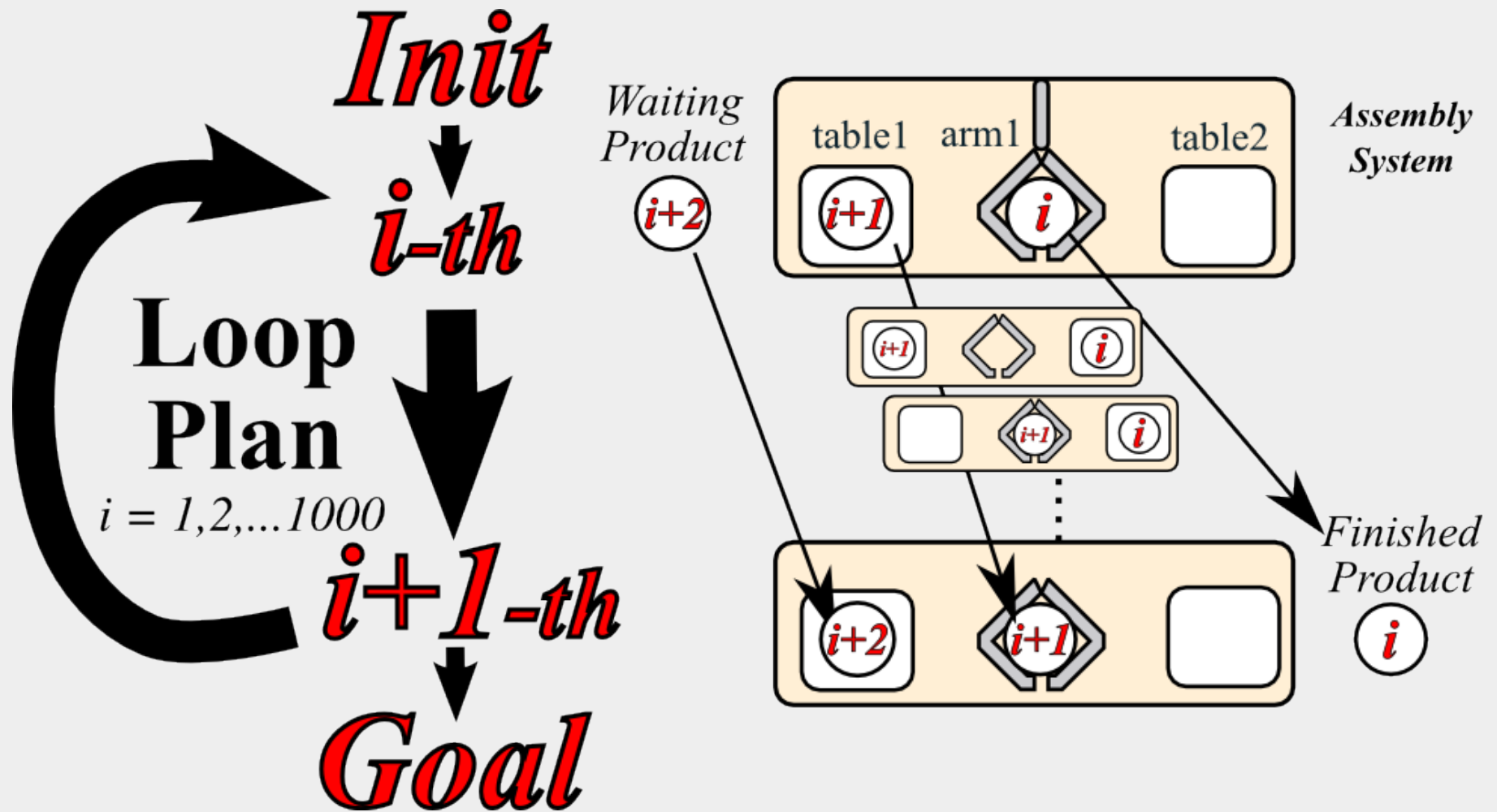
paint(A[1],table1)
weld(A[2],table2)
move(A[1],table1,table2)
...
weld(A[1000],table2)
move(A[999],table1,table2)
```

length : 3

```
for i in [0..1000]
  paint(p[i],table1)
  weld(p[i+1],table2)
  move(p[i],table1,table2)
```

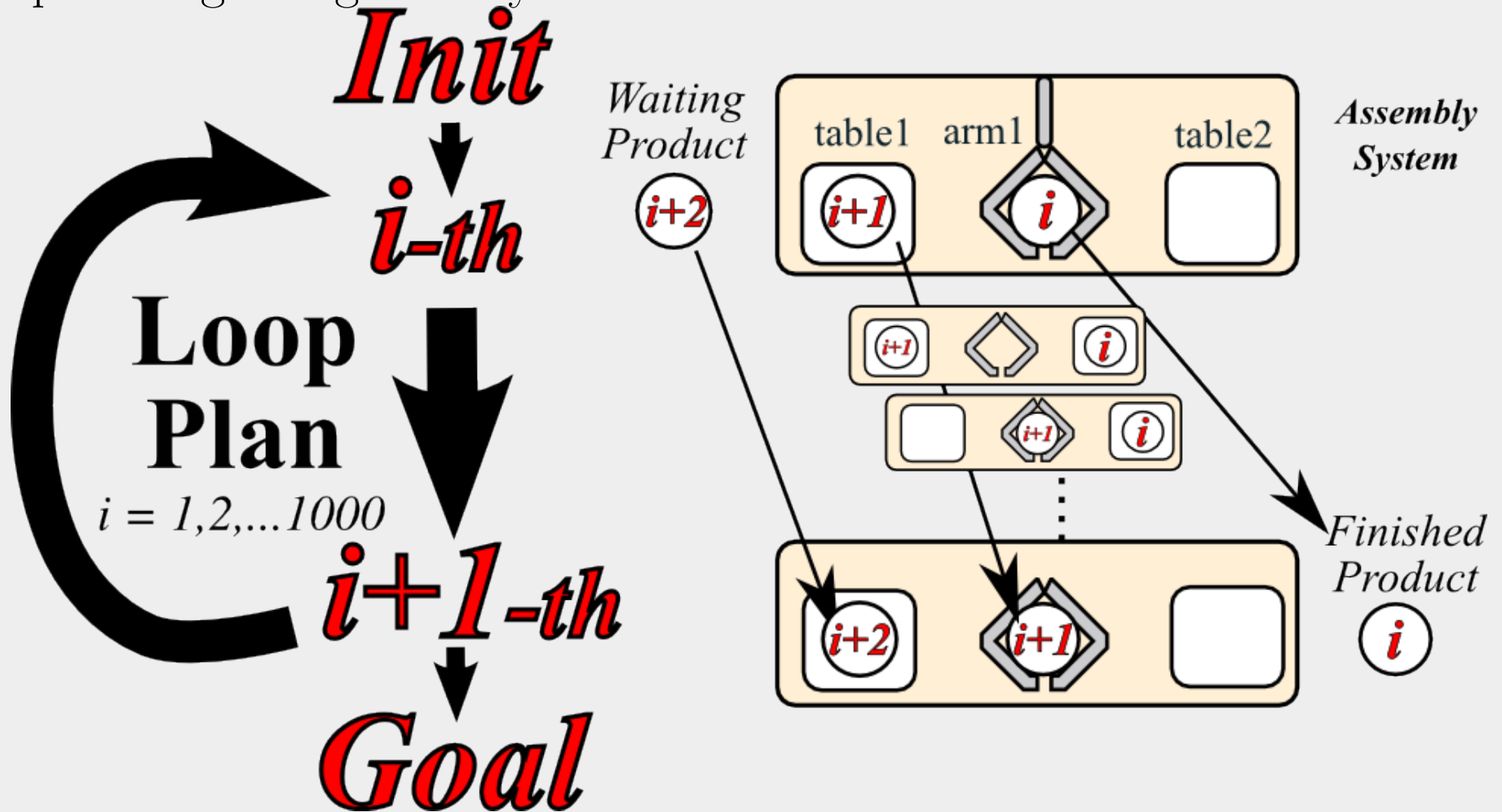
Planning is easier!

4.2 Interactions between subproblems



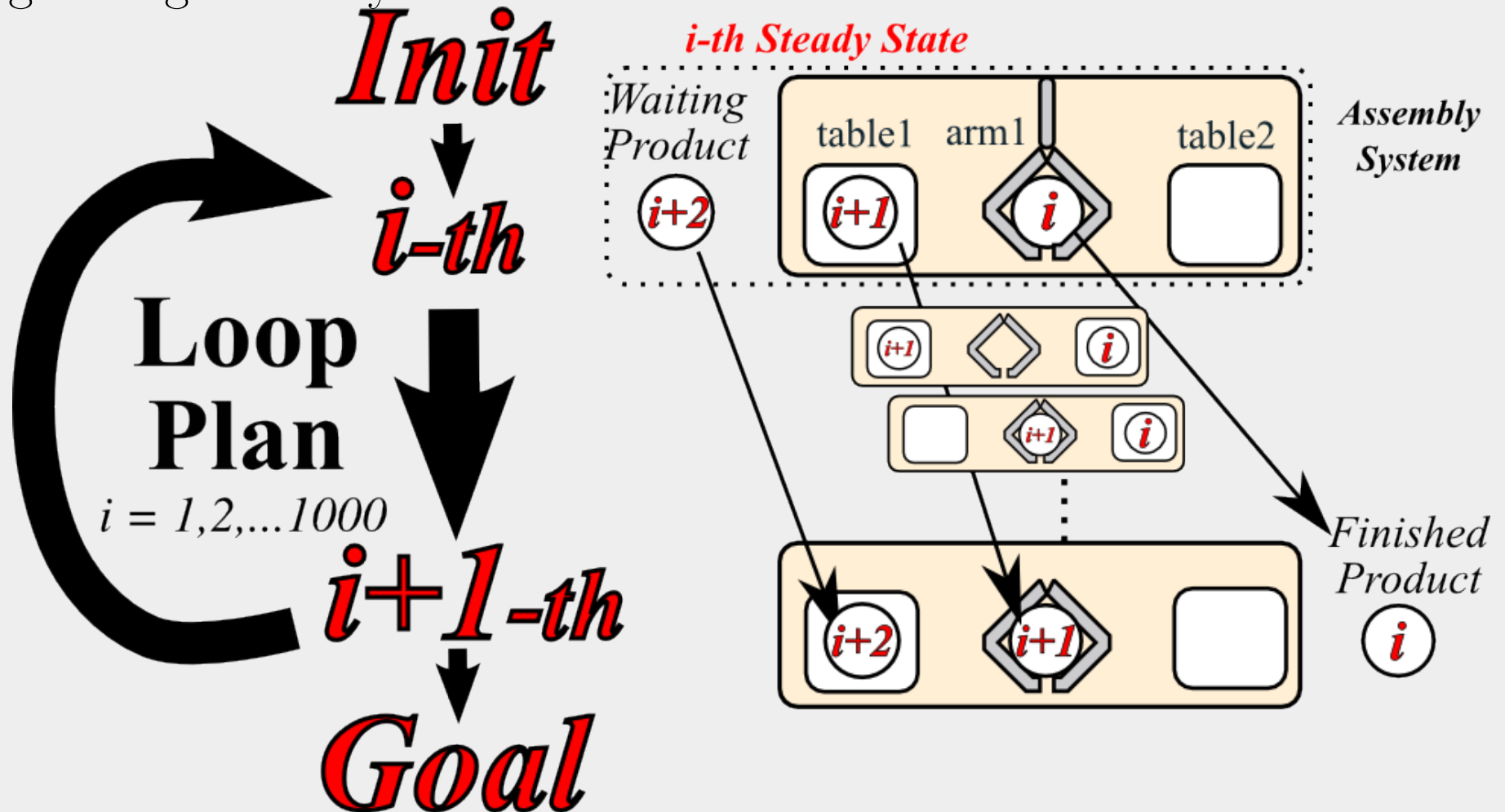
5 Loop : representation

Loop \Leftrightarrow Beginning of a Cycle

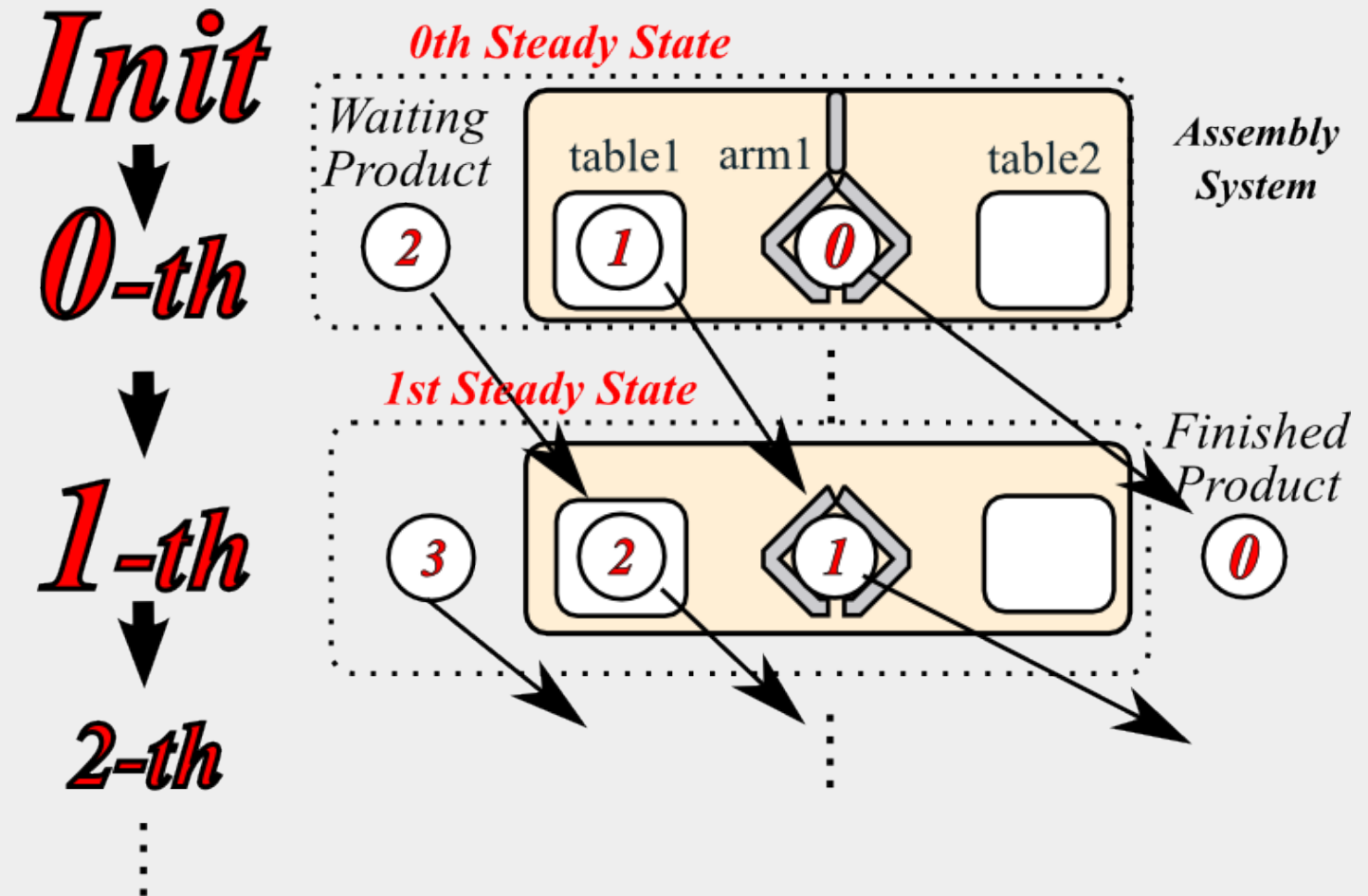


5.1 Steady State S_i

Beginneing of the cycle



5.2 Unrolling



6 Difficulty

It is not trivial to detect the information necessary for constructing an efficient loop.

6.1 Difficulty

WHICH STATES are the steady state? .

= which state can form a loop?

Checking ALL states from Init – impractical

Which SS yields a cost-efficient loop plan? .

Difficulty : (# of SSs) x (evaluation time for each SS)

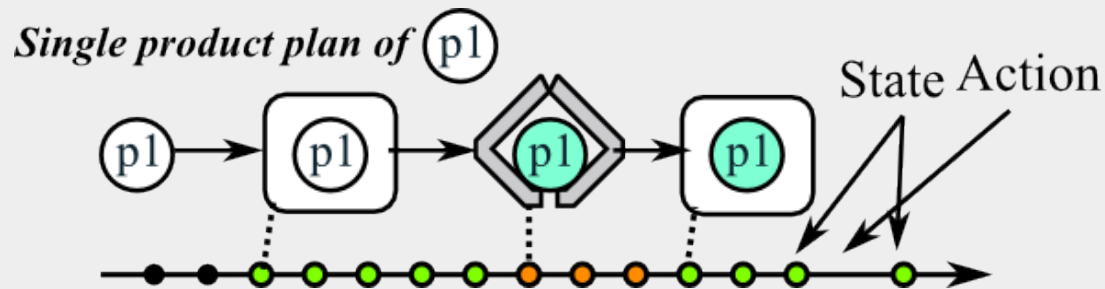
(# of SSs) – still exponentially large (e.g. 5×10^6)

(**expensive evaluation**) – loop plan : calls *FastDownward* each time

Even if we have a finite set of steady states,
checking ALL steady states is again impractical!

7 Process – to Enumerate SSs

Observation:



3 processes ($Attatch > \underline{Carry} > Attatch$)

There is at most **1** product in a process.

7.1 Process = Unit Capacity Resource

assemble @ table1 \approx use table1 (resource)

- $Process = \{\text{table, machine, painter}\} : \text{places}$
- *Which predicate specifies a resource?* is not trivial

It's *difficult even for human!* w/o analysing PDDL
semantically indistinguishable – (X A Y)

	seems like..
(color X red)	not a place
(at X table)	a place
(P X Y)	??????

7.2 Detection mechanism

Question Is a predicate $o = (P \ X \ Y \ Z \dots)$ a place?

Answer find a *lock* predicate l that satisfies **some condition** with o .

If it exists, it is a place.

Condition : for all action a ,

1 . If a occupies a place, a should check if the place is not in use, and a should acquire the lock.

2 . If a leaves the place, release the lock.

7.3 What's the difference?

```
(:action put-on
:parameters (?arm ?product ?table)
:precondition (and (in ?arm ?product)
                   (not-in-use ?table))
:effects (and (not (not-in-use ?table))
              (on ?product ?table)
              (clear ?arm)))
```

```
(:action paint
:parameters (?product ?table ?color)
:precondition (on ?product ?table)
:effects (color ?table ?color))
```

Puts a product in an arm onto a table

Paint a product

7.4 What's the difference?

```
(:action put-on
:parameters (?arm ?product ?table)
:precondition (and (in ?arm ?product)
                  (not-in-use ?table))
:effects (and (not (not-in-use ?table))
             (on ?product ?table)
             (clear ?arm)))
```

If an action **occupies a place**,
ensure the lock is free > **acquire the lock**.

```
(:action paint
:parameters (?product ?table ?color)
:precondition (on ?product ?table)
:effects (color ?table ?color))
```

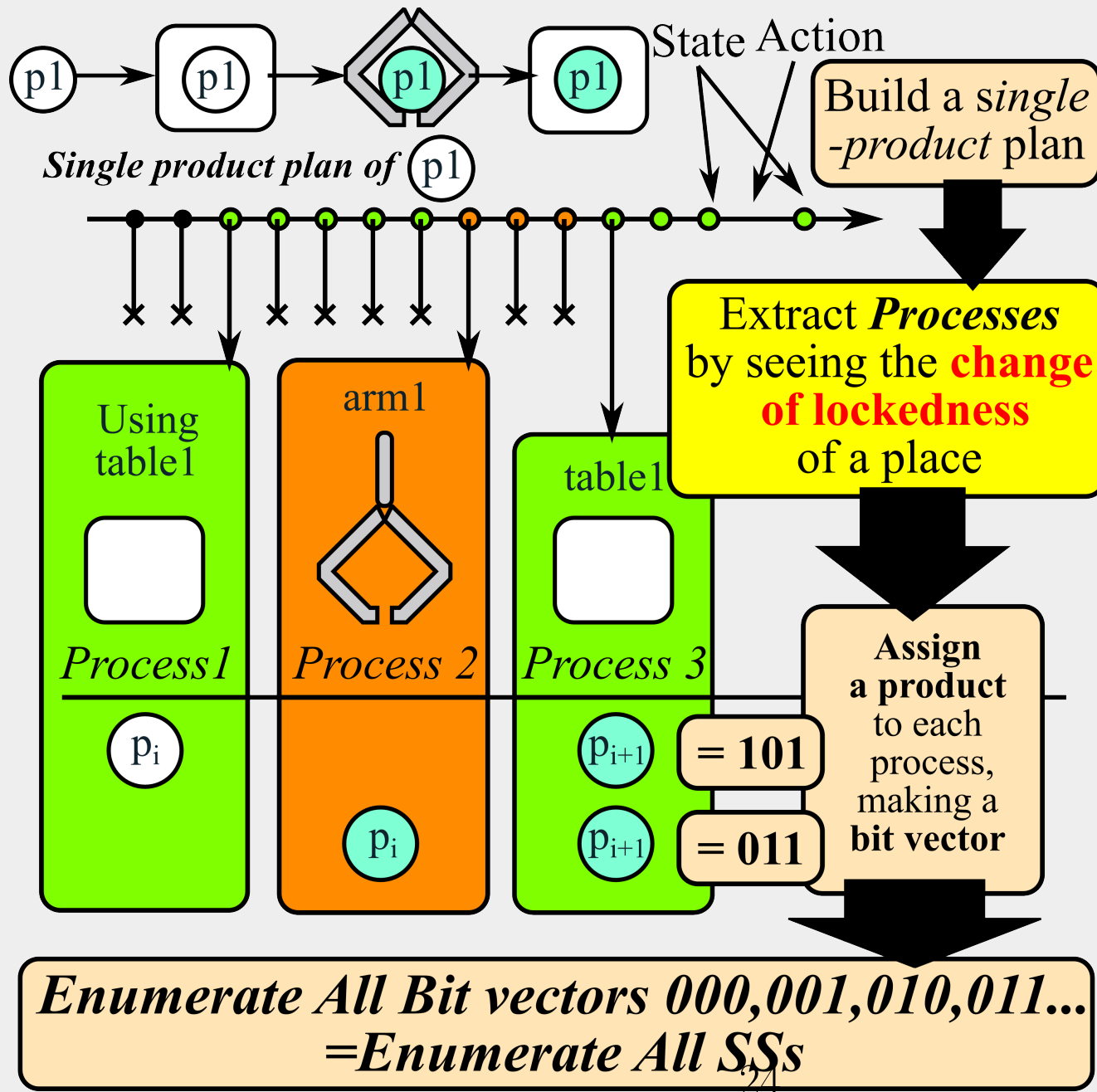
No such construct!

7.5 Detects ANY places possible

hilarious example, but...

(ტოქოლო ?აფტო ?嶺上開花 ?ឈ្នួប) – non-place

7.6 ACP : Automated Cyclic Planner



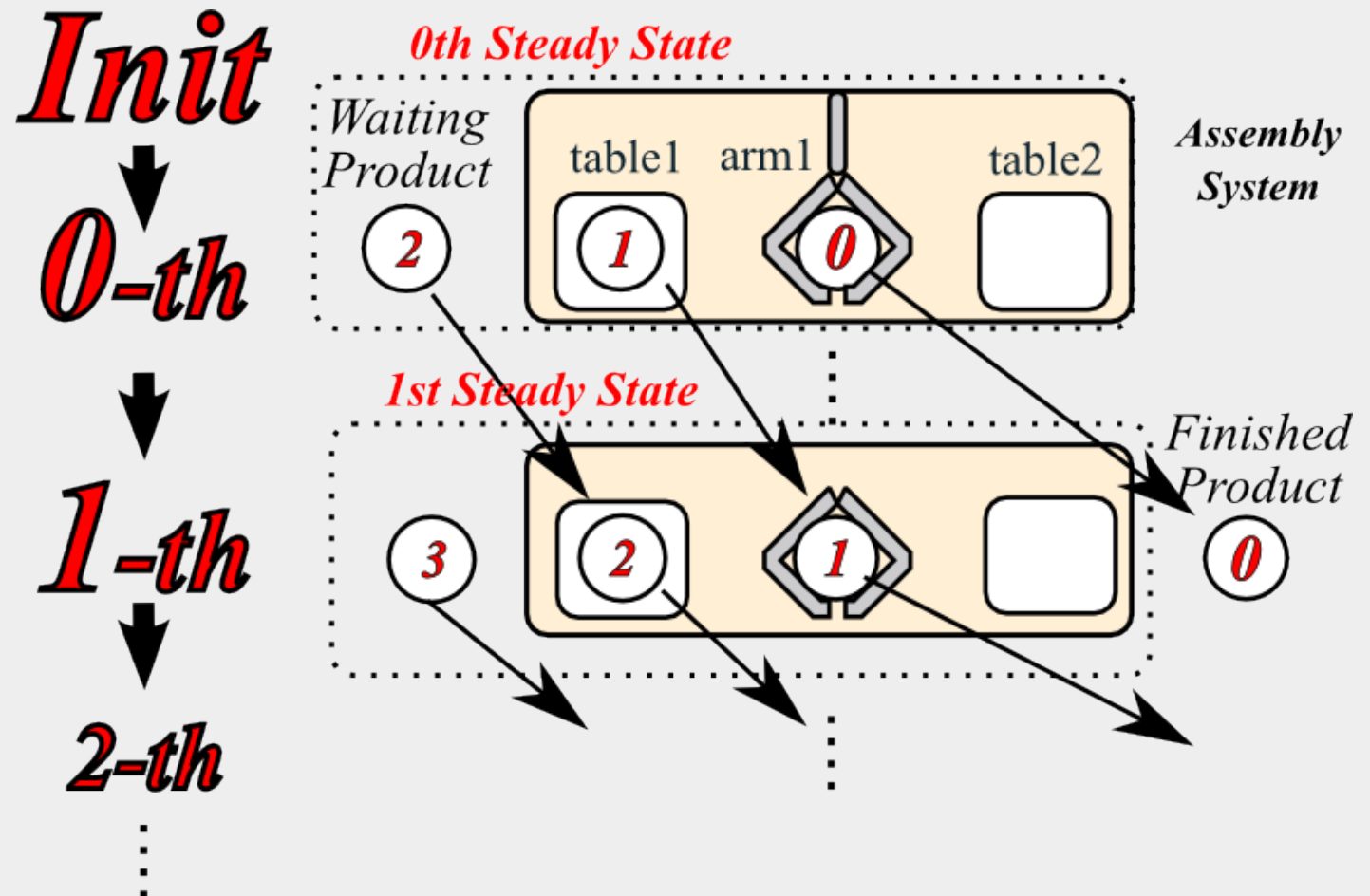
7.7 Infeasibility/Deadlock Detection

Using compact representation of steady state,

	table1	arm1	table1
Infeasible	1	0	1
Deadlock	1	1	0
Duplicated*	1	0	0
	0	1	0
(* They are the same loop)			

Reduced # : $5 \times 10^6 \rightarrow 677$

8 Unrolling



9 Experiments

- 5 CELL-ASSEMBLY problems (each with 4,16,64,256,1024 products)
- Woodworking (each with **same** 4,16,64,256,1024 parts)
- Barman (each with **same** 4,16,64,256,1024 cocktail)

9.1 ACP vs ...

- 5 temporal planners
 - FD/LAMA2011, FD/LMcut (+ min-slack scheduler)
 - yahsp2, DAEyahsp, CPT4
- best results of Simple Cyclic Planner (SCP)
 - 5 base planner x 9 configuration

9.2 Simple Cyclic Planner

Solve $K = \{1 \dots 9\}$ -products plan
with 5 planners (2FDs,yahsp,DAE,CPT)
get the average makespan per product ($K\text{-makespan}/K$)

9.3 All standard planners fails

Problem	# of products	Standard Planner				
		FD/LM _{cut} + scheduler	FD/LAMA + scheduler	yahsp	DAE	CPT4
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CELL-ASSEMBLY	4	fail	892	807	774	fail
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	16	fail	fail	fail	fail	fail
2b (1a, 5j)	64	fail	fail	fail	fail	fail
	256	fail	fail	fail	fail	fail
	1024	fail	fail	fail	fail	fail

similar results on other domains

9.4 Average makespan compared to SCP / lower bound

Problem	# of products	run-time	ACP makespan	makespan (per product)	SCP makespan (per product)	manual lbound	CPT(h2) lbound	gap (ACP / max. lbound)
	N	[sec]	c_{ACP}	c_{ACP}/N	c_{SCP}/K	l_m	l_{CPT}	
CELL- ASSEMBLY 1	4	1048	331	82.8	83 ($K = 2$)	156	176.3	1.9
	16	1049	1255	78.4	FD/LM _{cut}	624	460	2.0
	1024	1050	78871	77.0	(+ scheduler)	39936	fail	2.0
CELL- ASSEMBLY 2	4	34	246	61.5	62.3 ($K = 3$)	168	181.12	1.4
	16	33	978	61.1	FD/LM _{cut}	672	593	1.5
	1024	35	62466	61.0		43008	fail	1.5
CELL- ASSEMBLY 3	4	1893	660	165	171 ($K = 1$)	176	237	2.8
	16	1953	2352	147	FD/LAMA	704	345	3.3
	1024	1973	144480	141.1		45056	fail	3.2
CELL- ASSEMBLY 4	4	1163	318	79.5	81.3 ($K = 3$)	112	191	1.7
	16	1162	1074	67.1	FD/LM _{cut}	448	240	2.4
	1024	1165	64578	63.1		28672	fail	2.3
CELL- ASSEMBLY 5	4	1968	804	201	203 ($K = 1$)	172	335	2.4
	16	1856	2508	156.8	FD/LM _{cut}	688	532	3.6
	1024	1894	145644	142.2		44032	fail	3.3
WW product : parts	4	11	80	20	17.2 ($K = 9$)	60	80	1
	16	11	260	16.3	FD/LAMA	240	185	1.1
	1024	15	15380	15.0		15360	fail	1.0
Barman product : cocktail	4	331	35	8.8	6.3 ($K = 4$)	4	21	1.7
	16	332	179	11.2	FD/LM _{cut}	16	26	6.9
	1024	332	12275	12.0		1024	fail	12.0

9.5 Domain Independence

Problem	# of products	run-time	ACP makespan	makespan (per product)	SCP makespan (per product)	manual lbound	CPT(h2) lbound	gap (ACP / max. lbound)
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	1024	332	12275	12.0		1024	fail	12.0

Too much assumptions & strong requirements?

10 Summary of Contribution

Automated Framework to Form a Loop Structure .

First attempt: form and find the *best* cyclic plan

General domain/problem analysis method .

the basic method is applicable to other domains

Solved EXTREMELY large PDDL instances .

beyond state-of-the-art planners

Solve Large IPC problem (and variants) correctly

10.1 Is it useful?

No one doesn't even try to solve that large problems!

Dirty attempt – lessons might be learned

Global lock/owner in STRIPS – it may find a way to use

10.2 So, what's next?

- Categorizing the objects into identical groups (KEPS paper)
 - several mixed-orders becomes available (100 x A / 200 x B)
 - (100 loops A) + (200 loops B)
 - (100 loops AB) + (100 loops B)
- Categorization -> Checks serial decomposability / not.
 - check if a resource is released or not
 - consider the "release" action of the resources as an abstract action
- **Unit** capacity -> **arbitrary** capacity
 - Detect **numbers** in a problem, automatically?
 - up-converting STRIPS to ADL (opposite to the common strategy)

Thanks for listening!