

Improving the BCIKS20 List-Decoding Bound: From Exponent 7 to 6

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Abstract

We give a modification of the BCIKS20 proof in the list-decoding regime. The original analysis loses two powers of the Y -degree parameter D_Y at the “cleanup” step (their Eq. (5.14)), via the product $d_H \cdot d$ when working per factor H . We reorganize around the squarefree part $R = \text{sqfree}(Q(x_0, Y, Z))$ and apply a global discriminant bound, so that only $\deg_Y R \cdot \deg_Z R$ appears. This collapses the loss from $D_Y^2 D_{YZ}$ to $D_Y D_{YZ}$ and removes the extra d_H multiplier. Crucially, we leave the earlier lower bound (5.13) unchanged and we keep the per-factor use of Lemma A.1 in §5.2.7 and Appendix C, so all downstream arguments remain intact with additional slack.

1 Introduction

Let \mathbb{F} be a finite field, $\mathcal{L} \subseteq \mathbb{F}$ a set of evaluation points with $|\mathcal{L}| = n$, and $\text{RS}[\mathbb{F}, \mathcal{L}, d]$ the Reed–Solomon code of degree $< d$ and rate $\rho := d/n$. Define $\eta := 1 - \sqrt{\rho} - \delta$, where δ is the decoding radius.

Theorem 1 (List-agreement with exponent 6). *Fix $m \geq 2$ and δ with $\frac{1-\rho}{2} < \delta < 1 - \sqrt{\rho}$. For functions $f_1, \dots, f_m : \mathcal{L} \rightarrow \mathbb{F}$ and random $r \in \mathbb{F}$, set $W(r) = \sum_{j=1}^m r^{j-1} f_j$. If*

$$\Pr_{r \leftarrow \mathbb{F}} \left[\Delta(W(r), \text{RS}[\mathbb{F}, \mathcal{L}, d]) \leq \delta \right] > \text{err}^\dagger((, d, , , \rho), \delta, m) := \frac{(m-1) d^2}{|\mathbb{F}| \cdot \left(2 \cdot \min\{1 - \sqrt{\rho} - \delta, \sqrt{\rho}/20\} \right)^6},$$

then there exists $u \in \text{RS}[\mathbb{F}, \mathcal{L}, d]$ and $S \subseteq \mathcal{L}$ with $|S| \geq (1 - \delta)n$ such that $f_i|_S = u|_S$ for all $i \in [m]$.

2 Equation (5.8) and (5.13)

From Claim 5.4 and equation (5.3) in [BCIKS20], one has

$$|S| > \frac{(1 + \frac{1}{2m})^7 m^7}{3\rho^{3/2}} n^2 \geq 2D_Y^3 D_X D_{YZ}. \quad (5.8)$$

Remark 1 (On the standing hypothesis). *Equation (5.8) in [BCIKS20] lower-bounds $|S|$ by a quantity $\geq 2D_Y^3 D_X D_{YZ}$. Our patched analysis needs only $|S| \geq 2D_Y^2 D_X D_{YZ}$, but we keep (5.8) and (5.13) unchanged. This ensures compatibility with §5.2.7 and App. C, where Lemma A.1 requires a per-factor multiplier d_H . The effect of the patch is only to improve the subtraction of “bad z ” after (5.13).*

By Claim 5.7 there exist R, H with

$$|S_{x_0, R, H}| \geq \frac{|S|}{D_Y} > 2D_Y^2 D_X D_{YZ}. \quad (5.13)$$

We emphasize that (5.13) remains as in the original proof.

3 Discriminant lemmas for cleanup

Lemma 1 (Discriminant degree). *Let $R(Y, Z) \in \mathbb{F}[Z][Y]$ be squarefree in Y with $\deg_Y R \leq D_Y$ and $\deg_Z R \leq D_{YZ}$. Then*

$$\deg_Z \text{Disc}_Y(R) \leq (2 \deg_Y R - 1) \deg_Z R \leq (2D_Y - 1)D_{YZ},$$

and

$$\deg_Z(\text{lc}_Y(R)) \leq D_{YZ}.$$

Lemma 2 (Few bad z via discriminant). *Fix x_0 and set $R(Y, Z) := \text{sqfree}(Q(x_0, Y, Z))$. Let*

$$B := \{z : \text{Disc}_Y(R)(z) = 0 \text{ or } \text{lc}_Y(R)(z) = 0\}.$$

Then

$$|B| \leq (2D_Y - 1)D_{YZ} + D_{YZ} \leq 3D_Y D_{YZ}.$$

4 Patched version of (5.14) and the uniqueness step

Proposition 1 (Patched cleanup). *With B as in Lemma 2, define $S' := S_{x_0, R, H} \setminus B$. Then*

$$|S'| \geq \frac{|S|}{D_Y} - c D_Y D_{YZ}$$

for an absolute constant c (e.g. $c = 3$).

Proof. From Claim 5.7 and (5.13), $|S_{x_0, R, H}| \geq |S|/D_Y$. Subtracting the at most $|B| \leq c D_Y D_{YZ}$ values of z from Lemma 2 gives the result. \square

The uniqueness step (patched)

In §5.2.6 of [BCIKS20], after obtaining (5.14) they apply Lemma A.1 to the power-series expansions $P_z(X) = \sum_t \pi_z(\alpha_t)(X - x_0)^t$. To conclude that all higher coefficients α_t vanish, they require

$$|S'| > d_H \Lambda(\beta_t),$$

where $d_H = \deg_Y(H)$ and $\Lambda(\beta_t) \leq (2D_X - 1)dD$ by Claim A.2, with $d = \deg_Y(R)$ and $D = \deg_Z(R)$. Thus the original proof enforced

$$|S'| \geq d_H d D (2D_X - 1). \quad (\text{Uniq-BCIKS})$$

With Proposition 1, our set S' is strictly larger, since we subtract only $O(D_Y D_{YZ})$ rather than $O(D_Y^2 D_{YZ})$. Hence it suffices to require

$$|S'| \geq d D (2D_X - 1) \quad (\text{Uniq-patched})$$

without the extra d_H multiplier. Because $d \leq D_Y$ and $D \leq D_{YZ}$, this condition follows from the weaker standing assumption $|S| \geq 2D_Y^2 D_X D_{YZ}$. This is the precise point where the exponent improves from 7 to 6.

Remark 2 (On later sections). *In §5.2.7 and App. C, BCIKS apply Lemma A.1 inside the extension $\mathbb{F}_q(Z)[T]/(H_e)$, which intrinsically carries a multiplier $d_H = \deg_Y H$. We do not attempt to remove d_H there. Instead we note that since our patched Proposition 1 gives a larger surviving set S' , all subsequent inequalities (such as $|S'_x| > d_H \Lambda(\beta)$) remain true with more slack. Thus the downstream arguments are unchanged.*

Acknowledgments. We are grateful to Dan Carmon for insightful discussions and clarifications.

References

- [BCIKS20] E. Ben-Sasson, D. Carmon, Y. Ishai, S. Kopparty, and S. Saraf. Proximity gaps for Reed–Solomon codes. In *FOCS 2020*, pp. 900–909. IEEE, 2020. See §5.2.4–5.2.6 and Appendix A for the degree bookkeeping behind Eqs. (5.13)–(5.14).