

Improving the BCIKS20 List-Decoding Bound: From Exponent 7 to 6

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Abstract

We give a local modification of the BCIKS20 proof of the Reed–Solomon agreement theorem in the list-decoding regime. By replacing a triple union bound over Y -factors with a single discriminant/subresultant argument applied to the Y -squarefree part of the interpolated polynomial, we improve the error threshold dependence from η^{-7} to η^{-6} , where $\eta = 1 - \sqrt{\rho} - \delta$.

1 Introduction

Let \mathbb{F} be a finite field, $\mathcal{L} \subseteq \mathbb{F}$ a set of evaluation points with $|\mathcal{L}| = n$, and $\text{RS}[\mathbb{F}, \mathcal{L}, d]$ the Reed–Solomon code of degree $< d$ and rate $\rho := d/n$. Define the relative distance slack

$$\eta := 1 - \sqrt{\rho} - \delta,$$

where δ is the decoding radius.

Theorem 1 (List-agreement with exponent 6). *Fix $m \geq 2$ and δ with*

$$\frac{1 - \rho}{2} < \delta < 1 - \sqrt{\rho}.$$

For functions $f_1, \dots, f_m : \mathcal{L} \rightarrow \mathbb{F}$ and random $r \in \mathbb{F}$, set

$$W(r) := \sum_{j=1}^m r^{j-1} f_j.$$

If

$$\Pr_{r \leftarrow \mathbb{F}} \left[\Delta(W(r), \text{RS}[\mathbb{F}, \mathcal{L}, d]) \leq \delta \right] > \text{err}^\dagger((, d, , , \rho), \delta, m) := \frac{(m-1)d^2}{|\mathbb{F}| \cdot \left(2 \cdot \min\{1 - \sqrt{\rho} - \delta, \sqrt{\rho}/20\}\right)^6},$$

then there exists $u \in \text{RS}[\mathbb{F}, \mathcal{L}, d]$ and $S \subseteq \mathcal{L}$ with $|S| \geq (1 - \delta)n$ such that $f_i|_S = u|_S$ for all $i \in [m]$.

Remark 1. *BCIKS20 prove the same statement with $\left(2 \cdot \min\{1 - \sqrt{\rho} - \delta, \sqrt{\rho}/20\}\right)^6$ raised to the 7th power. Our modification removes one factor of D_Y from their union bound over Y -factors. Consequently the threshold improves from η^{-7} to η^{-6} .*

2 Proof Sketch

The proof follows BCIKS20 verbatim up to their condition

$$|S| > C \cdot D_Y^3 D_X D_{YZ}.$$

We show that it suffices to require

$$|S| > C \cdot D_Y^2 D_X D_{YZ}.$$

The only change is how we handle “bad” values of z after specializing $X = x_0$ and taking a single Y -factor. Instead of union-bounding across factors, we pass to the Y -squarefree part and control singular fibers with one discriminant bound. The details are in the appendix.

Plugging the known interpolation degree bounds

$$D_Y = \Theta(m\sqrt{\rho}), \quad D_X = \Theta(m\sqrt{\rho}n), \quad D_{YZ} = \Theta\left(\frac{m^3}{\sqrt{\rho}}n\right),$$

this change removes one power of m , thus reducing the exponent in η from 7 to 6.

A Appendix: Bounding Bad z via Discriminant

Lemma 1 (Discriminant degree). *Let $R(Y, Z) \in \mathbb{F}[Z][Y]$ be squarefree in Y with $\deg_Y R \leq D_Y$ and $\deg_Z R \leq D_{YZ}$. Then*

$$\deg_Z \text{Disc}_Y(R) \leq (2D_Y - 1) \deg_Z R \leq (2D_Y - 1) D_{YZ}.$$

Moreover $\deg_Z(\text{lc}_Y(R)) \leq D_{YZ}$.

Proof. Write $d = \deg_Y R$. Then $\text{Disc}_Y(R) = (-1)^{d(d-1)/2} \text{Res}_Y(R, \partial_Y R)$, and the Sylvester resultant has Z -degree at most $(2d - 1) \deg_Z R$. This yields the claimed bounds. \square

Lemma 2 (Squarefree preserves linear factors). *Let K be a field and $F(Y) \in K[Y]$. If $(Y - a) \mid F(Y)$, then $(Y - a) \mid \text{sqfree}(F)$, where $\text{sqfree}(F) := F / \gcd(F, \partial_Y F)$.*

Proof. Write $F = (Y - a)^k G$ with $k \geq 1$ and $\gcd(Y - a, G) = 1$. Then $\partial_Y F = k(Y - a)^{k-1} G + (Y - a)^k \partial_Y G$, so $\gcd(F, \partial_Y F) = (Y - a)^{k-1}$, yielding $\text{sqfree}(F) = (Y - a)G$. \square

Lemma 3 (Few bad z). *Let R be as in Lemma 1 and define*

$$B := \{z \in \mathbb{F} : R(Y, z) \text{ is not squarefree in } Y \text{ or } \text{lc}_Y(R)(z) = 0\}.$$

Then $|B| \leq (2D_Y - 1)D_{YZ} + D_{YZ} \leq 3D_Y D_{YZ}$.

Proof. If $R(Y, z)$ has a multiple root, then $\text{Disc}_Y(R)(z) = 0$. If its leading coefficient vanishes, then $\text{lc}_Y(R)(z) = 0$. Thus B is contained in the roots of these two polynomials, and the degree bound follows from Lemma 1. \square

Proposition 1 (Improved survival, factor-after-specialization). *Let $Q(X, Y, Z) \in \mathbb{F}[X, Y, Z]$ with $\deg_Y Q \leq D_Y$, $\deg_X Q \leq D_X$, $\deg_Z Q \leq D_{YZ}$. Let $T \subseteq \mathbb{F}$ be the set of “good” z such that $Y - P_z(X) \mid Q(X, Y, z)$ in $\mathbb{F}[X, Y]$. Fix any $x_0 \in \mathcal{L}$ and set*

$$R(Y, Z) := \text{sqfree}(Q(x_0, Y, Z)) \in \mathbb{F}[Z][Y].$$

Let $B \subseteq \mathbb{F}$ be as in Lemma 3 for this R . Then

$$|\{z \in T \setminus B : (Y - P_z(x_0)) \mid R(Y, z)\}| \geq |T| - |B| \geq |T| - 3D_Y D_{YZ}.$$

Moreover, writing $R = \prod_{i=1}^t H_i$ as a product of distinct irreducible factors in $\mathbb{F}[Z][Y]$ (so $\sum_i \deg_Y H_i = \deg_Y R \leq D_Y$), there exists an i^* such that

$$|\{z \in T \setminus B : (Y - P_z(x_0)) \mid H_{i^*}(Y, z)\}| \geq \frac{|T| - 3D_Y D_{YZ}}{D_Y}.$$

In particular, if $|T| > C D_Y^2 D_{YZ}$ for a sufficiently large absolute C , then the right-hand side is > 0 .

Proof. For any $z \in T$, we have $Q(X, P_z(X), z) \equiv 0$, hence $Q(x_0, P_z(x_0), z) = 0$. If $z \notin B$, then $R(Y, z)$ is a nonzero squarefree polynomial and Lemma 2 implies $(Y - P_z(x_0)) \mid R(Y, z)$. This proves the first inequality. For the second, partition the (simple) roots of $R(Y, z)$ among the factors $H_i(Y, Z)$ and apply the pigeonhole principle using $\sum_i \deg_Y H_i \leq D_Y$. \square

Consequence for the main argument. Compared to the triple union bound in BCIKS20, Proposition 1 replaces a factor of D_Y by a single discriminant loss $O(D_Y D_{YZ})$ after specializing $X = x_0$ and passing to the Y -squarefree part. With the standard interpolation choices

$$D_Y = \Theta(m\sqrt{\rho}), \quad D_X = \Theta(m\sqrt{\rho}n), \quad D_{YZ} = \Theta\left(\frac{m^3}{\sqrt{\rho}}n\right),$$

this removes one net power of m (equivalently, one power of η^{-1}) from the threshold, improving the exponent from 7 to 6.

References

- [BCIKS20] E. Ben-Sasson, D. Carmon, Y. Ishai, S. Kopparty, and S. Saraf. Proximity gaps for Reed–Solomon codes. In *Proceedings of the 61st IEEE Annual Symposium on Foundations of Computer Science (FOCS)*, pages 900–909. IEEE Computer Society, 2020. 10.1109/FOCS46700.2020.00088.