

University of Toronto Scarborough

CSCC24 Winter 2018 Midterm Test

Duration - 1 hour 30 minutes

Aid: 1 crib sheet (letter size, double-sided).

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1: 7/10

2: 8/10

3: 15/20

4: 10/10

5: 20/20

Total: 60/70

1. [10 marks] The *takeWhile* function takes two parameters: a predicate *p*, a list *lst*. It "takes elements from the list as long as the predicate holds"; formally, it returns the longest prefix of *lst* such that every element in the returned list satisfies the predicate. Example:

*takeWhile* even [2, 6, 4, 1, 4, 8]  
= [2, 6, 4]

- (a) [3 marks] Implement your own version in Scheme. Use your own recursion. Avoid using other list functions.

(define (take-while p lst) 2  
 (match lst  
 ['() '()]  
 [(cons hd tl)   
   (if (p hd) (cons hd (take-while p tl))  
       (take-while p tl))] ))

- (b) [3 marks] Implement your own version in Haskell. Use your own recursion. Avoid using other list functions.

takeWhile p [] = [] 2  
 takeWhile p (hd : tl) | p hd = hd : takeWhile p tl  
                           otherwise = takeWhile p tl p

- (c) [4 marks] Implement as a *foldr*. Do not use your own recursion. Write in Scheme or Haskell. 3

(define (take-while p lst)  
 (foldr lst))

OR Assuming  $\text{foldr} :: (a \rightarrow b \rightarrow b) \rightarrow b \rightarrow [a] \rightarrow b$   
 $p :: a \rightarrow \text{Bool}$

takeWhile p lst =

foldr ~~(\x y -> if (p y) (y : x) x)~~ binop [] lst  
 where binop  $\begin{cases} y : x & \text{if } p y \\ x & \text{otherwise} \end{cases}$

2. [10 marks] Standard lists in Haskell and Scheme take  $\Theta(n)$  time to compute lengths. Someone then comes up with the idea of a user-defined list type that stores lengths at list nodes, so that asking for lengths take  $O(1)$  time. In Haskell and Scheme:

```
data MyList a = End | More Int a (MyList a)
(struct More (len head tail))
; Still use '()' for the empty list in the Scheme version.
```

Example in Haskell: `More 3 x (More 2 y (More 1 z End))`

Example in Scheme: `(More 3 x (More 2 y (More 1 z '())))`

Clearly, we would not require users to use `More` directly. We would provide a decent API such as:

- (a) [5 marks] *mycons* adds an element to the beginning of a list. Implement in both Haskell and Scheme.

*mycons* :: a -> MyList a -> MyList a

*mycons* a End = More 1 a (End)

*mycons* a lst@(More n hd tl) = More (n+1) a lst

```
(define (mycons a lst)
```

```
  (match lst
```

```
    ['()
```

```
      (More 1 a '()) (More 1 a '())
```

```
      [(More n hd tl) (More new (+ n 1) a (More n hd tl) lst)])
```

- (b) [5 marks] *myappend* concatenates two lists. Use your own recursion. Maximize re-use of *mycons*. Implement in both Haskell and Scheme.

*myappend* :: MyList a -> MyList a -> MyList a

*myappend* End lst2 = lst2

*myappend* (More n hd tl) lst2 = *mycons* hd (*myappend* tl lst2)

```
(define (myappend lst1 lst2)
```

```
  (match lst1
```

```
    ['()
```

```
      [(More n hd tl)
```

```
      ]
```

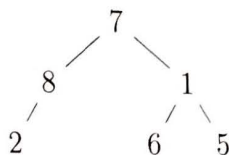
```
      ]))
```

3. [20 marks] A type of binary trees (not necessarily binary search trees) that store elements at internal nodes has been defined in Haskell and Scheme:

```
data BT a = Nil | Node (BT a) a (BT a)
(struct nil ())
(struct node (left val right))
```

Example in code and in picture:

```
exBT = Node (Node (Node Nil 2 Nil) 8 Nil)
          7
          (Node (Node Nil 6 Nil) 1 (Node Nil 5 Nil))
(define exBT
  (node (node (node (nil) 2 (nil)) 8 (nil))
        7
        (node (node (nil) 6 (nil)) 1 (node (nil) 5 (nil)))))
```



- (a) [5 marks] Implement the *btFoldl* function. It takes 3 parameters: a binary operator, a value, a binary tree. Its return value is as though you applied list *foldl* to the in-order traversal. **Important:** This only specifies external correctness, not internal implementation strategy. Example:

$$\begin{aligned}
 & \text{btFoldl } (-) \ 9 \ \text{exBT} \\
 &= ((((((9 - 2) - 8) - 7) - 6) - 1) - 5) \quad \text{(btFoldl leftchild, val)} \\
 & \quad \text{binop} \quad \text{btFoldl binop init rchild}
 \end{aligned}$$

The function type in Haskell would be  
 (b -> a -> b) -> b -> BT a -> b

Your implementation shall do it directly with recursion and avoid intermediate data structures and mutable variables. You may choose Scheme or Haskell.

~~btFoldl binop init~~

btFoldl - init Nil - = ~~init~~

btFoldl binop init (Node lchild val rchild) = btFoldl binop (binop (btFoldl binop init ~~lchild~~) val) rchild

- (b) [5 marks] Implement the *btFoldr* function. It takes 3 parameters: a binary operator, a value, a binary tree. Its return value is as though you applied list *foldr* to the in-order traversal. **Important:** This only specifies external correctness, not internal implementation strategy. Example:

$$\begin{aligned} & \text{btFoldr } (-) \ 9 \ \text{exBT} \\ &= 2 - (8 - (7 - (6 - (1 - (5 - 9)))))) \end{aligned}$$

The function type in Haskell would be

(a -> b -> b) -> b -> BT a -> b

Your implementation shall do it directly with recursion and avoid intermediate data structures and mutable variables. You may choose Scheme or Haskell.

*btFoldr* - ~~init~~ *NIL* - = ~~init~~ *init* ✓  
*btFoldr* binop init (Node (lchild val rchild)) = ~~binop~~ ~~leftside~~ ~~rightside~~ *btFoldr* binop leftside rchild  
 where leftside = binop val (*btFoldr* binop init lchild)

- (c) [5 marks] Use *btFoldr* to convert a binary tree to the list of its elements in in-order. You may choose Scheme or Haskell. Avoid writing more than minimum code.

(4) (define (toList tree) (*btFoldr* ~~binop~~ ~~init~~ tree))

OR

toList :: BT a -> [a] (*\x y -> x:y*)  
 toList tree = *btFoldr* ~~binop~~ ~~init~~ [] tree

- (d) [5 marks] Use *btFoldl* to count the number of elements in a binary tree. You may choose Scheme or Haskell. Avoid writing more than minimum code.

(5) (define (size tree)  
 (*btFoldl* ~~binop~~ ~~init~~ tree))

OR

(b -> a -> b) -> b -> BT a -> b

size :: BT a -> Integer  
 size tree = *btFoldl* (*\x y -> x+1*) 0 tree ✓



4. [10 marks] Implement in Haskell this function

`mwalk :: [Maybe a] -> Maybe [a]`

If the input list contains one or more `Nothing`, the answer is `Nothing`; otherwise, the answer factors out `Just` from the input list, e.g.,

$$mwalk [Just\ 1, Just\ 2] = Just\ [1, 2]$$

- (a) [5 marks] Use your own recursion and pattern matching on both list and `Maybe`. In the empty list case, do the one thing that makes the recursive case the simplest.

`mwalk [] = Just []`

`mwalk (Nothing : _) = Nothing`

`mwalk (Just a : tl) = case (mwalk tl) of`  
`Nothing -> Nothing`  
`Just list -> Just (a : list)`

- (b) [5 marks] Use your own recursion and pattern matching on list. But not pattern matching on `Maybe`—recall `Maybe` is an instance of `Applicative` with:

`fmap f Nothing = Nothing`

`fmap f (Just a) = Just (f a)`

`pure a = Just a`

`Just f <*> Just a = Just (f a)`

`_ <*> _ = Nothing`

So use `fmap`, `pure`, and `<*>` instead.

`mwalk [] = Just []`

`mwalk (hd : tl) =`

`fmap (\x -> \y -> x <*> y) hd <*> mwalk tl`

(This means `mwalk` is generalizable from `Maybe` to all `Applicative` instances:

`Applicative f => [f a] -> f [a]`

)

5. [20 marks] In this question, a row vector is represented by a non-empty list of numbers, and a matrix is represented by a non-empty list of row vectors. Example:

$$\begin{pmatrix} 4 & 1 & 6 \\ 9 & 0 & 5 \\ 6 & 4 & 7 \end{pmatrix} \text{ is represented as } [[4, 1, 6], [9, 0, 5], [6, 4, 7]]$$

In this question, it will be most useful to recall:

- In Scheme, *map* can apply a function to the elements of a list:  
`(map abs '(-1 -2)) = '(1 2)`  
 and can apply a binary operator to the respective elements of two lists:  
`(map + '(1 2) '(10 20)) = '(11 22)`

Summing a list can be done by `(apply + lst)`.

- In Haskell, the corresponding examples become:

`map abs [-1, -2]`

`zipWith (+) [1,2] [10,20]`

Summing a list can be done by `sum lst`.

- (a) [4 marks] Implement the function *absMat* that applies *abs* to every number in a matrix, e.g.,

$$\text{absMat } [[-1, -2], [-4, -3]] = [[1, 2], [4, 3]]$$

Use the least code; do not write your own recursion. Use Haskell.

*Handwritten:*  $\text{absMat} = \text{map } (\text{map } \text{abs})$

- (b) [4 marks] The dot product of two row vectors is defined as

$$\text{dot } [a_1, \dots, a_k] [b_1, \dots, b_k] = a_1 \times b_1 + \dots + a_k \times b_k$$

Implement this function with the least code; do not write your own recursion. You may assume that the two lists have the same length. You may use Scheme or Haskell.

*Handwritten:*  $\text{dot } a \text{ b} = \text{sum } (\text{zipWith } (*) \text{ a b})$

$$\begin{bmatrix} 9 & 0 & 5 \\ 6 & 4 & 7 \end{bmatrix}^T = \begin{bmatrix} 9 & 6 \\ 0 & 4 \\ 5 & 7 \end{bmatrix}$$

- (c) [6 marks] The transpose of a matrix switches the roles between rows and columns. Formally, given a matrix  $M$ , the  $i$ th row of  $M$  becomes the  $i$ th column of the transpose of  $M$ . Examples:

$$\begin{bmatrix} 4 & 6 \\ 9 & 0 & 5 \\ 6 & 4 & 7 \end{bmatrix}^T = \begin{bmatrix} 4 & 9 & 6 \\ 6 & 0 & 4 \\ 5 & 5 & 7 \end{bmatrix}$$

$$\text{transpose} [[9, 0, 5], [6, 4, 7]] = [[9, 6], [0, 4], [5, 7]]$$

$$\begin{aligned} \text{transpose} [[4, 1, 6], [9, 0, 5], [6, 4, 7]] &= [[4, 9, 6], [1, 0, 4], [6, 5, 7]] \\ &= [4 : [9, 6], 1 : [0, 4], 6 : [5, 7]] \end{aligned}$$

That last equation looks really like applying  $(:)$  as a binary operator to the respective elements of two lists:  $[4, 1, 6]$  and  $[[9, 6], [0, 4], [5, 7]]$ . Wait, those two lists look familiar...

Implement this function. Do not assume square matrices. You may assume that all rows have the same length. You may use Scheme or Haskell.

```
(define (transpose matrix)
  (match matrix
    [(cons row1 '())
     ]
    [(cons row1 more)
     (map cons
           ]))
```

OR

$\text{transpose} (\text{row1} : []) = \text{map } (\lambda x \rightarrow x : []) \text{ row1}$  ✓ 3

$\text{transpose} (\text{row1} : \text{more}) = \text{zipWith } (:) \text{ row1 } \text{more} (\text{transpose more})$  ✓ 3



*map dot A*

*A.B*

*A T*

~~*zipWith dot A (transpose B)*~~

(d) [6 marks] Matrix multiplication of two matrices  $A$  and  $B$  is defined as:

$[[\text{dot}(A\text{'s row } 1)(T\text{'s row } 1), \dots, \text{dot}(A\text{'s row } 1)(T\text{'s row } n)]$

$\dots,$

$[\text{dot}(A\text{'s row } m)(T\text{'s row } 1), \dots, \text{dot}(A\text{'s row } m)(T\text{'s row } n)]$

6

where  $T = \text{transpose } B$ .

The Prof has implemented this in a super-slick Haskell one-liner that doesn't need its own recursion but just takes advantage of *map*, *dot*, *transpose*, and a lambda. The Prof is about to present it in a lecture, but then...

Our most esteemed guest the much awaited Code Mangler finally enters! He deletes all parentheses, sorts the words on the RHS, and if a word occurs twice on the RHS, he deletes the second occurrence. (Fortunately, every word on the RHS appears at most twice in the Prof's correct code.) The code is mangled to:

```
mul matA matB = -> \ arowi dot map matA matB transpose
```

Help the Prof restore the correct one-liner. Remember: Some words should occur twice on the RHS, and you have to put back parentheses.

```
mul matA matB = map (\row -> map (dot row) (transpose matB)) matA
```

(End of questions.)

Blank page for sketch work.

[dot row 1, dot row 2 ... - ]

\ x →

map(\ row → map (row dot') (transpose mat B) ) mat A

[  
map (dot A's row 1) (transpose mat B)  
map (dot A's row 2) (transpose mat B)  
]