

# Pathfinding in 3D Space: A\*, Theta\*, Lazy Theta\* in Octree Structure

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# Introduction

Pathfinding addresses the problem of finding shortest paths from source to destination avoiding obstacles. It is an essential component of Machine Intelligence and finds many applications in the fields of robotics, logistics and video games. In particular, pathfinding in 3D space may be useful for drone navigation and real-time 3D strategy games. There exist different algorithms to solve exact shortest paths on graphs (Dijkstra) or 3D surfaces (exact geodesics). However, finding exact Euclidean shortest paths in three or higher dimensions is an NP-hard **problem**. Standard methods for pathfinding in a 2D plane could be extended to 3D, but they are either inefficient in time or memory.

The **objective** of this study is to find approximate shortest paths efficiently in 3D space with obstacles with reasonable memory consumption.



Screenshot from Homeworld

Multisource: reuse information

## Methods

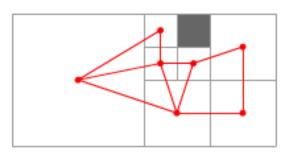
The general pipeline of our method is as follows. First, we construct an octree representing the 3D space with obstacles. Second, we construct a graph from the octree. Third, we inject source and destination in the graph. Finally, we apply a graph-based search algorithm to find approximate shortest paths.

#### Octree construction:

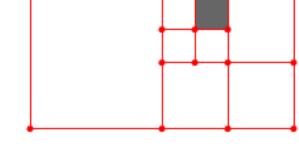
We subdivide the octree recursively wherever the space it represents intersects an obstacle, until a fixed lowest level. Since the obstacles are normally objects represented by triangle meshes, we implement a fast method to detect trianglecube intersection [1].

A progressive octree is an octree with an additional constraint: we require that the difference of levels between neighbouring octree leaves be no more than 1. This constraint can potentially reduce approximation error while applying graphbased pathfinding algorithms.

#### **Graph construction:**



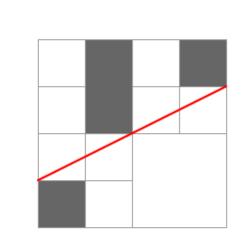
Dual graph



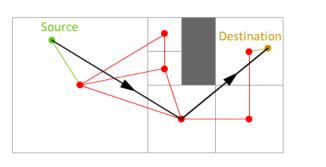
Edge-corner graph

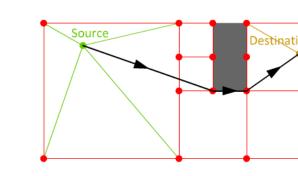
integer-based, robust

Line of sight:

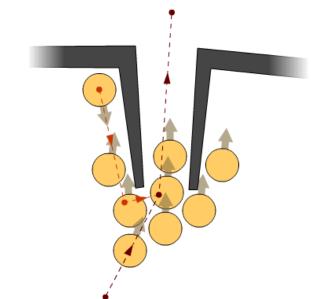


Injection of source and destination:





**Avoid exhaustive search:** precompute connectivity using union-find

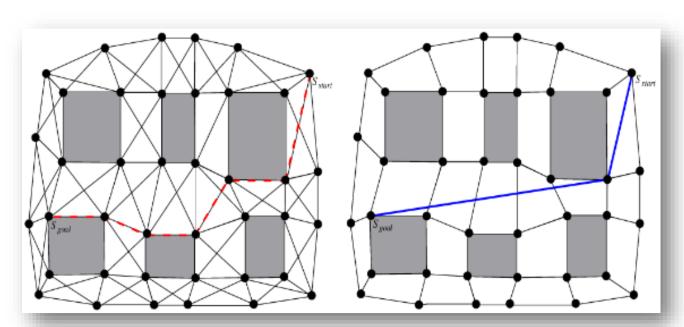


**Multi-agent** pathfinding: waypoints, repulsive force, replanning

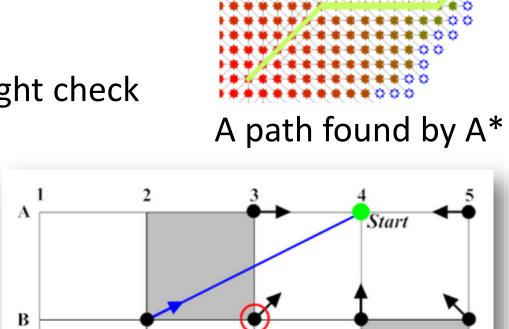
# Algorithms

## **Graph-based search algorithms:**

A\* (1968 Hart [2]) – generalization of Dijkstra Theta\* (2007 Nash [3]) – line-of-sight check Lazy Theta\* (2010 Nash [4]) – delay of line-of-sight check



Theta\* allows path outside of the edges by performing a line-of-sight check.

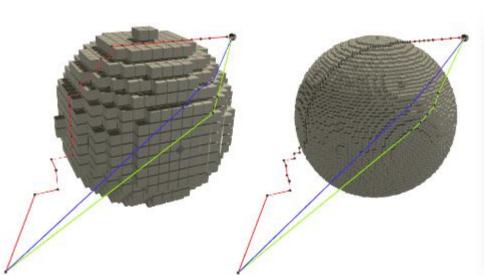


Theta\* checks whether the parent of a node is in its line of sight only before exploring its neighbours.

# Results

#### Comparison of data structures and algorithms:

Red: A\*, Green: Theta\*, Blue: Lazy Theta\*



Data Structure	Algorithm	distance	time cost	distance	time cost
Octree	A*	121.38%	$1.5 \mathrm{ms}$	121.03%	$28.0 \mathrm{ms}$
	Theta*	104.59%	$6.7 \mathrm{ms}$	102.71%	$236.1 \mathrm{ms}$
	Lazy Theta*	104.65%	$4.2 \mathrm{ms}$	102.34%	$114.8 \mathrm{ms}$
Progressive Octree	A*	127.43%	$3.3 \mathrm{ms}$	126.88%	$55.6 \mathrm{ms}$
	Theta*	103.49%	$6.3 \mathrm{ms}$	101.23%	$229.4 \mathrm{ms}$
	Lazy Theta*	103.68%	$3.2 \mathrm{ms}$	101.16%	$108.8 \mathrm{ms}$

Data Structure	Algorithm	distance	time cost	distance	time cost
Octree	A*	3.2472	$9.5 \mathrm{ms}$	3.3302	$5.3 \mathrm{ms}$
	Theta*	2.4108	$23.9 \mathrm{ms}$	2.4600	$45.5 \mathrm{ms}$
	Lazy Theta*	2.4135	$9.5 \mathrm{ms}$	2.4592	$16.3 \mathrm{ms}$
Progressive Octree	A*	3.3949	14.1ms	3.3222	7.15ms
	Theta*	2.4009	$17.59 \mathrm{ms}$	2.4158	$43.1 \mathrm{ms}$
	Lazy Theta*	2.4057	$7.78 \mathrm{ms}$	2.4205	$14.6 \mathrm{ms}$
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#### Table 2: Comparison - A complex scene - Left: edge-corner graph/ Right: dual graph

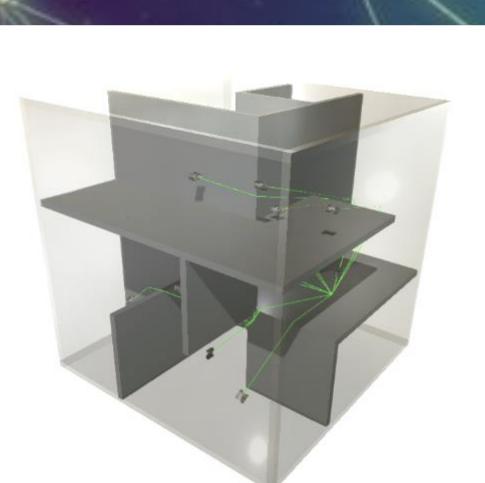
# Conclusion

**Best combination:** Lazy Theta\* + progressive octree + edge-corner graph

We proposed methods to find approximate short paths in 3D space efficiently. We achieved short paths in acceptable error (<2%) for reasonable resolution in all of our test cases with known exact solution.

#### **Future works / Possible improvements:**

Distributed computation at each frame Other possible heuristics in A\*-family algorithms Post-processing with local optimisation



### For more information:

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**Unity demo** 

References

[1] Tomas Akenine-Möller. Fast 3d triangle-box overlap testing. In ACM SIGGRAPH 2005 Courses, page 8. ACM, 2005.

[2] Nilsson Hart and Raphael. A formal basis for the heuristic determination of minimum cost paths. IEEE Transactions on Systems, Science, and Cybernetics, SSC-4(2):100–107, 1968.

[3] Alex Nash, Kenny Daniel, Sven Koenig, and Ariel Felner. Theta\*: Any-angle path planning on grids. Proceedings of the National Conference on Artificial Intelligence, 22(2):1177, 2007.

[4] Alex Nash, Sven Koenig, and Craig Tovey. Lazy theta\*: Any-angle path planning and path length analysis in 3d. In Proceedings of the National Conference on Artificial Intelligence, 2010.