Pathfinding in 3D Space

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Outline

- ▶ Introduction
- ▶ I. State of the art
- ▶ II. Algorithms
- ▶ III. Implementation in 3D space
- ▶ IV. Results
- ▶ Conclusion

Introduction

- ▶ Goal: Find the shortest paths efficiently in 3D space
- Applications: video games, drone navigation

► Homeworld (1999):

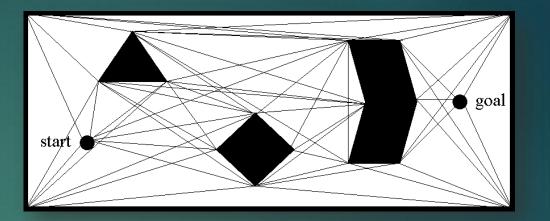
First famous real-time strategy game with movement in 3D space

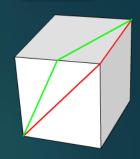




- Shortest paths in a graph
 - ▶ Dijkstra (single source)
 - ► O((|V|+|E|)log(|V|))
 - Bellman-Ford (single source, weighted directed graph)
 - ► O(|V||E|)
 - Floyd-Warshall (for all pairs of vertices, weighted graph, no negative cycle)
 - ► O(|V|³)
 - ► A* (single source, single destination)
 - \triangleright O(n), n = length of the solution path => O(|E|)

- ▶ 2D exact
 - Visibility graph
 - ► Anya (2D grid)
- ▶ 2D approximate
 - Waypoints
 - ► Navigation mesh + tunnel
 - ► Family of Theta*



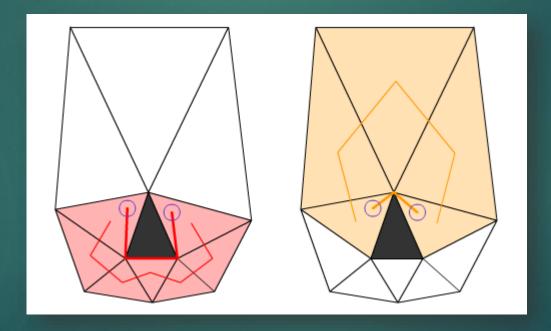






Non-optimality

Navigation mesh + tunnel



path found VS true shortest path

- ▶ 3D surface exact
 - Windows (Fast exact and approximate geodesics on meshes 2005 Surazhsky)
- ▶ 3D surface approximate
 - ► Heat (Geodesics in heat 2013 Crane)
 - ► Fast-marching (1996 Sethian)

- World representation
 - ► Tetrahedralization
 - ► Convex decomposition
 - ▶ Grid
 - Octree

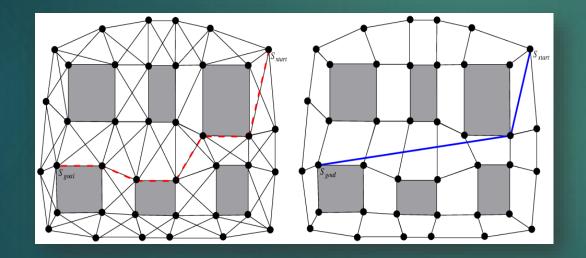
► A* (1968 Hart)

h admissible if no over-estimation and h(y) <= h(x) + d(x, y)

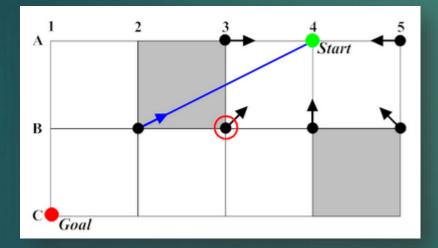


```
1 Main()
       open := closed := \emptyset;
       g(s_{start}) := 0;
       parent(s_{start}) := s_{start};
       open.Insert(s_{start}, g(s_{start}) + h(s_{start}));
        while open \neq \emptyset do
            s := open.Pop();
 7
            if s = s_{\text{goal}} then
                return "path found";
            closed := closed \cup \{s\};
10
            foreach s' \in nghbr_{vis}(s) do
11
                if s' \not\in \text{closed then}
12
                     if s' \notin \text{open then}
13
                         g(s') := \infty;
14
                         parent(s') := NULL;
15
                     UpdateVertex(s, s');
16
       return "no path found";
17
18 end
   UpdateVertex(s, s')
       g_{old} := g(s');
20
       ComputeCost(s, s');
21
       if g(s') < g_{old} then
            if s' \in \text{open then}
23
                open.Remove(s');
24
            open.Insert(s', g(s') + h(s'));
25
26 end
   ComputeCost(s, s')
        /* Path 1 */
       if g(s) + c(s, s') < g(s') then
29
            parent(s') := s;
30
            g(s') := g(s) + c(s, s');
31
32 end
```

▶ Theta* (2007 Nash)



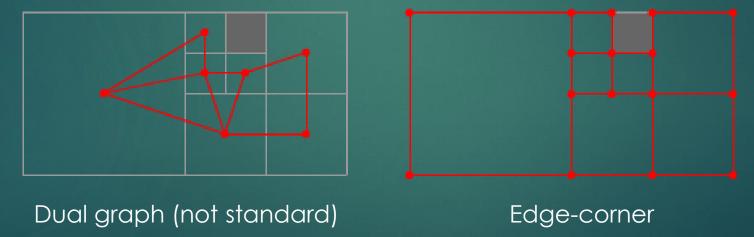
► Lazy Theta* (2010 Nash)



```
1 Main()
       open := closed := \emptyset;
       g(s_{start}) := 0;
       parent(s_{start}) := s_{start};
       open.Insert(s_{start}, g(s_{start}) + h(s_{start}));
        while open \neq \emptyset do
            s := open.Pop();
            SetVertex(s);
            if s = s_{\text{goal}} then
                return "path found";
            closed := closed \cup \{s\};
11
            foreach s' \in nghbr_{vis}(s) do
12
               if s' \not\in \text{closed then}
13
                    if s' \notin \text{open then}
                        g(s') := \infty;
                        parent(s') := NULL;
                   UpdateVertex(s, s');
17
       return "no path found";
19 end
20 UpdateVertex(s, s')
       q_{old} := g(s');
       ComputeCost(s, s');
       if g(s') < g_{old} then
           if s' \in \text{open then}
24
               open.Remove(s');
25
           open.Insert(s', g(s') + h(s'));
26
27 end
28 ComputeCost(s, s')
        /* Path 2 */
       if g(parent(s)) + c(parent(s), s') < g(s') then
           parent(s') := parent(s);
31
           g(s') := g(parent(s)) + c(parent(s), s');
33 end
34 SetVertex(s)
       if NOT lineofsight(parent(s), s) then
            /* Path 1*/
36
           parent(s) :=
37
            argmin_{s' \in nghbr_{vis}(s) \cap \boldsymbol{closed}}({}^{(g}(s') + c(s',s));
            g(s) :=
           min_{s' \in nghbr_{vis}(s) \cap closed}(g(s') + c(s', s));
39 end
```

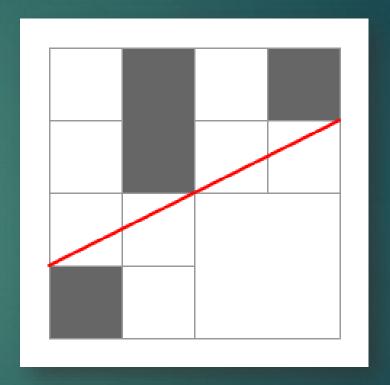
III. Implementation

- Octree construction
 - ► Triangle-cube intersection
 - ► Progressive octree
- Graph construction



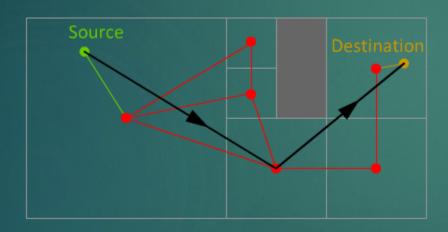
III. Implementation

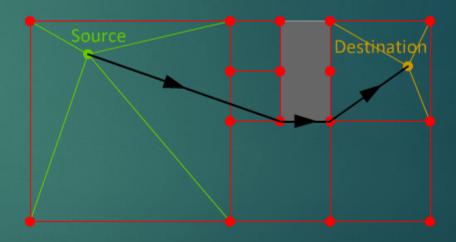
- ▶ Line of sight
 - ▶ Fast
 - ▶ Robust



III. Implementation

▶ Injection of source and destination





III. Implementation - Optimisation

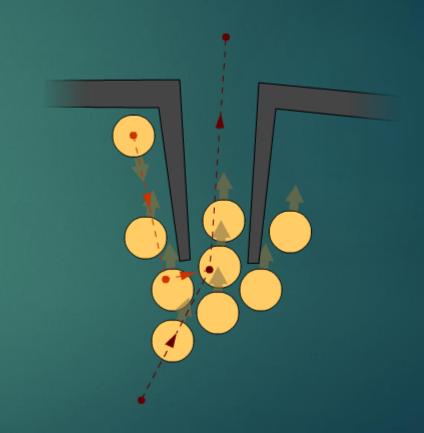
- Avoid exhaustive search
 - Precompute the connectedness of the graph nodes

III. Implementation - Optimisation

- Multisource
 - ▶ Reuse information

III. Implementation - Extension

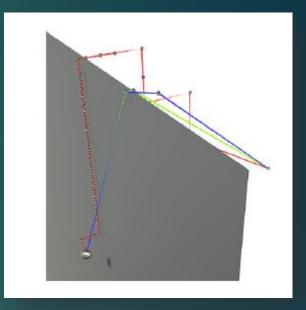
- Application in video games
 - Waypoints
 - ► Repulsive force
 - ► Replanification



IV. Results

Data Structure	Algorithm	distance	time cost	distance	time cost
Octree	A*	125.18%	$1.6 \mathrm{ms}$	124.56%	$4.5 \mathrm{ms}$
	Theta*	101.58%	$11.8 \mathrm{ms}$	109.24%	$53.0 \mathrm{ms}$
	Lazy Theta*	101.73%	$6.6 \mathrm{ms}$	109.48%	$22.1 \mathrm{ms}$
Progressive Octree	A*	125.19%	$4.6 \mathrm{ms}$	126.27%	$5.8 \mathrm{ms}$
	Theta*	101.58%	$21.5 \mathrm{ms}$	101.63%	$42.7 \mathrm{ms}$
	Lazy Theta*	101.39%	$10.6 \mathrm{ms}$	102.27%	$18.1 \mathrm{ms}$

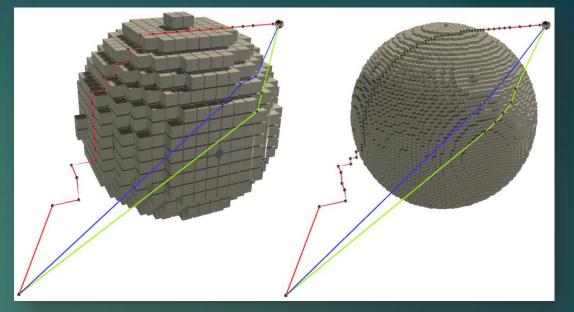
Table 1: Comparison - A single wall - Left: source sparse/ Right: destination sparse



Red: A* Green: Theta* Blue: Lazy Theta*

IV. Results

Red: A* Green: Theta* Blue: Lazy Theta*

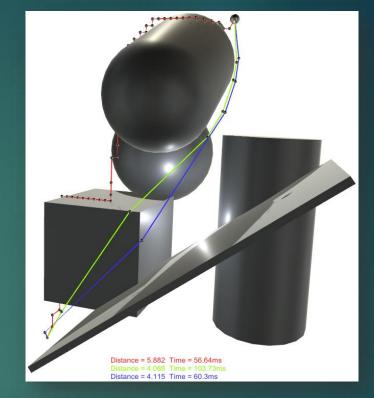


Data Structure	Algorithm	distance	time cost	distance	time cost
Octree	A*	121.38%	$1.5 \mathrm{ms}$	121.03%	$28.0 \mathrm{ms}$
	Theta*	104.59%	$6.7 \mathrm{ms}$	102.71%	$236.1 \mathrm{ms}$
	Lazy Theta*	104.65%	$4.2 \mathrm{ms}$	102.34%	114.8ms
Progressive Octree	A*	127.43%	$3.3 \mathrm{ms}$	126.88%	$55.6 \mathrm{ms}$
	Theta*	103.49%	$6.3 \mathrm{ms}$	101.23%	$229.4 \mathrm{ms}$
	Lazy Theta*	103.68%	$3.2 \mathrm{ms}$	101.16%	$108.8 \mathrm{ms}$

Table 2: Comparison - A single sphere - Left: level = 7/ Right: level = 9

IV. Results

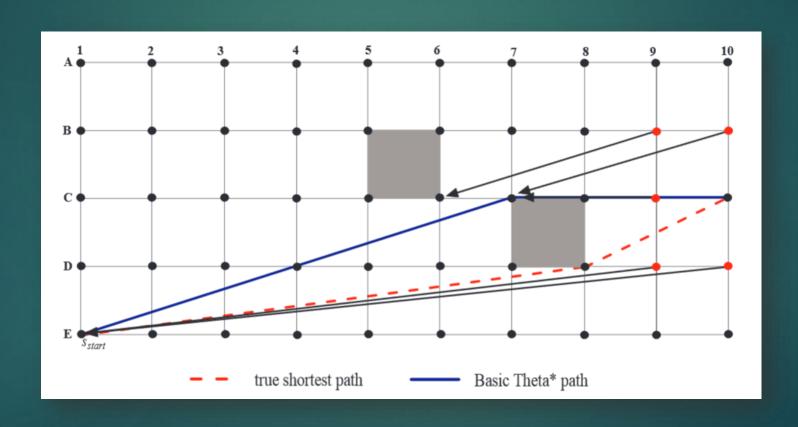
Red: A* Green: Theta* Blue: Lazy Theta*



Data Structure	Algorithm	distance	time cost	distance	time cost
Octree	A*	3.2472	$9.5 \mathrm{ms}$	3.3302	$5.3 \mathrm{ms}$
	Theta*	2.4108	$23.9 \mathrm{ms}$	2.4600	$45.5 \mathrm{ms}$
	Lazy Theta*	2.4135	$9.5 \mathrm{ms}$	2.4592	$16.3 \mathrm{ms}$
Progressive Octree	A*	3.3949	14.1ms	3.3222	7.15ms
	Theta*	2.4009	$17.59 \mathrm{ms}$	2.4158	$43.1 \mathrm{ms}$
	Lazy Theta*	2.4057	$7.78 \mathrm{ms}$	2.4205	14.6ms

Table 3: Comparison - A complex scene - Left: edge-corner graph/ Right: dual graph

Non-optimality of Theta*



Demo!

- ▶ Demo!
 - ▶ Demo!
 - ▶ Demo!
 - ▶ Demo!
 - ▶ Demo!
 - ▶ Demo!
 - ▶ Demo!
 - Demo



Conclusion

- Exploration in a new domain
- Our proposition: Lazy Theta * + Progressive Octree + Edge-corner graph
- Possible Improvements
 - ▶ Distribution of computation at each frame
 - ▶ Other possibilities of h
 - ▶ Post-processing

