

3)

a)

provide m^i and calculate $H(m^i)$.

$$k = H(m^i) \oplus MAC_K(m^i)$$

Since key is retrieved, MAC is insecure

b) is secure

c)

represent $m = m_1 || m_2 || \dots || m_n || b$, $MAC_k(m) = H(k || m_1 || m_2 || \dots || m_n || b) \oplus H(m_1 || m_2 || \dots || m_n || b)$

provide m^1 , receive $MAC_K(m^1)$

Conduct a length extension attack where $m' = m || m_{n+1} || b$

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