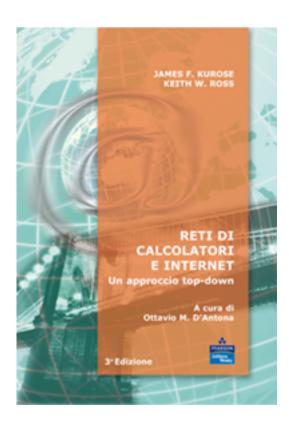
Reti di calcolatori: introduzione (Capitolo 1 Kurose-Ross)

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1 Marzo 2023

(Capitolo 1 Kurose-Ross)



Reti di calcolatori e Internet: Un approccio top-down

3ª edizione Jim Kurose, Keith Ross Pearson Education Italia ©2005

Delay in packet-switched networks

nodal processing: packets experience delay on end-to-end path check bit errors determine output link four sources of delay queueing at each hop Deterministic. time waiting at output link for transmission la portate del cono mon à sufficiente per extere code depends on congestion Sovgente dipendans dalle level of router transmission quanto of methors i bit and ~ 200,000 km/s Destinatione 8 mb/s si aspetta 1 secondo quento tenos ci processing queueing mette 11 ruter a accodamento proassare e decidere Co Voriobile, nou colcolosile a priori Jove speaking il paechetto Dipende dal router, Introduction < a) ms traffico dati

Kurose, Ross, Internet e reti di calcolatori

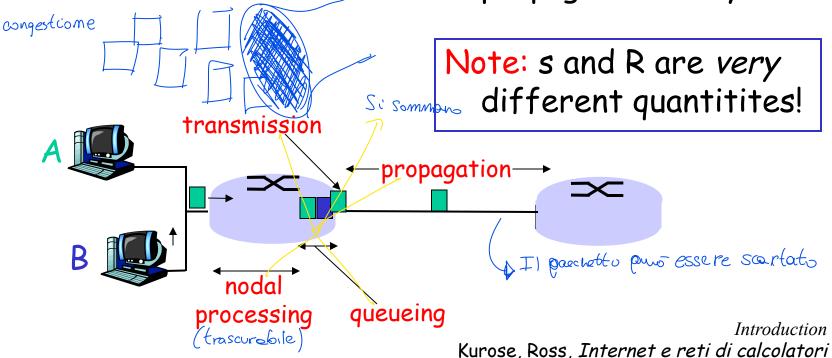
Delay in packet-switched networks

Transmission delay:

- R=link bandwidth (bps)
- □ L=packet length (bits)
- time to send bits into link = L/R

Propagation delay:

- d = length of physical link
- s = propagation speed in medium (~2×108 m/sec)
- propagation delay = d/s

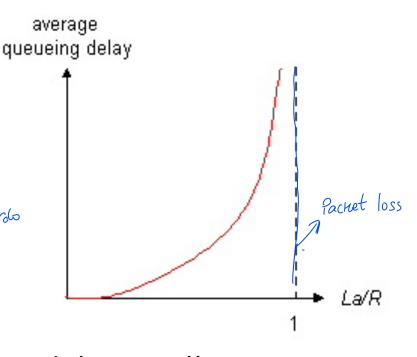


Queueing delay (revisited)

- R=link bandwidth (bps)
- □ L=packet length (bits)
- □ a=average packet arrival rate

p13 accords, p10 aumente il ritordo

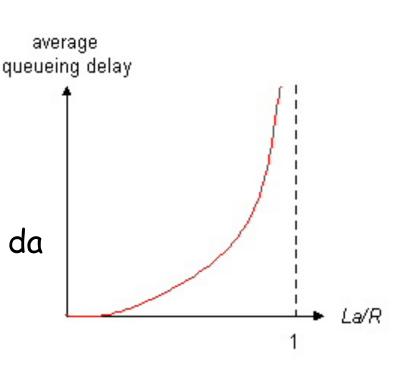




- □ La/R ~ 0: average queueing delay small
- □ La/R -> 1: delays become large
- □ La/R > 1: more "work" arriving than can be serviced, average delay infinite! File droppato

Packet Loss

- Tra l'altro: non si perderebbero pacchetti se la coda fosse di dimensione infinita!
- □ Purtroppo le code hanno dimensione finita, e dunque da un certo punto in poi i pacchetti possono essere persi: packet loss!
- Usate il programma trace route (RFC 1393) per fare un po' di prove su internet vera dove i pacchetti si accodano su N diversi router lungo una route



traffic intensity = L*a/R

Protocol "Layers"

Protocollo a livell; Accatastamento di Protocolli

Networks are complex!

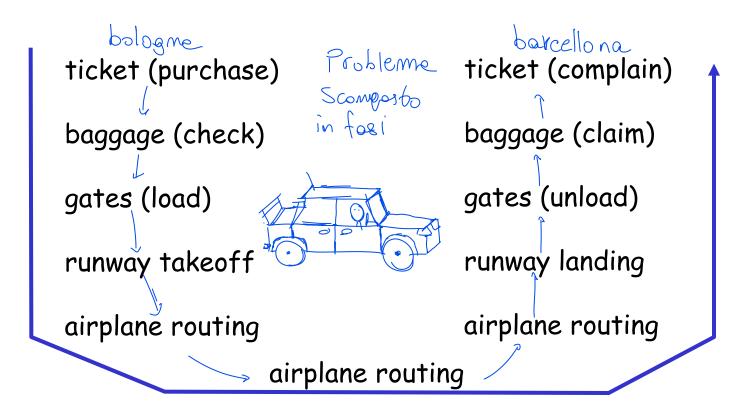
- □ many "pieces":
 - hosts
 - o routers
 - links of various media
 - applications
 - o protocols
 - hardware,software

Question:

Is there any hope of organizing structure of network?

Or at least our discussion of networks?

Organization of air travel



a series of steps

Organization of air travel: a different view

| ticket (purchase) | ticket (complain) |
|-------------------|-------------------|
| baggage (check) | baggage (claim) |
| gates (load) | gates (unload) |
| runway takeoff | runway landing |
| airplane routing | airplane routing |
| airplane routing | |

Layers: each layer implements a service

- o via its own internal-layer actions
- o relying on services provided by layer below

Layered air travel: services

Counter-to-counter delivery of person+bags

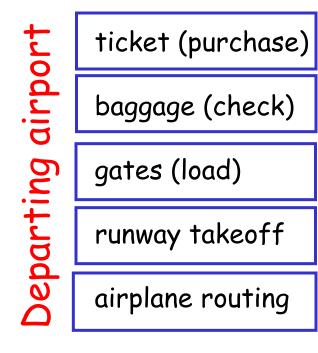
baggage-claim-to-baggage-claim delivery

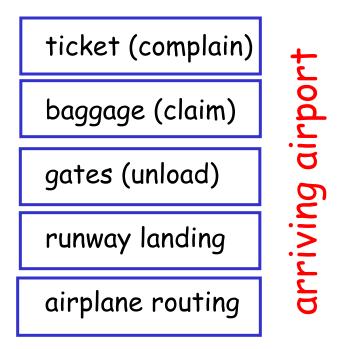
people transfer: loading gate to arrival gate

runway-to-runway delivery of plane

airplane routing from source to destination

<u>Distributed</u> implementation of layer functionality





intermediate air traffic sites

airplane routing

airplane routing

airplane routing

Ogni livello deve basars, su quello precedente/successivo

Why layering?

Dealing with complex systems:

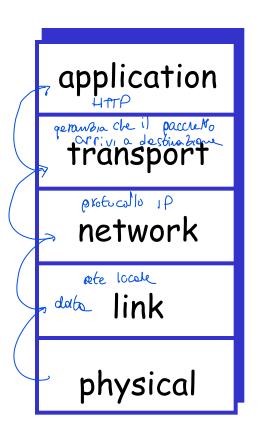
- explicit structure allows identification, relationship of complex system's pieces
 - layered reference model for discussion
- modularization eases maintenance, updating of system
 - change of implementation of layer's service transparent to rest of system
 - e.g., change in gate procedure doesn't affect rest of system
- layering considered harmful? Beware of duplications

How to place functions among layers?

- End-to-end argument: functions placed at low levels of a system may be redundant or of little value when compared with the cost of providing them at that low level.
- moving functions upward in a layered system, closer to the application that uses the function
- Higher-level layers, more specific to an application, are free to (and thus expected to) organize lower-level network resources to achieve application-specific design goals efficiently. (application autonomy)
- Lower-level layers, which support many independent applications, should provide only resources of broad utility across applications, while providing to applications usable means for effective sharing of resources and resolution of resource conflicts. (network transparency)

Internet protocol stack

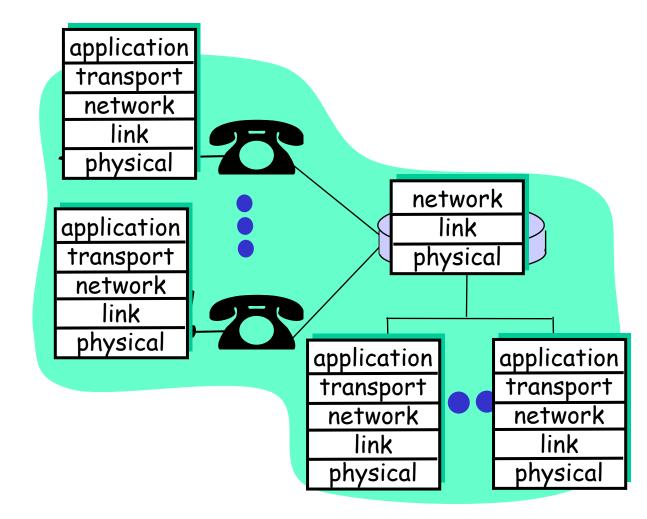
- application: supporting network applications (data)
 - o ftp, smtp, http
- transport: host-host data transfer (segment)
 - o tcp, udp
- network: routing of datagrams from source to destination
 - o ip, routing protocols
- □ link: data (frame) transfer between neighboring network elements
 - o ppp, ethernet
- physical: bits "on the wire"



Layering: logical communication

Each layer:

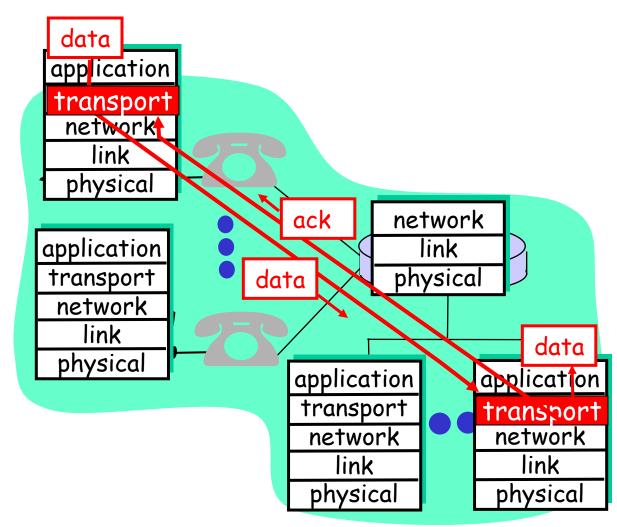
- distributed
 "entities"
 implement
 layer functions
 at each node
- entities
 perform
 actions,
 exchange
 messages with
 peers



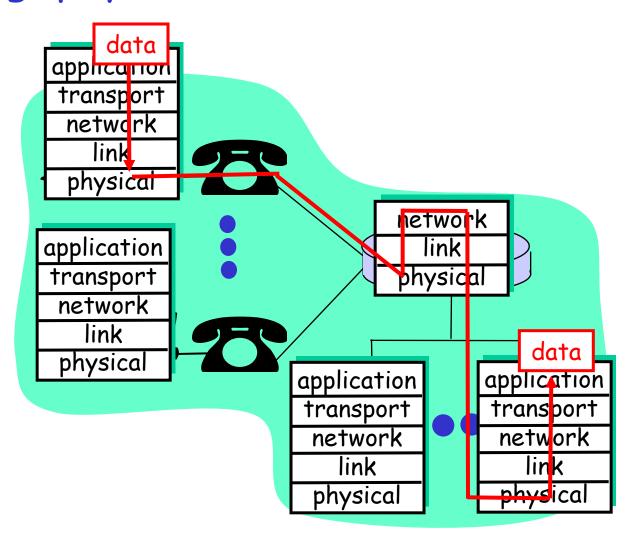
Layering: logical communication

E.g.: transport

- take data from app
- add addressing, reliability check info to form "datagram"
- send datagram to peer
- wait for peer to ack receipt
- analogy: post office



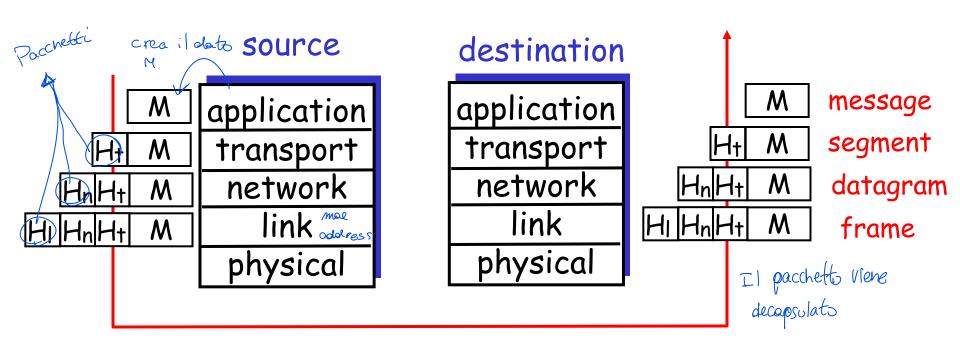
Layering: physical communication



Protocol layering and data encapsulation

Each layer takes data from above

- adds header information to create new data unit
- passes new data unit to layer below



1961-1972: Early packet-switching principles

- 1961: Kleinrock queueing theory shows effectiveness of packetswitching
- 1964: Baran packetswitching in military nets
- 1967: ARPAnet conceived by Advanced Reearch Projects Agency
- 1969: first ARPAnet node operational, Labor day
- □ Supervisionato da Kleinrock, 4 nodi UCLA, S. Barbara, SRI e Utah, primo messaggio LO

J 1972:

- ARPAnet demonstrated publicly
- NCP (Network Control Protocol) first host-host protocol
- o first e-mail program
- ARPAnet has 15 nodes

1972-1980: Internetworking, new and proprietary nets

- 1970: ALOHAnet satellite network in Hawaii
- □ 1973: Metcalfe's PhD thesis proposes Ethernet
- □ 1974: Cerf and Kahn architecture for interconnecting networks (TCP)
- □ late70's: proprietary architectures: DECnet, SNA, XNA
- □ late 70's: switching fixed length packets (ATM precursor)
- □ 1979: ARPAnet has 200 nodes

Cerf and Kahn's internetworking principles:

- minimalism, autonomy no internal changes required to interconnect networks
- best effort service model
- o stateless routers
- decentralized control

define today's Internet architecture

1980-1990: new protocols, a proliferation of networks

- ☐ 1983: deployment of TCP/IP
- □ 1982: smtp e-mail protocol defined
- 1983: DNS defined for name-to-IPaddress translation
- 1985: ftp protocol defined
- □ 1988: TCP congestion control

- new national networks:
 Csnet, BITnet,
 NSFnet, Minitel
- □ 100,000 hosts connected to confederation of networks

1990's: commercialization, the WWW

- □ Early 1990's: ARPAnet decomissioned
- □ 1991: NSF lifts restrictions on commercial use of NSFnet (decommissioned, 1995)
- □ early 1990s: WWW
 - hypertext [Bush 1945, Nelson 1960's]
 - HTML, http: Berners-Lee
 - 1994: Mosaic, laterNetscape
 - late 1990's: commercialization of the WWW

Late 1990's:

- est. 50 million computers on Internet
- □ est. 100 million+ users
- backbone links runnning at 1 Gbps

ATM: Asynchronous Transfer Mode nets

Internet:

 today's de facto standard for global data networking

1980's:

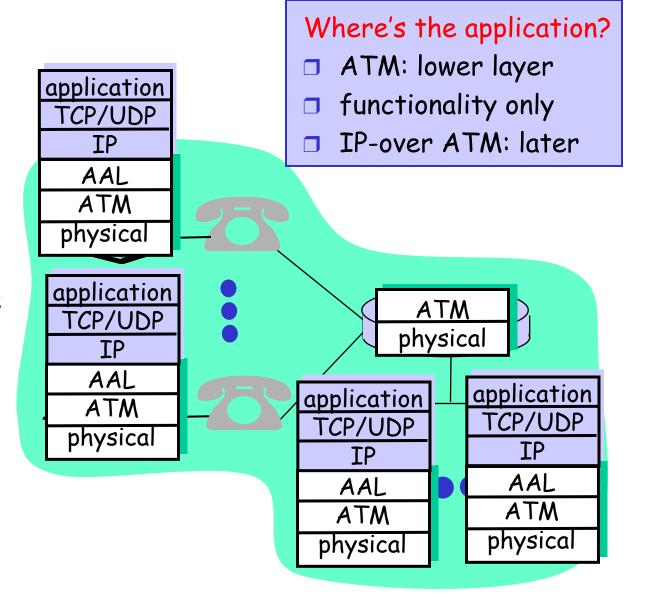
- telco's develop ATM: competing network standard for carrying high-speed voice/data
- standards bodies:
 - ATM Forum
 - O ITU

ATM principles:

- small (48 byte payload, 5 byte header) fixed length cells (like packets)
 - o fast switching
 - o small size good for voice
- virtual-circuit network: switches maintain state for each "call"
- well-defined interface between "network" and "user" (think of telephone company)

ATM layers

- ATM Adaptation Layer (AAL): interface to upper layers
 - o end-system
 - segmentation/rea ssembly
- ATM Layer: cell switching
- Physical



Chapter 1: Summary

Covered a "ton" of material!

- Internet overview
- what's a protocol?
- network edge, core, access network
- performance: loss, delay
- layering and service models
- backbones, NAPs, ISPs
- history
- ATM network

You now hopefully have:

- context, overview,
 "feel" of networking
- more depth, detail later in course