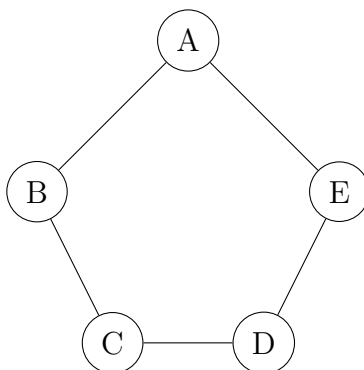


# Homework 2

Team F  
CSC 565 - Graph Theory

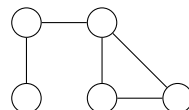
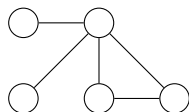
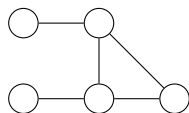
January 29, 2018

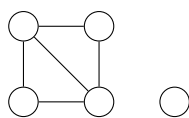
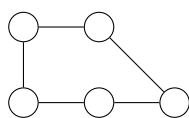
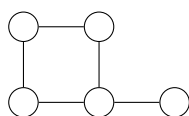
**Question 1.** Find a graph  $G$  with 5 vertices which has neither a clique of size 3 nor an independent set of size 3.



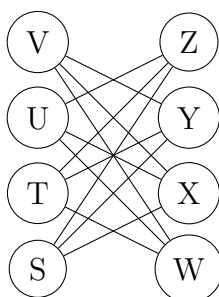
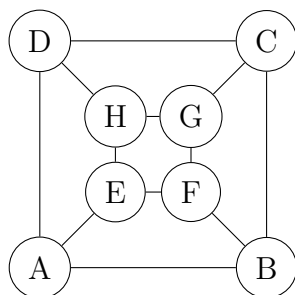
$C_5$  has a maximal clique of 2 and a maximal independent set of 2.

**Question 2.** Find all metamorphism classes of simple graphs with 5 vertices and 5 edges.

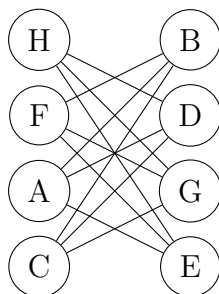




**Question 3.** Determine which pairs of graphs below are bipartite.



The first and second graph are isomorphic. The first can be relabeled to look like the second:



The third graph is not isomorphic to the first two graphs, because it is not a bipartite graph like the first two graphs (its non bipartite nature is easy discernible- it contains an odd cycle of 5, and by Theorem 1.2.18, a graph is bipartite iff it contains no odd cycles)

**Question 4.** Let  $A$  be the adjacency matrix of  $K_n$ . If  $i = j$ , then  $A^2[i, j]$ ----- . Otherwise,  $A^2[i, j] =$ -----.

- If  $i = j$ , then  $A^2[i, j] = (n-1)$
- Otherwise,  $A^2[i, j] = (n-2)$

**Question 5.** Prove or disprove: if  $G$  is a disconnected graph, then  $\overline{G}$ , the complement of  $G$ , is connected.

i. Let  $u, v$  be on different components in  $G$ . Since there is no  $uv$  edge in  $G$ , there must be a  $uv$  edge in  $\overline{G}$ . Thus  $u, v$  are connected in  $\overline{G}$

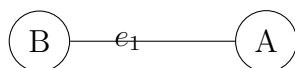
ii. Let  $u, v$  be on the same component in  $G$ . Let  $w$  be a vertex on a separate component in  $G$ . There must be a  $uw$  edge in  $\overline{G}$ . There must also be a  $vw$  edge in  $\overline{G}$ .

Because there is a  $uw$  path, and a  $vw$  path, by Lemma 1.2.5 there must also be a  $uv$  path. Thus  $u$  and  $v$  are connected in  $\overline{G}$

Because there exists a path between any two vertices in  $\overline{G}$  it is a connected graph.

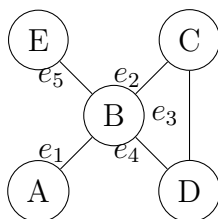
**Question 6.** Determine whether  $K_4$  contains the following (give an example or a proof of non-existence);

a) A walk that is not a trail.



The walk  $B, e_1, A, e_1, B$  is a walk that is not a trail; it repeats the  $e_1$  edge.

b) A trail that is not closed and is not a path.



The walk  $A, e_1, B, e_2, C, e_3, D, e_4, B, e_5, E$  is a trail that is not closed and is not a path, as the vertex  $B$  is repeated.

c) A closed trail that is not a cycle.

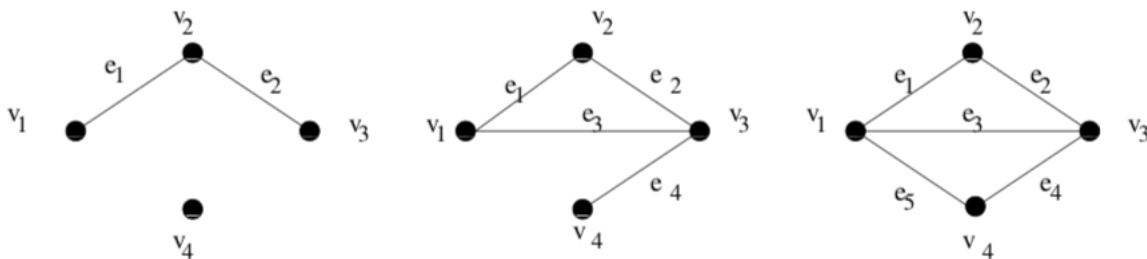


The closed trail of A alone is a closed trail that is not a cycle. The length of the walk is 0, so it is not a cycle.  $v_0 = v_n$  so it is closed. It also does not repeat any edges, therefore it is a closed trail that is not a cycle.

**Question 7.** A graph is called chordal if it has no induced subgraph isomorphic to  $C_n$  for any  $n \leq 4$ . Which of the following 10-vertex graphs are chordal?

- $K_{10}$  : Yes
- $K_{5,5}$  : No -  $K_{5,5}$  contains  $C_4$  as an induced subgraph
- $K_{1,9}$  : Yes -  $K_{1,9}$  is trivially chordal, since it does not contain any cycles
- $C_{10}$  : No -  $C_{10}$  contains itself as an induced subgraph
- $P_{10}$  : Yes -  $P_{10}$  is trivially chordal, since it does not contain any cycles
- Petersen Graph: No - The Petersen graph has  $C_5$  as an induced subgraph

**Question 8.** a. For each graph  $G$  below, find the incidence matrix  $M(G)$ . (See definition 1.1.17 and Example 1.1.19.)



Far Left

	$e_1$	$e_2$
$v_1$	1	0
$v_2$	1	1
$v_3$	0	1
$v_4$	0	0

Middle

	e <sub>1</sub>	e <sub>2</sub>	e <sub>3</sub>	e <sub>4</sub>
v <sub>1</sub>	1	0	1	0
v <sub>2</sub>	1	1	0	0
v <sub>3</sub>	0	1	1	1
v <sub>4</sub>	0	0	0	1

Far Right

	e <sub>1</sub>	e <sub>2</sub>	e <sub>3</sub>	e <sub>4</sub>	e <sub>5</sub>
v <sub>1</sub>	1	0	1	0	1
v <sub>2</sub>	1	1	0	0	0
v <sub>3</sub>	0	1	1	1	0
v <sub>4</sub>	0	0	0	1	1

b. Recall that the transpose of a  $p \times q$  matrix  $M$  is the  $q \times p$  matrix  $M^T$  with  $M^T[i, j] = M[j, i]$ . For each graph  $G$  above in (a), find the matrix product  $M(G) \cdot (M(G))^T$ .

Far Left

1	1	0	0
1	2	1	0
0	1	1	0
0	0	0	0

Middle

2	1	1	0
1	2	1	0
1	1	3	1
0	0	1	1

Far Right

3	1	1	1
1	2	1	0
1	1	3	1
1	0	1	2