

Relationale Datenbanken

Dinosaurier oder Dauerbrenner?

Folien



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- Aufgaben:
 - Bildverarbeitung
 - Backend & Verteilte Systeme
 - Desktop Apps
 - Produktionsanlagennahe Software



Werkzeuge



Werkzeuge



Werkzeuge



Werkzeuge

Java



C#



C++

Werkzeuge

MS SQL



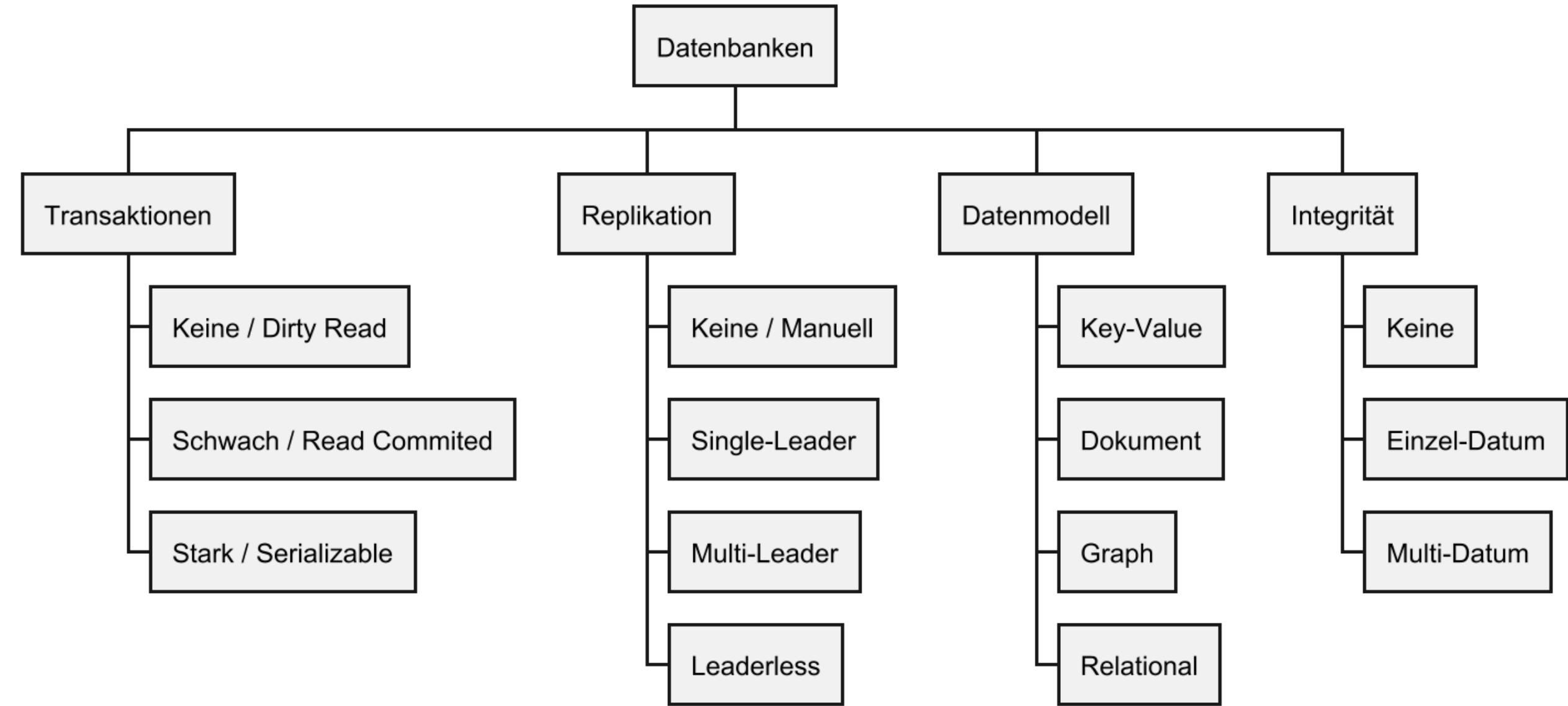
PostgreSQL



Oracle

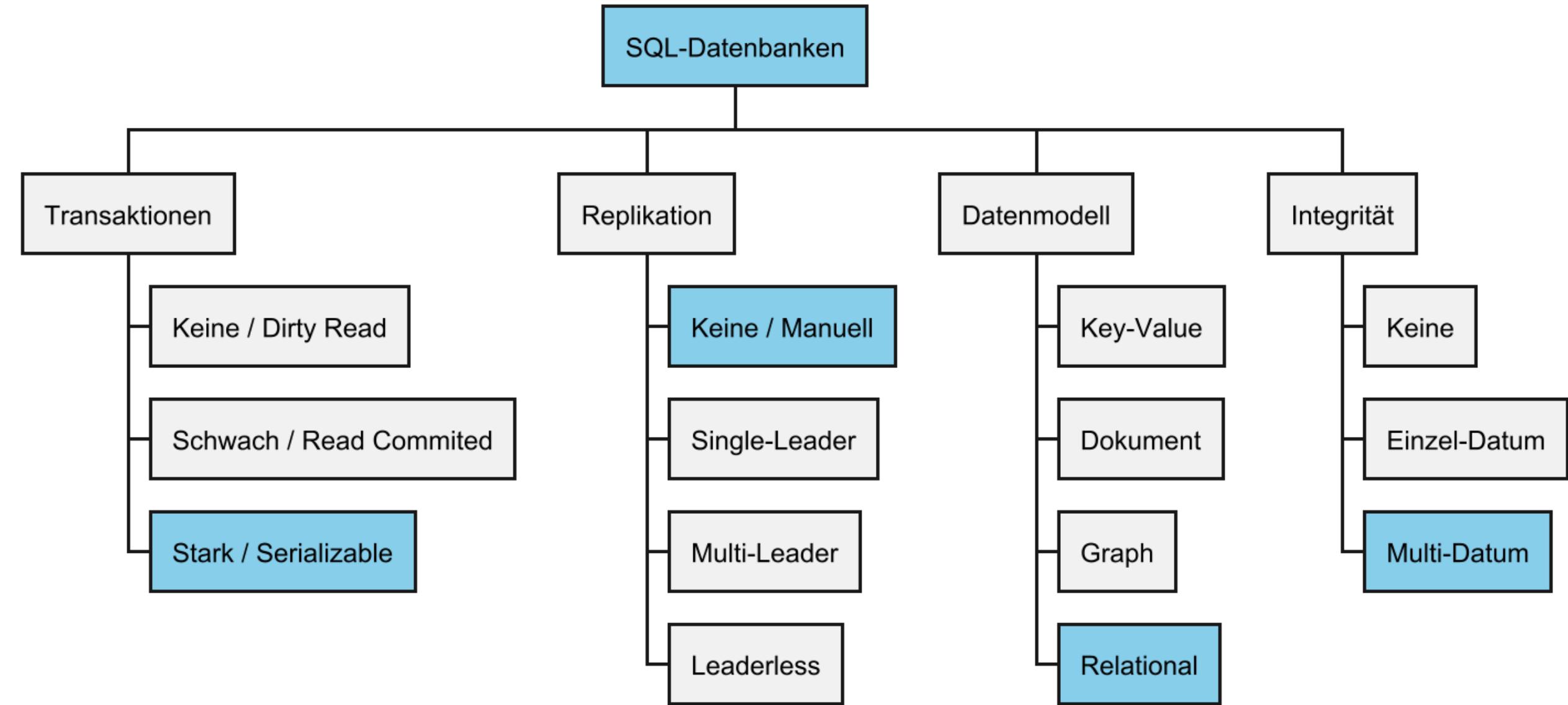
Das "Ich bau mir eine Datenbank"-Spiel

Wähle 1x aus jedem Ast

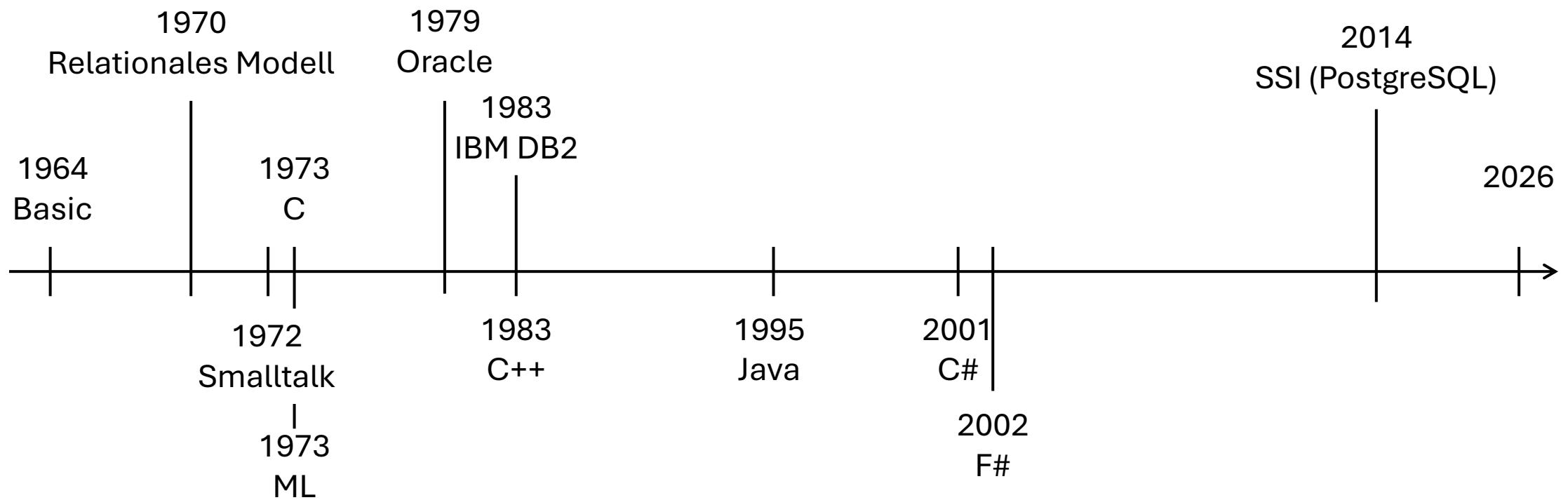


Das "Ich bau mir eine Datenbank"-Spiel

Wähle 1x aus jedem Ast



Geschichte



Lösung für was?

Future users of large data banks must be protected from having to know how the data is organized in the machine (the internal representation). A prompting service which supplies such information is not a satisfactory solution. Activities of users at terminals and most application programs should remain unaffected when the internal representation of data is changed and even when some aspects of the external representation are changed. Changes in data representation will often be needed as a result of changes in query, update, and report traffic and natural growth in the types of stored information.

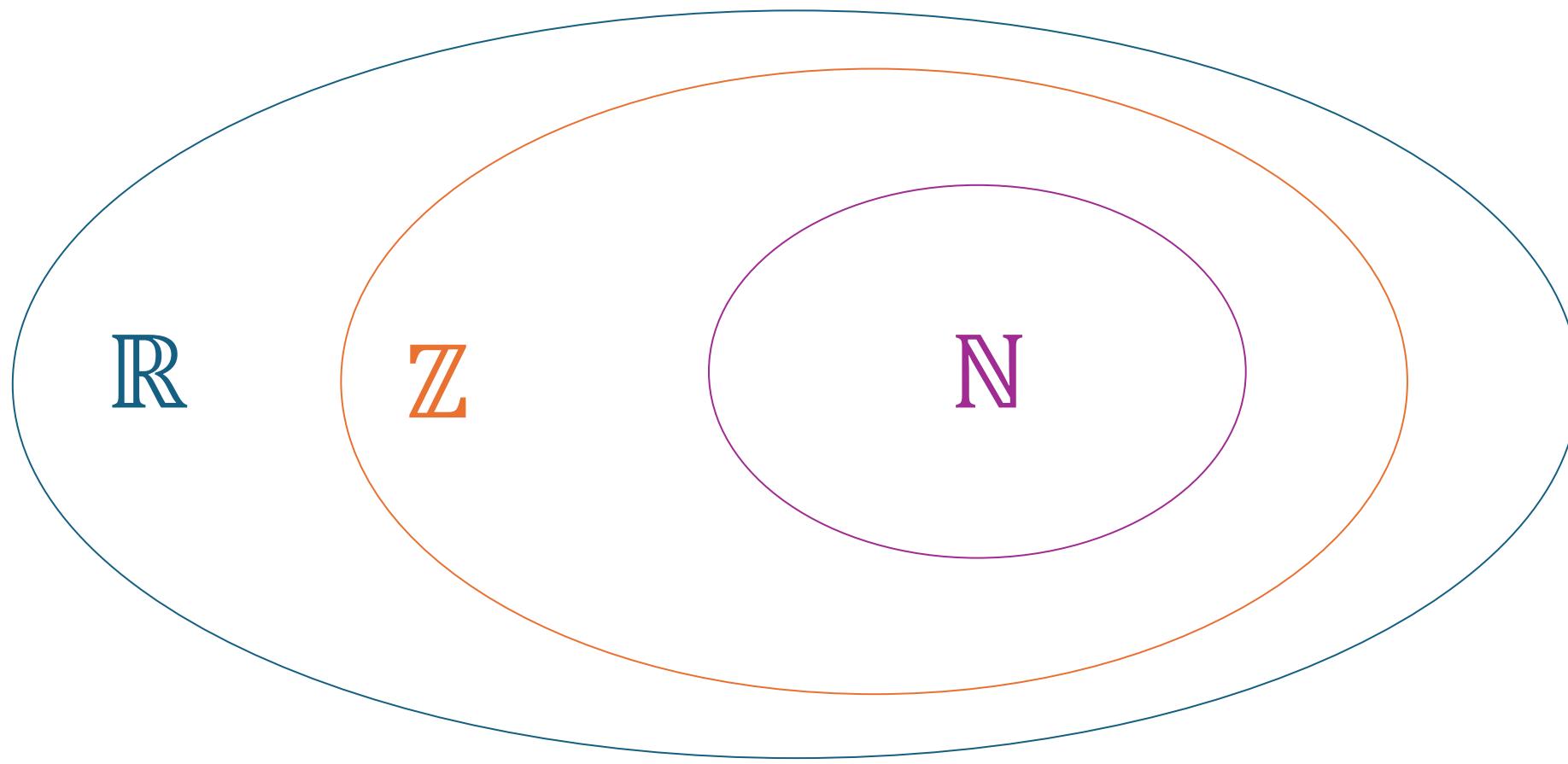
Existing noninferential, formatted data systems provide users with tree-structured files or slightly more general network models of the data. In Section 1, inadequacies of these models are discussed. A model based on n-ary relations, a normal form for data base relations, and the concept of a universal data sublanguage are introduced. In Section 2, certain operations on relations (other than logical inference) are discussed and applied to the problems of redundancy and consistency in the user's model.

Abstract von “A relational model of data for large shared data banks”
E.F. Codd, 1. Juli 1970

Warum ist Struktur so wichtig?

- 052knai0111D160129Erec727711ap208ne55S001061016
- 100SpenserE2106198026150160PiankaD110117727250
- 100 Spenser E21 06-19-80 26150 160 Pianka D11
10-11-77 27250
- 100 Spenser E21 06-19-80 26150
160 Pianka D11 10-11-77 27250
- “BUSINESS MODELING FOR DATABASE DESIGN: Formalizing the Informal (PRACTICAL DATABASE FOUNDATIONS Book 1)”, Fabian Pascal, 2015

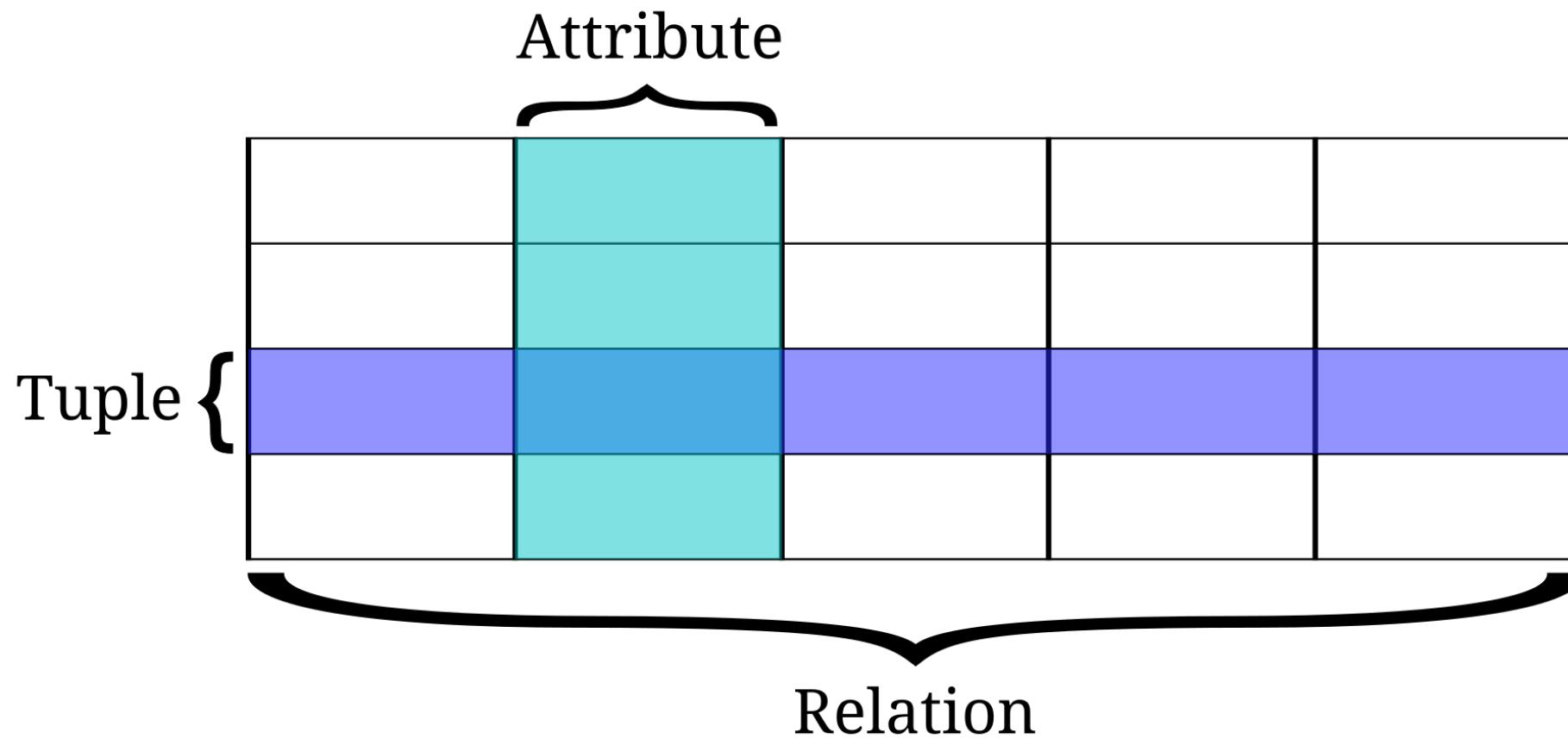
Mengentheorie



Logik

- 100 Spenser E21 06-19-80 26150
160 Pianka D11 10-11-77 27250
- Ein Mitarbeiter wird identifiziert durch Nummer (EMP#), hat den Namen (ENAME), arbeitet in der Abteilung mit der Nummer (DEPT#), wurde eingestellt zum Datum (HIREDATE) und verdient jährlich (SALARY) €.
- EMP# ENAME DEPT# HIREDATE SALARY
100 Spenser E21 06-19-80 26150
160 Pianka D11 10-11-77 27250

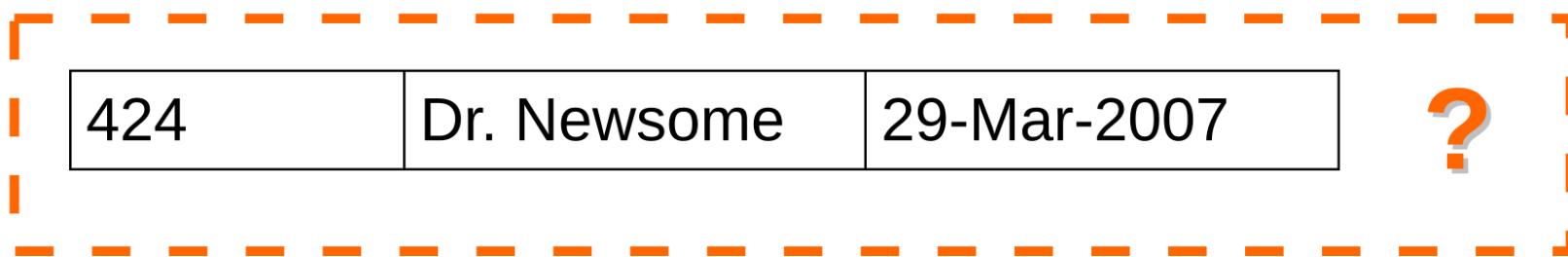
Relationales Modell



Normalformen: insert anomaly

Faculty and Their Courses

Faculty ID	Faculty Name	Faculty Hire Date	Course Code
389	Dr. Giddens	10-Feb-1985	ENG-206
407	Dr. Saperstein	19-Apr-1999	CMP-101
407	Dr. Saperstein	19-Apr-1999	CMP-201



Normalformen: Update anomaly

Employees' Skills

Employee ID	Employee Address	Skill
426	87 Sycamore Grove	Typing
426	87 Sycamore Grove	Shorthand
519	94 Chestnut Street	Public Speaking
519	96 Walnut Avenue	Carpentry

Normalformen: delete anomaly

Faculty and Their Courses

Faculty ID	Faculty Name	Faculty Hire Date	Course Code
389	Dr. Giddens	10-Feb-1985	ENG-206
407	Dr. Saperstein	19-Apr-1999	CMP-101
407	Dr. Saperstein	19-Apr-1999	CMP-201



DELETE

Normalformen

- 0 Normalform:
 - nicht normalisiert, Redundanzen
- 3 Normalform:
 - fast keine Redundanzen mehr in einer “Tabelle” (Funktionale Abhangigkeiten)
- 3,5 / Boyce-Codd Normalform:
 - Keine Redundanzen mehr in einer “Tabelle”
 - Normalerweise ausreichend
- 5 Normalform:
 - Keine Redundanzen mehr zwischen mehreren “Tabellen” / (Join Abhangigkeiten). Weitere Zerlegung in kleinere “Tabellen” fuhrt zu Datenverlust.

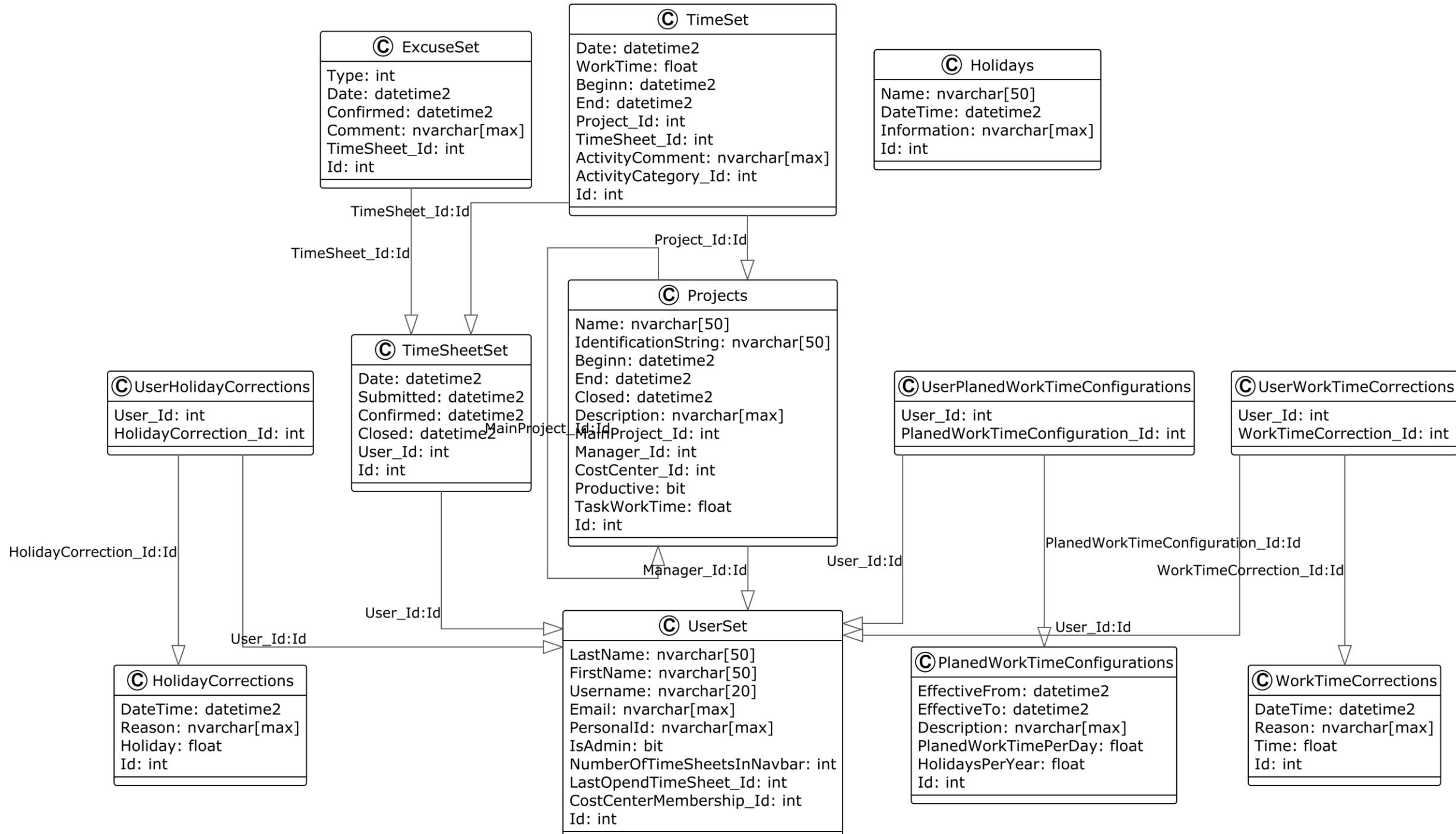
Warum SQL so scheiön ist

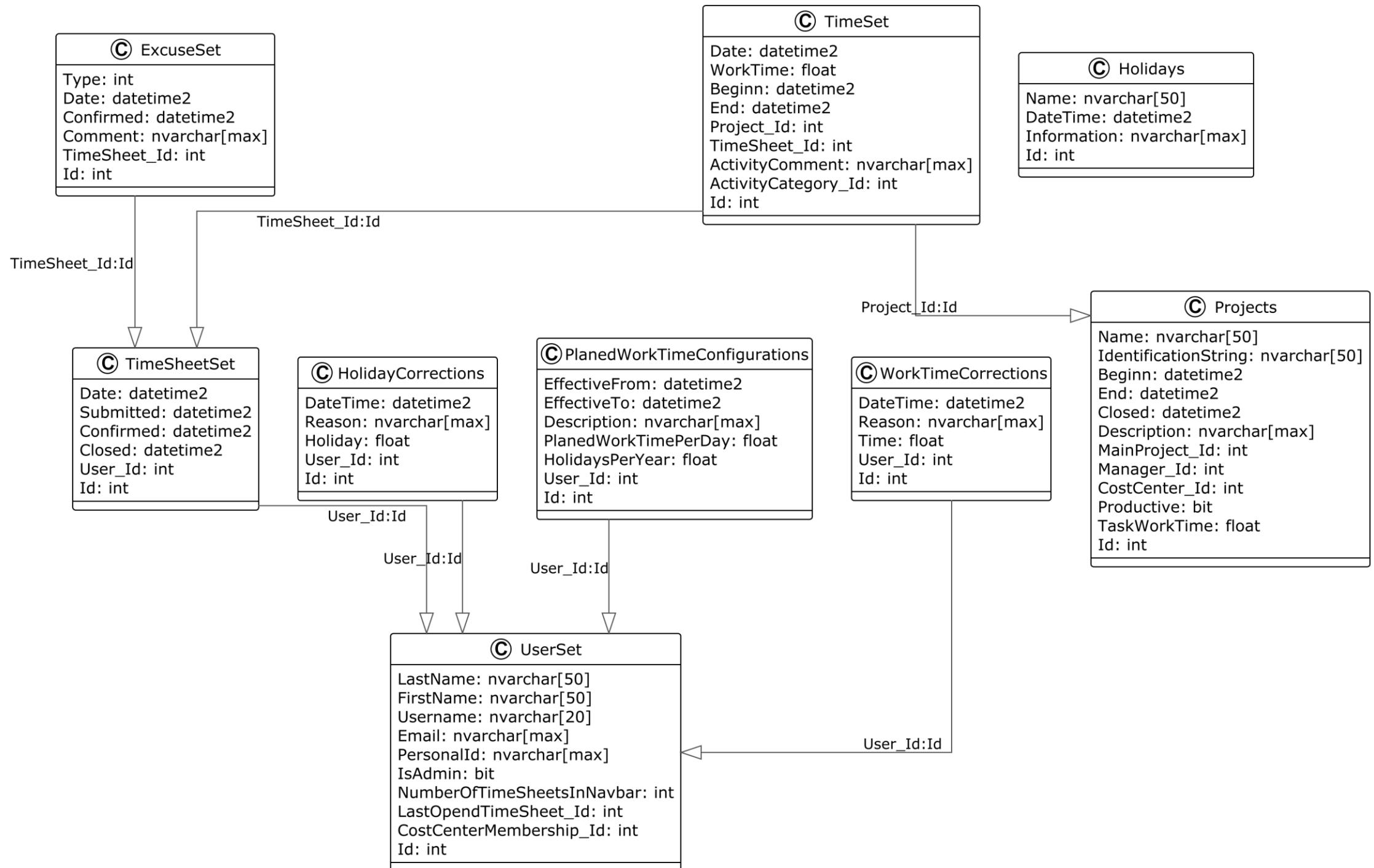
- Katastrophale Syntax: select ... from ... where
- Keine Trennung von logischem und physischem Schema
- WAHR, FALSCH, NULL → Dreiwertlogik

```
SELECT * FROM MyTable WHERE MyColumn != NULL (0 Results)
SELECT * FROM MyTable WHERE MyColumn <> NULL (0 Results)
SELECT * FROM MyTable WHERE MyColumn IS NOT NULL (568 Results)
```

- SQL erlaubt doppelte Einträge
- SQL erlaubt keine Constraints über mehrere Einträge oder Spalten
 - CREATE ASSERTION ...

Beispiel: Sollarbeitszeitberechnung





Literatur

- Data-Oriented Design, Richard Fabian, 2018
- Designing Data-Intensive Applications, Martin Kleppmann, 2017
- An Introduction to Relational Database Theory, Darwin Hugh, 2014
 - SQL: A Comparative Survey, Darwin Hugh, 2017
- Applied Mathematics for Database Professionals, Lex de Haan, Toon Koppelaars, 2007

Werkzeuge

