

Basic Graph Theory

Binary trees:

- a special case of graphs
- start with basic graph theory -
- a pictorial way to represent information.

1

References:

- Discrete Mathematics by Ross & Wright
- Discrete Mathematics by R. Johnsonbaugh
- Discrete Mathematics and Applications by S. Epp
- Schaum's Outline Series in Discrete Mathematics

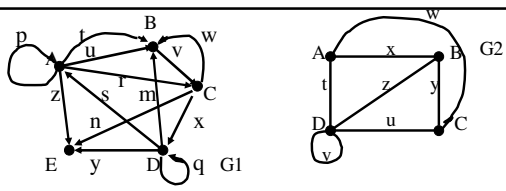
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Definition: Graph

A graph $G=(V,E)$, an ordered pair, where V is a set of vertices (nodes), & E is a set of edges on the vertices.

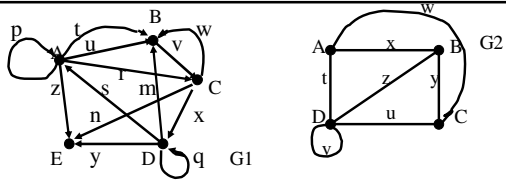
- V is not empty and has size n .
- E may be empty.
- Graphs are a modeling tool

3



- G1 has 5 vertices and 13 edges.
 $V = \{A, B, C, D, E\}$ & $E = \{m, n, p, q, r, s, t, u, v, w, x, y, z\}$
- G2 has 4 vertices and 7 edges.
 $V = \{A, B, C, D\}$ and $E = \{t, u, v, w, x, y, z\}$

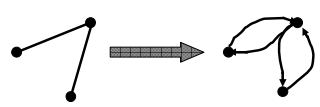
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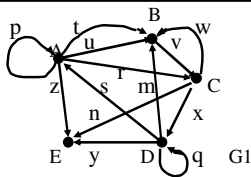
- Edge and vertex labels are arbitrary and optional.
- In G1 and G2, the vertices and edges are each labeled.

5

- G1 is a directed graph (arrowheads)
- G2 is an undirected graph (no arrows)
- All edges on a graph must be the same.
- To convert undirected graphs to directed graphs, replace each undirected edge with a pair of opposing edges.



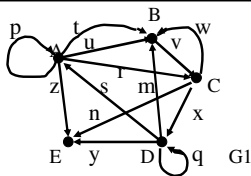
6



Edges can be named by the vertex pair they connect, e.g. in G1, AE, BC, DA are examples of edges. They are the same as edges **z**, **v**, and **s**, respectively.

- Vertex B in G1 is said to be adjacent to vertex A. Vertex A is NOT adjacent to B
- B is also adjacent to vertices C and D.
- Vertex E is adjacent to vertices A,C and D.

7

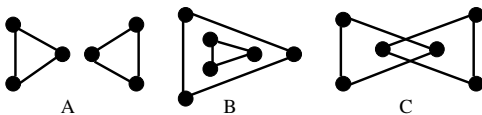


Graphs - Special edges

- Loops: on vertices A and D in G1, edge v in G2
- Parallel Edges: edges t and u in G1
- Note: edges v and w in G1 are not parallel

8

- Are graphs A, B and C
- all the same?
- all different?
- or is one different from the other two?



- This is an example of a disconnected graph with connected components.
- The way a graph is drawn is not important as long as the salient physical characteristics are preserved.

9

A **path** of length n is sequence of n edges from one vertex to another.

uvxy	length = 4
uvn	length = 3
z	length = 1
rxxy	length = 3
ABCDE	length = 4

are examples of paths in graph G_1 .
The first and last path are the same.

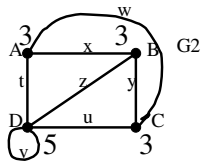
10

- As a rule interesting paths are **simple**, without repeated vertices or edges.
- A **cycle** or **closed path** begins and ends at the same vertex.
- uvxs is a cycle of length 4 in G_1
- vxm is a cycle of length 3 in G_1
- ACDA is a cycle of length 3 in G_1
- xyut is a cycle of length 4 in G_2
- An **acyclic** graph has no cycles.

11

- It is possible to attach additional information to an edge,
- e.g. time units required
cost info in \$
capacity (e.g. gallons per minute).
- These values are called **weights**. All edges are weighted. If it is not relevant to the problem, then all the weights are 1 and are not shown.

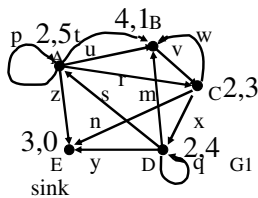
12



- The **degree** of a vertex is the number of incident edges.

- In G2, vertices A, B, & C have degree 3.
- In G2, vertex D has degree 5.

13

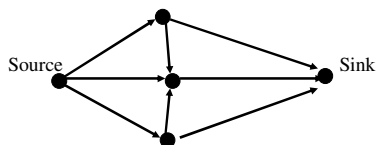


- In a directed graph, we distinguish between in-degree and out-degree.

- A has in-degree 2 & out-degree 5
- In G1, B has in-degree 4 & out-degree 1

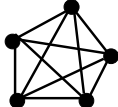
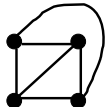
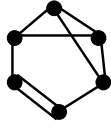
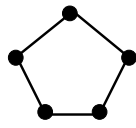
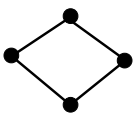
14

- Vertex E is a sink. It has out-degree 0.
- A vertex with in-degree zero and positive out-degree is called a source.



15

- If the degree of each vertex is the same, the graph is regular.



K1

K2

K3

K4

K5₁₆
