

Problem 1 (r is power of 2)

(a)

• if exist $r \in \mathbb{Z}_2^*$ s.t. $f(r) = 5^r \pmod{39} \equiv 1$

$$\text{then, } \begin{cases} f(x-r) = 5^x \cdot 5^{-r} \pmod{39} \equiv 5^x \pmod{39} = f(x) \\ f(x+r) = 5^x \cdot 5^r \pmod{39} \equiv 5^x \pmod{39} = f(x) \end{cases} \Rightarrow f(x) \text{ is periodic.}$$

• Furthermore, $f(4) = 5^4 = 625 \equiv 1 \pmod{39}$ (and $f(1), f(2), f(3) \not\equiv 1$), $r=4$.

(b) $0, 4, \dots, 2044 \rightarrow 1$
 $x \rightarrow f(x): 1, 5, \dots, 2045 \rightarrow 5$ $\Rightarrow f(x) = \frac{1}{2}(11\rangle + 15\rangle + 125\rangle + 18\rangle)$ #
 $2, 6, \dots, 2046 \rightarrow 25$
 $3, 7, \dots, 2047 \rightarrow 8$

(c) if $Z=8$, $|x\rangle = \frac{1}{\sqrt{2^9}} (|3\rangle + |7\rangle + |11\rangle + \dots + |2047\rangle)$
 $= \frac{1}{\sqrt{2^9}} \sum_{j=0}^{2^9-1} |3+j \cdot 4\rangle$, $r=4$.

(d) QFT $_{2^n}: |x\rangle \rightarrow \frac{1}{\sqrt{N}} \sum_{y \in \mathbb{Z}_N} W^{xy} |y\rangle$, $N \triangleq 2^n$

$$\begin{aligned} \frac{1}{\sqrt{2^9}} \sum_{j=0}^{2^9-1} |3+4j\rangle &\rightarrow \frac{1}{\sqrt{2^9}} \frac{1}{\sqrt{N}} \sum_{j=0}^{2^9-1} \sum_{y \in \mathbb{Z}_N} W^{(3+4j)y} |y\rangle \\ &= \frac{1 \cdot 1}{\sqrt{2^9} \cdot \sqrt{N}} \sum_{y \in \mathbb{Z}_N} W^{3y} \left(\sum_{j=0}^{2^9-1} W^{4jy} \right) |y\rangle \\ &= \frac{1}{\sqrt{2^9} \sqrt{N}} \sum_{l=0}^3 W^{3 \cdot 2^9 l} \cdot 2^9 \cdot |2^9 l\rangle \\ &= \frac{1}{2} \sum_{l=0}^3 W^{3 \cdot 2^9 l} |2^9 l\rangle \rightarrow \text{measure result: } \begin{cases} 0: 25\% \\ 2^9: 25\% \\ 2^{10}: 25\% \\ 3 \cdot 2^9: 25\% \end{cases} \end{aligned}$$

(e) $\frac{2^9 \cdot l}{2^n} = \frac{l}{4} \rightarrow \boxed{50\%}$: if $l=1, 3$, $\frac{l}{4}$ has denominator 4, $f(x) = f(x+4)$, $r=4$
 $\downarrow 50\%$: if $l=0, 2$, $\frac{l}{4}$ has denominator 1/2, $f(x) \neq f(x+1), f(x+2)$, failed.

(f) $5^4 - 1 = (5^2 - 1)(5^2 + 1) \pmod{39} \equiv 0$

$$\because 5^2 + 1 \pmod{39} \neq 0 \quad \therefore \text{Shor's algo holds, } \gcd(5^2 - 1, 39) = 3, \frac{3 \cdot 13 = 39}{\gcd(5^2 + 1, 39) = 13} \text{ 6}$$

Problem 2 (r is not power of 2)

(a) $t=9, 2^t=512$

(b) $\frac{c}{N} = \frac{427}{512} = \frac{1}{1+\frac{85}{427}} = \frac{1}{1+\frac{1}{5+\frac{2}{85}}} = \frac{1}{1+\frac{1}{5+\frac{1}{42+\frac{1}{2}}}} = [1, 5, 42, 2].$

• Restriction: $\left| \frac{c}{N} - \frac{l}{r} \right| \leq \frac{1}{2N^2} = \frac{1}{882} \rightarrow 0.83285 \leq \frac{l}{r} \leq 0.83512$

$[1]=1, \times \times$

$[1, 5] = \frac{1}{1+\frac{1}{5}} = \frac{5}{6} = 0.83333, \text{ WV } \Rightarrow \underline{\underline{r=6}}$

$[1, 5, 42] = \frac{1}{1+\frac{1}{5+\frac{1}{42}}} = \frac{47}{89}, r > N = 21, \times.$

(c) $f(x) = 2^6 = 64 \pmod{21} \equiv 1$

$f(x+r) = f(x) = f(x-r) \rightarrow \text{periodic}$

$\because f(1), f(2), f(3), f(4), f(5) \not\equiv 1 \quad \therefore \text{period is } \underline{\underline{r=6}}.$

$$2^6 - 1 = (2^3 + 1)(2^3 - 1) \pmod{21} = 0$$

$$\therefore 2^3 + 1 \pmod{21} \not\equiv 0$$

$$\therefore \text{Shor's algo holds, } \begin{aligned} \gcd(2^3 + 1, 21) &= 3, \quad \underline{3 \mid 7 = 21}, \\ \gcd(2^3 - 1, 21) &= 7 \end{aligned} \quad \#$$

Problem 3.

(a) let $y = \log_g(x)$, $g^y = x$ and $k = \log_g(c)$, $g^k = c$, then

$$f(a, b) = c \Leftrightarrow g^a \cdot g^{-by} \equiv g^k \pmod{p} \Leftrightarrow g^{a-(by+k)} \equiv 1 \equiv g^{p-1} \pmod{p} \Leftrightarrow a - (by+k) \equiv 0 \pmod{p}$$

(b)

$$\frac{1}{p-1} \sum_{a, b \in \mathbb{Z}_{p-1}} |a\rangle |b\rangle |f(a, b)\rangle \xrightarrow[\text{measure result}]{c_0} \frac{1}{\sqrt{p-1}} \sum_{b \in \mathbb{Z}_{p-1}} |by+k\rangle |b\rangle \xrightarrow[\text{if } d \text{ is min s.t. } x^d \equiv 1 \pmod{p}]{\substack{\text{if } d \text{ is min s.t. } x^d \equiv 1 \pmod{p} \\ \text{then } x^d = g^{dy} = 1 \equiv g^{p-1} \pmod{p} \\ y \mid p-1, m = \frac{p-1}{y}}} \frac{1}{\sqrt{y}} \sum_{b \in \mathbb{Z}_m} |by+k\rangle |b\rangle.$$

(c) QFT_{p-1}: $|x\rangle \rightarrow \frac{1}{\sqrt{p-1}} \sum_{u \in \mathbb{Z}_{p-1}} w^{xu} |u\rangle$.

$$\begin{aligned} |1\rangle &\rightarrow \frac{1}{\sqrt{y}} \sum_{b \in \mathbb{Z}_m} \left(\frac{1}{\sqrt{p-1}} \sum_{u \in \mathbb{Z}_{p-1}} w^{(by+k)u} |u\rangle \right) \left(\frac{1}{\sqrt{p-1}} \sum_{v \in \mathbb{Z}_{p-1}} w^{bv} |v\rangle \right) \\ &= \frac{1}{\sqrt{y}} \left(\frac{1}{\sqrt{p-1}} \sum_{u \in \mathbb{Z}_{p-1}} w^{ku} \left(\sum_{b \in \mathbb{Z}_m} w^{byu} \right) |u\rangle \right) \left(\frac{1}{\sqrt{p-1}} \sum_{v \in \mathbb{Z}_{p-1}} \left(\sum_{b \in \mathbb{Z}_m} w^{bv} \right) |v\rangle \right) \\ &= \left(\frac{1}{\sqrt{y}} \sum_{k=0}^{y-1} w^{kem} |km\rangle \right) (|0\rangle) \end{aligned}$$

$\xrightarrow{\text{measure result}}$ $(0, km)$ with $k = 0, 1, \dots, y-1$; each probability = $\frac{1}{y} = \frac{m}{p-1}$

(d). we have: $\frac{lm}{p-1} = \frac{y}{f}$... obtain denominator y in $O(\log^3(p))$ time by Continued fraction algo.

- Choosing g , g must be co-prime to $p-1$

\rightarrow we must run $O(\log \log(p))$ times to achieve a constant level prob q ... [week 8 ppt, P28]

$\xrightarrow{\text{Given } \epsilon_{\text{ex}}}$ For n repetition, successful prob = $(1-q)^n \geq (1-\epsilon)$, take $n = \lceil \log_{1-q} (1-\epsilon) \rceil$. #

This part need $O(\log \log(p) \cdot \lceil \log_{1-q} (1-\epsilon) \rceil)$ times.

QFT_{p-1} need $O(\text{poly}(\log p))$ times.

\Rightarrow total time is $O(\text{poly}(\log p))$ times. #

Problem 4.

(a) PVM

- $\langle \Pr(H_0=1|q_0) \rangle = \frac{1}{2}$, let $\Pi_0 = \begin{pmatrix} a & b \\ c & d \end{pmatrix}$, $\Pi_1 = \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}$ by unambiguous discrimination
 $\Pr(H_1=1|q_1) = \frac{1}{2}$
- $\langle \Pr(\text{outcome is } 0 | H_0=1|q_0) \rangle = \langle 0 | \Pi_0 | 0 \rangle \cdot \frac{1}{2} = \frac{1}{2} |a\rangle$
 $\Pr(\text{outcome is } 1 | H_1=1|q_1) = \langle + | \Pi_1 | + \rangle \cdot \frac{1}{2} = 0$
- $4 \cdot P_S = a + 0 = a \leq 1 \Rightarrow P_S \leq \frac{1}{4}$

(b) POVM

- $R = \frac{1}{2} |0\rangle\langle 0| + \frac{1}{2} |+\rangle\langle +| = \begin{pmatrix} \frac{1}{2} & 0 \\ 0 & 0 \end{pmatrix} + \begin{pmatrix} \frac{1}{4} & \frac{1}{4} \\ \frac{1}{4} & \frac{1}{4} \end{pmatrix} = \begin{pmatrix} \frac{3}{4} & \frac{1}{4} \\ \frac{1}{4} & \frac{1}{4} \end{pmatrix}$
- $\Pr(\text{outcome } 0) = \text{Tr}(R\Pi_0)$ $\Pr(\text{outcome } 1) = \text{Tr}(R\Pi_1)$ $\Pr(\text{outcome } *) = \text{Tr}(R\Pi_*)$
 $\Pr(\text{outcome } 0) + \Pr(\text{outcome } 1) + \Pr(\text{outcome } *) = \text{Tr}(R) = 1.$
- Take $\Pi_* = \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix}$, $\Pi_0 = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}$, $\Pi_1 = \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}$; $\Pr(\text{outcome } *) = \frac{1}{4}$, $P_S = \frac{1 - \frac{1}{4}}{2} = \frac{3}{8}$

(c) Since $P_S = \frac{1 - \text{Tr}(R\Pi_*)}{2}$ and $\text{Tr}(R\Pi_*) \geq 0$. When $\Pi_* = \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}$, P_S has max value = $\frac{1}{2}$

(d) Yes! In 4-dim, $\Pi_0 = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$, $\Pi_1 = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$, $\Pi_* = I - \Pi_0 - \Pi_1$.

$$\text{Then, } P_S = \frac{1}{4} (\langle 0 | \Pi_0 | 0 \rangle + \langle + | \Pi_1 | + \rangle) \\ = \frac{1}{4} \left(1 + \frac{1}{2} \right) = \frac{1}{4} \cdot \frac{3}{2} = \frac{3}{8} > \frac{1}{4}.$$

Problem 5.

(a) $\text{Tr}(I \otimes \sigma_x) = \text{Tr}(I \otimes \sigma_y) = \text{Tr}(I \otimes \sigma_z) = 0$
 $\text{Tr}(\sigma_x \otimes \sigma_y) = \text{Tr}\left(\begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}\right) = 0, \quad \text{Tr}(\sigma_x \otimes \sigma_z) = \text{Tr}\left(\begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix}\right) = 0, \quad \text{Tr}(\sigma_y \otimes \sigma_z) = \text{Tr}\left(\begin{pmatrix} 0 & 1 \\ -1 & 0 \end{pmatrix}\right) = 0.$

$\cdot \text{Tr}(\rho \sigma_\lambda) = \text{Tr}\left(\frac{1}{2} r_\lambda \sigma_\lambda \cdot \sigma_\lambda\right) = \frac{r_\lambda}{2} \text{Tr}(I) = r_\lambda, \quad \lambda = x, y, z. \#$

(b) $R = \frac{1}{2} \left[\left(\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} + \begin{pmatrix} 0 & r_x \\ r_x & 0 \end{pmatrix} + \begin{pmatrix} 0 & -r_y \\ r_y & 0 \end{pmatrix} + \begin{pmatrix} 1 & 0 \\ 0 & -r_z \end{pmatrix} \right) \right] = \frac{1}{2} \begin{pmatrix} 1+r_z & r_x-r_y \\ r_x+r_y & 1-r_z \end{pmatrix}$

4. $\det(R - \pi I) = (1-2\pi+r_z)(1-2\pi-r_z) - (r_x+\pi r_y)(r_x-\pi r_y)$
 $= (1-2\pi)^2 - r_z^2 - r_x^2 - r_y^2 = (1-2\pi)^2 - \|R\|^2 = 0 \Rightarrow \pi = \frac{1 \pm \|R\|}{2}$

R is density operator \Leftrightarrow eigenvalue $\pi_0, \pi_1 \geq 0 \Leftrightarrow \frac{1+\|R\|}{2}, \frac{1-\|R\|}{2} \geq 0 \Leftrightarrow \|R\| \leq 1$. #

(c) $\text{Tr}(\rho \sigma) = \frac{1}{2} \left[\text{Tr}(I \sigma) + r_x \text{Tr}(\sigma_x \sigma) + r_y \text{Tr}(\sigma_y \sigma) + r_z \text{Tr}(\sigma_z \sigma) \right]$
 $= \frac{1}{2} (1 + r_x s_x + r_y s_y + r_z s_z) = \frac{1}{2} (1 + \vec{r} \cdot \vec{s}) . \#$

$\cdot R$ is orthogonal to $\sigma \Leftrightarrow 0 = \text{Tr}(\rho \sigma) = \frac{1}{2} (1 + \vec{r} \cdot \vec{s}) \Leftrightarrow \vec{r} \cdot \vec{s} = -1 \dots$ 向量平行，反方向
球面上的两端，组成直径

(d) $P(\text{outcome is } 0) = \text{Tr}(\rho | \psi_0 \rangle \langle \psi_0 |) \xrightarrow{\text{by (c)}} \frac{1}{2} (1 + \vec{r} \cdot \vec{r}_0)$
 $P(\text{outcome is } 1) = \text{Tr}(\rho | \psi_1 \rangle \langle \psi_1 |) \xrightarrow{\text{by (c)}} \frac{1}{2} (1 + \vec{r} \cdot \vec{r}_1) . \#$

(e) $\sum_{a,b} \Pi_{(a,b)} = \frac{I + G_x + G_y + G_z - \sigma_x - \sigma_y - \sigma_z}{6} = I .$

$\cdot \Pi_{(a,b)}$ are positive semi-definite Hermitian matrices:

$$\text{Tr}(\Pi_{(a,b)}) = \frac{1}{6} [2 + (-1)^b \cdot 0] = \frac{1}{3}$$

$$\langle \frac{\Pi_{(a,b)}}{\Pi_{(a,b)}^T} \rangle = \frac{I + (-1)^b \overline{G_a}}{6} = \frac{I + (-1)^b G_a}{6} = \Pi_{(a,b)} . \#$$

(f) $P_{(a,b)} = \text{Tr}(\rho \Pi_{(a,b)}) = \frac{1}{2} [\text{Tr}(I \Pi_{(a,b)}) + r_x \text{Tr}(\sigma_x \Pi_{(a,b)}) + r_y \text{Tr}(\sigma_y \Pi_{(a,b)}) + r_z \text{Tr}(\sigma_z \Pi_{(a,b)})]$
 $= \frac{1}{2} \left[\frac{2}{6} + r_a \text{Tr}(\sigma_a \cdot \frac{I + (-1)^b G_a}{6}) \right]$
 $= \frac{1}{2} \left[\frac{2}{6} + r_a \text{Tr}(\frac{G_a + (-1)^b I}{6}) \right] = \frac{1}{2} \left[\frac{2}{6} + r_a \cdot 2 \cdot (-1)^b \right] = \frac{1 + (-1)^b r_a}{6} . \#$

(g) Given each copy ρ , we can use $\Pi_{(a,0)}$ to get $\frac{1+r_a}{6} \triangleq V$, and $r_a = 6V - 1$ to get r_a .
In this case, ρ is 2×2 matrix, has r_x, r_y, r_z ; we need 3 copy to obtain $V = (r_x, r_y, r_z)$.
[If ρ is $n \times n$ matrix, we need $n^2 - 1$ copy $\Leftrightarrow n^2 - 1$ freedom].