

# MARKING SCHEME DO NOT COPY

# **School of Computer Science and Engineering**

Module: Reasoning about Programs Module Code: 6SENG001W, 6SENG003C

Module Leader: Klaus Draeger

Date: 9<sup>th</sup> January 2019

Start: 10:00 Time allowed: 2 Hours

#### **Instructions for Candidates:**

You are advised (but not required) to spend the first ten minutes of the examination reading the questions and planning how you will answer those you have selected.

Answer ALL questions in Section A and TWO questions from Section B.

Section A is worth a total of 50 marks. Each question in section B is worth 25 marks.

In section B, only the TWO questions with the HIGHEST MARKS will count towards the FINAL MARK for the EXAM.

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## Section A

Answer ALL questions from this section.

## Question 1

```
(a) BakerlooStations \cap VictoriaStations = \{Oxford\_Circus\}
[1 mark]
```

- (b) CentralStations-VictoriaStations = { Bond\_Street, Tottenham\_Court\_Road } [1 mark]
- (c) card(CircleStations) = 3 [1 mark]
- (d)  $Baker\_Street \mapsto Central \in onLine = FALSE$  [1 mark]
- (e) ran(onLine) = TubeLines [1 mark]
- (f) colour(Central) = red [1 mark]
- (g)  $\bigcup$  ( { { Bond\_Street, Euston\_Square }, {}, { Warren\_Street } } ) = {Bond\_Street, Euston\_Square, Warren\_Street } [2 marks]
- (h)  $BakerlooStations \subseteq dom(onLine) = TRUE$  [2 marks]
- (i)  $CentralStations \cap dom(onLine) = CentralStations$  [2 marks]
- (j)  $\mathbb{P}(BakerlooStations)$

```
= { {}, { Baker_Street }, { Regents_Park }, { Oxford_Circus }, { Baker_Street, Regents_Park}, { Baker_Street, Oxford_Circus }, { Regents_Park, Oxford_Circus }, { Baker_Street, Regents_Park, Oxford_Circus } }
```

[3 marks]

[QUESTION Total 15]

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#### Question 2

See the *TubeSystem* B machine of the London Underground System, that is given in an Appendix of the Exam paper.

(a) The relation onLine in cannot be defined as a function because it contains maplets that map one element of the domain to more than one element in the range, e.g.

```
= \{ \ldots, Baker\_Street \mapsto Bakerloo, Baker\_Street \mapsto Circle, \\ Oxford\_Circus \mapsto Bakerloo, Oxford\_Circus \mapsto Central, \\ Oxford\_Circus \mapsto Victoria, \ldots \}
```

[2 marks]

[PART Total 2]

**(b)** Yes, the *colour* function can be defined as a total injective function,

```
colour \in TubeStation \rightarrow Colour
```

Because the domain equals the target set (total) & each element in the domain is mapped to a different element in the range, i.e. it is a total injective function. [3 marks]

[PART Total 3]

- **(c)** Evaluate the following expressions:
  - (i) onLine[ { Baker\_Street, Oxford\_Circus } ] = { Circle, Central, Bakerloo, Victoria } [2 marks]
  - (ii)  $CircleStations \triangleleft onLine$ 
    - $= \{ Baker\_Street \mapsto Circle, \ Baker\_Street \mapsto Bakerloo, \\ Great\_Portland\_Street \mapsto Circle, \ Euston\_Square \mapsto Circle \}$

[2 marks]

(iii)  $onLine > \{ Victoria \}$ =  $\{ Oxford\_Circus \mapsto Victoria, Warren\_Street \mapsto Victoria \}$ 

[2 marks]

(iv) onLine ⇒ { Bakerloo, Central, Victoria }
 = { Baker\_Street → Circle, Great\_Portland\_Street → Circle,
 Euston\_Square → Circle }
 [2 marks]

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(v) 
$$colour \Leftrightarrow \{ Circle \mapsto red, Central \mapsto yellow \}$$
  
=  $\{ Bakerloo \mapsto brown, Circle \mapsto red,$   
 $Central \mapsto yellow, Victoria \mapsto lightblue \}$  [2 marks]

(vi) 
$$colour^{-1}$$
  
= {  $brown \mapsto Bakerloo, \ yellow \mapsto Circle, \ red \mapsto Central, \ lightblue \mapsto Victoria$  } [2 marks]

[3 marks]

[QUESTION Total 20]

## Question 3

(a) An Abstract Machine is similar to the programming concepts of: modules, class definition (e.g. Java) or abstract data types. [1 mark]

An Abstract Machine is a specification of what a system should be like, or how it should behave (operations); but not how a system is to be built, i.e. no implementation details. [1 mark]

The main logical parts of an Abstract Machine are its: *name*, *local state*, represented by "encapsulated" variables, *collection of operations*, that can access & update the state variables. [3 marks]

[PART Total 5]

- **(b)** Explain the purpose of the following B Machine *clauses*:
  - VARIABLES declare state variable identifiers. [1 mark]
  - INVARIANT define the *state invariant* for the system, including the types of the variables & any additional constraints on them. [2 marks]
  - INITIALISATION initialise all the state variable with values that satisfy the state invariant. [2 marks]

[PART Total 5]

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(c) To add a tube passenger's current *location*, (station and line), to the *TubeSystem* B machine, need to add the following clauses:

```
\begin{tabular}{ll} VARIABLES & station, line \\ INVARIANT & station \in Stations \ \land \ line \in TubeLine \ \land \\ & station \mapsto line \ \in \ onLine \\ INITIALISATION & station := Oxford\_Circus \ \mid \ \mid \\ & line := Victoria \\ [5 marks] \end{tabular}
```

Any pair of station & line can be used to initialise the two variables, as long as they satisfy the invariant, i.e. are *consistent* with onLine.

[PART Total 5]

[QUESTION Total 15]

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# Section B

Answer TWO questions from this section.

#### Question 4

The LuggageRack B machine; basically its a **stack**. So would expect something similar to the following machine, but a student's version is unlikely to include all the details included in this *solution*, as long as its a stack & has the main parts will not penalise for minor errors, e.g. syntax errors, etc.

#### (a) MACHINE LuggageRack

```
SETS
 LUGGAGE = { case1, case2, case3, case4, case5,
               bag1, bag2, bag3, bag4, bag5,
               null_bag } ;
 ANSWER = \{ Yes, No \} ;
 MESSAGE = { Luggage_Added,
                              ERROR_No_Space_Left,
              Luggage_Removed, ERROR_No_Luggage_Left }
CONSTANTS
 MaxItemsOfLuggage, No_Luggage, EMPTY_LuggageRack
PROPERTIES
 MaxItemsOfLuggage : NAT1
                                   & MaxItemsOfLuggage = 5
                   : LUGGAGE
                                   & No_Luggage = null_bag
 No_Luggage
 EMPTY_LuggageRack : seq(LUGGAGE) & EMPTY_LuggageRack = []
VARIABLES
 luggageRack
INVARIANT
 luggageRack : seq( LUGGAGE ) &
 size( luggageRack ) <= MaxItemsOfLuggage</pre>
INITIALISATION
 luggageRack := EMPTY_LuggageRack
```

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```
Roughly award: SETS [2\ marks], CONSTANTS & PROPERTIES [3\ marks], VARIABLES & INVARIANT [3\ marks], INITIALISATION [1\ mark].
```

### [PART Total 9]

(ii)

(b) (i) In this version the "top" of the luggage rack (stack) is the front of the sequence, but okay if the end. Might use strings for reporting rather than MESSAGE.

**OPERATIONS** 

```
report <-- AddLuggage( luggage ) =</pre>
      PRE
          report : MESSAGE & luggage : LUGGAGE
      THEN
          IF ( size(luggageRack) < MaxItemsOfLuggage )</pre>
          THEN
              luggageRack := luggage -> luggageRack ||
                           := Luggage_Added
              report
          ELSE
              report := ERROR_No_Space_Left
          END
      END ;
[PART Total 6]
  report, topcase <-- RemoveLuggage =
      PRE
```

```
report, topcase <-- RemoveLuggage =
    PRE
        report : MESSAGE & topcase : LUGGAGE
    THEN
        IF ( luggageRack /= EMPTY_LuggageRack )
        THEN
            luggageRack := tail(luggageRack) ||
                 report := Luggage_Removed ||
                      topcase := first(luggageRack)
        ELSE
            report := ERROR_No_Luggage_Left ||
                     topcase := No_Luggage
        END
        END
        END
        [PART Total 7]</pre>
```

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```
(iii)
             answer <-- AnyLuggageLeft =</pre>
                 PRE
                     answer : ANSWER
                 THEN
                           ( luggageRack /= EMPTY_LuggageRack )
                     IF
                     THEN
                           answer := Yes
                     ELSE
                          answer := No
                     END
                 END
          END /* LuggageRack */
          [PART Total 3]
     [PART Total 16]
[QUESTION Total 25]
```

# **Question 5**

See Exam paper Appendix for the BirthdayBook B machine.

- (a) The B type of DATE is an element of the Cartesian product of days & months, i.e. an ordered pair (or maplet). [1 mark]

  Today's date is represented:  $9 \mapsto Jan$  (exam date). [1 mark]

  [PART Total 2]
- (b) The B type of known is a set of names. Example value is any subset of NAME, e.g. known = { Jim, Sue, Mon, Zoe }. [1 mark] [PART Total 1]
- (c) birthday is a partial function because the birthday book will not always hold everyone's birthday & everyone has just one birthday. [1 mark] It is not:
  - a relation because no one has more than one birthday, (excluding the Queen). [1 mark]
  - a total function since the birthday book does not always hold everyone's birthday. [1 mark]

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- an injective function since several people can have the same birthday. [1 mark]
- a surjective function since not every possible date will be recorded as someone's birthday. [1 mark]
- a bijections, since its neither an injective, surjective or total function. [1 mark]

#### [PART Total 6]

- (d) birthday's constraints:
  - card( birthday ) <= maximum
     the birthday book has a maximum limit on how many birthdays it can
     record. [1 mark]</li>
  - known = dom( birthday )
     the known people are those that have a birthday recorded.
    [1 mark]
  - NonDate /: ran( birthday )
     no one can have a birthday recorded as being on the 31<sup>st</sup> February.
     [1 mark]

### [PART Total 3]

- **(e)** *Pre-conditions* for the following operations:
  - (i) AddBirthday
    The name must be a proper name. [1 mark] The date must be valid date & it cannot be 31<sup>st</sup> February. [2 marks] There must be room in the birthday book to record at least one more birthday. [1 mark]

[SUBPART Total 4]

(ii) Reminder

The date must be valid date & it cannot be  $31^{st}$  February. [1 mark] Only a collection (set) of real names can be output. [1 mark]

[SUBPART Total 2]

(iii) NumKnownBirthdays

Since there is no explicit pre-condition for this operation, it is implicitly just "TRUE". [2 marks]

[PART Total 2]

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#### [PART Total 8]

- (f) AddBirthday updates the state as follows:
  - Provided the person's name is not already in the birthday book it updates the state by adding the name to the known people & adds the name and date pair to the birthday book. [1 mark]
  - If the person's name is already in the birthday book it does not add them again, i.e. it does not update the state. [1 mark]

#### [PART Total 2]

(g) The state to produce the ProB output: report = Birthdays\_On\_Date & cards = { Jim, Sue } for "Reminder( 10 |-> Jul )" is any pair of known & birthday that satisfy:

```
\{Jim, Sue\} \subseteq known
\{Jim \mapsto (10 \mapsto Jul), Sue \mapsto (10 \mapsto Jul)\} = birthday \triangleright \{10 \mapsto Jul\}
```

[3 marks]

[PART Total 3]

[QUESTION Total 25]

# Question 6

Marking Scheme for Hoare Logic & Program Verification.

(a) The Hoare triple

$$[x = 0] \ y := z \ [z = x + y]$$

means that executing the instruction y := z (i.e. assigning the value of z to y), starting from a state in which x is 0, leads to a state in which z equals x + y. [2  $\mathbf{marks}$ ]

#### [SUBPART Total 2]

(b) (i)  $[x < y] \ y := 0 \ [x < 0]$  is invalid.  $[\mathbf{1} \ \mathbf{mark}]$  Counterexample: Starting in a state with x = 1, y = 2 leads to a state with x = 1, y = 0.  $[\mathbf{1} \ \mathbf{mark}]$  [SUBPART Total 2]

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- (ii)  $[x < y] \ y := y+1 \ [x < y+1]$  is valid: the pre-condition is x < y+2, which follows from x < y. [2 marks]  $[\mathbf{SUBPART\ Total\ 2}]$
- (iii)  $[x < y] \ y := y 1 \ [x < y 1]$  is invalid.  $[\mathbf{1} \ \mathbf{mark}]$  Counterexample: Starting in a state with x = 1, y = 2 leads to a state with x = 1, y = 0.  $[\mathbf{1} \ \mathbf{mark}]$  [SUBPART Total 2]
- (iv) [true] x := 0 [true] is valid, since any post-state satisfies true. [2 marks] [SUBPART Total 2]

[PART Total 10]

- (c) The intermediate assertions are
  - 1. 0 < z [2 marks]
  - 2. 0 < y + z [2 marks]
  - 3. x < y + z [2 marks]

[PART Total 6]

$$\begin{array}{lll} \textbf{(d)} & [y=10] \\ & x:=0; \\ & [x+y=10 \ \& \ y \geq 0] & \textbf{[3 marks]} \\ & WHILE \ y > 0 \ DO \\ & [x+y=10 \ \& \ y > 0] & \textbf{[2 marks]} \\ & x:=x+1; \\ & [x+y=11 \ \& \ y > 0] & \textbf{[2 marks]} \\ & y:=y-1 \\ & [x+y=10 \ \& \ y \geq 0] & \textbf{[2 marks]} \\ & END \\ & [x=10] \end{array}$$

[PART Total 9]

[QUESTION Total 25]