CODEPATH*ORG

Welcome to class!

Turn on your camera 👝

Rename yourself to include your pod # at the beginning of your Zoom name - ie 6 - Emily Garvey (she/her)

Answer the following question in the chat: What is your favorite TV show | book | thing?



Agenda

Check In 5 mins

Core Data Structures: 25 mins

Stacks and Queues

Breakout Sessions 25 mins

BREAK 5 mins

Breakout Sessions 30 mins

Go Over Problem 1 25 mins

Wrap Up 5 mins

Announcements

☐ TA Office Hours



SE103 TA Office Hours Logistics

- When are Office Hours held, and who hosts?
 - Mondays from 10-11am PT / 1-3pm ET with Emily Crowl
 - Thursdays from 5-6pm PT / 8-9pm ET with Suresh Venkatesan
- How do I join the meeting?
 - You can find the zoom link in the <u>"Schedule" page</u> of the course portal
- Are they mandatory?
 - No, but we would love for you to join, even if you don't have a specific question in mind
- Are they recorded?
 - Yes! The playlist link is also in the <u>"Schedule" page</u>, and the <u>"Resources"</u>



TA Office Hours Tips

- What kinds of things will the TAs go over during Office Hours?
 - TAs will usually prepare questions from the current or previous week -- including ones in the Session tabs, Resources tab, or from the assignments.
- How can I make sure the TA goes over the topic I want them to cover?
 - Ask them ahead of time! If you DM the TA before Office Hours, you can ensure they
 will prepare the question you want to discuss.
- How can I make the most of Office Hours?
 - Attend!! Even if you think you don't have a question, make it a routine to show up.
- What if I can't attend Office Hours?
 - You can post any questions you have in the #help slack channel. Our instructors, TAs, mentors, and fellow students all may be able to help. Posting code snippets, links, or specific questions are encouraged (just don't give away HackerRank answers before the due date!)

Core Data Structures: Stacks and Queues | 30 mins



A quack guide to collections





What is a Stack?

Stacks stores objects in a last in, first out (LIFO) fashion

Think of it in terms of stacks of plates





How does a stack work?

- The last thing added (pushed) is the first thing to be retrieved (popped)
- A stack is a sequence of items that are accessible at only one end of the sequence
- Operations that can be performed on a stack:
 - push: add an item to the top of the stack
 - pop: remove the item at the top of the stack
 - top/peek: get (but do not remove) the item at the top of the stack



Stack operations

```
import java.util.Stack;

Stack<Object> stack = new Stack<>();
stack.add(1);
stack.add(2);
System.out.println(stack.pop()); // 2
System.out.println(stack.peek()); // 1
if (!stack.isEmpty()) {
    stack.pop(); // 1
}
```

```
import collections

stack = collections.deque()
stack.append(1)
stack.append(2)
print(stack.pop()) # 2
print(stack[-1]) # 1
if stack:
    stack.pop() # 1
```



Postfix Notation

- Infix expression is the form AOB
 - A and B are numbers or also infix expression
 - O is operator (+, -, *, /)
- Postfix expression is the form ABO
 - A and B are numbers or also postfix expression
 - O is operator (+, -, *, /)



Expression notations

Infix notation

- A < op> B
- Examples

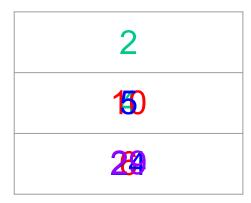
Postfix notation

- AB <op>
- Examples



Expression notations





Postfix notation

- A B <op>
- Examples



Expression notations



Postfix notation

- A B <op>
- Examples

- Advantage over infix
 - Parentheses never needed
 - Easier to evaluate



What is a Queue?

Queues store objects in a first in, first out (FIFO) fashion.

Think of it like waiting in line to cast your ballot.





Queue operations

```
import java.util.Queue;

Queue<Object> queue = new LinkedList<>();
queue.add(1);
queue.add(2);
queue.remove(); // 1
if (!queue.isEmpty()) {
   obj = queue.remove(); // 2
}
```

```
import collections
queue = collections.deque()
queue.append(1)
queue.append(2)
queue.popleft() # 1
if queue:
    print(queue.popleft()) # 2
```



Queues: Key Takeaways

- Useful when ordering of the data matters as it preserves that ordering.
- Make sure you know the right classes and methods for your language.



Tips on what data structure to use

- Learning what data structure to use takes practice.
- When reviewing a data structure:
 - How does this compare to other data structures
 I know?
 - What does this data structure do well?
 - What does this data structure do poorly at?
- When you read a solution to a coding problem, spend time analyzing why a data structure was optimal for the problem.



	finding max/min value	lookup	ordering by time
Неар			
Hash			
Stack			
Queue			



	finding max/min value	lookup	ordering by time
Неар	+	-	-
Hash			
Stack			
Queue			



	finding max/min value	lookup	ordering by time
Неар	+	-	-
Hash			
Stack			
Queue			



	finding max/min value	lookup	ordering by time
Неар	+	-	-
Hash	-	+	-
Stack			
Queue			



	finding max/min value	lookup	ordering by time
Неар	+	-	-
Hash	-	+	-
Stack			
Queue			



	finding max/min value	lookup	ordering by time
Неар	+	-	-
Hash	-	+	-
Stack	-	-	LIFO
Queue			



	finding max/min value	lookup	ordering by time
Неар	+	-	-
Hash	-	+	-
Stack	-	-	LIFO
Queue			



	finding max/min value	lookup	ordering by time
Неар	+	-	-
Hash	-	+	-
Stack	-	-	LIFO
Queue	-	-	FIFO



Comparison: the fine print

	finding max/min value	lookup	ordering by time
Неар	+	-	-
Hash	-	+	-
Stack	-	-	LIFO
Queue	-	-	FIFO

- There are types of hash sets/tables that keep track of:
 - o relative ordering of entries (TreeMap and TreeSet in Java)
 - when items were added (LinkedHashSet in Java)
 - as of python 3.7, dictionaries maintain order of when items were added

In Class Walkthrough | 15 mins



https://leetcode.com/problems/remove-outermost-parentheses/

A valid parentheses string is either:

- empty ""
- "(" + A + ")"
- or A + B

where A and B are valid parentheses strings, and + represents string concatenation.

For example:

"", "()", "(())()", and "(()(()))" are all **valid parentheses strings**.



https://leetcode.com/problems/remove-outermost-parentheses/

A valid parentheses string s is primitive if it is nonempty, and there does not exist a way to split it into s = A + B, with A and B nonempty valid parentheses strings.

Given a valid parentheses string s, consider its primitive decomposition:

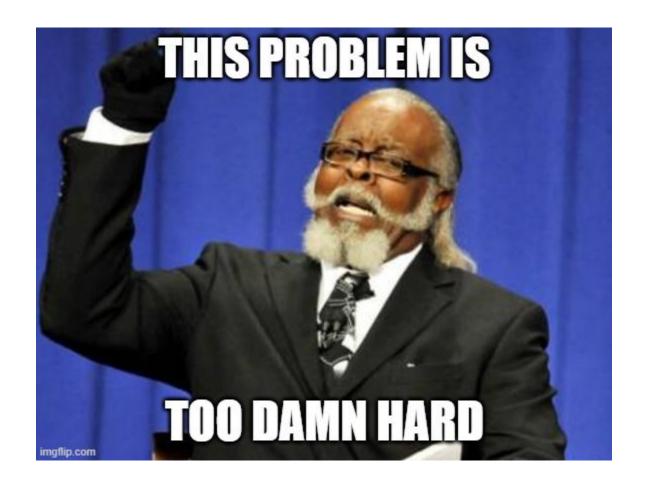
s = P1 + P2 + ... + Pk, where Pi are primitive valid parentheses strings.

Return s after removing the outermost parentheses of every primitive string in the primitive decomposition of s

Constraints:

1 <= s.length <= 105 s[i] is either '(' or ')' s is a valid parentheses string







https://leetcode.com/problems/remove-outermost-parentheses/

Example 1:

Input: s = "(()())(())"

Output: "()()()"

Explanation:

The input string is "(()())(())", with primitive decomposition "(()())" + "(())"After removing outer parentheses of each part, this is "()()" + "()" = "()()()"



https://leetcode.com/problems/remove-outermost-parentheses/

Example 2:

Input: s = "(()())(()(()()))"

Output: "()()()(())"

Explanation:

The input string is "(()())(()(()))", with primitive decomposition "(()())" + "(()())" + "(()(()))".

After removing outer parentheses of each part, this is "()()" + "()" + "()(())" = "()()()(())"



https://leetcode.com/problems/remove-outermost-parentheses/

Example 3:

Input: s = "()()"

Output: ""

Explanation:

The input string is "()()", with primitive decomposition "()" + "()"

After removing outer parentheses of each part, this is "" + "" = ""



M-atch

Make use of a stack and a set

- Use a stack to operate the string
- Or... instead of using set to store the index to remove, we can directly use list slice to append the string to answer
- Skipping first and last element of the valid sub-parentheses requires some extra attention.



P-lan:

- 1. Use a stack to operate the string
- 2. Append "(" when there is "(" and pop "(" there is meet ")"



P-lan:

- 3. We also need a set to store the indexes that need to remove
- 4. After pop, once the stack is empty we add the last value index and current value index into set.
- 5. Then we iterate the string again
- 6. Skip those index in set and concatenate others character to the answer.



I-mplement

```
def removeOuterParentheses(self, S: str) -> str:
        stack = []
        index remove=set()
        res=""
        for i, v in enumerate(S):
            if v == "(":
                stack.append(i)
            else:
                left index = stack.pop()
                if not stack:
                    index remove.add(left index)
                    index remove.add(i)
        for i, v in enumerate(S):
            if i not in index remove:
                res+=v
         return res
```



R-eview:

- Trace through your code with an input to check for the expected output.
- Catch possible edge cases and off-by-one errors.



E-valuate

Time complexity

O(2n), since we iterate the list twice so the total cost will be roughly

Space complexity

• O(n), since we need to store items in stack



Breakout Sessions | 60 mins



Problem #1: Brick Wall

Original Problem on Leetcode: Brick Wall

There is a rectangular brick wall in front of you with n rows of bricks. The ith row has some number of bricks each of the same height (i.e., one unit) but they can be of different widths. The total width of each row is the same.

Draw a vertical line from the top to the bottom and cross the least bricks. If your line goes through the edge of a brick, then the brick is not considered as crossed. You cannot draw a line just along one of the two vertical edges of the wall, in which case the line will obviously cross no bricks.

Given the 2D array wall that contains the information about the wall, return the minimum number of crossed bricks after drawing such a vertical line.

Example:



```
Input: wall = [[1,2,2,1],[3,1,2],[1,3,2],[2,4],[3,1,2],[1,3,1,1]]
Output: 2
```



Activity | Breakout Groups



Directions

- IceBreaker: You get a million dollars but you must legally change your first name to Buttergunt for the rest of your life. Do you do it? Why?
- Work together through the UMPIRE steps of the problems in the Course Portal under Week
 Session 2 Brick Wall
- 3. If your pod needs help, post a message on the slack channel and tag @se103-tas

Reminders

Don't forget to turn on your cameras!

Take a 5 min break sometime during your breakout session!



Breakout Problem Review | 15 mins

Don't peek at these until I present them!



U-nderstand

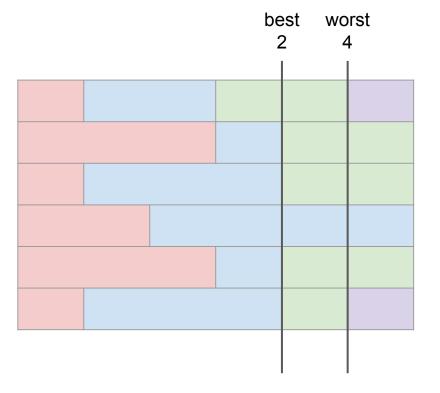
```
Input: wall = [
    [1,2,2,1],
    [3,1,2],
    [2,4],
    [3,1,2],
    [1,3,1,1]
]
Output: 2
```

1	2		2		1	
3			1	2		
1		3		2		
2			4			
3			1	2		
1		3		1	1	



U-nderstand

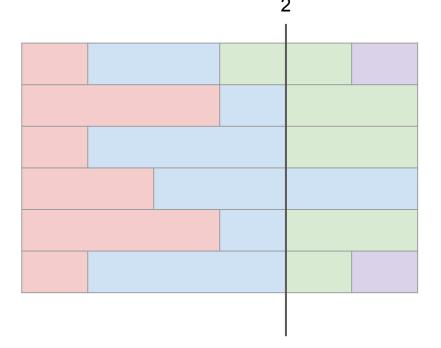
```
Input: wall = [
    [1,2,2,1],
    [3,1,2],
    [2,4],
    [3,1,2],
    [1,3,1,1]
]
Output: 2
```





U-nderstand: Test Case 1

```
Input: wall = [
    [1,2,2,1],
    [3,1,2],
    [2,4],
    [3,1,2],
    [1,3,1,1]
]
Output: 2
```

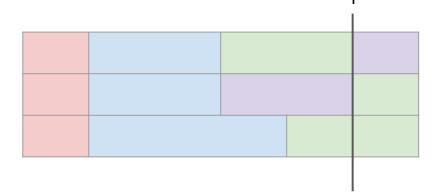


best



U-nderstand: Test Case 2

```
Input: wall = [
    [1,2,2,1],
    [1,2,2,1],
    [1,3,2],
]
Output: 1
```

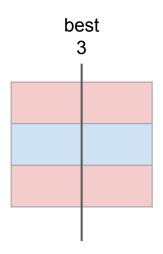


best



U-nderstand: Test Case 3 (edge case)

```
Input: wall = [
   [2],
   [2],
   [2],
]
Output: 3
```



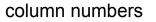


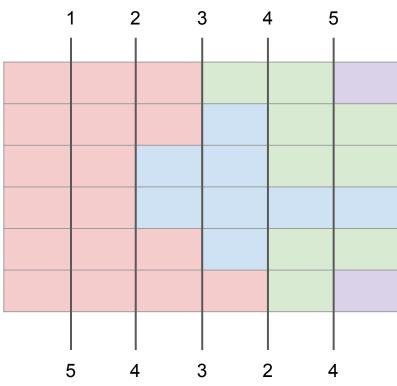
U-nderstand: Test Case 4 (edge case)

```
Input: wall = []
Output: 0
```



M-atch





crossed bricks

- 1. Build a representation of the wall.
- 2. Find the column with the fewest bricks.



M-atch: Representation

column numbers

Constraints: n == wall.length $1 <= n <= 10^4$ 1 <= wall[i].length <= 104 1 <= sum(wall[i].length) <= 2 * 104 sum(wall[i]) is the same for each row i. $1 \le \text{wall[i][j]} \le 2^{31} - 1$

crossed bricks

We'll need to keep count of a number of bricks spanning each column.

How about an array?





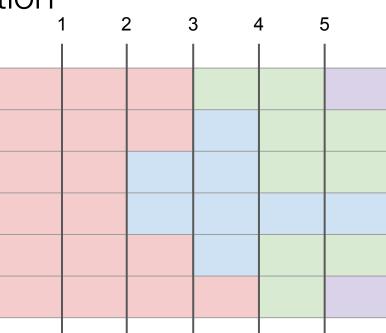
```
Input: wall = [
  [3,2,1],
  [3,1,2],
  [2,2,2],
  [2,4],
  [3,1,2],
  [4,1,1]
Output: 2
```

We'll need to keep count of a number of bricks spanning each column.

How about an array? No

Consider a hash table instead.

That's still a lot of entries. Maybe there will be fewer if we keep track of gaps instead.



column numbers

crossed bricks

3

5

4





Input: wall = [[3,2,1],[3,1,2],[2,2,2], [2,4], [3,1,2],[4,1,1]

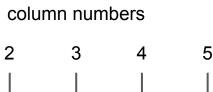
Output: 2

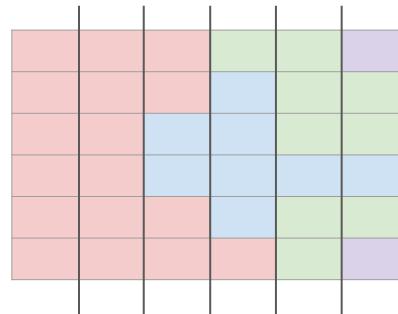
We'll need to keep count of a number of bricks spanning each column.

How about an array? No

Consider a hash table instead.

That's still a lot of entries. Maybe there will be fewer if we keep track of gaps instead.





crossed gaps

3



P-lan

- 1. Build a hash table to keep track of gaps per column.
- 2. For each row in rows:
 - a. Set x to 0
 - b. For each brick in the row:
 - i. x = x + brick.width
 - ii. Increment gaps[x]
- 3. Find the maximum number of gaps (max value in gaps).
- 4. Return {number of rows} {max number of gaps} to get fewest bricks crossed.



I-mplementation

```
import java.util.*;
public int leastBricks(List<List<Integer>> wall) {
   Map<Integer, Integer> map = new HashMap<>();
    for (List<Integer> row : wall) {
        int x = 0:
        for (int i = 0; i < row.size() - 1; i++) {
            x += row.get(i);
            int gaps = map.getOrDefault(x, 0) + 1;
            map.put(x, gaps);
    return wall.size() - Collections.max(map.values());
```

```
import collections
def leastBricks(self, wall):
    gaps = defaultdict(int)
    for bricks in wall:
        x = 0
        for brick in bricks[:-1]:
            x += brick
            gaps[x] += 1
   maxGaps = max(gaps.values())
    return len(wall) - maxGaps
```



R-eview

- Test case 1: pass
- Test case 2: pass
- Test case 3: fail
- Test case 4: fail

```
Test case 3
Input: wall = [
    [2],
    [2],
    [2],
]
Output: 3
```

```
Test case 4
Input: wall = [ ]
Output: 0
```



I-mplementation (corrected)

```
public int leastBricks(List<List<Integer>> wall) {
   Map<Integer, Integer> map = new HashMap<>();
    for (List<Integer> row : wall) {
        int x = 0;
        for (int i = 0; i < row.size() - 1; i++) {
           x += row.get(i);
            int gaps = map.getOrDefault(x, 0) + 1;
           map.put(x, gaps);
    if (map.size() > 0) {
        return wall.size() -
           Collections.max(map.values());
    } else {
        return wall.size();
```

```
def leastBricks(self, wall):
    gaps = defaultdict(int)

for bricks in wall:
    x = 0
    for brick in bricks[:-1]:
     x += brick
     gaps[x] += 1

maxGaps = max(gaps.values()) if gaps else 0
    return len(wall) - maxGaps
```



R-eview (2nd time)

- Test case 1: pass
- Test case 2: pass
- Test case 3: pass
- Test case 4: pass



E-valuate: time

- 1. Build a hash table to keep track of gaps per column.
- 2. For each row in rows:
 - a. Set x to 0

O(n), where n is the number of bricks.

- b. For each brick in the row:
 - i. x = x + brick.width
 - ii. Increment gaps[x]
- 3. Find the maximum number of gaps (max value in gaps).
- 4. Return {number of rows} {max number of gaps} to get fewest bricks crossed.



E-valuate: space

- 1. Build a hash table to keep track of gaps per column.
- 2. For each row in rows:
 - a. Set x to 0
 - b. For each brick in the row:
 - i. x = x + brick.width
 - ii. Increment gaps[x]
- 3. Find the maximum number of gaps (max value in gaps).
- 4. Return {number of rows} {max number of gaps} to get fewest bricks crossed.

O(g), where g is the number of columns with gaps



Wrap Up | 10 mins



Activity | Exit Ticket





Answer the 3 multiple choice in the Zoom Poll.



Exit tickets are a great way for us to check your understanding of the topics we discussed in class and identify for you what you should review after class.



Exit Ticket: Question #1

What are **heaps** best for?

- a) Keeping track of the maximum or minimum value
- b) Quickly checking whether they contain a specific value
- c) Retrieving values based on when they were added



Exit Ticket: Question #2

What are **hash tables** best for?

- a) Keeping track of the maximum or minimum value
- b) Quickly checking whether they contain a specific value
- c) Retrieving values based on when they were added



Exit Ticket: Question #2

What are **stacks and queues** best for?

- a) Keeping track of the maximum or minimum value
- b) Quickly checking whether they contain a specific value
- c) Retrieving values based on when they were added



Shout Outs!



Take a moment to shoutout someone today who helped you out.

Alternatively, drop in the chat something new that you were excited to learn about today!



HackerRank Reminders

- You must use the same email you provided CodePath during application for your submission to be tracked
- → HRs are due the day before Session #1 of the following week
 - Wed/Sat class: Tuesdays at 11:59pm PDT
- ☐ You are allowed 2 missing/submitted late HRs, no questions asked

At CodePath we collect your scores as a way to track your progress and provide you feedback throughout the program. You are not graded in this course, these assignments are made solely for practice only.

All information is on the **assignment tab** on the course portal.



Before you Leave |



- Complete the Session Survey [5 min]
- ☐ Next session is Wed 09/28 5pm pst
- Complete your HackerRank

assessment by 1 day before the next

session at 11:59pm PDT

Appendix



Comparison: the fine print

	finding max/min value	lookup	ordering by time
Неар	+	-	-
Hash	-	+	-
Stack	-	-	+ (LIFO)
Queue	-	-	+ (FIFO)

