

# Theory of Computation

## (Introduction)

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# Hand-axe



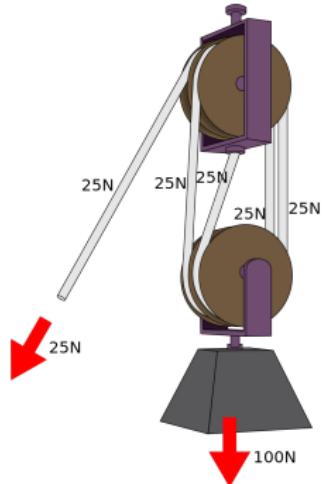
Longest-used tool in human history (1.5 million years)

# Wheel

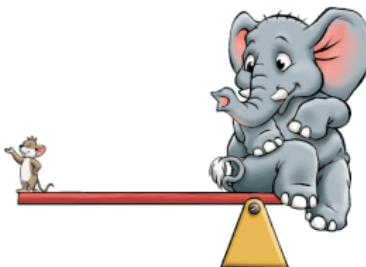


Idea behind **transportation revolution**  
(Other uses: potter's wheel, steering wheel, flywheel)

# Simple machines



Pulley



Lever



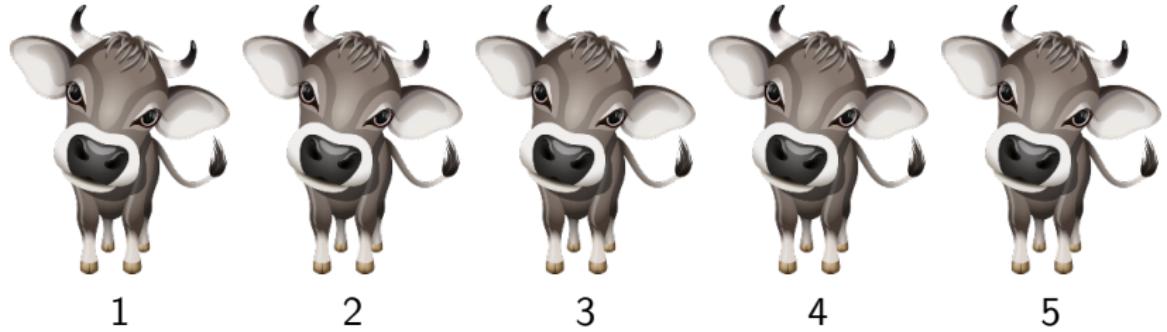
Screw

# Machines



# Computing

# Counting cattle

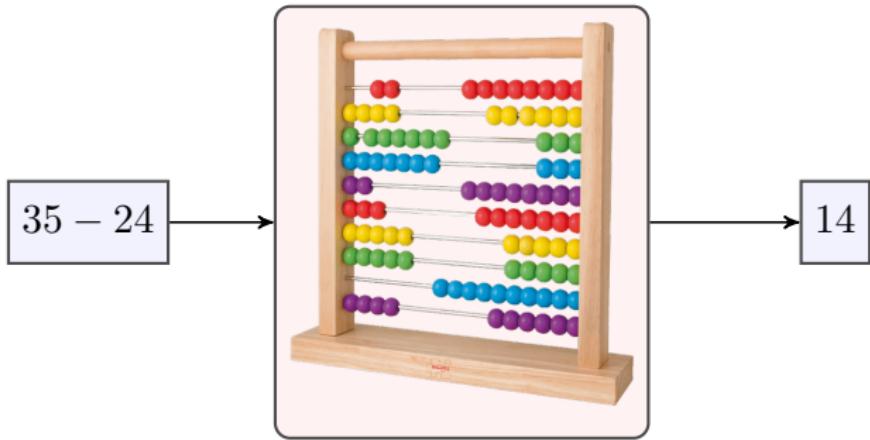


# Machine for computing



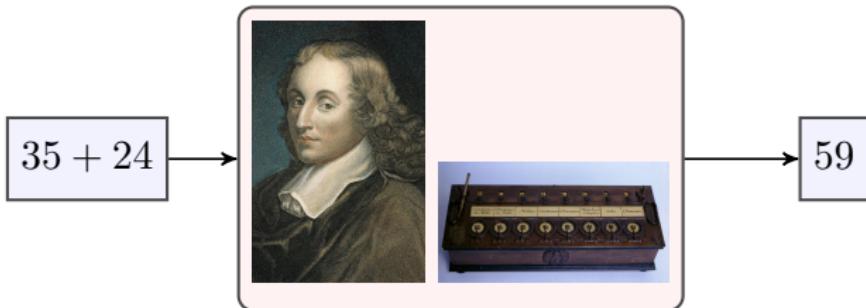
- How do humans compute or calculate or solve problems?
- Is it possible to build a **computing machine** that can mechanically (i.e., without thinking) simulate the computations performed by a human brain like that of Galileo or Newton or Einstein?
- If so, what problems can or cannot be solved by such a computing machine?

# < 2000 BC: Abacus



- Not automatic
- Operations: **Addition, subtraction, multiplication, and division**

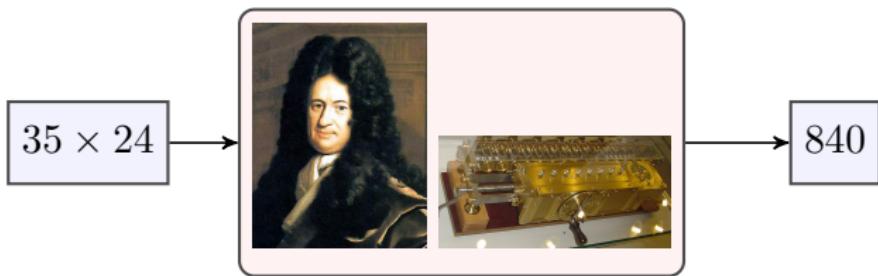
# 1643: Pascal's calculator (Pascaline)



Source: Computer Museum History Center

- Inventor: Blaise Pascal
- Operations: **Addition and subtraction**
- World's first mechanical calculator

# 1694: Leibniz' calculator (Step reckoner)



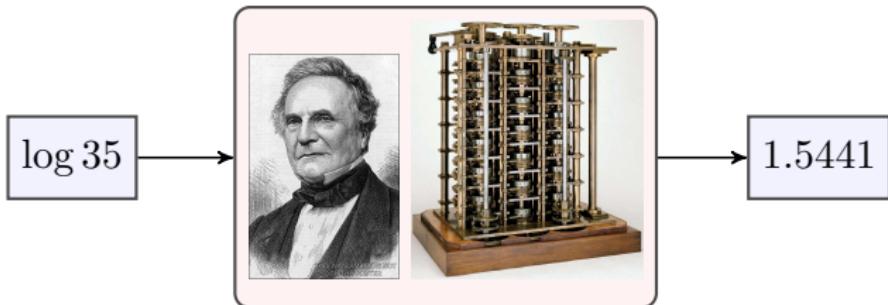
- Inventor: Gottfried Wilhelm Leibniz
- Operations: Addition, subtraction, multiplication, and division

# 1820: Colmar's calculator (Arithmometer)



- Inventor: Thomas de Colmar
- Operations: Addition, subtraction, multiplication, division, square root, involution, resolution of triangles, etc
- Applications: Financial organizations

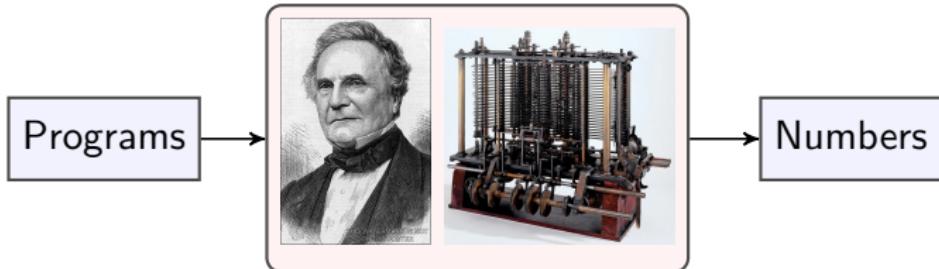
# 1822: Babbage's calculator (Difference engine)



Source: Science Museum London

- Designer: Charles Babbage
- The system was never built due to conflicts and insufficient funding
- Operations: **Addition, subtraction, multiplication, division, logarithmic, trigonometric functions, etc**

# 1833: Babbage's computer (Analytical engine)



Source: Science Museum London

- Designer: Charles Babbage
- The system was never built due to conflicts and insufficient funding
- World's first general-purpose computer (Turing-complete)
- Components: arithmetic logic unit, control flow in the form of conditional branching and loops, and integrated memory

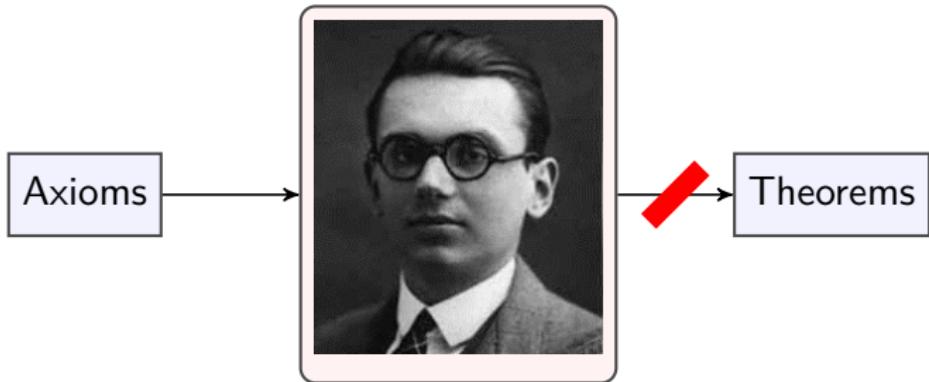
# 1843: Lovelace's algorithm



Pic by: Antoine Claudet

- Designer: Ada Lovelace
- World's first programmer
- Published the first algorithm to be implemented on a computer
- The algorithm was used to compute Bernoulli numbers

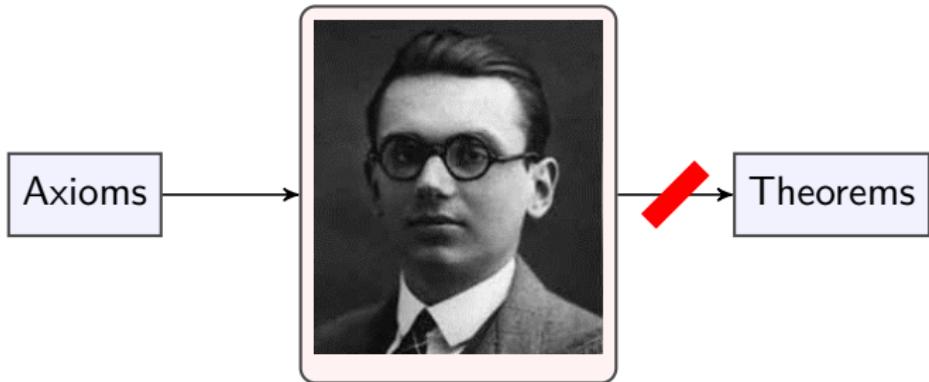
# 1931: Gödel's proof



Source: geni.com

- Discoverer: Kurt Gödel
- Some mathematical truths cannot be proved

# 1931: Gödel's proof



Source: geni.com

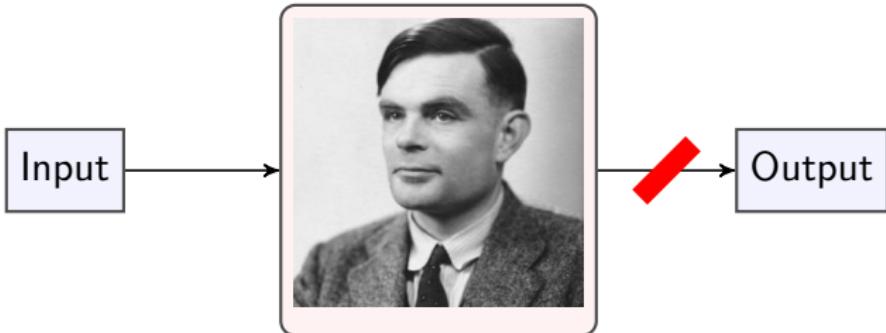
- Discoverer: Kurt Gödel
- Some mathematical truths cannot be proved
  - (If you cannot prove a mathematical statement, then how do you know that the statement is true?)

# 1936: Turing machine



- Discoverer: Alan Mathison Turing
- Creator of computer science
- Turing machine – the simplest, the most intuitive, the most generic, and the most powerful mathematical model of a computing human brain and a computer
- Algorithm and computation

# 1936: Turing's proof



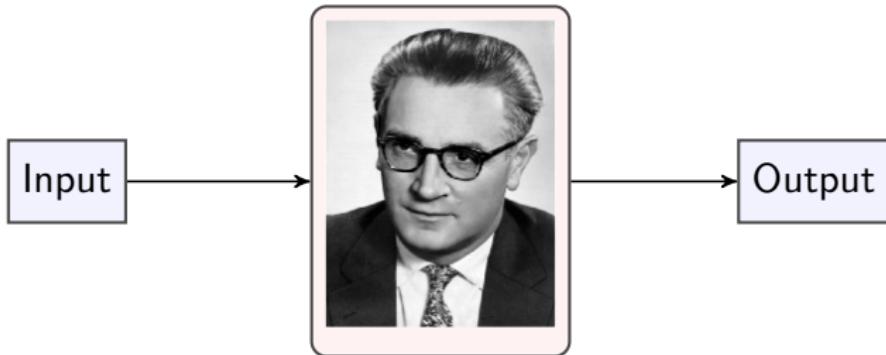
- Discoverer: Alan Mathison Turing
- Some computational problems cannot have algorithms

# 1936: Turing's proof



- Discoverer: Alan Mathison Turing
- Some computational problems cannot have algorithms  
(If you cannot mechanically compute a computational problem, then why is it called a computational problem?)

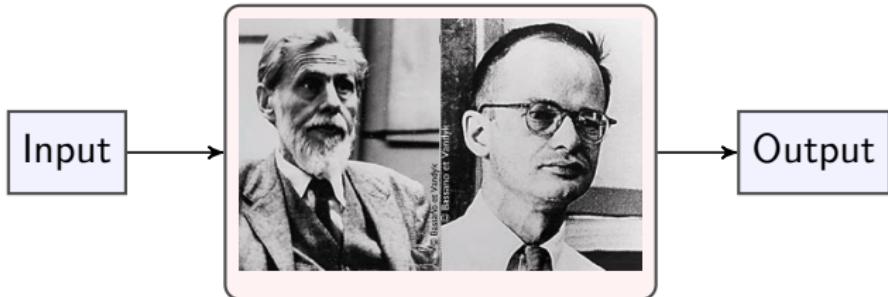
# 1941: Zuse's Z3



Source: <http://www.horst-zuse.homepage.t-online.de/>

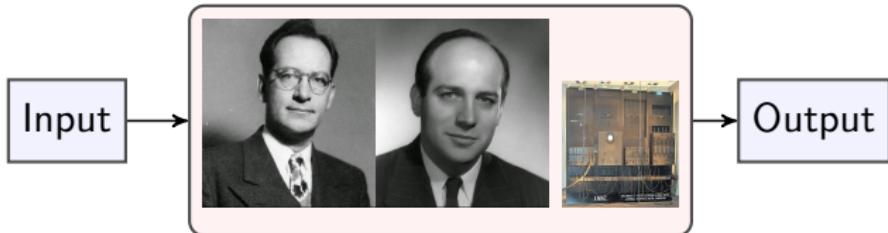
- Designer: Konrad Zuse
- World's first working programmable, fully automatic digital computer (Turing-complete)

# 1943: McCulloch and Pitts' finite automata



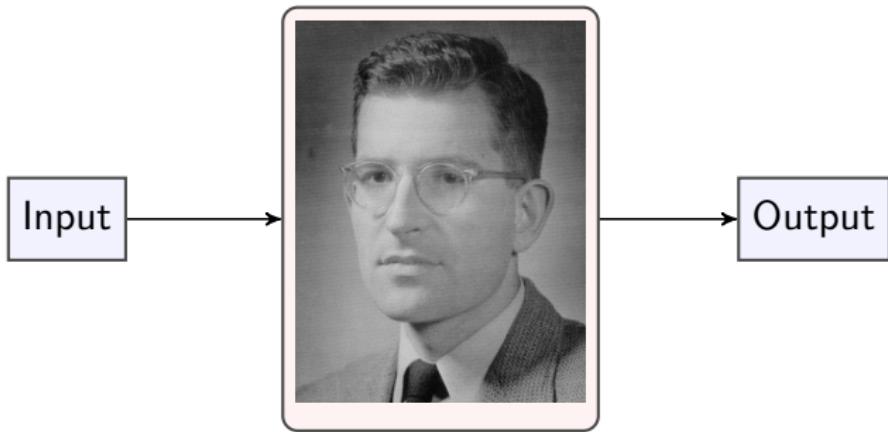
- Designers: Warren McCulloch and Walter Pitts
- **Finite automata** – simple model of computation

# 1945: Mauchly and Eckert's ENIAC



- Designers: John Mauchly, J. Presper Eckert
- World's first electronic general-purpose computer  
(Turing-complete)

# 1957: Chomsky's grammars



- Designer: Noam Chomsky
- Context-free grammar and context-sensitive grammar – models of computation

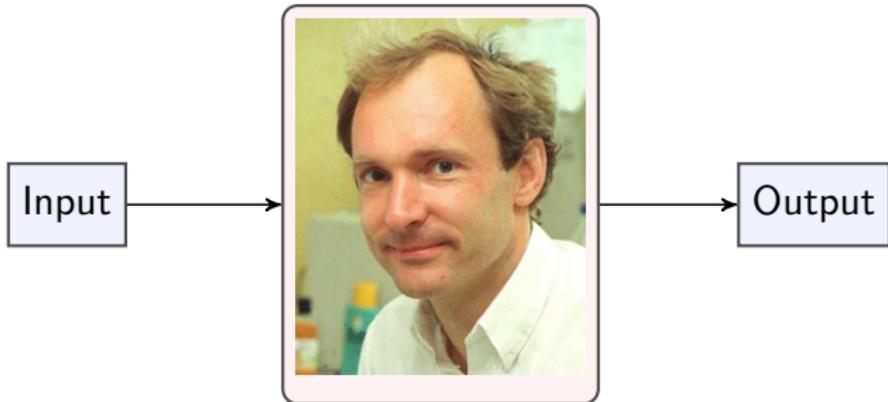
# 1985: Deutsch's quantum machine



Source: twitter

- Discoverer: David Deutsch
- Quantum model of computation
- Model based on quantum physics and not classical physics
- Exponentially faster than classical computing for some problems

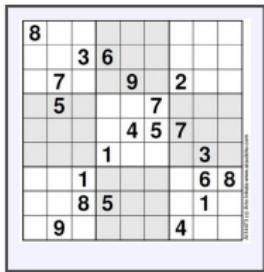
# 1989: Lee's world wide web



Source: CERN

- Designer: Tim Berners Lee
- World wide web – led to Internet revolution

# What is a computer/computation/algorithim?



# What is a computer/computation/algorithm?



# What is an alphabet?

## Definition

- An **alphabet**, denoted by  $\Sigma$ , is a finite, non-empty set of symbols.

## Examples

- $\Sigma = \{a, b\}$
- Unary alphabet  $\Sigma = \{1\}$
- Binary alphabet  $\Sigma = \{0, 1\}$
- English alphabet  $\Sigma = \{a, \dots, z, A, \dots, Z\}$
- Alphanumeric alphabet  $\Sigma = \{a-z, A-Z, 0-9\}$
- Morse code alphabet  $\Sigma = \{\text{dot}, \text{dash}, \text{pause}\}$
- DNA alphabet  $\Sigma = \{A, C, G, T\}$
- Java programming language alphabet  
 $\Sigma = \{a-z, A-Z, 0-9, (,), \{, \}, \dots, ;\}$
- $\{1, 2, 3, \dots\}$  is not an alphabet as the set is not finite

# Powers of an alphabet

## Definition

- $\Sigma$  = Some alphabet
- $\Sigma^k$  = Set of all strings of length  $k$  over  $\Sigma$
- $\Sigma^* = \Sigma^0 \cup \Sigma^1 \cup \Sigma^2 \cup \dots$  = Set of all strings over  $\Sigma$   
 $\Sigma^*$  is the universal set of all strings
- $\Sigma^+ = \Sigma^1 \cup \Sigma^2 \cup \Sigma^3 \cup \dots$  = Set of nonempty strings over  $\Sigma$

## Examples

- Let  $\Sigma = \{a, b\}$
- $\Sigma^0 = \{\epsilon\}$
- $\Sigma^1 = \{a, b\}$
- $\Sigma^2 = \{aa, ab, ba, bb\}$
- $\Sigma^* = \{\epsilon, a, b, aa, ab, ba, bb, \dots\}$

This ordering is called **canonical ordering**, which is different from lexicographic ordering

- $\Sigma^+ = \{a, b, aa, ab, ba, bb, \dots\}$

# What is a string?

## Definition

- A **string** or word is a finite sequence of symbols chosen from  $\Sigma$ .  
A string  $x \in \Sigma^*$ . An empty string is denoted by  $\epsilon$ .
- $|x|$  = length of string  $x$
- $n_\sigma(x)$  = #occurrences of symbol  $\sigma \in \Sigma$  in the string  $x$

## Examples

- $x = abaabb$  from  $\Sigma = \{a, b\}$
- $x = 111$  from  $\Sigma = \{0, 1\}$
- $x = \epsilon$  from  $\Sigma = \{a, \dots, z, A, \dots, Z\}$
- $x = Bond007$  from  $\Sigma = \{a - z, A - Z, 0 - 9\}$
- $x = CGGTCCGC$  from  $\Sigma = \{A, C, G, T\}$
- $x =$  a simple hello world C program from  
 $\Sigma = \{\text{if}, \text{main}, \text{return}, \text{for}, (\,), \{, \}, \dots, ;\}$

# What is a language?

## Definition

- A **language** over  $\Sigma$  is a subset of  $\Sigma^*$ .

## Examples

- The empty language  $\phi$ .
- $\{\epsilon, a, aab\}$  - a finite language.
- Language of palindromes over  $\{a, b\}$
- $\{x \in \{a, b\}^* \mid n_a(x) > n_b(x)\}$ .
- $\{x \in \{a, b\}^* \mid |x| \geq 2 \text{ and } x \text{ begins and ends with } b\}$

# What is a language?

## Examples (continued)

- Language of **your favorite quotations**
- Language of legal **Java identifiers**
- Language of legal **algebraic expressions** involving the identifier  $a$ , the binary operations  $+$  and  $*$ , and parentheses  
(strings:  $a$ ,  $a + a * a$ , and  $(a + a * (a + a))$ )
- Language of **balanced strings of parentheses**.  
(strings:  $\epsilon$ ,  $()()$ , and  $((((()))))$ )
- Language of **numeric “literals”** in Java (e.g:  $-41$ ,  $0.03$ ,  $5.0E - 3$ ).
- Language of legal **Java programs**.
- Language of **theorems** (true statements) in arithmetic
- Language of **theorems** (true statements) in geometry

# How can we represent information?

## Representation

- **Strings** can be used to represent all types of information
- Strings can **encode information** about names, numbers, dates, text documents, images, videos, and literally any type of data
- **Binary strings** are the simplest type of strings that can encode any information
- Binary strings can also be viewed as **numbers**
- Hence, **numbers** can also be used to represent all types of information

# Three major concepts in Theory of Computation

Concept	Meaning
Model of computation	An <b>abstract</b> but physically realistic machine that does computation
Language	Set of all <b>strings</b> that the computational model accepts
Grammar	Set of <b>rules</b> to derive any string from the language

# Core idea of Theory of Computation

Computation model	Language	Grammar
Finite automaton	Regular language	Regular grammar
Pushdown automaton	Context-free language	Context-free grammar
Linear-bounded automaton	Context-sensitive language	Context-sensitive grammar
Turing machine	Recursively enumerable language	Unrestricted grammar
No computer or no algorithm	Undecidable language	?

- We will spend an entire semester for this course trying to understand this table.

# Three major topics of Theory of Computation

Covered topic	Questions
Automata theory	What can be computed with extremely limited space?
Computability theory	What can be computed? Can a computer solve all computational problems, given enough (finite) time and space?
Complexity theory	How fast can we solve a problem? How small space can we use to solve a problem?
Not covered topic	Questions
Algorithms	How can a given computational problem be solved efficiently (less time and space)?

# What can be computed?

Problem	DFA	PDA	TM
Draw money from ATM	✓	✓	✓
Check if a string is present in another string	✓	✓	✓
Linux regular expressions	✓	✓	✓
Parse if-else blocks and for loops in C/C++/Java programs	✗	✓	✓
Parse nested arithmetic expressions	✗	✓	✓
Parse markup languages such as HTML	✗	✓	✓
Multiply two integers	✗	✗	✓
Factorize an integer into two integers	✗	✗	✓
Find a shortest path between two cities	✗	✗	✓
Check if a computer program halts or terminates	✗	✗	✗
Check if a computer program crashes	✗	✗	✗
Check if a computer program is correct	✗	✗	✗

- DFA: Deterministic Finite Automaton
- PDA: Pushdown Automaton
- TM: Turing Machine

# Applications of Theory of Computation

Topic	Applications
Finite automaton	<b>Regular expressions</b> Traffic signals, Vending machines, ATMs String matching Lexical analysis in a compiler Combination/sequential digital logic circuits Spell checkers
Pushdown automaton	<b>Stack applications</b> Balanced parentheses Syntax analysis in a compiler Evaluating arithmetic expressions
Linear-bounded automaton	Variable declaration and definition in a compiler Genetic programming
Turing machine	Understanding computation Mother of classical computers and algorithms Grandmother of quantum computers
Complexity theory	<b>Cryptography</b>

# Turing-complete systems

Time	Turing-complete system	Designer
1830s	Analytical engine	Charles Babbage
1930s	Recursive functions $\lambda$ -calculus Turing machine	Stephen Kleene Alonzo Church Alan Turing
—	Unrestricted grammar	—
1940s	Z3 Tag systems	Konrad Zuse Emil Leon Post
1960s	Markov's algorithms Unlimited register machines	Andrey Markov, Jr. John Shepherdson, Howard Sturgis
1970s	C Game of life	Dennis Ritchie John Conway
1980s	Rule 110 Quantum computers	Stephen Wolfram David Deutsch

# What can be computed?

## Problems

- [Halting program]

Write a computer program that takes a computer program  $P$  as input and outputs whether  $P$  halts (i.e., terminates) or not.

- [Correctness program]

Write a computer program that takes a computer program  $P$  and a specification  $s$  for  $P$  as input and outputs whether  $P$  is correct or not (i.e., if  $P$  follows the input-output specification  $s$  or not).

- [Equivalence program]

Write a computer program that takes two computer programs  $P_1$  and  $P_2$  as input and outputs whether  $P_1$  is functionally equivalent to  $P_2$  or not.

- [Self-replicating program]

Write a computer program that does not take any input and outputs its own source code.

# What can be computed?

## Problems

- [Halting program]

▷ Impossible

Write a computer program that takes a computer program  $P$  as input and outputs whether  $P$  halts (i.e., terminates) or not.

- [Correctness program]

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- [Self-replicating program]

▷ Possible

Write a computer program that does not take any input and outputs its own source code.