Algorithms

(Arrays and Lists)

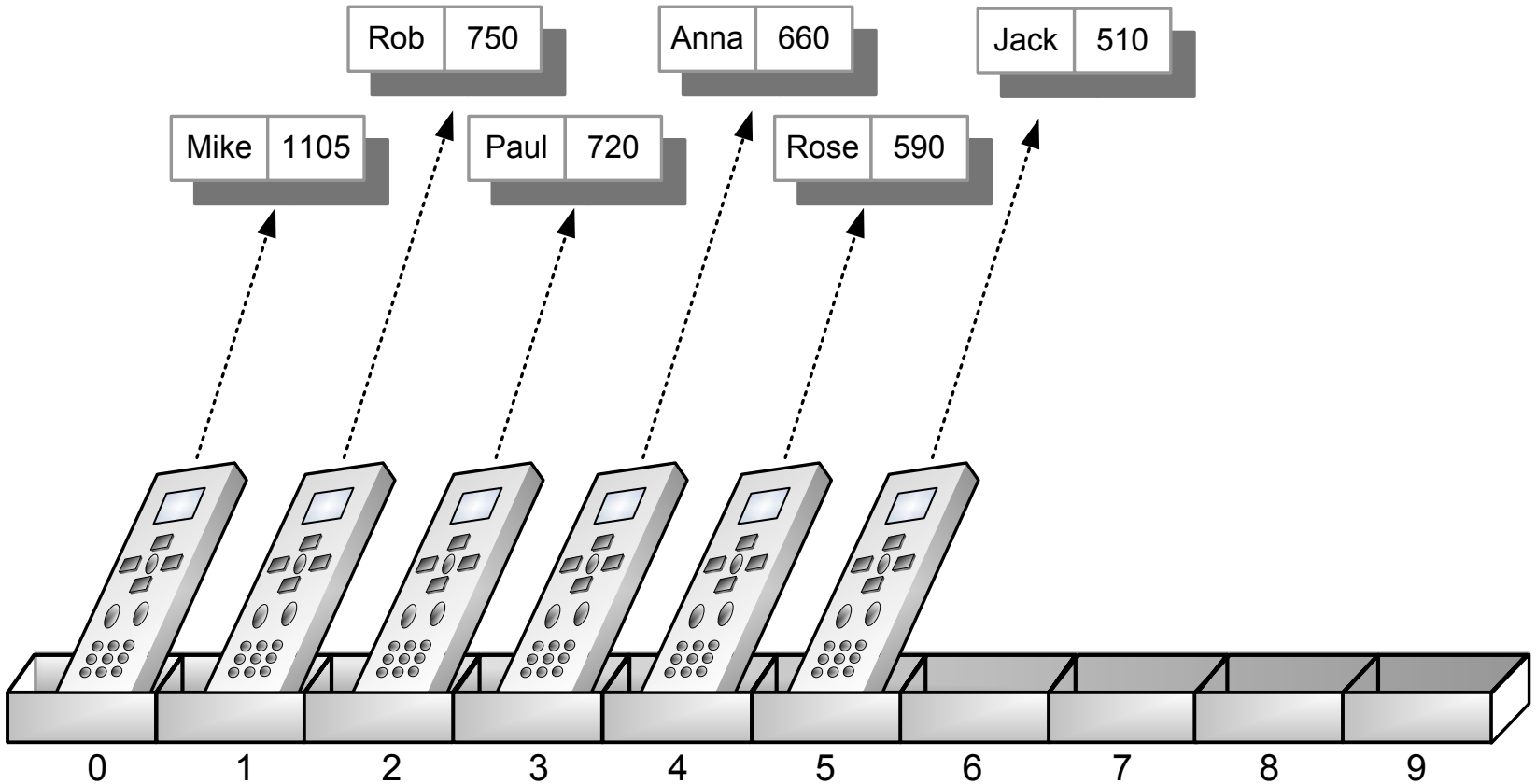
Pramod Ganapathi
Department of Computer Science
State University of New York at Stony Brook

December 29, 2020

Contents

Arrays

Scoreboard: Storing game entries in an array



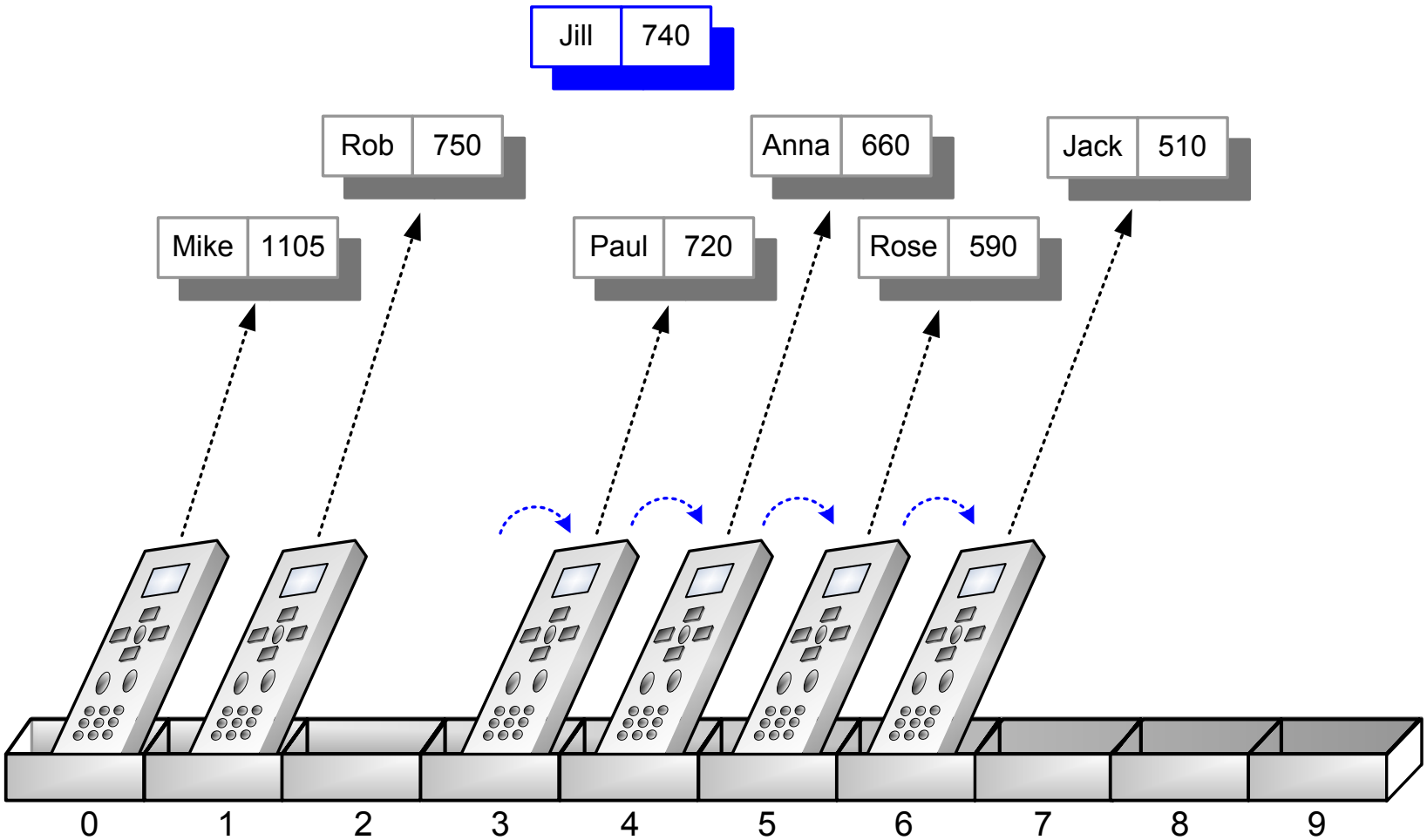
Scoreboard: High score entry

```
1.  public class GameEntry {
2.    private String name;    // name of the person earning this score
3.    private int score;      // the score value
4.
5.    /** Constructs a game entry with given parameters.. */
6.    public GameEntry(String n, int s) { name = n; score = s; }
7.    /** Returns the name field. */
8.    public String getName() { return name; }
9.    /** Returns the score field. */
10.   public int getScore() { return score; }
11.   /** Returns a string representation of this entry. */
12.   public String toString() { return "(" + name + ", " + score + ")"; }
13. }
```

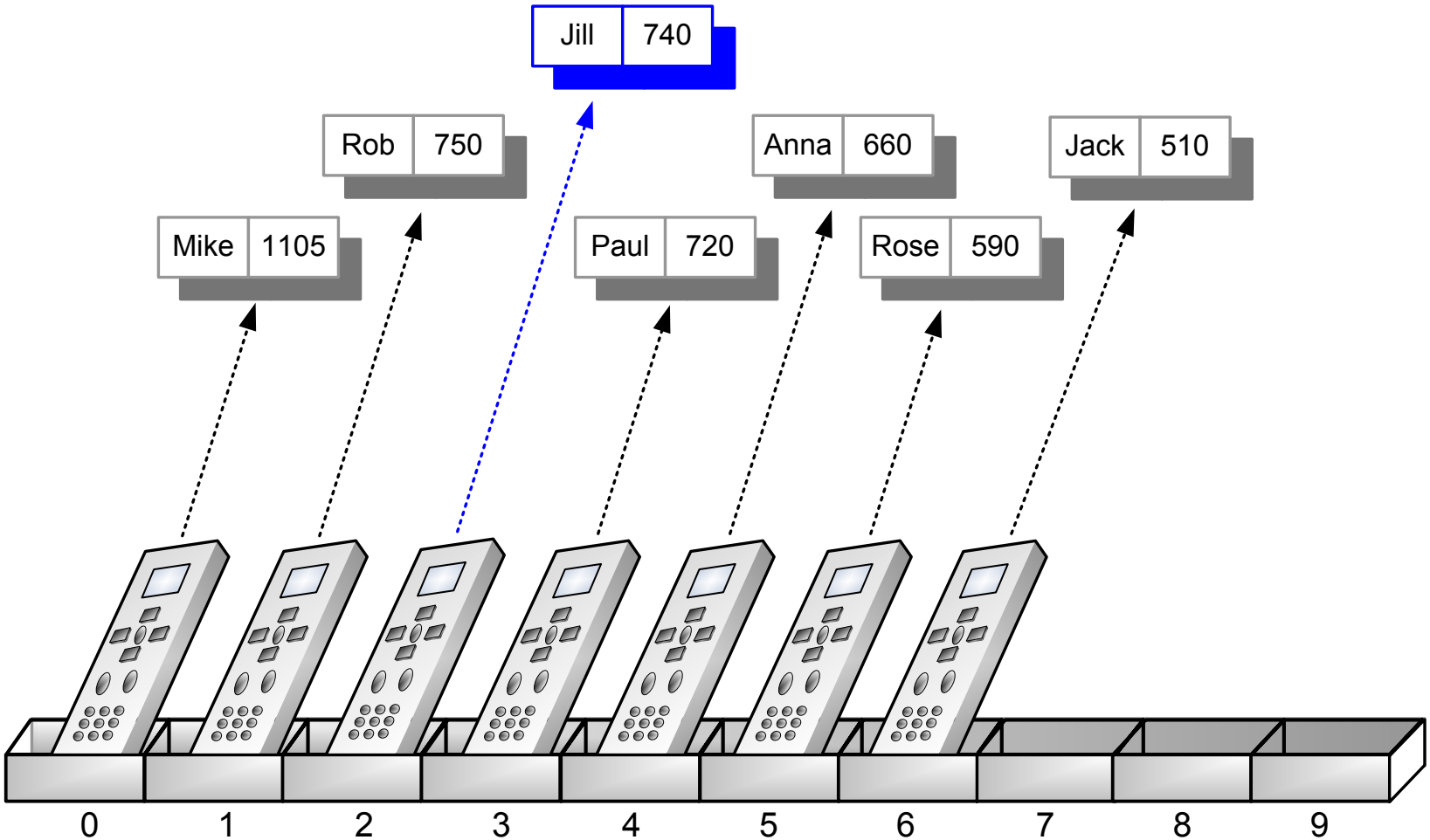
Scoreboard: Board of high scores

```
1.  /** Class for storing high scores in an array in nondecreasing order. */
2.  public class Scoreboard {
3.    private int numEntries = 0;                // number of actual entries
4.    private GameEntry[] board;                 // array of game entries
5.
6.    /** Constructs an empty scoreboard with the given capacity. */
7.    public Scoreboard(int capacity) { board = new GameEntry[capacity]; }
8.    /** Attempt to add a new high score to the collection. */
9.    public void add(GameEntry e) {...}
10.   /** Remove and return the high score at index i. */
11.   public GameEntry remove(int i) throws IndexOutOfBoundsException {...}
12.   /** Returns a string representation of the high scores list. */
13.   public String toString() {...}
14.
15.   public static void main(String[] args) {...}
16. }
```

Scoreboard: Add an entry



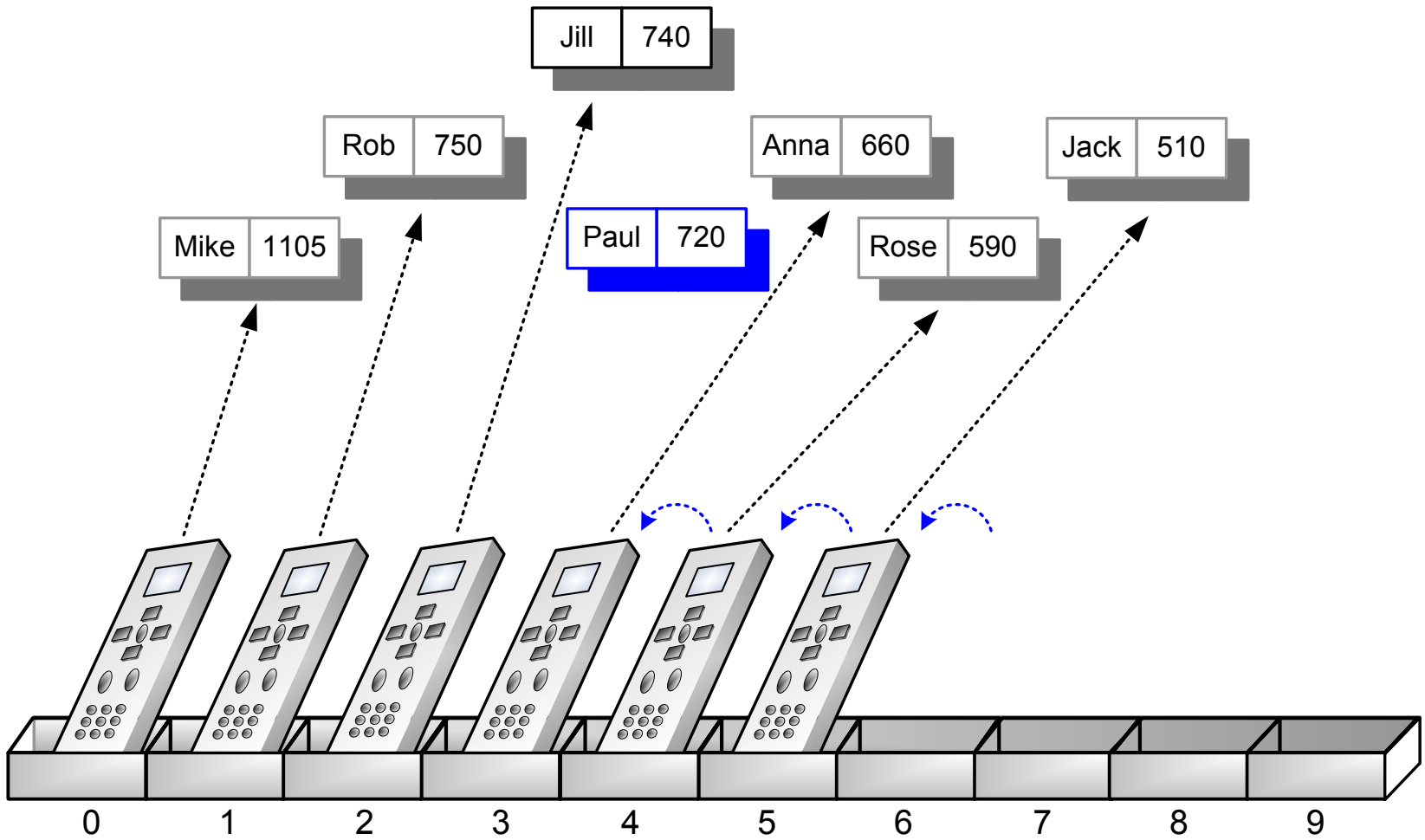
Scoreboard: Add an entry



Scoreboard: Add an entry

```
1.  /** Attempt to add a new score to the collection */
2.  public void add(GameEntry e) {
3.    int newScore = e.getScore();
4.
5.    // is the new entry e really a high score?
6.    if (numEntries < board.length || newScore > board[numEntries-1].getScore()) {
7.      if (numEntries < board.length)        // no score drops from the board
8.        numEntries++;                       // so overall number increases
9.      // shift any lower scores rightward to make room for the new entry
10.     int j = numEntries - 1;
11.     while (j > 0 && board[j-1].getScore() < newScore) {
12.       board[j] = board[j-1];              // shift entry from j-1 to j
13.       j--;                                // and decrement j
14.     }
15.     board[j] = e;                         // when done, add new entry
16.   }
17. }
```

Scoreboard: Remove an entry



Scoreboard: Remove an entry

```
1.  /** Remove and return the high score at index i. */
2.  public GameEntry remove(int i) throws IndexOutOfBoundsException {
3.    if (i < 0 || i >= numEntries)
4.      throw new IndexOutOfBoundsException("Invalid index: " + i);
5.    GameEntry temp = board[i];                  // save the object to be removed
6.    for (int j = i; j < numEntries - 1; j++)    // count up from i (not down)
7.      board[j] = board[j+1];                    // move one cell to the left
8.    board[numEntries - 1] = null;               // null out the old last score
9.    numEntries--;
10.   return temp;                                // return the removed object
11. }
```

Scoreboard: toString function

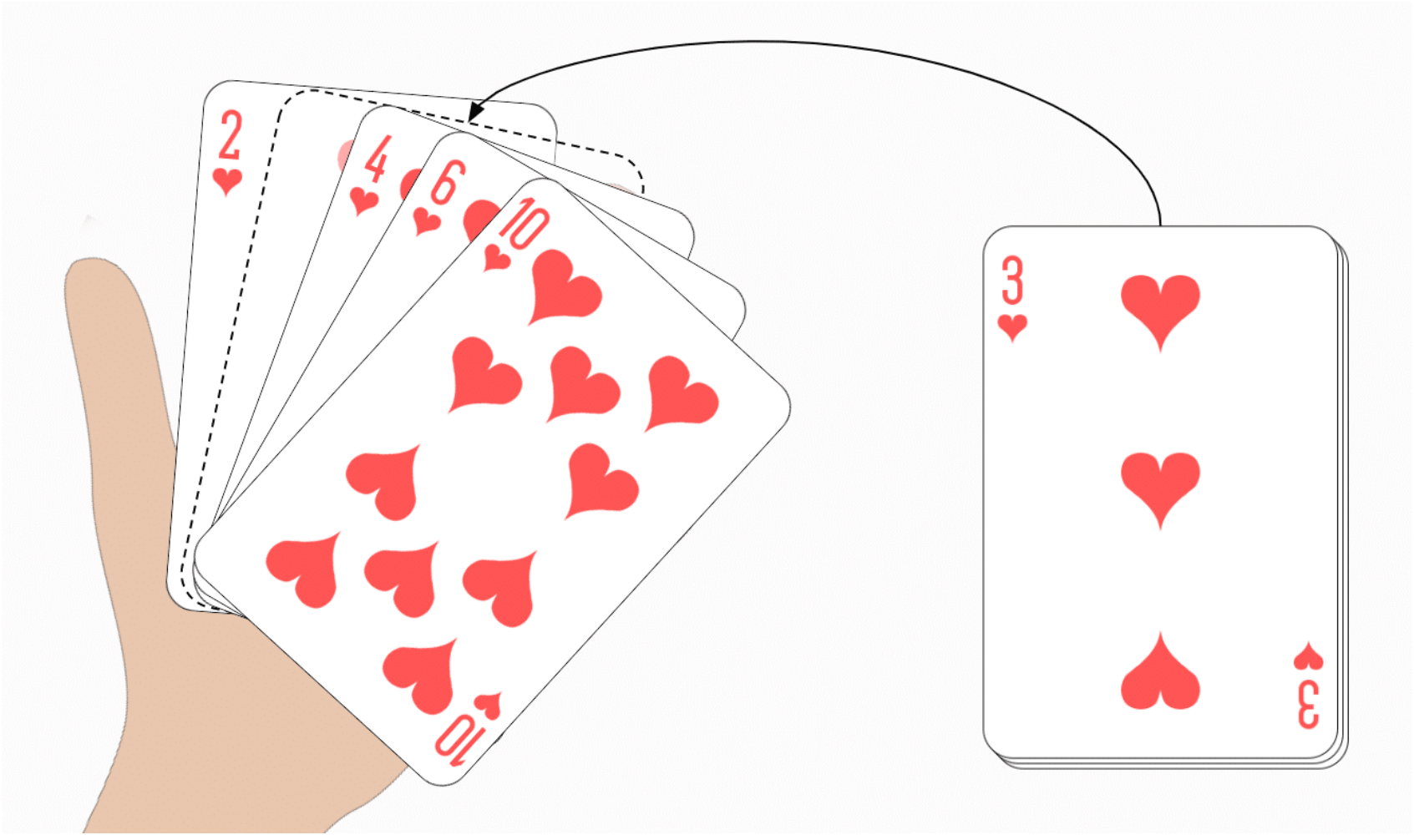
- Print the board consisting of high scores:
  `[board[0], board[1], ..., board[numEntries-1]]`

```
1.  /** Returns a string representation of the high scores list. */
2.  public String toString() {
3.    StringBuilder sb = new StringBuilder("[");
4.
5.    for (int j = 0; j < numEntries; j++) {
6.      if (j > 0) sb.append(", ");                    // separate entries by commas
7.      sb.append(board[j]);
8.    }
9.    sb.append("]");
10.   return sb.toString();
11. }
```

Scoreboard: main function

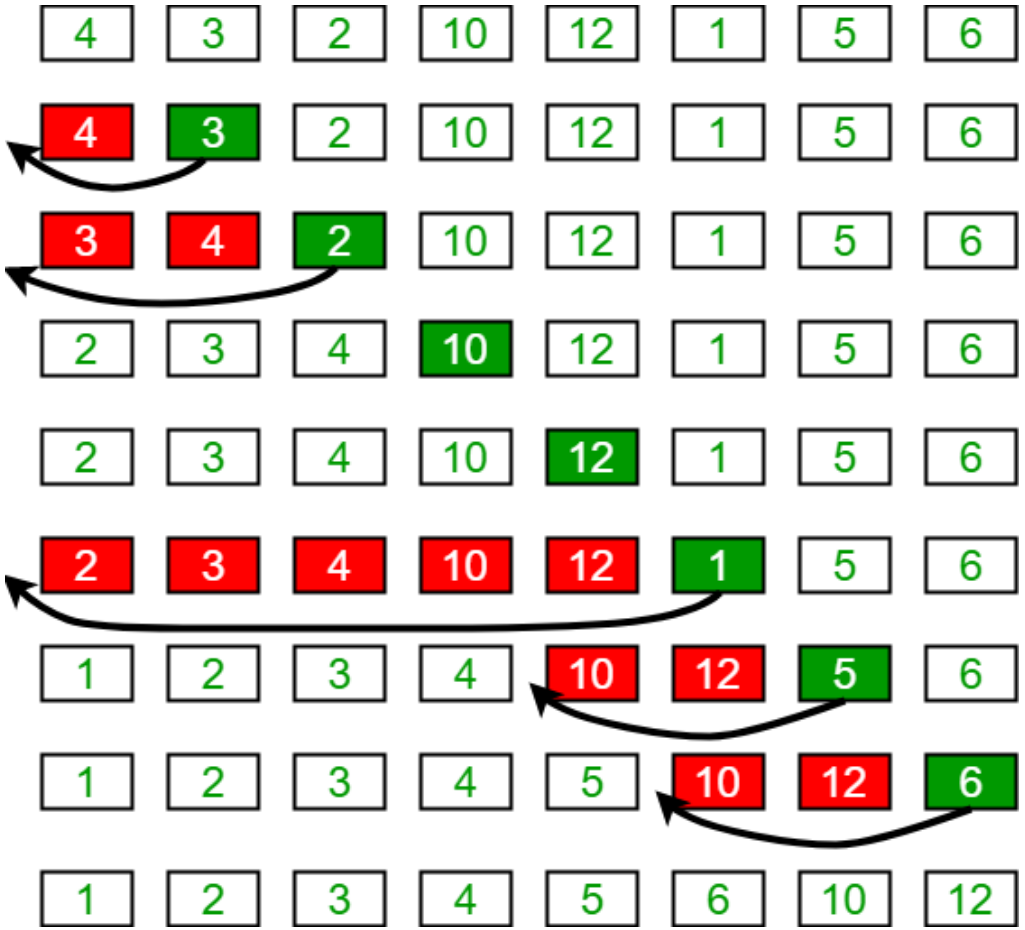
```
1.  public static void main(String[] args) {
2.    Scoreboard highscores = new Scoreboard(5);
3.    String[] names = {"Rob","Mike","Rose","Jill","Jack","Anna","Paul","Bob"};
4.    int[] scores = {750, 1105, 590, 740, 510, 660, 720, 400};
5.
6.    for (int i = 0; i < names.length; i++) {
7.      GameEntry ge = new GameEntry(names[i], scores[i]);
8.      System.out.println("Adding " + ge);
9.      highscores.add(ge);
10.     System.out.println(" Scoreboard: " + highscores);
11.   }
12.   System.out.println("Remove score at index " + 3); highscores.remove(3);
13.   System.out.println(highscores);
14.   System.out.println("Remove score at index " + 0); highscores.remove(0);
15.   System.out.println(highscores);
16.   System.out.println("Remove score at index " + 1); highscores.remove(1);
17.   System.out.println(highscores);
18.   System.out.println("Remove score at index " + 1); highscores.remove(1);
19.   System.out.println(highscores);
20.   System.out.println("Remove score at index " + 0); highscores.remove(0);
21.   System.out.println(highscores);
22. }
```

Sorting: Insertion sort



https://www.happycoders.eu/wp-content/uploads/2020/05/Insertion_Sort_Playing_Card_Example.png

Sorting: Insertion sort



Source: `https://media.geeksforgeeks.org/wp-content/uploads/insertionsort.png`

Sorting: Insertion sort

INSERTION-SORT($A[0..n-1]$)

Input: An array $A[0..n-1]$ of n orderable elements
Output: Array $A[0..n-1]$ sorted in nondecreasing order

1. **for** $i \leftarrow 1$ **to** $n-1$ **do**
2. $\quad v \leftarrow A[i]$
3. $\quad j \leftarrow i-1$
4. $\quad$ **while** $j \geq 0$ **and** $A[j] > v$ **do**
5. $\quad\quad A[j+1] \leftarrow A[j]$
6. $\quad\quad j \leftarrow j-1$
7. $\quad A[j+1] \leftarrow v$

Built-in methods for java.util.Arrays class

Method	Functionality
`equals(A, B)`	Compares arrays A and B.
`fill(A, x)`	Stores x in every cell of array A.
`copyOf(A, n)`	Returns n-sized array where the first $k = \min\{n, A.\text{length}\}$ elements are copied from A. If $n > A.\text{length}$, then the remaining elements are padded with 0 or null.
`toString(A)`	Returns string representation of array A.
`sort(A)`	Sorts array A based on natural ordering.
`binarySearch(A)`	Searches the sorted array A for value x.

Pseudorandom numbers

Linear congruential generator

$$X_i = \begin{cases} \text{seed} & \text{if } i = 0 \\ (a \times X_{i-1} + b) \ \% \ n & \text{if } i \geq 1 \end{cases}$$

Example

- Suppose seed $= 467$, $a = 17$, $b = 1$, $n = 1$ million. Then

  $X_0 = 467$

  $X_1 = (17 \times 467 + 1) \ \% \ 10^6 = 7940$

  $X_2 = (17 \times 7940 + 1) \ \% \ 10^6 = 134981$

  $X_3 = (17 \times 134981 + 1) \ \% \ 10^6 = 294678$

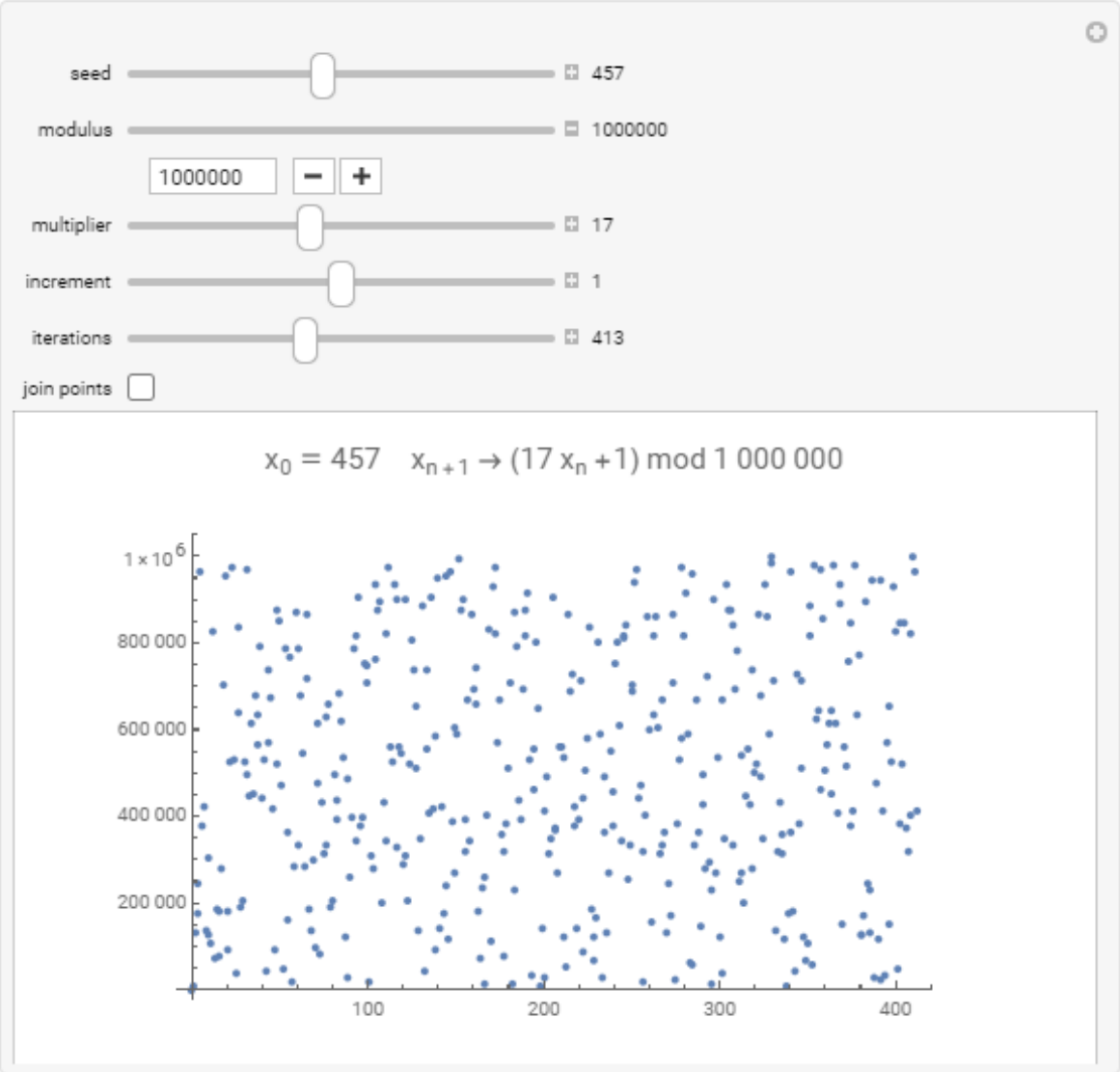
  $X_4 = (17 \times 294678 + 1) \ \% \ 10^6 = 9527$

  $X_5 = (17 \times 9527 + 1) \ \% \ 10^6 = 161960$

  $X_6 = (17 \times 161960 + 1) \ \% \ 10^6 = 753321$

  $X_7 = (17 \times 753321 + 1) \ \% \ 10^6 = 806458$

Pseudorandom numbers



https://demonstrations.wolfram.com/LinearCongruentialGenerators/
https://asecuritysite.com/encryption/linear

Built-in methods for java.util.Random class

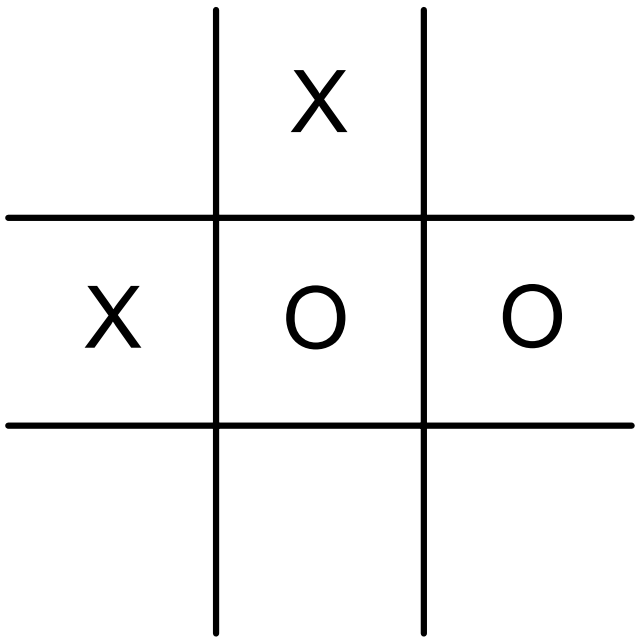
Method	Functionality
`nextBoolean()`	Returns the next pseudorandom boolean value.
`nextDouble()`	Returns the next pseudorandom double value in the range $[0.0, 1.0]$
`nextInt()`	Returns the next pseudorandom int.
`nextInt(n)`	Returns the next pseudorandom int in the range $[0, n)$.
`setSeed(s)`	Sets the seed of the generator to the long s.

2-D Arrays

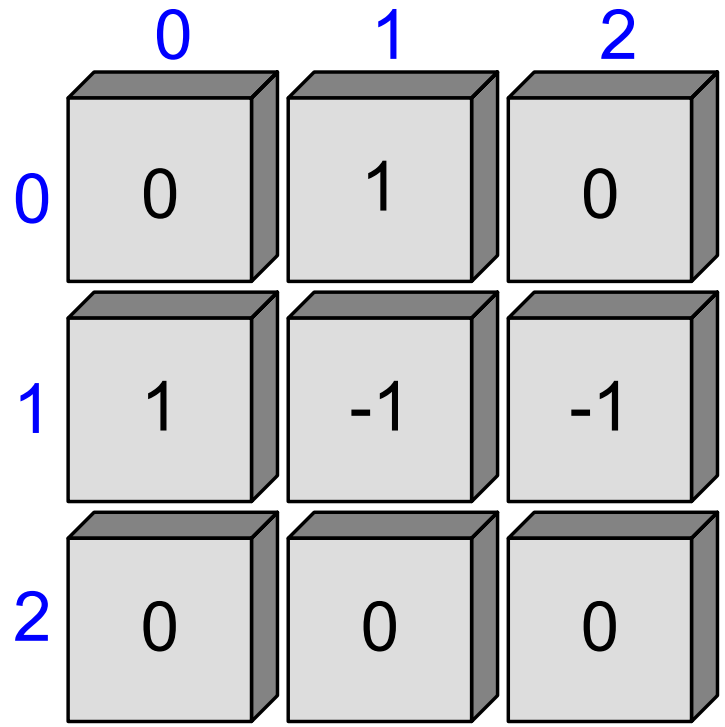
- Definition.
  A 2-D array in Java is created as array of arrays
- Declaration.
  ```
  int[][] data = new int[8][10];
  ```
- Valid uses.
  ```
  data[i][i+1] = data[i][i] + 3;
  j = data.length;          // j is 8
  k = data[4].length;       // k is 10
  ```

	0	1	2	3	4	5	6	7	8	9
0	22	18	709	5	33	10	4	56	82	440
1	45	32	830	120	750	660	13	77	20	105
2	4	880	45	66	61	28	650	7	510	67
3	940	12	36	3	20	100	306	590	0	500
4	50	65	42	49	88	25	70	126	83	288
5	398	233	5	83	59	232	49	8	365	90
6	33	58	632	87	94	5	59	204	120	829
7	62	394	3	4	102	140	183	390	16	26

2-D Arrays: Tic-Tac-Toe



playing board



board array

2-D Arrays: Tic-Tac-Toe

```
1.  /** Simulation of a Tic-Tac-Toe game (does not do strategy). */
2.  public class TicTacToe {
3.    public static final int X = 1, O = -1;      // players
4.    public static final int EMPTY = 0;          // empty cell
5.    private int board[][] = new int[3][3];      // game board
6.    private int player;                         // current player
7.
8.    /** Constructor */
9.    public TicTacToe() { clearBoard(); }
10.   /** Clears the board */
11.   public void clearBoard() {...}
12.   /** Puts an X or O mark at position i,j. */
13.   public void putMark(int i, int j) throws IllegalArgumentException {...}
14.   /** Checks whether the board configuration is a win for the given player. */
15.   public boolean isWin(int mark) {...}
16.   /** Returns the winning player's code, or 0 to indicate a tie.*/
17.   public int winner() {...}
18.   /** Returns a simple character string showing the current board. */
19.   public String toString() {...}
20.   /** Test run of a simple game */
21.   public static void main(String[] args) {...}
22. }
```

2-D Arrays: Tic-Tac-Toe

```
1.  /** Clears the board */
2.  public void clearBoard() {
3.    for (int i = 0; i < 3; i++)
4.      for (int j = 0; j < 3; j++)
5.        board[i][j] = EMPTY;                    // every cell should be empty
6.    player = X;                                 // the first player is 'X'
7.  }
```

```
1.  /** Puts an X or O mark at position i,j. */
2.  public void putMark(int i, int j) throws IllegalArgumentException {
3.    if ((i < 0) || (i > 2) || (j < 0) || (j > 2))
4.      throw new IllegalArgumentException("Invalid board position");
5.    if (board[i][j] != EMPTY)
6.      throw new IllegalArgumentException("Board position occupied");
7.    board[i][j] = player;         // place the mark for the current player
8.    player = - player;            // switch players (uses fact that O = - X)
9.  }
```

2-D Arrays: Tic-Tac-Toe

```
1.  /** Checks whether the board configuration is a win for the given player. */
2.  public boolean isWin(int mark) {
3.    return ((board[0][0] + board[0][1] + board[0][2] == mark*3)   // row 0
4.         || (board[1][0] + board[1][1] + board[1][2] == mark*3)   // row 1
5.         || (board[2][0] + board[2][1] + board[2][2] == mark*3)   // row 2
6.         || (board[0][0] + board[1][0] + board[2][0] == mark*3)   // column 0
7.         || (board[0][1] + board[1][1] + board[2][1] == mark*3)   // column 1
8.         || (board[0][2] + board[1][2] + board[2][2] == mark*3)   // column 2
9.         || (board[0][0] + board[1][1] + board[2][2] == mark*3)   // diagonal
10.        || (board[2][0] + board[1][1] + board[0][2] == mark*3)); // rev diag
11. }
```

```
1.  /** Returns the winning player's code, or 0 to indicate a tie.*/
2.  public int winner() {
3.    if (isWin(X))          return(X);
4.    else if (isWin(O))     return(O);
5.    else                   return(0);
6.  }
```

2-D Arrays: Tic-Tac-Toe

```
1.  /** Returns a simple character string showing the current board. */
2.  public String toString() {
3.    StringBuilder sb = new StringBuilder();
4.    for (int i = 0; i < 3; i++) {
5.      for (int j = 0; j < 3; j++) {
6.        switch (board[i][j]) {
7.          case X:      sb.append("X"); break;
8.          case O:      sb.append("O"); break;
9.          case EMPTY:  sb.append(" "); break;
10.       }
11.       if (j < 2) sb.append("|");              // column boundary
12.     }
13.     if (i < 2) sb.append("\n-----\n");        // row boundary
14.   }
15.   return sb.toString();
16. }
```

2-D Arrays: Tic-Tac-Toe

```
1.  /** Test run of a simple game */
2.  public static void main(String[] args) {
3.    TicTacToe game = new TicTacToe();
4.    /* X moves: */           /* O moves: */
5.    game.putMark(1,1);       game.putMark(0,2);
6.    game.putMark(2,2);       game.putMark(0,0);
7.    game.putMark(0,1);       game.putMark(2,1);
8.    game.putMark(1,2);       game.putMark(1,0);
9.    game.putMark(2,0);
10.   System.out.println(game);
11.   int winningPlayer = game.winner();
12.   String[] outcome = {"O wins", "Tie", "X wins"};  // rely on ordering
13.   System.out.println(outcome[1 + winningPlayer]);
14. }
```

```
O|X|O
-----
O|X|X
-----
X|O|X
Tie
```

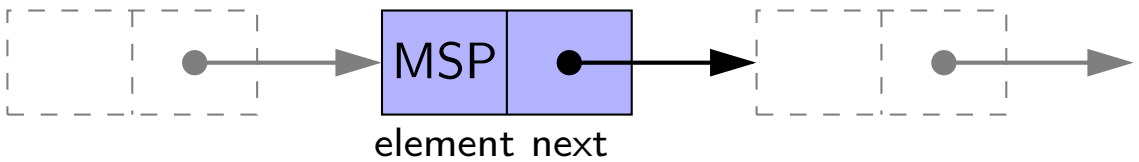
Advantages and disadvantages

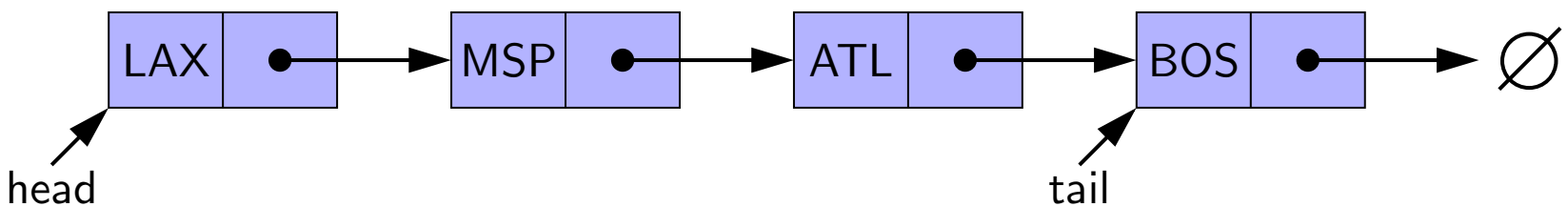
Operation	Time complexity
Fast operations	
Access/modify	$\mathcal{O}(1)$
Insert last	$\mathcal{O}(1)$
Slow operations	
Insert	$\mathcal{O}(n)$
Delete	$\mathcal{O}(n)$
Increase size	–

Singly Linked Lists

Singly linked lists

- A singly linked list, an alternative of array, is a linear sequence of nodes.
- E.g.: A singly linked list of airport codes.





Node class

```
1.  //---------------- nested Node class ----------------
2.  /** Node of a singly linked list, which stores a reference to its
3.      element and to the subsequent node in the list (or null if this
4.      is the last node). */
5.  private static class Node<E> {
6.    private E element;          // reference to the element stored at this node
7.    private Node<E> next;       // reference to the subsequent node in the list
8.
9.    /** Creates a node with the given element and next node. */
10.   public Node(E e, Node<E> n) { element = e; next = n; }
11.   /** Returns the element. */
12.   public E getElement() { return element; }
13.   /** Returns the node that follows this one (or null if no such node). */
14.   public Node<E> getNext() { return next; }
15.   /** Sets the node's next reference to point to Node n. */
16.   public void setNext(Node<E> n) { next = n; }
17. } //----------- end of nested Node class -----------
```

SinglyLinkedList class

Method	Functionality
`size()`	Returns the number of elements in the list.
`isEmpty()`	Returns true if the list is empty, and false otherwise.
`first()`	Returns the first element in the list.
`last()`	Returns the last element in the list.
`addFirst(e)`	Adds a new element to the front of the list.
`addLast(e)`	Adds a new element to the end of the list.
`removeFirst()`	Removes and returns the first element of the list.

SinglyLinkedList class

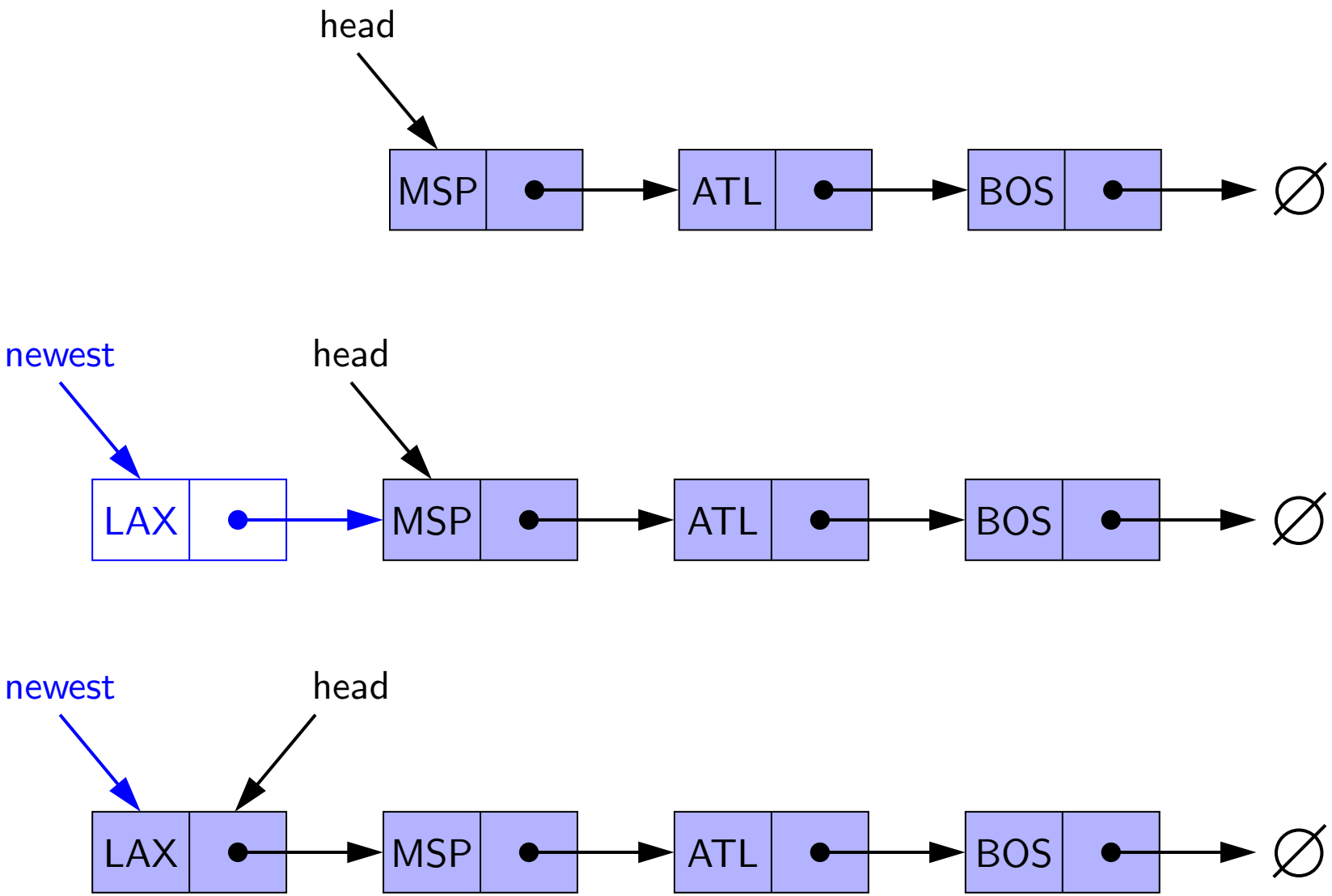
```
1.  public class SinglyLinkedList<E> {
2.    private static class Node<E> {...}
3.
4.    private Node<E> head = null;      // head node of the list
5.    private Node<E> tail = null;      // last node of the list
6.    private int size = 0;             // number of nodes in the list
7.
8.    public SinglyLinkedList() { }     // constructs an initially empty list
9.
10.   // access methods
11.   public int size() { return size; }
12.   public boolean isEmpty() { return size == 0; }
13.   public E first() {...}            // returns the first element
14.   public E last() {...}             // returns the last element
15.
16.   // update methods
17.   public void addFirst(E e) {...}  // adds element e to the front of the list
18.   public void addLast(E e) {...}   // adds element e to the end of the list
19.   public E removeFirst() {...}     // removes and returns the first element
20. }
```

Head and the tail

```
1. public E first() {     // returns (but does not remove) the first element
2.   if (isEmpty()) return null;
3.   return head.getElement();
4. }
```

```
1. public E last() {      // returns (but does not remove) the last element
2.   if (isEmpty()) return null;
3.   return tail.getElement();
4. }
```

Insert an element at the head



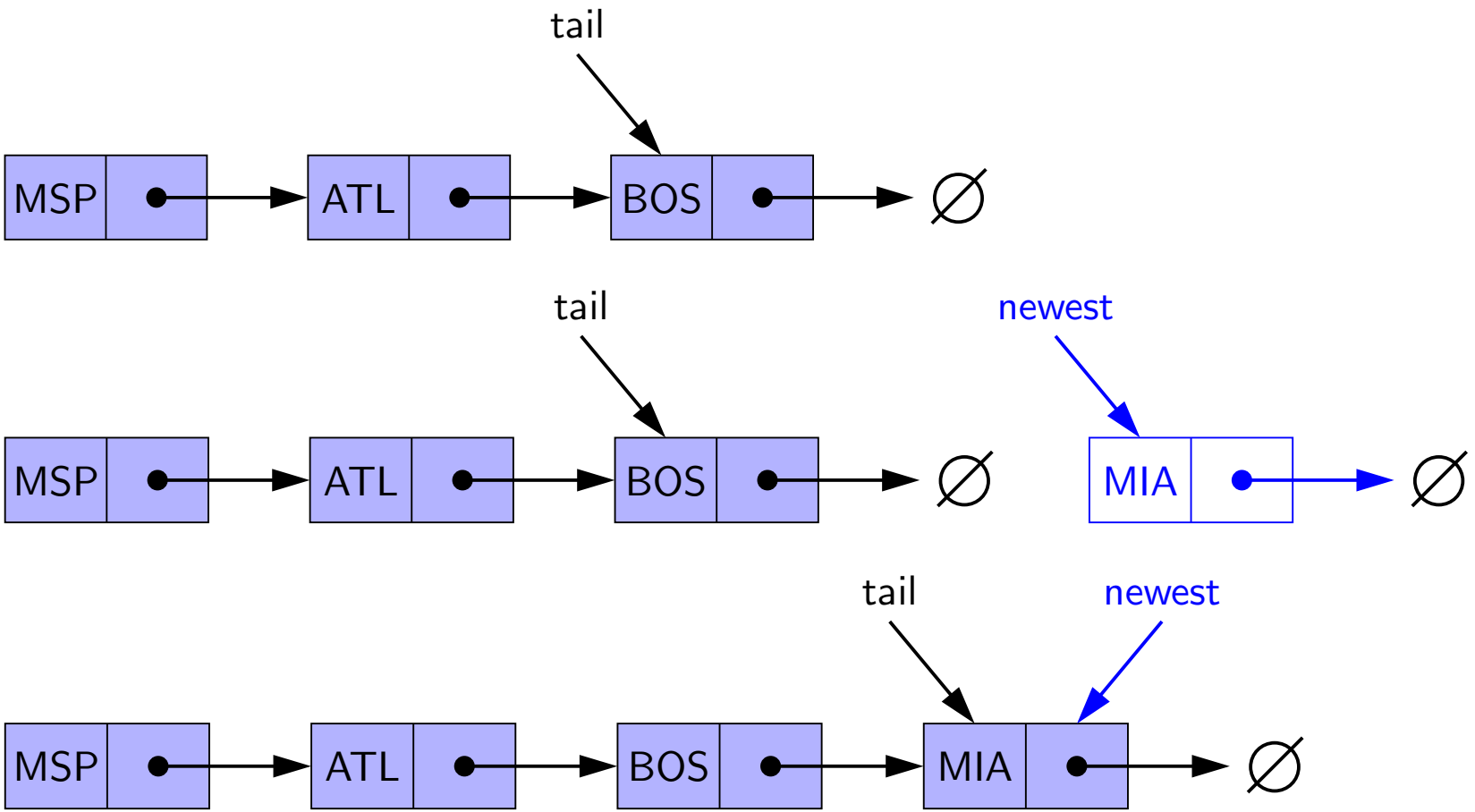
Insert an element at the head

ADD-FIRST(e)

1. $newest \leftarrow \text{NODE}(e)$
2. $newest.next \leftarrow head$
3. $head \leftarrow newest$
4. $size \leftarrow size + 1$

```
1.  public void addFirst(E e) {    // adds element e to the front of the list
2.    head = new Node<>(e, head);  // create and link a new node
3.    if (size == 0)
4.      tail = head;               // special case: new node becomes tail also
5.    size++;
6.  }
```

Insert an element at the tail



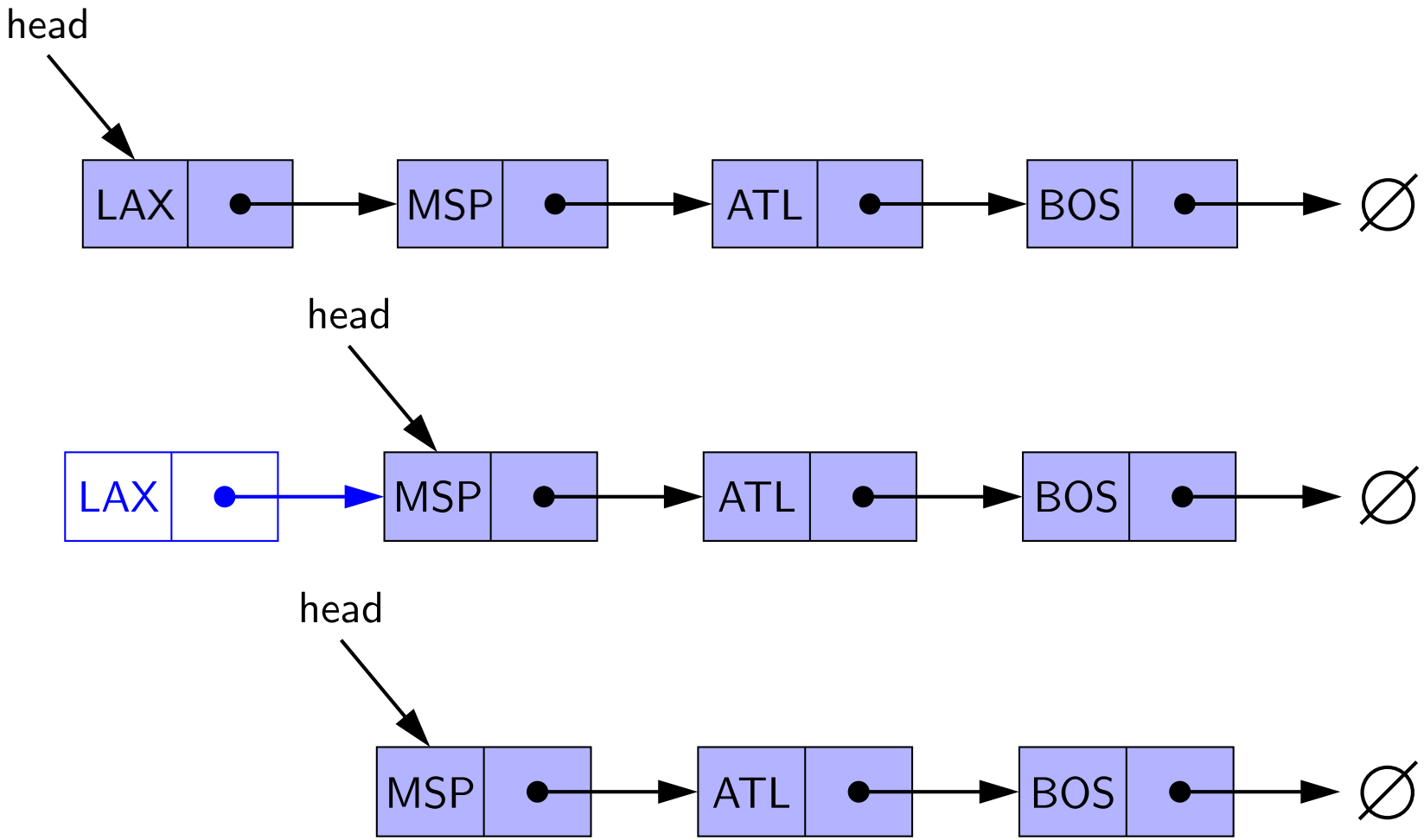
Insert an element at the tail

ADD-LAST(e)

1. $newest \leftarrow \text{NODE}(e)$
2. $newest.next \leftarrow null$
3. $tail.next \leftarrow newest$
4. $tail \leftarrow newest$
5. $size \leftarrow size + 1$

```
1. public void addLast(E e) {                  // adds element e to the end of the list
2.   Node<E> newest = new Node<>(e, null); // node will eventually be the tail
3.   if (isEmpty())
4.     head = newest;                          // special case: previously empty list
5.   else
6.     tail.setNext(newest);                   // new node after existing tail
7.   tail = newest;                            // new node becomes the tail
8.   size++;
9. }
```

Remove an element at the head



Remove an element at the head

REMOVE-FIRST()

1. **if** $head = null$ **then**
2. the list is empty
3. $head \leftarrow head.next$
4. $size \leftarrow size - 1$

```
1.  public E removeFirst() {          // removes and returns the first element
2.    if (isEmpty()) return null;     // nothing to remove
3.    E answer = head.getElement();
4.    head = head.getNext();          // will become null if list had only one node
5.    size--;
6.    if (size == 0)
7.      tail = null;                  // special case as list is now empty
8.    return answer;
9.  }
```

Advantages and disadvantages

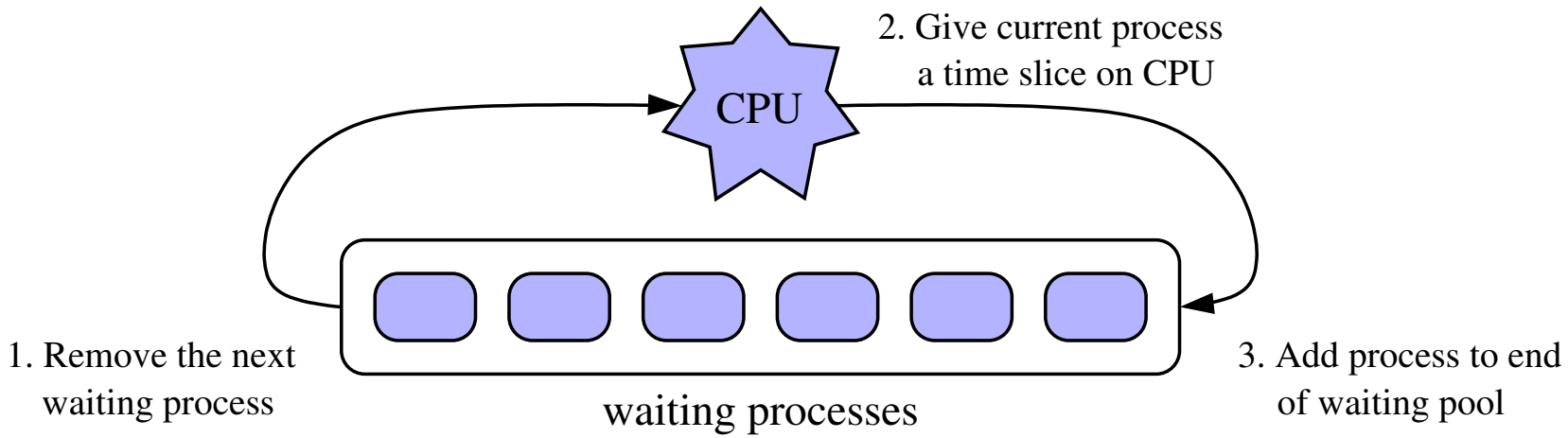
Operation	Time complexity
Fast operations	
Insert first	$\mathcal{O}(1)$
Insert last	$\mathcal{O}(1)$
Delete first	$\mathcal{O}(1)$
Increase size	$\mathcal{O}(1)$
Slow operations	
Delete last	$\mathcal{O}(n)$
Access/modify	$\mathcal{O}(n)$
Insert	$\mathcal{O}(n)$
Delete	$\mathcal{O}(n)$

Circularly Linked Lists

Applications requiring cyclic order

- Operating system
  Round-robin scheduling of processes/jobs
- Multiplayer games
  Scheduling of player turns
- Buses and subways
  Scheduling of stops in a continuous loop

Round-robin scheduling of processes



Round-robin scheduler can be implemented using a singly linked list L by repeatedly performing:

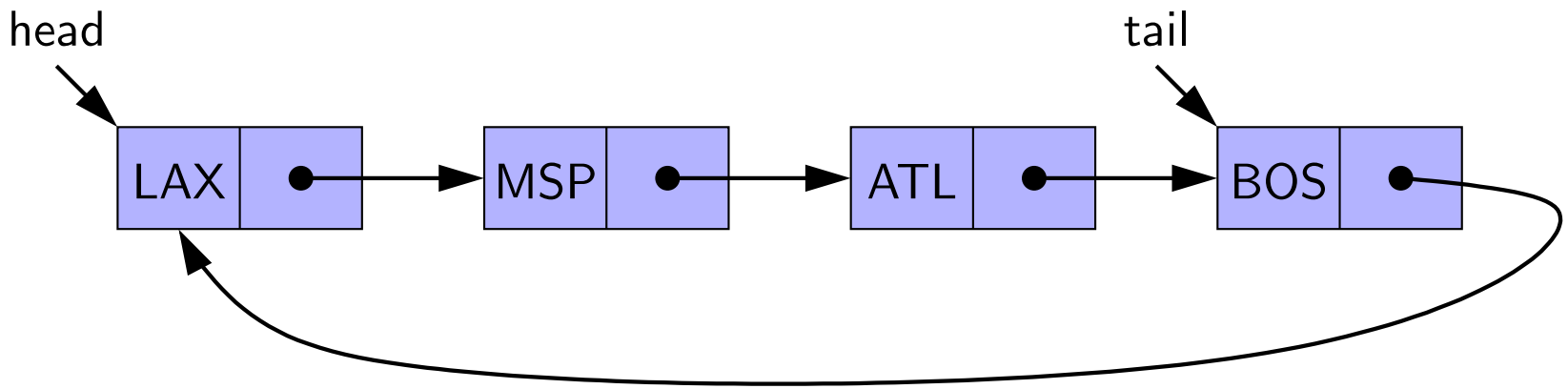
1. process $p = L$.removeFirst()
2. Give a time slice to process p
3. L.addLast(p)

Designing a CircularlyLinkedList class

CircularlyLinkedList =

SinglarlyLinkedList + (tail.next $\leftarrow$ head) + rotate() method

(rotate() moves the first element to the end of the list)



Round-robin scheduler can be implemented using a circularly linked list C by repeatedly performing:

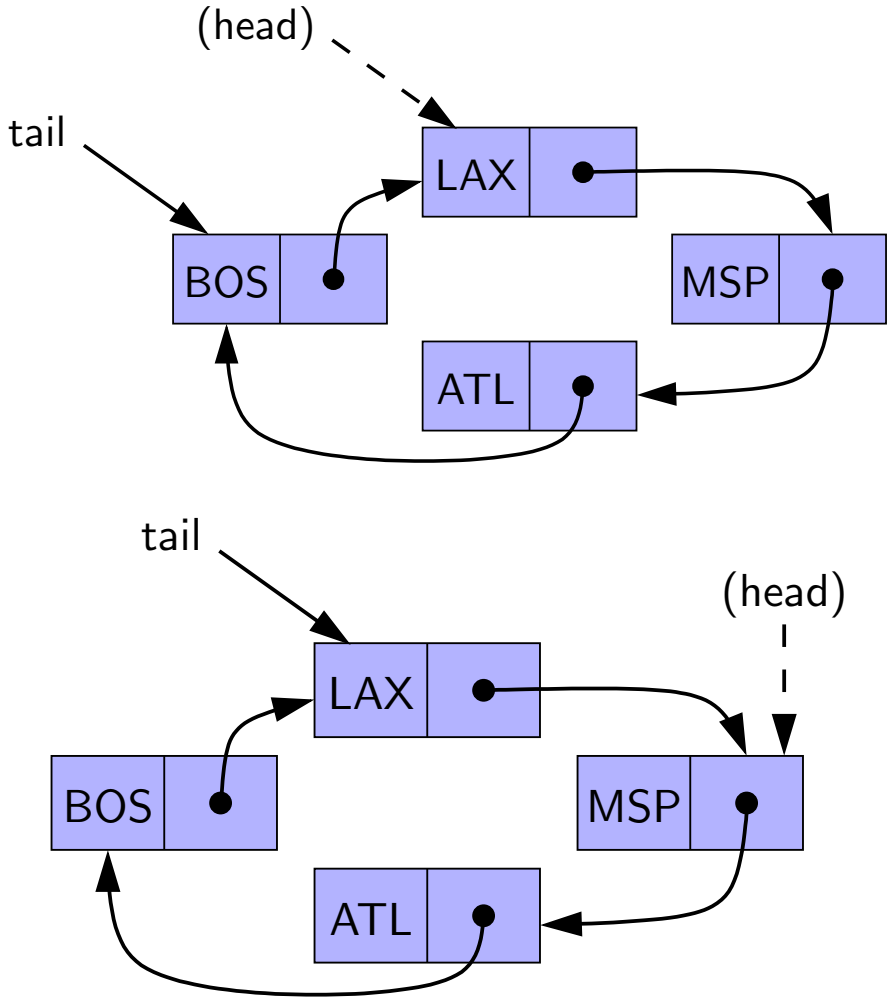
1. Give a time slice to process C.first()
2. C.rotate()

Additional optimization

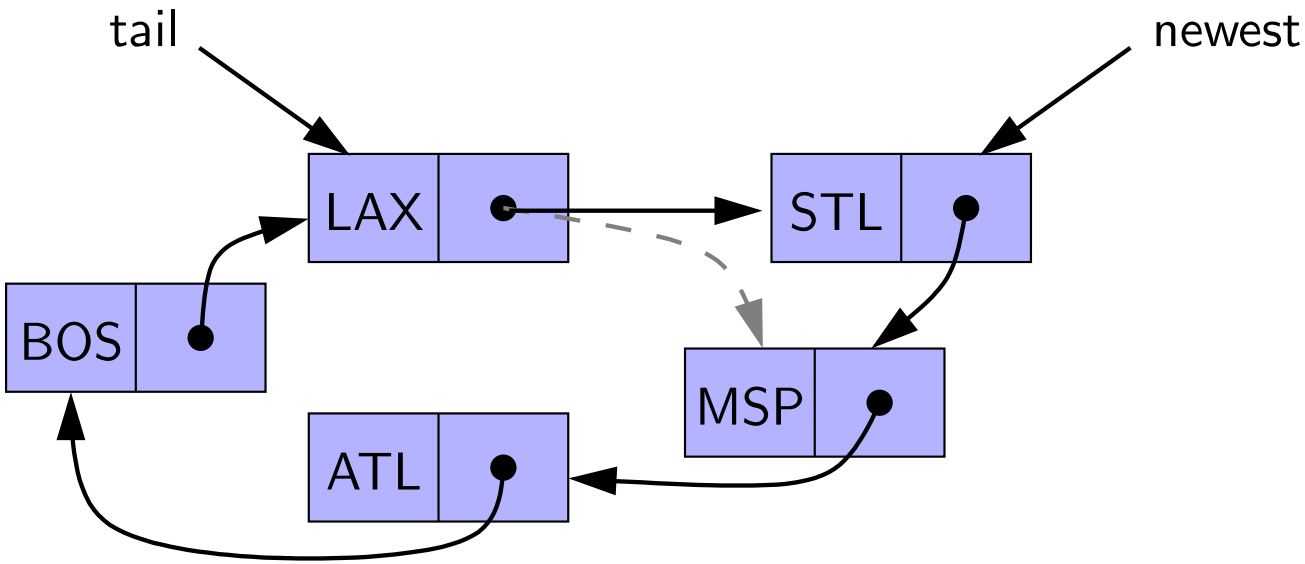
- Head reference is no longer required.
  Head can be accessed as tail.getNext()
- Maintaining only the tail reference is simpler, time-, and space-efficient.
- This implementation is superior to singly linked list implementation.

Rotate operation

- Simply advance the tail reference to its next node.



Insert element at head



CircularlyLinkedList class

```
1.  public class CircularlyLinkedList<E> {
2.    // nested node class identical to that of the SinglyLinkedList class
3.    private static class Node<E> {...}
4.
5.    private Node<E> tail = null;        // we store tail (but not head)
6.    private int size = 0;               // number of nodes in the list
7.    public CircularlyLinkedList() { }  // constructs an initially empty list
8.
9.    // access methods
10.   public int size() { return size; }
11.   public boolean isEmpty() { return size == 0; }
12.   public E first() {...}              // returns the first element
13.   public E last() {...}               // returns the last element
14.
15.   // update methods
16.   public void rotate() {...}          // rotate the first element to the last
17.   public void addFirst(E e) {...}     // adds element e to the front
18.   public void addLast(E e) {...}      // adds element e to the end
19.   public E removeFirst() {...}        // removes and returns the first element
20. }
```

Access methods

```
1. public E first() {                          // returns the first element
2.   if (isEmpty()) return null;
3.   return tail.getNext().getElement();       // the head is after the tail
4. }
```

```
1. public E last() {                           // returns the last element
2.   if (isEmpty()) return null;
3.   return tail.getElement();
4. }
```

Update methods

```
1.  public void rotate() {              // rotate the first element to the last
2.    if (tail != null)                 // if empty, do nothing
3.      tail = tail.getNext();          // the old head becomes the new tail
4.  }
```

```
1.  public void addFirst(E e) {         // adds element e to the front of the list
2.    if (size == 0) {
3.      tail = new Node<>(e, null);
4.      tail.setNext(tail);             // link to itself circularly
5.    } else {
6.      Node<E> newest = new Node<>(e, tail.getNext());
7.      tail.setNext(newest);
8.    }
9.    size++;
10. }
```

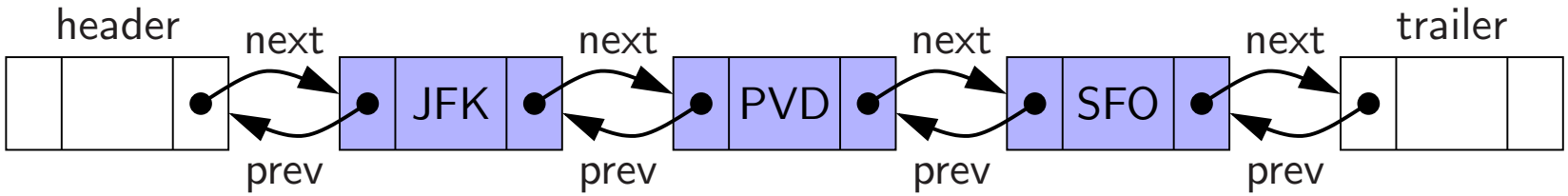
Update methods

```
1.  public void addLast(E e) {              // adds element e to the end of the list
2.    addFirst(e);                          // insert new element at front of list
3.    tail = tail.getNext();                // now new element becomes the tail
4.  }
```

```
1.  public E removeFirst() {                // removes and returns the first element
2.    if (isEmpty()) return null;           // nothing to remove
3.    Node<E> head = tail.getNext();
4.    if (head == tail) tail = null;        // must be the only node left
5.    else tail.setNext(head.getNext());    // removes "head" from the list
6.    size--;
7.    return head.getElement();
8.  }
```

Doubly Linked Lists

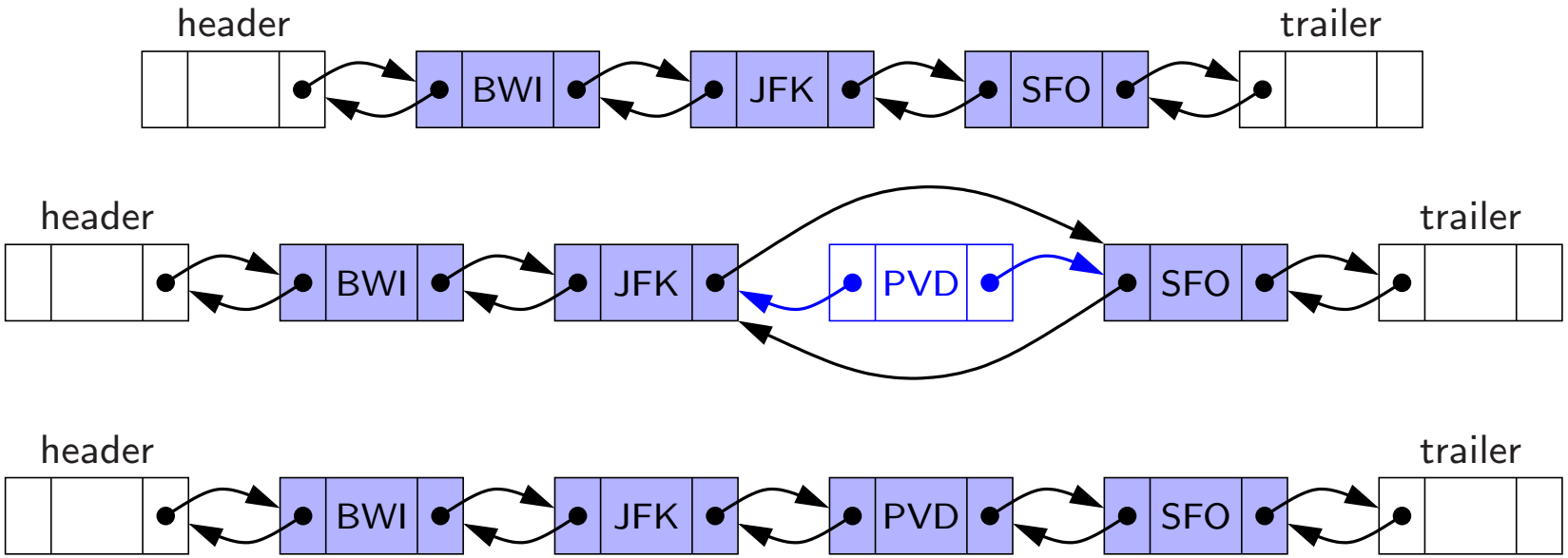
Doubly linked lists



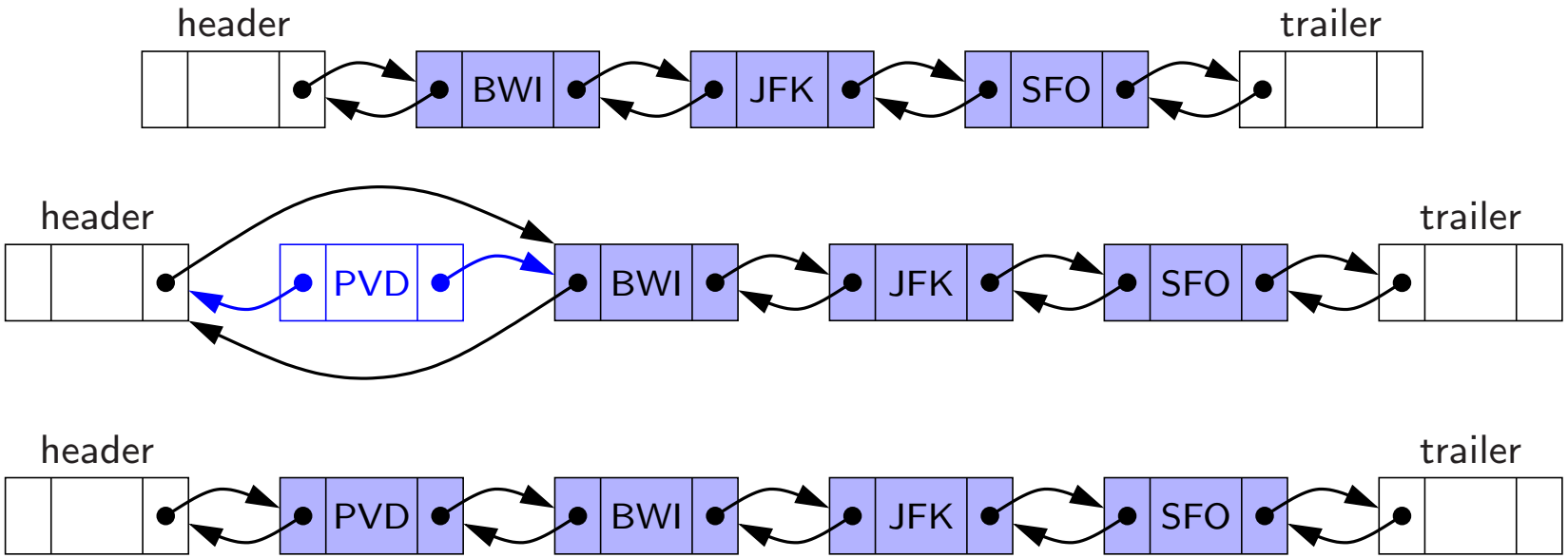
Advantages of using sentinels:

- Header and trailer nodes never change, only the nodes between them change
- Insertions and deletions can be handled in a unified manner

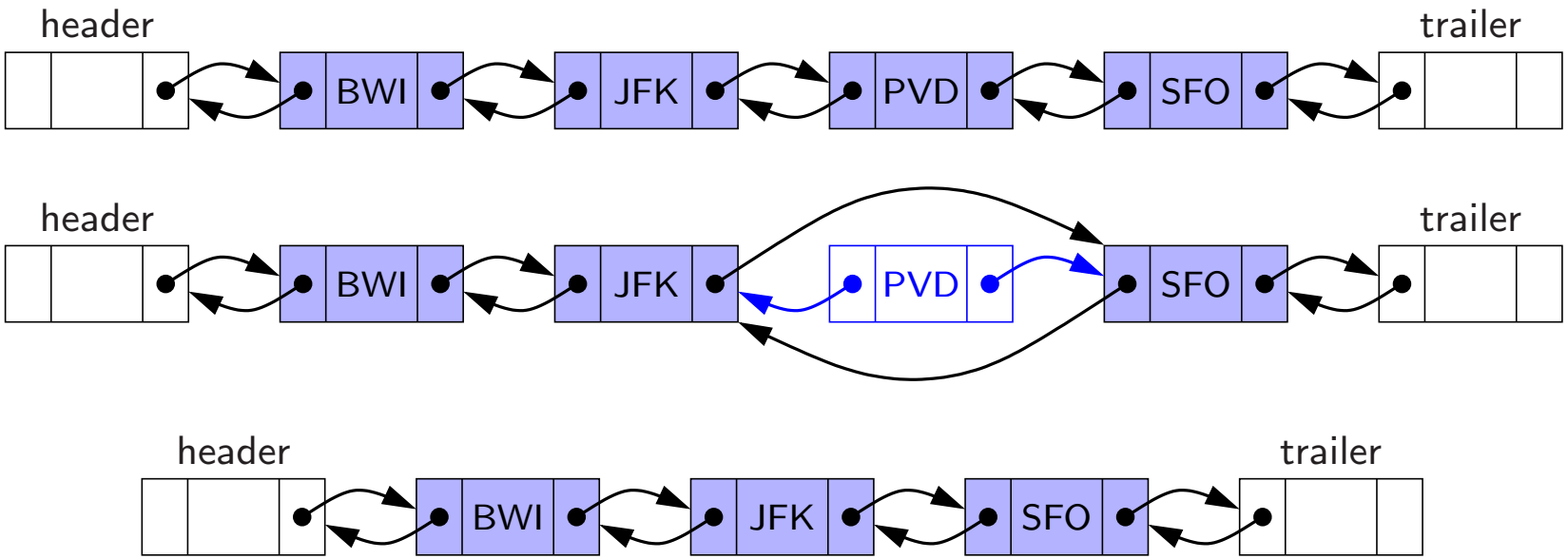
Inserting a node



Inserting at the front



Deleting a node



Methods for DoublyLinkedList class

Method	Functionality
`size()`	Returns the number of elements in the list.
`isEmpty()`	Returns true if the list is empty, and false otherwise.
`first()`	Returns the first element in the list.
`last()`	Returns the last element in the list.
`addFirst(e)`	Adds a new element to the front of the list.
`addLast(e)`	Adds a new element to the end of the list.
`removeFirst()`	Removes and returns the first element of the list.
`removeLast()`	Removes and returns the last element of the list.

DoublyLinkedList class

```
1.  public class DoublyLinkedList<E> {
2.    // nested Node class
3.    private static class Node<E> {...}
4.    private Node<E> header;                    // header sentinel
5.    private Node<E> trailer;                   // trailer sentinel
6.    private int size = 0;
7.
8.    // access methods
9.    public DoublyLinkedList() {...}
10.   public int size() {...}
11.   public boolean isEmpty() {...}
12.   public E first() {...}
13.   public E last() {...}
14.
15.   // update methods
16.   public void addFirst(E e) {...}
17.   public void addLast(E e) {...}
18.   public E removeFirst() {...}
19.   public E removeLast() {...}
20.   // private update methods
21.   private void addBetween(E e, Node<E> predecessor, Node<E> successor) {...}
22.   private E remove(Node<E> node) {...}
23. }
```

Node class

```
1.  //---------------- nested Node class ----------------
2.  private static class Node<E> {
3.    private E element;          // reference to the element stored at this node
4.    private Node<E> prev;       // reference to the previous node in the list
5.    private Node<E> next;       // reference to the subsequent node in the list
6.
7.    public Node(E e, Node<E> p, Node<E> n) {
8.      element = e; prev = p; next = n;
9.    }
10.   public E getElement() { return element; }
11.   public Node<E> getPrev() { return prev; }
12.   public Node<E> getNext() { return next; }
13.   public void setPrev(Node<E> p) { prev = p; }
14.   public void setNext(Node<E> n) { next = n; }
15. } //----------- end of nested Node class -----------
```

Access methods

```
1.  public DoublyLinkedList() {
2.    header = new Node<>(null, null, null);      // create header
3.    trailer = new Node<>(null, header, null);   // trailer is preceded by header
4.    header.setNext(trailer);                    // header is followed by trailer
5.  }
```

```
1.  // public access methods
2.  public int size() { return size; }
```

```
1.  public boolean isEmpty() { return size == 0; }
```

```
1.  public E first() {
2.    if (isEmpty()) return null;
3.    return header.getNext().getElement();       // first element is beyond header
4.  }
```

```
1.  public E last() {
2.    if (isEmpty()) return null;
3.    return trailer.getPrev().getElement();      // last element is before trailer
4.  }
```

Private update methods

```
1.  /* Adds an element to the linked list in between the given nodes. */
2.  private void addBetween(E e, Node<E> predecessor, Node<E> successor) {
3.    // create and link a new node
4.    Node<E> newest = new Node<>(e, predecessor, successor);
5.    predecessor.setNext(newest);
6.    successor.setPrev(newest);
7.    size++;
8.  }
```

```
1.  /* Removes the given node from the list and returns its element. */
2.  private E remove(Node<E> node) {
3.    Node<E> predecessor = node.getPrev();
4.    Node<E> successor = node.getNext();
5.    predecessor.setNext(successor);
6.    successor.setPrev(predecessor);
7.    size--;
8.    return node.getElement();
9.  }
```

Update methods

```
1.  /* Adds an element to the front of the list. */
2.  public void addFirst(E e) {
3.    addBetween(e, header, header.getNext());   // place just after the header
4.  }
```

```
1.  /* Adds an element to the end of the list. */
2.  public void addLast(E e) {
3.    addBetween(e, trailer.getPrev(), trailer); // place just before the trailer
4.  }
```

```
1.  /* Removes and returns the first element of the list. */
2.  public E removeFirst() {
3.    if (isEmpty()) return null;                // nothing to remove
4.    return remove(header.getNext());           // first element is beyond header
5.  }
```

```
1.  /* Removes and returns the last element of the list. */
2.  public E removeLast() {
3.    if (isEmpty()) return null;                // nothing to remove
4.    return remove(trailer.getPrev());          // last element is before trailer
5.  }
```

Comparison table of linear data structures

Operation	Dyn. array	SLL/CLL	DLL
Insert first	$\mathcal{O}(n)$	$\mathcal{O}(1)$	$\mathcal{O}(1)$
Insert last	$\mathcal{O}(1)$	$\mathcal{O}(1)$	$\mathcal{O}(1)$
Insert between	$\mathcal{O}(n)$	$\mathcal{O}(1)$	$\mathcal{O}(1)$
Insert at index	$\mathcal{O}(n)$	$\mathcal{O}(n)$	$\mathcal{O}(n)$
Delete first	$\mathcal{O}(n)$	$\mathcal{O}(1)$	$\mathcal{O}(1)$
Delete last	$\mathcal{O}(1)$	$\mathcal{O}(n)$	$\mathcal{O}(1)$
Delete at index	$\mathcal{O}(n)$	$\mathcal{O}(n)$	$\mathcal{O}(n)$
Access at index	$\mathcal{O}(1)$	$\mathcal{O}(n)$	$\mathcal{O}(n)$
Modify at index	$\mathcal{O}(1)$	$\mathcal{O}(n)$	$\mathcal{O}(n)$
Size	unlimited	unlimited	unlimited

List ADT

List ADT

Method	Functionality
`size()`	Returns the number of elements in the list.
`isEmpty()`	Returns a boolean indicating whether the list is empty.
`get(i)`	Returns the element of the list having index i; an error condition occurs if i is not in range [0,size() - 1].
`set(i, e)`	Replaces the element at index i with e, and returns the old element that was replaced; an error condition occurs if i is not in range [0,size() - 1].
`add(i, e)`	Inserts a new element e into the list so that it has index i, moving all subsequent elements one index later in the list; an error condition occurs if i is not in range [0,size()].
`remove(i)`	Removes and returns the element at index i, moving all subsequent elements one index earlier in the list; an error condition occurs if i is not in range [0,size() - 1]

Operations on a list

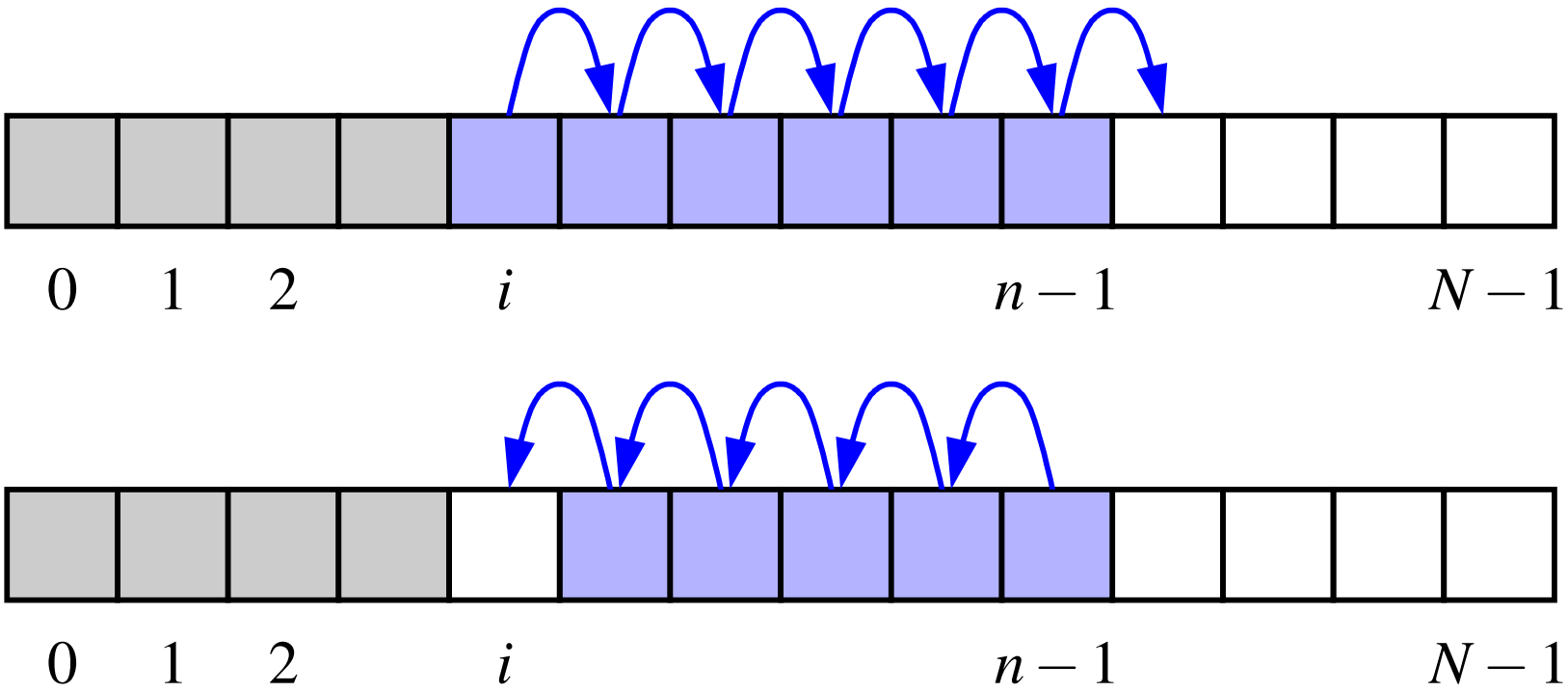
Method	Return value	List contents
add(0, A)	-	(A)
add(0, B)	-	(B, A)
get(1)	A	(B, A)
set(2, C)	error	(B, A)
add(2, C)	-	(B, A, C)
add(4, D)	error	(B, A, C)
remove(1)	A	(B, C)
add(1, D)	-	(B, D, C)
add(1, E)	-	(B, E, D, C)
get(4)	error	(B, E, D, C)
add(4, F)	-	(B, E, D, C, F)
set(2, G)	D	(B, E, G, C, F)
get(2)	G	(B, E, G, C, F)

Simplified java.util.List interface

```
1.  /* A simplified version of the java.util.List interface. */
2.  public interface List<E> {
3.    int size();
4.    boolean isEmpty();
5.
6.    /* Returns (but does not remove) the element at index i. */
7.    E get(int i) throws IndexOutOfBoundsException;
8.
9.    /* Replaces element at index i with e, and returns replaced element. */
10.   E set(int i, E e) throws IndexOutOfBoundsException;
11.
12.   /* Inserts e to be at index i, shifting subsequent elements later. */
13.   void add(int i, E e) throws IndexOutOfBoundsException;
14.
15.   /* Removes the element at index i, shifting subsequent elements earlier. */
16.   E remove(int i) throws IndexOutOfBoundsException;
17. }
```

Array Lists

- Implement the List ADT using an array.
- Get/set methods are fast, but add/remove methods are slow.



Simple ArrayList implementation

```
1.  public class ArrayList<E> implements List<E> {
2.    public static final int CAPACITY=16;     // default array capacity
3.    private E[] data;                        // generic array used for storage
4.    private int size = 0;                    // current number of elements
5.
6.    public ArrayList() { this(CAPACITY); }   // constructs list with default cap.
7.    public ArrayList(int capacity) { data = (E[]) new Object[capacity]; }
8.
9.    public int size() { return size; }
10.   public boolean isEmpty() { return size == 0; }
11.   public E get(int i) throws IndexOutOfBoundsException {...}
12.   public E set(int i, E e) throws IndexOutOfBoundsException {...}
13.   public void add(int i, E e) throws IndexOutOfBoundsException {...}
14.   public E remove(int i) throws IndexOutOfBoundsException {...}
15.
16.   // utility methods
17.   /** Checks whether the given index is in the range [0, n-1]. */
18.   protected void checkIndex(int i, int n) throws IndexOutOfBoundsException {...}
19.   /** Resizes internal array to have given capacity >= size. */
20.   protected void resize(int capacity) {...}
21. }
```

Access methods

```
1.  /* Returns (but does not remove) the element at index i. */
2.  public E get(int i) throws IndexOutOfBoundsException {
3.    checkIndex(i, size);
4.    return data[i];
5.  }
```

```
1.  /* Replaces the element at the specified index, and
2.   * returns the element previously stored. */
3.  public E set(int i, E e) throws IndexOutOfBoundsException {
4.    checkIndex(i, size);
5.    E temp = data[i];
6.    data[i] = e;
7.    return temp;
8.  }
```

Update methods

```
1.  /* Inserts the given element at the specified index of the list, shifting all
2.   * subsequent elements in the list one position further to make room. */
3.  public void add(int i, E e) throws IndexOutOfBoundsException {
4.    checkIndex(i, size + 1);
5.    if (size == data.length)                  // not enough capacity
6.      resize(2 * data.length);                // so double the current capacity
7.    for (int k=size-1; k >= i; k--)           // start by shifting rightmost
8.      data[k+1] = data[k];
9.    data[i] = e;                              // ready to place the new element
10.   size++;
11. }
```

```
1.  /* Removes and returns the element at the given index, shifting all subsequent
2.   * elements in the list one position closer to the front. */
3.  public E remove(int i) throws IndexOutOfBoundsException {
4.    checkIndex(i, size);
5.    E temp = data[i];
6.    for (int k=i; k < size-1; k++)            // shift elements to fill hole
7.      data[k] = data[k+1];
8.    data[size-1] = null;                      // help garbage collection
9.    size--;
10.   return temp;
11. }
```

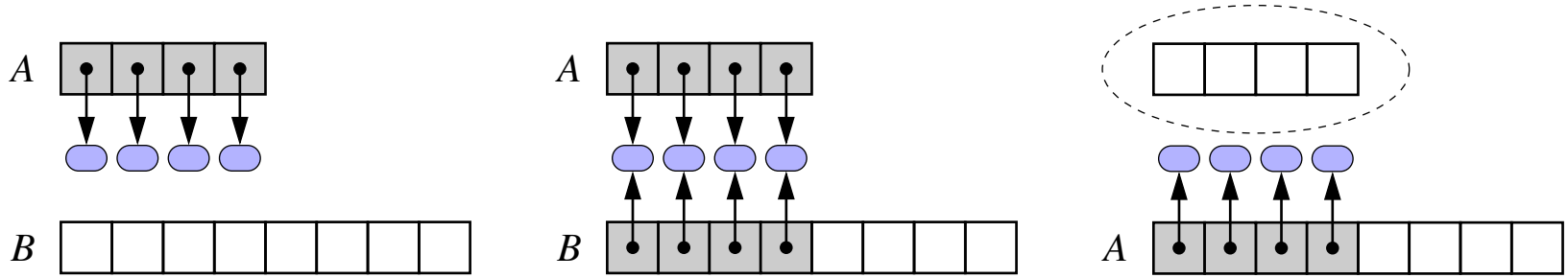
Utility methods

```
1.  /* Checks whether the given index is in the range [0, n-1]. */
2.  protected void checkIndex(int i, int n) throws IndexOutOfBoundsException {
3.    if (i < 0 || i >= n)
4.      throw new IndexOutOfBoundsException("Illegal index: " + i);
5.  }
```

```
1.  /* Resizes internal array to have given capacity >= size. */
2.  protected void resize(int capacity) {
3.    E[] temp = (E[]) new Object[capacity];     // safe cast
4.    for (int k=0; k < size; k++)
5.      temp[k] = data[k];
6.    data = temp;                               // start using the new array
7.  }
```

Dynamic array

- Adding elements leads to the overflow problem
- The overflow problem can be handled by growing the array



GROW-ARRAY(A, n)

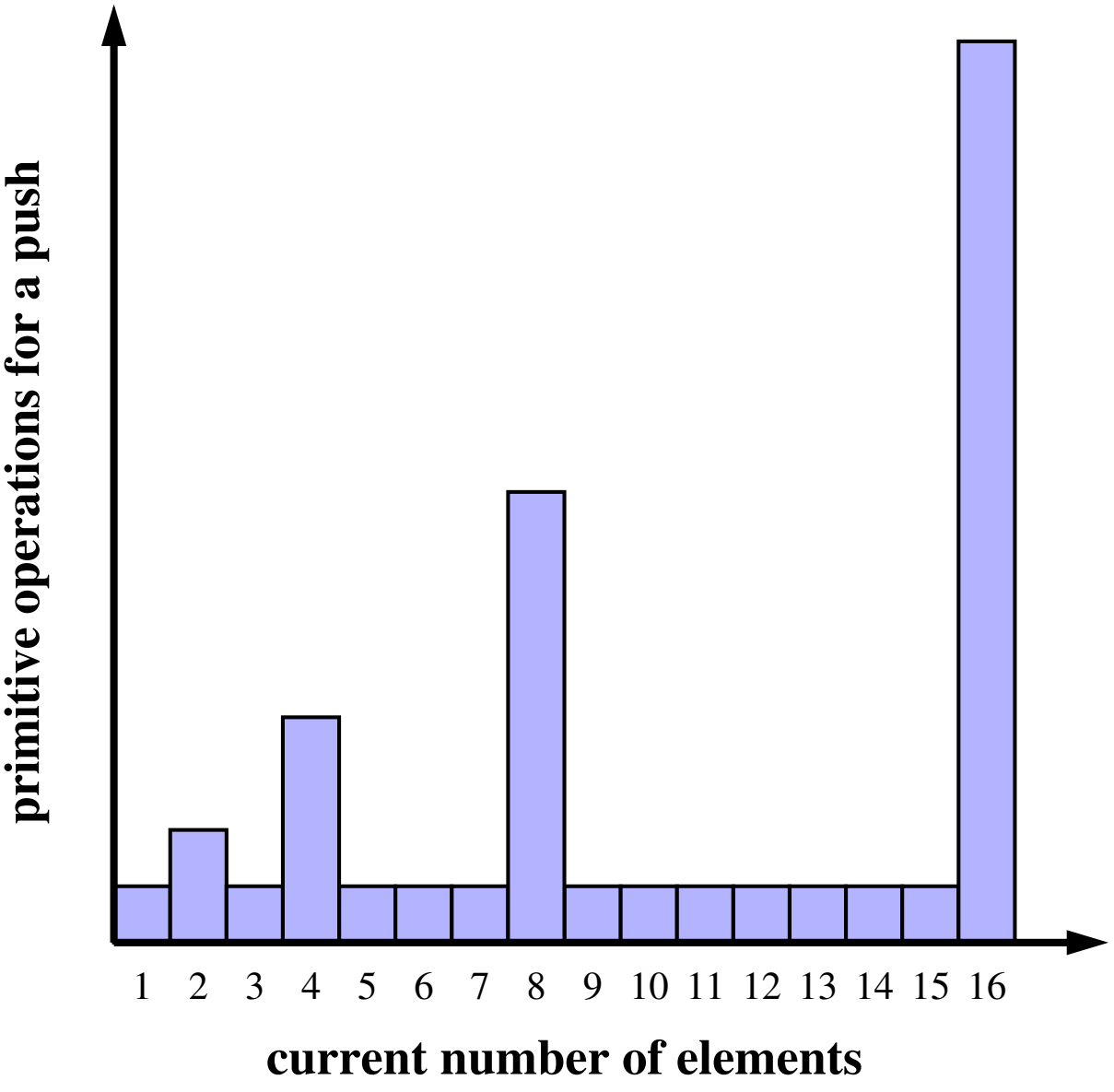
1. Allocate a new array B with larger capacity.
2. Set $B[k] = A[k]$, for $k \leftarrow 0, \ldots, n-1$, where n denotes current number of items.
3. Set $A \leftarrow B$, that is, we henceforth use the new array to support the list.
4. Leave the old array to be garbage collected.

Dynamic array using array doubling

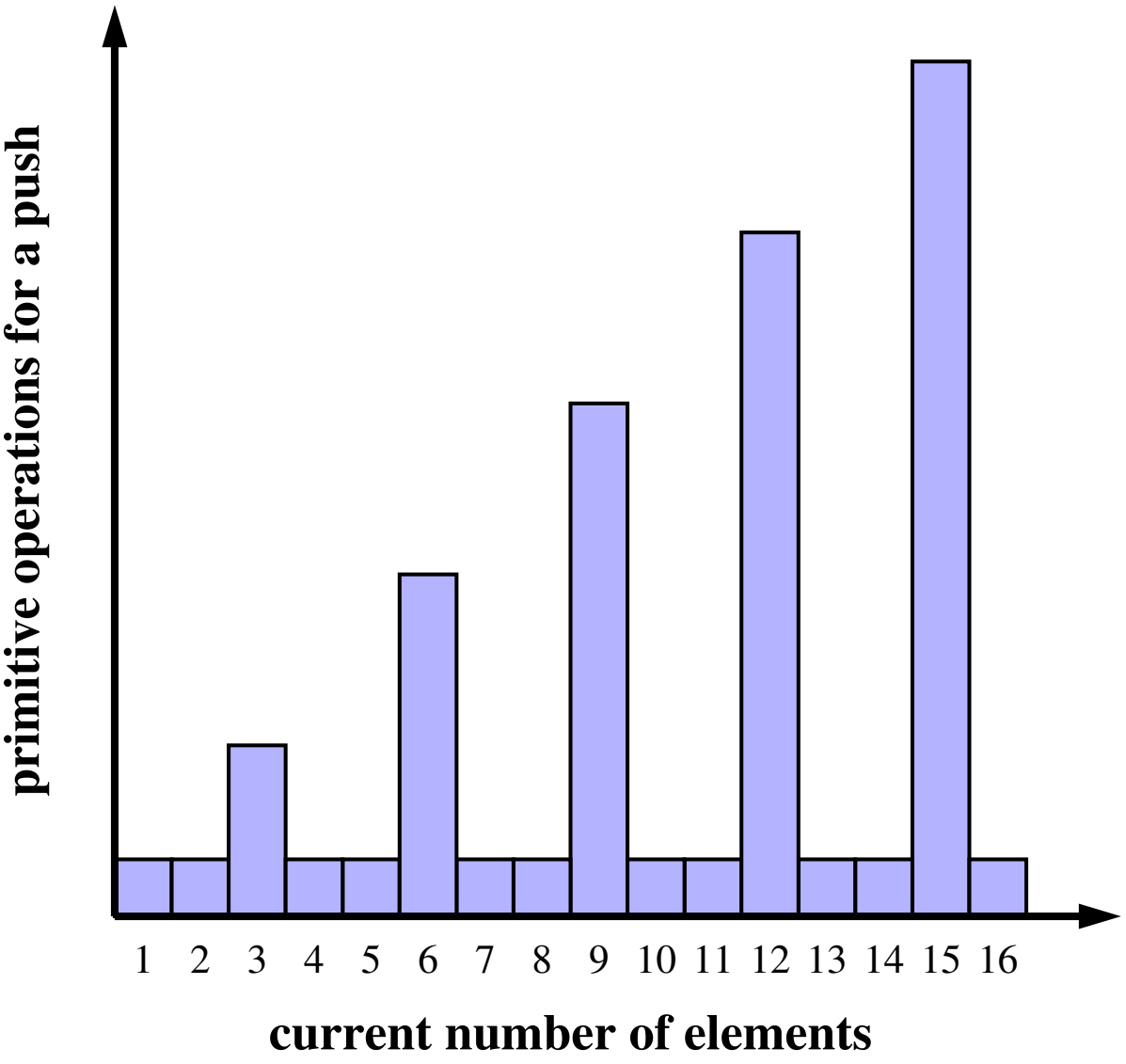
```
1.  /* Resizes internal array to have given capacity >= size. */
2.  protected void resize(int capacity) {
3.    E[] temp = (E[]) new Object[capacity];     // safe cast
4.    for (int k=0; k < size; k++)
5.      temp[k] = data[k];
6.    data = temp;                               // start using the new array
7.  }
```

```
1.  /* Inserts the given element at the specified index of the list, shifting all
2.   * subsequent elements in the list one position further to make room. */
3.  public void add(int i, E e) throws IndexOutOfBoundsException {
4.    checkIndex(i, size + 1);
5.    if (size == data.length)                  // not enough capacity
6.      resize(2 * data.length);                // so double the current capacity
7.    // rest of the method
8.  }
```

Functions for growing/resizing arrays



Functions for growing/resizing arrays



Amortized analysis of dynamic arrays

- Amortized analysis.
  Show that performing a sequence of operations is quite efficient
- Core idea.
  Instead of considering worst-case time taken per operation, consider the average time taken per operation.

Amortized analysis of dynamic arrays

- Use geometric progressions
  $\langle a, ar, ar^2, \ldots \rangle$, such that $r \in \mathbb{R}$ and $r > 1$
  Total time to perform n add operations is $\Theta(n)$
  The value r chosen depends on the trade-off between runtime efficiency and memory usage
- Do not use arithmetic progressions
  $\langle a, a + d, a + 2d, \ldots \rangle$, such that $d \in \mathbb{N}$
  Total time to perform n add operations is $\Theta(n^2)$

Shrinking the dynamic array

- What if you repeatedly remove elements from an arbitrarily large array?
- What if there is an oscillation between growing and shrinking the underlying array?
- The array capacity is halved when the number of elements falls below 1/4th of the capacity

String vs. StringBuilder class

```
1. public String repeat1(char c, int n) {
2.   String answer = "";
3.   for (int j=0; j < n; j++)
4.     answer += c;
5.   return answer;
6. }
```

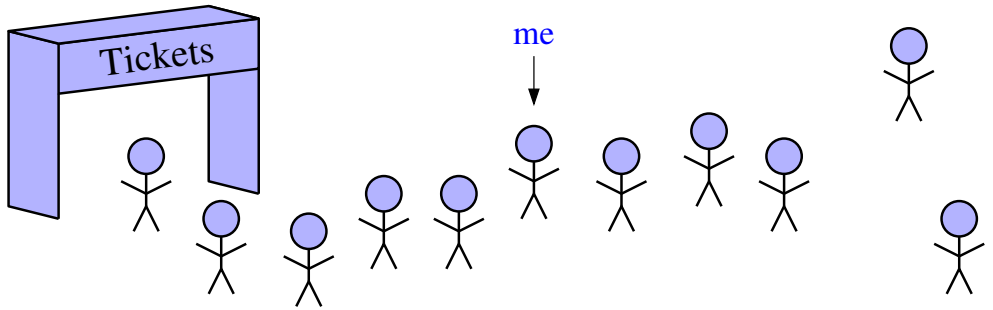
- Static array
- Resize every time
- Time for n adds is $\Theta\left(n^2\right)$

```
1. public String repeat2(char c, int n) {
2.   StringBuilder sb = new StringBuilder();
3.   for (int j=0; j < n; j++)
4.     sb.append(c);
5.   return sb.toString();
6. }
```

- Dynamic array
- Resize a few times
- Time for n adds is $\Theta\left(n\right)$

Positional Lists

Location/position in a sequence



- Position in a queue = node reference

```
1  adfsa
2  fdggfh
3  hgh
4  sfg
5  gh
6  df
7  gdfa adf sdfdsf
8  gsdf
9  gf
10 gfh
11 adfg
```

- Cursor in a text editor = node reference

Positional list ADT

Method	Functionality
`getElement()`	Returns the element stored at this position.

- The position of an element does not change even its index changes due to insertions/deletions in the list

Accessor methods in a positional list

Method	Functionality
`first()`	Returns the position of the first element of L (or null if empty).
`last()`	Returns the position of the last element of L (or null if empty).
`before(p)`	Returns the position of L immediately before position p (or null if p is the first position).
`after(p)`	Returns the position of L immediately after position p (or null if p is the last position).
`isEmpty()`	Returns true if list L does not contain any elements.
`size()`	Returns the number of elements in list L.

A traversal of a positional list

```
1. Position<String> cursor = guests.first();
2. while (cursor != null) {
3.   System.out.println(cursor.getElement());
4.   cursor = guests.after(cursor);    // advance to the next position (if any)
5. }
```

Update methods in a positional list

Method	Functionality
`addFirst(e)`	Inserts a new element e at the front of the list, returning the position of the new element.
`addLast(e)`	Inserts a new element e at the back of the list, returning the position of the new element.
`addBefore(p,e)`	Inserts a new element e in the list, just before position p, returning the position of the new element.
`addAfter(p,e)`	Inserts a new element e in the list, just after position p, returning the position of the new element.
`set(p,e)`	Replaces the element at position p with element e, returning the element formerly at position p.
`remove(p)`	Removes and returns the element at position p in the list, invalidating the position.

Operations on a positional list

Method	Return value	List contents
addLast(8)	p	(8p)
first()	p	(8p)
addAfter(p, 5)	q	(8p, 5q)
before(q)	p	(8p, 5q)
addBefore(q, 3)	r	(8p, 3r, 5q)
getElement()	3	(8p, 3r, 5q)
after(p)	r	(8p, 3r, 5q)
before(p)	null	(8p, 3r, 5q)
addFirst(9)	s	(9s, 8p, 3r, 5q)
remove(last())	5	(9s, 8p, 3r)
set(p, 7)	8	(9s, 7p, 3r)
remove(q)	error	(9s, 7p, 3r)

Java interface for positional list ADT

```
1.  public interface Position<E> {
2.    E getElement() throws IllegalStateException;
3.  }
```

```
1.  public interface PositionalList<E> {
2.    int size();
3.    boolean isEmpty();
4.    Position<E> first();
5.    Position<E> last();
6.    Position<E> before(Position<E> p) throws IllegalArgumentException;
7.    Position<E> after(Position<E> p) throws IllegalArgumentException;
8.    Position<E> addFirst(E e);
9.    Position<E> addLast(E e);
10.   Position<E> addBefore(Position<E> p, E e) throws IllegalArgumentException;
11.   Position<E> addAfter(Position<E> p, E e) throws IllegalArgumentException;
12.   E set(Position<E> p, E e) throws IllegalArgumentException;
13.   E remove(Position<E> p) throws IllegalArgumentException;
14. }
```

Positional list using DLL

```
1.  public class LinkedPositionalList<E> implements PositionalList<E> {
2.    private static class Node<E> implements Position<E> {...}
3.    private Node<E> header;              // header sentinel
4.    private Node<E> trailer;             // trailer sentinel
5.    private int size = 0;
6.
7.    public LinkedPositionalList() {...}
8.    private Node<E> validate(Position<E> p) throws IllegalArgExcep {...}
9.    private Position<E> position(Node<E> node) {...}
10.   public int size() {...}
11.   public boolean isEmpty() {...}
12.   public Position<E> first() {...}
13.   public Position<E> last() {...}
14.   public Position<E> before(Position<E> p) throws IllegalArgExcep {...}
15.   public Position<E> after(Position<E> p) throws IllegalArgExcep {...}
16.   private Position<E> addBetween(E e, Node<E> pred, Node<E> succ) {...}
17.   public Position<E> addFirst(E e) {...}
18.   public Position<E> addLast(E e) {...}
19.   public Position<E> addBefore(Position<E> p, E e) throws IllegalArgExcep {...}
20.   public Position<E> addAfter(Position<E> p, E e) throws IllegalArgExcep {...}
21.   public E set(Position<E> p, E e) throws IllegalArgExcep {...}
22.   public E remove(Position<E> p) throws IllegalArgExcep {...}
23. }
```

Positional list using DLL

```
1.  //---------------- nested Node class ----------------
2.  private static class Node<E> implements Position<E> {
3.    private E element;          // reference to the element stored at this node
4.    private Node<E> prev;       // reference to the previous node in the list
5.    private Node<E> next;       // reference to the subsequent node in the list
6.    public Node(E e, Node<E> p, Node<E> n)
7.    { element = e; prev = p; next = n; }
8.
9.    // public accessor methods
10.   public E getElement() throws IllegalStateException {
11.     if (next == null)         // convention for defunct node
12.       throw new IllegalStateException("Position no longer valid");
13.     return element;
14.   }
15.   public Node<E> getPrev() { return prev; }
16.   public Node<E> getNext() { return next; }
17.   public void setElement(E e) { element = e; }
18.   public void setPrev(Node<E> p) { prev = p; }
19.   public void setNext(Node<E> n) { next = n; }
20. } //----------- end of nested Node class -----------
```

Constructor and private utilities

```
1. public LinkedPositionalList() {
2.   header = new Node<>(null, null, null);       // create header
3.   trailer = new Node<>(null, header, null);    // trailer is preceded by header
4.   header.setNext(trailer);                      // header is followed by trailer
5. }
```

```
1. private Node<E> validate(Position<E> p) throws IllegalArgumentException {
2.   if (!(p instanceof Node)) throw new IllegalArgumentException("Invalid p");
3.   Node<E> node = (Node<E>) p;      // safe cast
4.   if (node.getNext() == null)      // convention for defunct node
5.     throw new IllegalArgumentException("p is no longer in the list");
6.   return node;
7. }
```

```
1. private Position<E> position(Node<E> node) {
2.   if (node == header || node == trailer)
3.     return null;   // do not expose user to the sentinels
4.   return node;
5. }
```

Accessor methods

```
1. public Position<E> first() { return position(header.getNext()); }
```

```
1. public Position<E> last() { return position(trailer.getPrev()); }
```

```
1. public Position<E> before(Position<E> p) throws IllegalArgumentException {
2.   Node<E> node = validate(p);
3.   return position(node.getPrev());
4. }
```

```
1. public Position<E> after(Position<E> p) throws IllegalArgumentException {
2.   Node<E> node = validate(p);
3.   return position(node.getNext());
4. }
```

Update methods

```
1. private Position<E> addBetween(E e, Node<E> pred, Node<E> succ) {
2.   Node<E> newest = new Node<>(e, pred, succ);  // create and link a new node
3.   pred.setNext(newest); succ.setPrev(newest); size++;
4.   return newest;
5. }
```

```
1. public Position<E> addFirst(E e)
2. { return addBetween(e, header, header.getNext()); }
```

```
1. public Position<E> addLast(E e)
2. { return addBetween(e, trailer.getPrev(), trailer); }
```

```
1. public Position<E> addBefore(Position<E> p, E e)
2.                                   throws IllegalArgumentException {
3.   Node<E> node = validate(p);
4.   return addBetween(e, node.getPrev(), node);
5. }
```

```
1. public Position<E> addAfter(Position<E> p, E e)
2.                                throws IllegalArgumentException {
3.   Node<E> node = validate(p);
4.   return addBetween(e, node, node.getNext());
5. }
```

Update methods

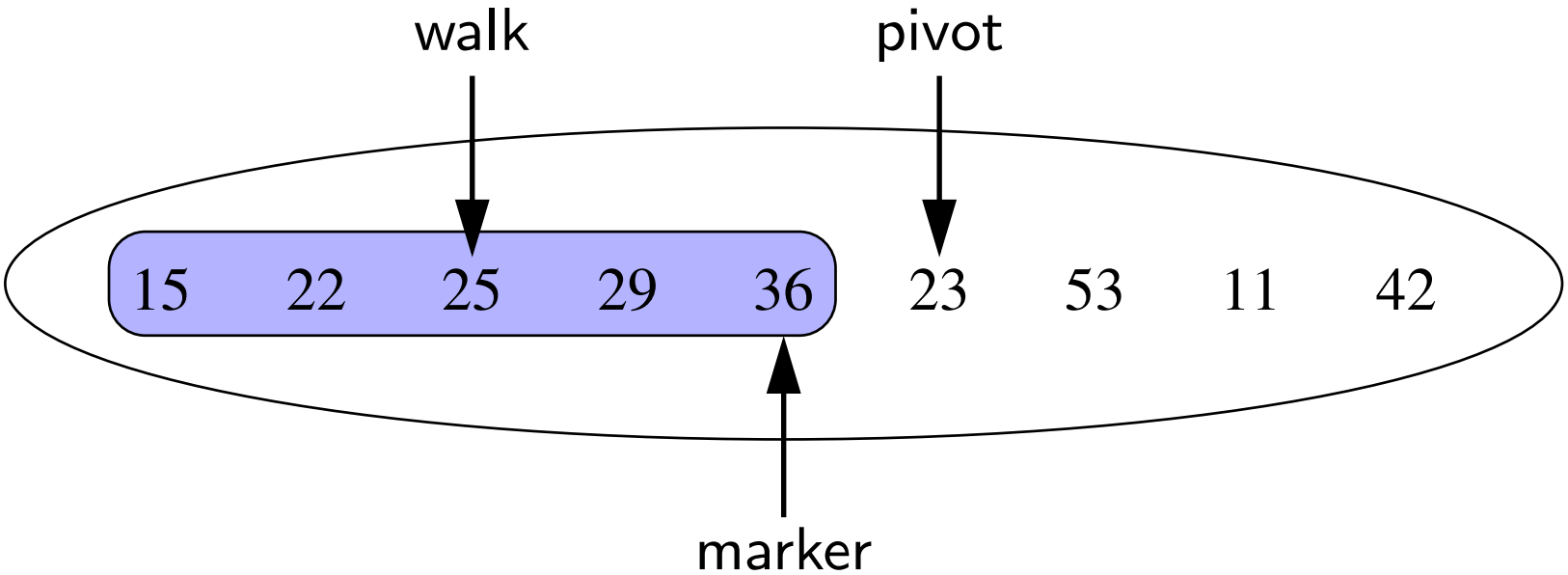
```
1.  public E set(Position<E> p, E e) throws IllegalArgumentException {
2.    Node<E> node = validate(p);
3.    E answer = node.getElement();
4.    node.setElement(e);
5.    return answer;
6.  }
```

```
1.  public E remove(Position<E> p) throws IllegalArgumentException {
2.    Node<E> node = validate(p);
3.    Node<E> predecessor = node.getPrev();
4.    Node<E> successor = node.getNext();
5.    predecessor.setNext(successor);
6.    successor.setPrev(predecessor);
7.    size--;
8.    E answer = node.getElement();
9.    node.setElement(null);            // help with garbage collection
10.   node.setNext(null);               // and convention for defunct node
11.   node.setPrev(null);
12.   return answer;
13. }
```

Performance of a linked positional list

Method	Running time
size()	$\mathcal{O}(1)$
isEmpty()	$\mathcal{O}(1)$
first(), last()	$\mathcal{O}(1)$
before(p), after(p)	$\mathcal{O}(1)$
addFirst(e), addLast(e)	$\mathcal{O}(1)$
addBefore(p, e), addAfter(p, e)	$\mathcal{O}(1)$
set(p, e)	$\mathcal{O}(1)$
remove(p)	$\mathcal{O}(1)$

Sorting a positional list



Sorting a positional list

```
1.  /** Insertion-sort of a positional list of integers into nondecreasing order */
2.  public static void insertionSort(PositionalList<Integer> list) {
3.    Position<Integer> marker = list.first(); // last position known to be sorted
4.    while (marker != list.last()) {
5.      Position<Integer> pivot = list.after(marker);
6.      int value = pivot.getElement(); // number to be placed
7.      if (value > marker.getElement()) // pivot is already sorted
8.        marker = pivot;
9.      else { // must relocate pivot
10.       Position<Integer> walk = marker; // find leftmost item greater than value
11.       while (walk != list.first() && list.before(walk).getElement() > value)
12.         walk = list.before(walk);
13.       list.remove(pivot); // remove pivot entry and
14.       list.addBefore(walk, value); // reinsert value in front of walk
15.     }
16.   }
17. }
```