

## Final term Examination

Spring 2021

CSE-210 Data structure and algorithms

Submitted by: Ashfaq Ahmad  
Registration No: 19PWCSE1795  
Class Section: B

"On my honor, as student of University of Engineering and Technology, I have neither given nor received unauthorized assistance on this academic work."

Student Signature: \_\_\_\_\_

Submitted to:  
**Prof: Nasruminullah**  
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**Department of Computer Systems Engineering  
University of Engineering and Technology, Peshawar**

(Q 1-4)

Question No \*

Application and limitation of all data structure:

→ Important application of all Data Structure:

→ Data structure help in efficient storage of Data in storage device.

→ Data structure usage provide convenience while retrieving the data from storage device.

→ Usage of proper data structure can help programmer save lots of time or processing time while operation such as storage retrieval or processing of Data.

→ Data structure usage can simply encourage reusability in long run as well.

→ manipulation of large amount of data can be carried out easily with the use good data structure approach.

P + T + O

## Limitation of All Data Structure:

- An application using data structure Requires highly qualified professional resource to manage the operation related to data structure.
- Bigger the application or data structure involved in creating and maintaining application more is the requirement of man power. This can increase maintaining of Data Structure.
- Designing your own Data Structure may involve complex algorithm and may require lot of time and testing to conclude they are fool-proof and ready to use for organization purpose. This again will come with more cost.

—xx —xx —xx —xx —xx

## Question No # 2:

### Complexities of Insertion function in Array:

#### i) Worst Case - O(N):

if we

want to insert an element to index 0, then we need to shift all elements to right.

→ In general if we have n-elements then we have need to shift all n elements.

So

Time Complexity =  $O(N)$

→ Best Case:  $O(1)$

if position value equal to size of array then below for loop will not work as the pos = size  
 $\text{for } (i = \text{size}; i > \text{pos}; i++)$   
 $\text{arr}[i] = \text{arr}[i-1];$

→ we can directly place the element in arr[size]

→ so best case is  $O(1)$ .

— XX — XX — XY — XX XY

⊗ ↑

\* Time complexities of insertion function - of linklist:

→ Insertion at beginning:

if we want to insert an node at the beginning

of Linklist then time  
complexity  $O(1)$

→ Insertion in middle  $O(1)$ :

if we want to insert in  
middle then Time Complexity  
is  $O(1)$  with iteration  $O(n)$ .

→ Insertion at the END:

```
void ins (int* list, int val, int in)
{
    int n = in;
    int *t;
    for (int i=0; i<n; i++)
    {
        t = list;
        list = list -> next;
    }
    int *a = new node;
    node -> value = val;
    nod -> next = list -> next;
    t -> next = node;
}
```

Time Complexity is  $O(N)$ .

—xx —xx —xx —xv

END → Question 2:

P T T F U

(Question No # 3)

Error Finding = ?

a)

~~list -> first -> last~~

list -> last = -1;

b)

Inorder (tree -> left, n+1);

inorder (tree -> right, n+1);

—xx —xx —xx — xx

(Question No 4)

8 main drive routine \*

LIST

WINDOW-TYPE w;

ELEMENT-TYPE e;

LIST-TYPE list;

int i;

empty (§ list);

print (§ list);

assign-element-values (&e, 1, "Syed");

wc = first (§ list);

insert (e, w, § list);

print (§ list);

P + T + Q

assign\_element\_values (Se, 2,  
"Ashfaq");

Insert (e, w, Slist);  
print (Slist);

assign\_element\_values (Se, 3,  
"Ahmad");

Insert (e, last (Slist), Slist);  
print (Slist);

### Output:

Syed Ashfaq Ahmad

—xx —xx —xx —x

(Question No (8))

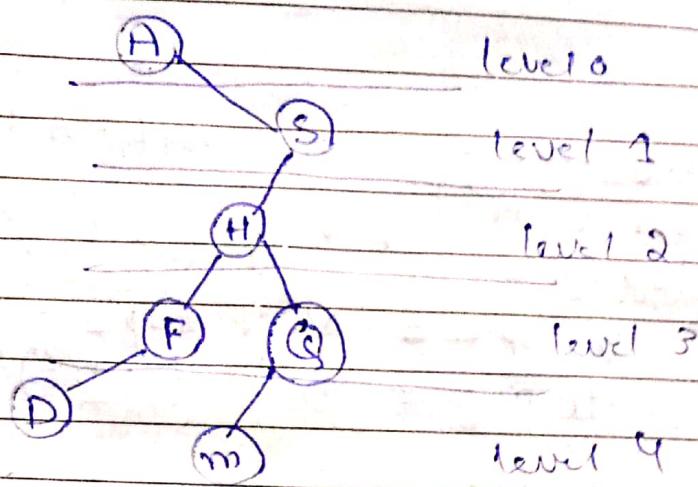
(a) BST by inserting  
Alphabets of full name.

Ans

my full name is

Astha S Ahmad

P + T + O

BST

Some elements are repeating in my name. So I ignored.

(b)

$$\text{no. of internal nodes} = 4 \quad (\text{S, H, F, Q})$$

$$\text{no. of external nodes} = 2 \quad (\text{D, m})$$

$$\text{level of tree} = 4$$

$$\text{height of tree} = 4$$

(c)

Right Complete binary tree.

(d)

Pre order (A, D, F, H, m, Q, S)

Post order (D, F, H, m, Q, S, A)

In order (A, D, F, H, m, Q, S)

— XX — XX — XX — XY

Question # 5Ans

my name is ASHFA ♀.

```

Push ('A');
if (top() == 'C' || top() == 'O'
    || top() == 'V' || top() == 'I'
    || top() == 'D') POP();

```

Push ('S');

→ Same if Statement will be used  
as used in case of Push ('A')

Push ('H');

→ Same case

Push ('F');

→ Same case

Push ('A');

→ Same case

Push ('Q');

→ Same case

↳ i will not use this Push  
function b/c A is already  
used before.

P → T + O

Diagram:

POP ↓ → last element



Push ↑ → first element

— xx — xx — xx — xx

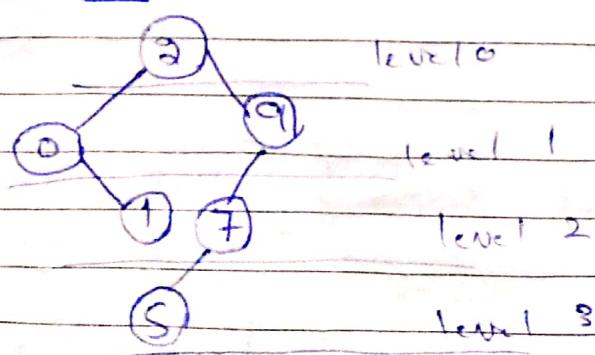
Answer No : 6

Ans

my Reg No = 19PWESE1798  
So

2019 1795

(b) BST Tree!



P → TF O

(c)

no. of internal node = 3 (0, 2, 7)

no. of external node = 2 (1, 5)

level of tree = 3

height of tree = 3

(d)

the tree generated is neither perfect nor complete since not all level are filled. it is degenerated tree in which each internal node has exactly one child.

(e)

Inorder (0, 1, 2, 5, 7, 9)

Preorder (2, 0, 1, 5, 7, 9)

Postorder (0, 1, 5, 7, 9, 2)

xx — xx — xk — xy

Question No # 7:Ans

Struct node ( int val;  
Node \* left, right )

Insert ( Node \* tree, int val )

```
{
    if (*tree == NULL)
        {
```

P → T → 0

```

Node * b = new Node;
b->left = NULL;
b->right = NULL;
b->val = val;
tree = b;
return;
}

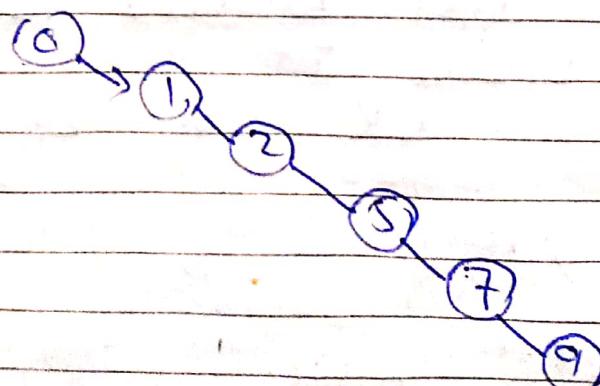
if (val < tree->val)
    insert (tree->left, val);
else
    if (val > tree->val)
        insert (tree->right, val);
}

```

```

int main () {
    Node * tree;
    insert (tree, 0);
    insert (tree, 1);
    insert (tree, 2);
    insert (tree, 5);
    insert (tree, 7);
    insert (tree, 9);
}

```

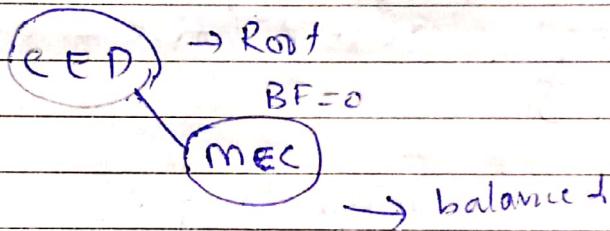


~~xx~~ — ~~xx~~ — ~~xx~~ — ~~xx~~  
 P → T + 0

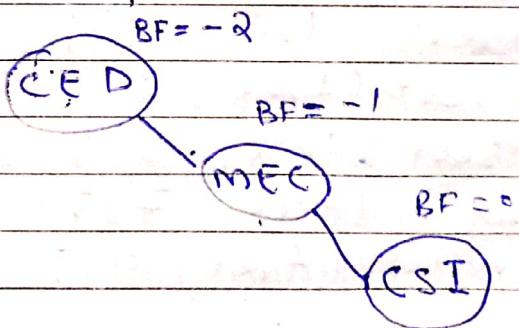
(Question No # 9)Ans:

Let I suppose given  
10 departments.

CED, MEC, CSI, EPE, MIN  
E&E, SED, TEC, RSD, AID  
BF=1

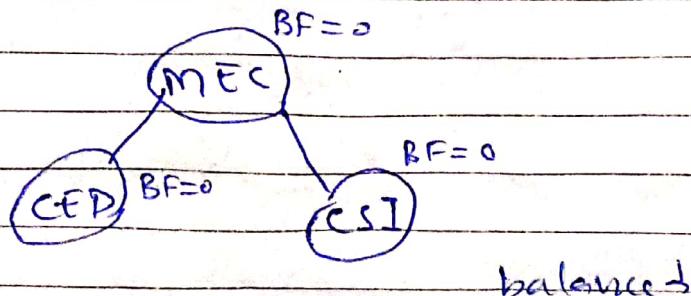


Insert CSI



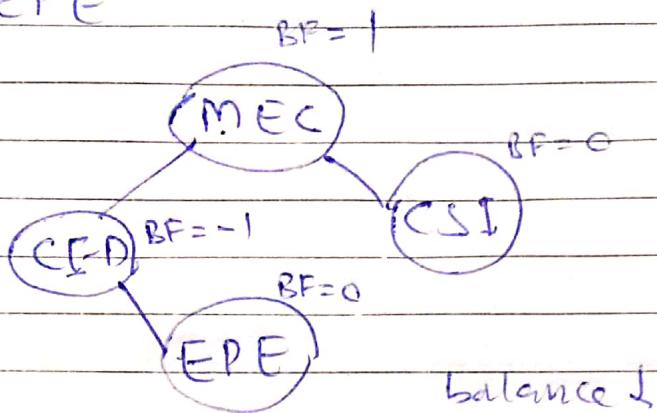
It is not balanced so first  
we balance it

it is RR so left rotate



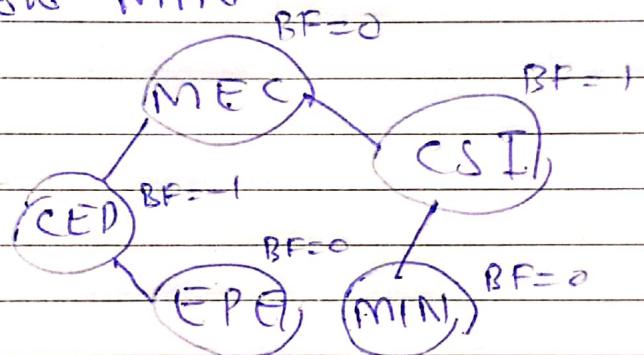
P → T + 0

EPE

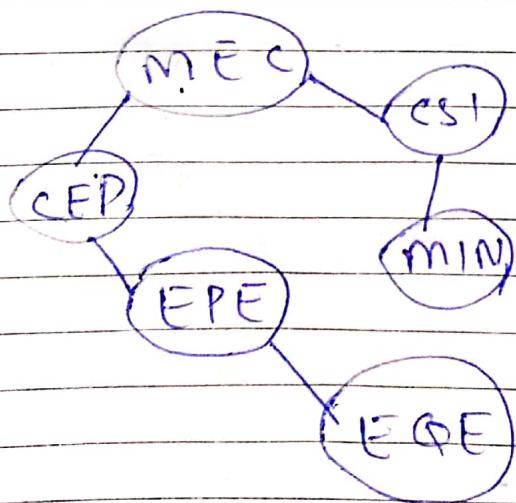


balance ↓

Now MIN



Now EQE



END

XIX