# Lecture 15

#### Topics:

- 1. Logic of Integer Representation
- 2. Decimal and Hexadecimal Representations
- 3. Binary Representation
- 4. Binary Addition Algorithm
- 5. Binary Multiplication Algorithm

#### 4.2 Integer Representations and Algorithms

Integers can be expressed using any integer greater than one as a base, as we will show in this section. Although we commonly use decimal (base 10), representations, binary (base 2), octal (base 8), and hexadecimal (base 16) representations are often used, especially in computer science. Given a base b and an integer n, we will show how to construct the base b representation of this integer. We will also explain how to quickly convert between binary and octal and between binary and hexadecimal notations.

#### 4.2.2 Representations of Integers

In everyday life we use decimal notation to express integers. In decimal notation, an integer n is written as a sum of the form  $a_k \cdot 10^k + a_{k-1} \cdot 10^{k-1} + \dots + a_1 \cdot 10 + a_0$ , where  $a_j$  is an integer with  $0 \le a_j \le 9$  for  $j = 0,1,\dots,k$ . For example, 965 is used to denote  $9 \cdot 10^2 + 6 \cdot 10 + 5$ . However, it is often convenient to use bases other than 10.

**THEOREM 1:** Let b be an integer greater than 1. Then if n is a positive integer, it can be expressed uniquely in the form  $n = a_k b_k + a_{k-1} b_{k-1} + \dots + a_1 b + a_0$ , where k is a nonnegative integer,  $a_0, a_1, \dots, a_k$  are nonnegative integers less than b, and  $a_k \neq 0$ .

**EXAMPLE 1:** What is the decimal expansion of the integer that has  $(1\,0101\,1111)_2$  as its binary expansion?

Solution: We have  $(1\ 0101\ 1111)_2 = 1 \cdot 2^8 + 0 \cdot 2^7 + 1 \cdot 2^6 + 0 \cdot 2^5 + 1 \cdot 2^4 + 1 \cdot 2^3 + 1 \cdot 2^2 + 1 \cdot 2^1 + 1 \cdot 2^0 = 351$ .

**EXAMPLE 2:** What is the decimal expansion of the number with octal expansion  $(7016)_8$ ? Solution: Using the definition of a base b expansion with b = 8 tells us that  $(7016)_8 = 7 \cdot 8^3 + 0 \cdot 8^2 + 1 \cdot 8 + 6 = 3598$ .

Sixteen different digits are required for hexadecimal expansions. Usually, the hexadecimal digits used are 0,1,2,3,4,5,6,7,8,9,A,B,C,D,E, and F, where the letters A through F represent the digits corresponding to the numbers 10 through 15 (in decimal notation).

**EXAMPLE 3:** What is the decimal expansion of the number with hexadecimal expansion  $(2AE0B)_{16}$ ?

<u>Solution</u>: Using the definition of a base *b* expansion with b = 16 tells us that  $(2AE0B)_{16} = 2 \cdot 16^4 + 10 \cdot 16^3 + 14 \cdot 16^2 + 0 \cdot 16 + 11 = 175627$ .

**EXAMPLE 4:** Find the octal expansion of  $(12345)_{10}$ .

Extra Solution: First, divide 12345 by 8 to obtain  $12345 = 8 \cdot 1543 + 1 \cdot 262$ . Successively dividing quotients by 8 gives  $1543 = 8 \cdot 192 + 7$ ,  $192 = 8 \cdot 24 + 0$ ,  $24 = 8 \cdot 3 + 0$ ,  $3 = 8 \cdot 0 + 3$ . The successive remainders that we have found, 1,7,0,0, and 3, are the digits from the right to the left of 12345 in base 8. Hence,  $(12345)_{10} = (30071)_8$ 

**EXAMPLE 5:** Find the hexadecimal expansion of  $(177130)_{10}$ .

Solution: First divide 177130 by 16 to obtain  $177130 = 16 \cdot 11070 + 10$ . Successively dividing quotients by 16 gives  $11070 = 16 \cdot 691 + 14$ ,  $691 = 16 \cdot 43 + 3$ ,  $43 = 16 \cdot 2 + 11$ ,  $2 = 16 \cdot 0 + 2$ . The successive remainders that we have found, 10,14,3,11,2, give us the digits from the right to the left of 177130 in the hexadecimal (base 16) expansion of  $(177130)_{10}$ . It follows that  $(177130)_{10} = (2B3EA)_{16}$ .

**EXAMPLE 6:** Find the binary expansion of  $(241)_{10}$ .

Solution: First divide 241 by 2 to obtain  $241 = 2 \cdot 120 + 1$ . Successively dividing quotients by 2 gives  $120 = 2 \cdot 60 + 0$ ,  $60 = 2 \cdot 30 + 0$ ,  $30 = 2 \cdot 15 + 0$ ,  $15 = 2 \cdot 7 + 1$ ,  $7 = 2 \cdot 3 + 1$ ,  $3 = 2 \cdot 1 + 1$ ,  $1 = 2 \cdot 0 + 1$ . The successive remainders that we have found, 1, 0, 0, 0, 1, 1, 1, 1, are the digits from the right to the left in the binary (base 2) expansion of  $(241)_{10}$ . Hence,  $(241)_{10} = (1111 \ 0001)_2$ .

**ALGORITHM 1:** Constructing Base *b* Expansions.

procedure base b expansion(n, b: positive integers with b > 1)

$$q := n$$

$$k \coloneqq 0$$

while  $q \neq 0$ 

```
a_k \coloneqq q \mod b, \ q \coloneqq q \operatorname{div} b, \ k \coloneqq k+1
```

return  $(a_{k-1}, \dots, a_1, a_0)$   $\{(a_{k-1} \dots a_1 a_0)_b \text{ is the base } b \text{ expansion of } n\}$ 

**ALGORITHM 2:** Addition of Integers.

```
procedure \operatorname{add}(a,b): positive integers) {the binary expansions of a and b are (a_{n-1}a_{n-2}\dots a_1a_0)_2 and (b_{n-1}b_{n-2}\dots b_1b_0)_2, respectively} c\coloneqq 0 for j\coloneqq 0 to n-1 d\coloneqq \left\lfloor (a_j+b_j+c)/2\right\rfloor s_j\coloneqq a_j+b_j+c-2d c\coloneqq d s_n\coloneqq c return (s_0,s_1,\dots,s_n) {the binary expansion of the sum is (s_ns_{n-1}\dots s_0)_2
```

**MULTIPLICATION ALGORITHM:** Next, consider the multiplication of two n-bit integers a and b. The conventional algorithm (used when multiplying with pencil and paper) works as follows.

Using the distributive law, we see that  $ab = a(b_0 2^0 + b_1 2^1 + \dots + b_{n-1} 2^{n-1}) = a(b_0 2^0) + a(b_1 2^1) + \dots + a(b_{n-1} 2^{n-1})$ 

**ALGORITHM 3:** Multiplication of Integers.

```
procedure multiply(a, b: positive integers) {the binary expansions of a and b are
(a_{n-1}a_{n-2}...a_1a_0)_2 and (b_{n-1}b_{n-2}...b_1b_0)_2, respectively
for j := 0 to n-1
if b_j = 1 then c_j := a shifted j places
      else c_i := 0 \ \{c_0, c_1, ..., c_{n-1} \text{ are the partial products}\}
     p \coloneqq 0
      for i := 0 to n-1
           p := \operatorname{add}(p, c_i)
return p \{ p \text{ is the value of } ab \}
```

**EXAMPLE 7:** Find the octal and hexadecimal expansions of  $(11\ 1110\ 1011\ 1100)_2$  and the binary expansions of  $(765)_8$  and  $(A8D)_{16}$ .

Solution: To convert (11 1110 1011 1100)<sub>2</sub> into octal notation we group the binary digits into blocks of three, adding initial zeros at the start of the leftmost block if necessary. These blocks, from left to right, are 011, 111, 010, 111, and 100, corresponding to 3, 7, 2, 7, and 4, respectively. Consequently,  $(11\ 1110\ 1011\ 1100)_2 = (37274)_8$ . To convert  $(11\ 1110\ 1011\ 1100)_2$  into hexadecimal notation we group the binary digits into blocks of four, adding initial zeros at the start of the leftmost block if necessary. These blocks, from left to right, are 0011, 1110, 1011, and 1100, corresponding to the hexadecimal digits 3, E, B, and C, respectively. Consequently,  $(11\ 1110\ 1011\ 1100)_2 = (3EBC)_{16}$ . To convert (765)<sub>8</sub> into binary notation, we replace each octal digit by a block of three binary digits. These blocks are 111,110, and 101. Hence,  $(765)_8 = (111110101)_2$ . To convert  $(A8D)_{16}$  into binary notation, we replace each hexadecimal digit by a block of four binary digits. These blocks are 1010, 1000, and 1101. Hence,  $(A8D)_{16} = (1010\ 1000\ 1101)_2$ .

# **Integer Operations**

```
ALGORITHM 2: Addition of Integers.
procedure add(a,b: positive integers) {the binary expansions of a and b are
(a_{n-1} a_{n-2} \dots a_1 a_0)_2 and (b_{n-1} b_{n-2} \dots b_1 b_0)_2
c \coloneqq 0
for j := 0 to n-1
  d \coloneqq |(a_i + b_i + c)/2|
   s_i \coloneqq a_i + b_i + c - 2d
   c \coloneqq d
   S_n := C
return (s_0, s_1, ..., s_n) {the binary expansion of the sum is (s_n s_{n-1} ... s_n)_2}
```

# Integer Operations (Continued)

**AEGORITHM 3:** Multiplication of Integers.

```
procedure multiply(a, b: positive integers) {the binary expansions of a and b are
(a_{n-1}a_{n-2}...a_1a_0)_2 and (b_{n-1}b_{n-2}...b_1b_0)_2
for i := 0 to n-1
  if b_i = 1 then c_i := a shifted j places
  else c_i := 0 {c_0, c_1, ..., c_{n-1} are the partial products}
p := 0
for j := 0 to n-1
  p := add(p, c_i)
return p \{ p \text{ is the value of } ab \}
```

#### 4.2.4 Modular Exponentiation

Before we present an algorithm for fast modular exponentiation based on the binary expansion of the exponent, first observe that we can avoid using large amount of memory if we compute  $b_n \mod m$  by successively computing  $b_k \mod m$  for k=1,2,...,n using the fact that  $b^{k+1} \mod m = b \left( b^k \mod m \right) \mod m$  (by Corollary 2 of Theorem 5 of Section 4.1). (Recall that  $1 \le b < m$ .) However, this approach is impractical because it requires n-1 multiplications of integers and n might be huge.

To motivate the fast modular exponentiation algorithm, we illustrate its basic idea. We will explain how to use the binary expansion of n, say  $n = (a_{k-1} \dots a_1 a_0)_2$ , to compute  $b_n$ . First, note that  $b_n = b^{a_{k-1} \cdot 2^{k-1} + \dots + a_1 \cdot 2 + a_0} = b^{a_{k-1} \cdot 2^{k-1}} \dots b^{a_1 \cdot 2} \cdot b^{a_0}$ 

For example, to compute 311 we first note that  $11 = (1011)_2$ , so that  $3^{11} = 383231$ . By successively squaring, we find that  $3^2 = 9$ ,  $3^4 = 9^2 = 81$ , and  $3^8 = (81)^2 = 6561$ .

Consequently,  $3^{11} = 3^8 3^2 3^1$ 

```
ALGORITHM 5: Fast Modular Exponentiation. procedure modular exponentiation(b: integer, n=(a_{n-1}a_{n-2}...a_1a_0)_2, m: positive integers) x\coloneqq 1 power \equiv b \bmod m for i\coloneqq 0 to k-1
```

 $power := (power \cdot power) \mod m$ 

if  $a_i = 1$  then  $x := (x \cdot \text{power}) \mod m$ 

return  $x \{x \text{ equals } b_n \mod m\}$ 

**EXAMPLE 12:** Use Algorithm 5 to find 3<sup>644</sup> mod 645.

<u>Solution:</u> Algorithm 5 initially sets x = 1 and power =  $3 \mod 645 = 3$ . In the computation of  $3^{644} \mod 645$ , this algorithm determines  $3^{2^j} \mod 645$  for j = 1,2,...,9 by successively squaring and reducing modulo 645. If  $a_j = 1$  (where  $a_j$  is the bit in the jth position in the binary expansion of 644, which is  $(1010000100)_2$ , it multiplies the current value of x by  $3^{2^j} \mod 645$  and reduces the result modulo 645. Here are the steps used:

```
i \bullet 0: Because a_0 = 0, we have x = 1 and power = 32 \mod 645 = 9 \mod 645 = 9;
i = 1: Because a_1 = 0, we have x = 1 and power = 9^2 \mod 645 = 81 \mod 645 = 81;
i = 2: Because a_2 = 1, we have x = 1 \cdot 81 \mod 645 = 81 and power = 81^2 \mod 645 = 81
6561 \mod 645 = 111;
i = 3: Because a_3 = 0, we have x = 81 and power = 111^2 \mod 645 = 12{,}321 \mod 645 = 66;
i = 4: Because a_4 = 0, we have x = 81 and power = 66^2 \mod 645 = 4356 \mod 645 = 486;
i = 5: Because a_5 = 0, we have x = 81 and power = 486^2 \mod 645 = 236,196 \mod 645 = 126:
i = 6: Because a_6 = 0, we have x = 81 and power = 12^{62} \mod 645 = 15,876 \mod 645 = 396;
i = 7: Because a_7 = 1, we find that x = (81 \cdot 396) \mod 645 = 471 and power = 39^{62} \mod 645 = 471
156,816 \mod 645 = 81;
i = 8: Because a_8 = 0, we have x = 471 and power = 81^2 \mod 645 = 6561 \mod 645 = 111;
i = 9: Because a_9 = 1, we find that x = (471 \cdot 111) \mod 645 = 36.
```